Predicate-based Testing

Input (X)

- · Predicates are conditions
 - Divides the input domain into partitions
 - Define the paths of the program
- · Program P
 - Input X; Predicate C
 - If outcome of C is incorrect,
 - Either C is incorrect,
 - · Or statement(s) executed before C
 - Most likely, P's output is incorrect
 - · Low probability of "coincidental correctness"
- Predicate-based testing
 - Require certain types of tests for each predicate in the program

Importance of Predicatebased Testing

- · Thorough testing of C used to
 - Detect faults in C.
 - Statements executed before C
 - Statements executed after C

Terms Defined

- · Predicate
 - Simple or compound predicate
- · Simple predicate
 - Boolean variable, or
 - Relational expression,
 - May have one or more NOT (¬) operators
- · Relational expression
 - E1 <rop> E2
 - E1 and E2 are arithmetic expressions
 - $\langle \mathsf{rop} \rangle \in \{\langle, \langle =, \rangle, \rangle =, /=, =\}$

Terms Defined (2)

- · Compound predicate
 - At least one "binary Boolean operator"
 - Two or more operands
 - Maybe NOT operators
 - Maybe parenthesis
- · Binary Boolean operators
 - OR (|) and AND (&)
- · Simple operand
 - Operand without binary Boolean operators
- · Compound operand
 - Operand with at least one binary Boolean operator

Terms Defined (3)

- · Boolean expression
 - Predicate with no relational expressions
- · Bi = Boolean expression
- Ei = Arithmetic expression
- · <rop> or <rop_i> = relational
 operator
- · <bop> or <bop_i> = binary Boolean operator

Assumptions

· Predicate has no syntactic faults

Types of Faults

- An "incorrect" predicate may have one or more of the following faults
 - Boolean operator fault
 - Incorrect AND/OR or missing/extra NOT
 - Boolean variable fault
 - · Incorrect Boolean variable
 - Parenthesis fault
 - Incorrect locationRelational operator fault
 - · Incorrect relational operator
 - Arithmetic expression fault
 - · Various types

Yet More Terms

- Existence of one/more faults is "detected by a test" T if an execution of C with T produces an incorrect outcome of C
- Test set T for C "guarantees the detection" of certain type of faults F in C if the existence of F in C can be detected by at least one element in T, provided C doesn't contain faults of other types

Yet More Terms (2)

- Assume that C* has the same set of variables as C and is not equivalent to C.
 Test set T "distinguishes" C from C* if C and C* produce different outcomes for T
- Assume that C contains faults and C" is the correct version of C. Test set T is "insensitive" to the faults in C if this test cannot distinguish C from C"

Testing Simple Predicates

- · Branch testing
 - TRUE and FALSE branches be executed at least once
- · Relational Operator Testing
 - Given E1 <rop> E2
 - Need 3 tests
 - E1 > E2: E1 < E2: E1 = E2
 - If only <rop> is incorrect and E1 and E2 are correct, then detection is guaranteed

Testing Compound Predicates

- · Complete branch testing
 - All TRUE and FALSE branches of each simple/compound operand in compound predicate C be executed at least once
- · Exhaustive branch testing
 - All combinations of TRUE and FALSE branches of simple operands in C be executed at least once
 - C has N Boolean Operators, then N+1 simple operands. Requires 2^(n+1) test cases

Testing Compound Predicates (2)

- · Complete relational operator testing
 - Relational operator testing for each relational expression in C
 - Let C# be (E1 = E2) & (E3 /= E4)
 - Assume T1 contains 3 tests
 - T11 makes E1 = E2 and E3 = E4
 - T12 makes E1 > E2 and E3 > E4
 - T13 makes E1 < E2 and E3 < E4
 - T1 satisfies relational operator testing for each simple operand of $\it C\#$
- If E1, E2, E3, and E4 are correct, what can we say about the correctness of operators?

