

1 Further Results

1.1 More APSP-Complete Problems

1. **MINIMUM WEIGHT CYCLE IN GRAPH OF NON-NEGATIVE EDGE WEIGHT:** Given a weighted graph with nonnegative edge weights, find the minimum weight cycle in the graph. Williams & Williams [2] showed that this problem is APSP-complete.
2. **SECOND SHORTEST SIMPLE PATH** is as follows: Given a weighted directed graph G , and two nodes s and t , find the second shortest simple path between s and t in G . Williams & Williams [2] showed that this problem is APSP-complete.
3. **CoDIAMETER:** Given a graph G , the goal of CoDiameter is to report a vertex which does not participate in an edge of length equal to the diameter of G . Boroujeni et al. [1] showed that this problem is APSP-complete by a reduction from APSP. Boroujeni also defines CoRADIUS, CoRADIUS, CoNEGATIVETRIANGLE, and CoMEDIAN are APSP-complete, and show them APSP-complete.
4. **APSPVERIFICATION:** Given a graph G and a matrix D , determine if D is the correct distance matrix for G . That is, check that

$$(\forall(i, j) \in E)[D_{i,j} = \text{dist}_G(i, j)].$$

It is known that this problem is APSP-complete.

1.2 Misc

1. **BOOLEAN MATRIX MULTIPLICATION (BMM):** If boolean matrix multiplication has a sub cubic combinatorial algorithm, then so does the triangle detection problem in graphs. This was shown by Williams & Williams [2]. All known algorithms for triangle detection take cubic time, hence using the hardness of triangle detection as an assumption is reasonable.
2. **CoAPSPVERIFICATION:** Given a graph G and a matrix D , either find a pair (i, j) such that $D_{i,j}$ is equal to the distance between vertices i and j in G , or determine that there is no such pair. Boroujeni et al. [1] gave a subcubic reduction from DIAM to it. Recall that DIAM is thought to require cubic time; however, we do not know if DIAM is APSP-hard.
3. **$\{-1, 0, 1\}$ – APSP IS AS FOLLOWS:** Given a weighted directed graph with edge weights in $\{-1, 0, 1\}$, compute the APSP. Despite the complication of having negative edge weights, this problem has a subcubic ($O(n^{2.52})$) algorithm given by Zwick [3]. This problem seemed to require cubic time but did not. Consider that a cautionary note.

References

- [1] M. Boroujeni, S. Dehghani, S. Ehsani, M. T. Hajiaghayi, and S. Seddighin. Subcubic equivalences between graph centrality measures and complementary problems, 2019.
<http://arxiv.org/abs/1905.08127>.
- [2] V. V. Williams and R. R. Williams. Subcubic equivalences between path, matrix, and triangle problems. *Journal of the Association of Computing Machinery (JACM)*, 65(5):27:1–27:38, 2018.
<https://doi.org/10.1145/3186893>.
- [3] U. Zwick. All pairs shortest paths using bridging sets and rectangular matrix multiplication. *Journal of the ACM*, 49(3):289–317, 2002.
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