Open Problems Column Edited by William Gasarch This Issue's Column!

This issue's Open Problem Column is by William Gasarch and is A Known Problem in Ramsey Theory: Ramsey Multiplicity.

Request for Columns!

I invite any reader who has knowledge of some area to contact me and arrange to write a column about open problems in that area. That area can be (1) broad or narrow or anywhere inbetween, and (2) really important or really unimportant or anywhere inbetween.

A Known Problem in Ramsey Theory: Ramsey Multiplicity

by William Gasarch

1 Introduction

In this column we state a class of open problems that are well known in Ramsey Theory but probably not to my non-Ramsey readers. Nothing I present is original.

The problem is as follows: Let G be a graph. Fill in the blank: For every 2-coloring of the edges of K_n (the complete graph on n vertices) there exist BLANK monochromatic copies of G. The key words to Google are Ramsey Multiplicity.

In Section 2 we state a motivating question. In Section 3 we look at K_n such that you always get 2 monochromatic copies of K_3 . In Section 4 we show that, for all 2-colorings of the edges of K_n there exist $\frac{n^3}{24} - O(n^2)$ monochromatic K_3 's. Our upper bound is precise and matches the lower bound; however we do not prove this. In Section 5 we state some of what is known. In Section 6 we state a class of open problems.

2 A Motivating Question

We abbreviate *monochromatic* by *mono* throughout.

Recall the first theorem one usually hears in Ramsey Theory:

For all 2-colorings of the edges of K_6 there is a mono triangle.

From this we obtain a trivial theorem:

For all 2-colorings of the edges of K_{12} there are 2 mono triangles.

How big does n have to be such that any 2-coloring of K_n has 2 mono triangles? The answer is in the next section and may surprise you.

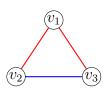
3 Two Monochromatic K_3

Theorem 3.1 For all 2-colorings of edges of K_6 there are 2 mono triangles.

Proof: Let COL be a 2-coloring of the edges of K_6 . Let R, B, M, be the sets of RED, BLUE, and MIXED (having both RED and BLUE edges) triangles, respectively. Clearly

$$|R| + |B| + |M| = {6 \choose 3} = 20.$$

We show that $|M| \le 18$, so $|R| + |B| \ge 2$. Let T be a mixed triangle. It looks like this:

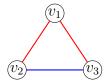


Note that there will be two vertices that have both a RED and a BLUE edge coming out of them.

- (v_2, v_1) is red, (v_2, v_3) is blue. View this as $(v_2, \{v_1, v_3\})$.
- (v_3, v_1) is red, (v_3, v_2) is blue. View this as $(v_3, \{v_1, v_2\})$.

Def 3.2 A Zan is an element $(v, \{u, w\}) \in V \times {V \choose 2}$ such that $v \notin \{u, w\}$ and $COL(v, u) \neq COL(v, w)$. ZAN is the set of all Zans.

Map ZAN to M by mapping $(v, \{u, w\})$ to triangle (v, u, w). This mapping is exactly 2-to-1: every element of M has two Zans mapping to it. The Zans that map to



are $(v_2, \{v_1, v_3\})$ and $(v_3, \{v_1, v_2\})$.

Since there is a 2-to-1 map from ZAN to M, |M| = |ZAN|/2. Now we want to bound |ZAN|. Look at how much each vertex can contribute to ZAN. Note that each vertex has degree 5.

Cases:

- 1. v has $\deg_R(v) = 5$ and $\deg_R(v) = 0$: v contributes 0.
- 2. v has $\deg_R(v) = 4$ and $\deg_R(v) = 1$: v contributes 4.
- 3. $v \text{ has } \deg_R(v) = 3 \text{ and } \deg_B(v) = 2$: v contributes 6. Max.

There are 6 vertices, each contribute ≤ 6 , $|M| \leq |ZAN|/2 \leq 6 \times 6/2 = 18$

$$|R| + |B| = 20 - |M| \ge 2$$

4 Many Mono Triangles

The following theorem was first proven by Goodman [5]; however, we give an easier proof given by Schwenk [9], as presented by Dorwart and Finkbeiner [2].

Theorem 4.1

- 1. Assume $n \equiv 1 \pmod{4}$. For all 2-colorings of the edges of K_n there are at least $\frac{n^3}{24} \frac{n^2}{4} + \frac{5n}{24}$ mono triangles.
- 2. Assume $n \equiv 3 \pmod{4}$. For all 2-colorings of the edges of K_n there are at least $\frac{n^3}{24} \frac{n^2}{4} + \frac{5n}{24} + \frac{1}{2}$ mono triangles.
- 3. Assume $n \equiv 0 \pmod{2}$. For all 2-colorings of the edges of K_n there are at least $\frac{n^3}{24} \frac{n^2}{4} + \frac{n}{3}$ mono triangles

Proof:

A porism to the proof of Theorem 3.1 is that

$$|R| + |B| = \binom{n}{3} - |ZAN|/2.$$

Hence we will upper bound |ZAN|.

