HW07: Solution to Problem 7

NP in DTIME PROBLEM

Prove $A \in NP \rightarrow \exists$ poly $r: A \in DTIME(O(2^{r(n)}))$.

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People did this three ways. I discuss all three.



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Time: $q(n) \times 2^{p(n)}$. This is the solution I had in mind.

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I do not like it since it uses Cook-Levin; however, it is correct.

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Some Teachers Disagree

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The class did not cover them, so for a student to use them is odd.

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Not going to waste time on a solution I disapprove of.