Complete Relational Operator Testing

- · Can the test cases T11, T12, and T13 distinguish between C# and
 - (E1 = E2) & (E3 < E4)
 - (E1 /= E2) & (E3 = E4)

BR-constraints

- · Given a predicate
 - (<opd₁> <bop₁> <opd₂> <bop₂> ... <opd_n> <bop_n>
 - · <opdi> is the ith simple operand
- · BR-constraint
 - (D1, D2, ..., Dn)
 - Each Di is a symbol specifying a constraint on the Boolean variable or relational expression in <pd;

BR-constraints (2)

- · Constraints for a Boolean variable B
 - The value of B is TRUE
 - The value of B is FALSE
 - No constraint
- · Symbols
 - 🕇
 - **f**
 - _ *

BR-constraints (2)

- Constraints for a relational expression (E1 <rop> E2)
 - · Value is TRUE t
 - · Value is FALSE f
 - · (E1 E2) > 0
 - (E1 E2) = 0
 - · (E1 E2) < 0
 - · No constraint *

Constraint Satisfaction

- · Definition
 - Constraint D on predicate C is covered (or satisfied) by a test if during the execution of C with this test, the value of each Boolean variable or relational expression in C satisfies the corresponding constraint in D
- E.g.,
 - (=, <)
- for ((E1 >= E2) | \neg (E3 > E4))
- Coverage requires that (E1 = E2) and (E3 < E4)

Constraint Satisfaction (2)

- · Definition
 - Set S of BR-constraints on predicate
 C is covered (or satisfied) by a test
 set T if each constraint in S is
 covered for C by at least one test in T

Terms Redefined

- · In terms of BR-constraints

 - - · {(†, *), (f, *), (*, †), (*, f)}

Terms Defined

- · Concatenation
 - Let $\mathbf{u}=(\mathbf{u}_1,\ \mathbf{u}_2,\ ...,\ \mathbf{u}_m)$ and $\mathbf{v}=(\mathbf{v}_1,\ \mathbf{v}_2,\ ...,\ \mathbf{v}_n)$ be two sequences
 - $-(u,v) = (u_1, u_2, ..., u_m, v_1, v_2, ..., v_n)$
- · Other terms
 - Let A and B be two sets
 - A\$B denotes the union of A and B
 - A*B is the product of A and B
 - |A| is the size of A
 - A%B is called the onto from A to B
 - Minimal set of (u,v) such that u ∈ A and every element in A appears in u at least once; v ∈ B and every element in B appears in v at least once

Terms Defined

- · Observations
 - $|A\%B| = \max(|A|, |B|)$
 - A%B may have several possible values
 - If $C = \{(a), (b)\}$ and $D = \{(c), (d)\}$
 - Then what is C%D
 - ((a,c),(b,d))
 - ((a,d),(b,c))
 - How about if E = {(a), (b)} and F = {(c), (d), (e)}

Expected Outcome

- · Let X be a constraint that contains "t", "f", ">", "<", and "=" for a predicate ${\it C}$
- Value produced by C on any input covering X; C(X)
- X covers the TRUE branch of C if C(X)=TRUE, and
- X covers the FALSE branch of C if C(X)=FALSE
- · Let S be a set of constraints for C
- Partition S into S_t and S_f
 - $S_{t}(C) = \{X \text{ in } S \mid C(X) = t\}$
 - $S_f(C) = \{X \text{ in } S \mid C(X) = f\}$

Lets Try Them Out

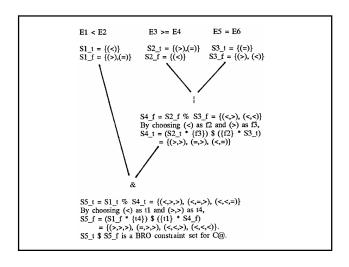
- E1 < E2
 - S1 = {(<), (>), (=)}
 - S1_t = {(<)}
 - S1_f = {(>), (=)}
- E3 >= E4
 - **52** = {(>), (=), (<)}
 - S2_t = {(>), (=)}
 - S2_f = {(<)}
- E5 = E6
 - S3 = {(=), (<), (>)}
 - **53_t** = {(=)}
 - $53_f = {(<), (>)}$

1&

- More complex predicates
 - -(E3 >= E4) | (E5 = E6)
 - 54_f = {(<, <), (<, >)}
 - -(E3 >= E4) & (E5 = E6)
 - · S9_t = {(>, =), (=, =)}
- · How about S4_t and S9_f?

Surprise Quiz

· How About S9_f?



What Next?

 Once all the constraints have been obtained, test cases may be generated