Case 1: $n \equiv 1 \pmod{2}$. The degree of each vertex is $n-1 \equiv 0 \pmod{2}$. To maximize |ZAN| we would, at each vertex, color half of the edges RED and half BLUE. So each vertex contributes $(\frac{n-1}{2})^2$, and there are n vertices, so we have $|ZAN| \leq \frac{(n-1)^2n}{4}$. Since $|ZAN|/2 = M \in \mathbb{N}$ we have

$$|\mathrm{ZAN}|/2 \le \left| \frac{(n-1)^2 n}{8} \right|$$

hence

$$|R| + |B| \ge \frac{n(n-1)(n-2)}{6} - \left| \frac{(n-1)^2 n}{8} \right|$$

Case 1a: $n \equiv 1 \pmod{4}$ so $(n-1)^2 \equiv 0 \pmod{16}$. Hence:

$$\left| \frac{(n-1)^2 n}{8} \right| = \frac{(n-1)^2 n}{8}.$$

Hence

$$|R| + |B| \ge \frac{n(n-1)(n-2)}{6} - \frac{(n-1)^2n}{8} = \frac{n^3}{24} - \frac{n^2}{4} + \frac{5n}{24}$$

Case 1b: $n \equiv 3 \pmod{4}$. One can easily show that

$$\left| \frac{(n-1)^2 n}{8} \right| = \frac{(n-1)^2 n - 4}{8}.$$

Hence

$$|R| + |B| \ge \frac{n(n-1)(n-2)}{6} - \frac{(n-1)^2n - 4}{8} = \frac{n^3}{24} - \frac{n^2}{4} + \frac{5n}{24} + \frac{1}{2}$$

Case 2: $n \equiv 0 \pmod{2}$. The degree of each vertex is $n-1 \equiv 1 \pmod{2}$. To maximize |ZAN| we would, at each vertex, color $\frac{n-2}{2}$ of the edges RED

and color $\frac{n}{2}$ of the edges BLUE. So each vertex contributes $\frac{n(n-2)}{4}$, and there are n vertices, so we have $|ZAN| \leq \frac{n^2(n-2)}{4}$. Since $|ZAN|/2 = M \in \mathbb{N}$ we have

$$|\mathrm{ZAN}|/2 \le \left| \frac{n^2(n-2)}{8} \right|.$$

Since $n \equiv 0 \pmod{2}$, $n^2(n-2) \equiv 0 \pmod{8}$, so the bound on |ZAN|/2 is always in N. Hence we do not need the floor. Hence

$$|\mathrm{ZAN}|/2 \le \frac{n^2(n-2)}{8}$$

$$|R| + |B| \ge \frac{n(n-1)(n-2)}{6} - \frac{n^2(n-2)}{8} = \frac{n^3}{24} - \frac{n^2}{4} + \frac{n}{3}.$$

Note 4.2 The bounds given in Theorem 4.1 are tight. This was proven by Goodman [5] and Savvy [8]. They view the problem differently. They asked given that you want t triangles, how big must n be?

5 What is Already Known

Def 5.1 R(k) is the least number n such that any 2-coloring of the edges of K_n has a mono K_k . By Ramsey's theorem for graphs, for all k, R(k) exists. It is known that R(3) = 6 and R(4) = 18; however, all that is known about R(5) is $43 \le R(5) \le 48$. It is a standard result that $R(k) \le 4^{k-c \log k}$ for some constant c.

- 1. Stanislaw Radziszowski and Konrad Piwakowski [6] proved the following: All 2-colorings of K_{18} have 9 mono K_4 's. This is tight. This was proven with the help of a computer.
- 2. Erdős [3] proved the following:

Theorem 5.2 Let $k \in \mathbb{N}$. Let R = R(k). Let N be large, so large that $R \ll N$. Let COL be a 2-coloring of the edges of K_N . Then there are $\geq \frac{N^k}{4(1+o(1))k^2}$ mono K_k 's.

We give the proof because we can.

Proof: Let $A_1, \ldots, A_{\binom{N}{R}}$ be a list of all the *R*-subsets of [N].

Note that by the definition of R and the A_i 's, every A_i has a mono K_k . We can't just say there are $\binom{N}{R}$ mono K_k 's since it may be that two A_i 's produce the same mono K_k .

We now produce many monochromatic K_k 's.

- (a) $X = \{A_i : 1 \le i \le {N \choose R} \}$
- (b) $Y = \emptyset$. Y will contain many mono K_k 's.
- (c) If $X \neq \emptyset$ do the following (else terminate). Let *i* be the least number such that $A_i \in X$. It has a mono K_k . Call it C.
 - i. Add C to Y (we will soon see that C is not already in Y).
 - ii. Remove from X all A_j 's that have C in them. Hence we are removing $\binom{N-k}{R-k}$ A_j 's.
 - iii. Goto Step c

In every iteration X loses $\binom{N-k}{R-k}$ A_j 's. Hence the number of mono K_k 's that this process produces is at least

$$\frac{\binom{N}{R}}{\binom{N-k}{R-k}} = \frac{N!}{R!(N-R)!} \times \frac{(N-R)!(R-k)!}{(N-k)!} = \frac{N!}{(N-k)!} \times \frac{(R-k)!}{R!}.$$

We need to lower bound this quantity. We use $\frac{(R-k)!}{R!} \ge \frac{1}{R^k} \ge \frac{1}{4^{k^2}}$. The other inequality we need will be more delicate.

$$\frac{N!}{(N-k)!} \ge (N-k+1)^k = \frac{N^k}{(N/(N-k+1))^k}$$

We look at the denominator

$$\left(\frac{N}{N-k+1}\right)^k = \left(1 + \frac{k-1}{N-k+1}\right)^k \sim e^{((k-1)k)/(N-k+1)} \sim 4^{(c(k-1)k)/(N-k+1)}$$

for some constant c. Hence the number of mono K_k 's is at least

$$\frac{N^k}{4^{k^2 + (c(k-1)k/(N-k+1))}} = \frac{N^k}{4^{k^2(1+o(1))}}.$$

- 3. David Conlon [1] proved the following: Fix t. For n large, for any 2-coloring of K_n there are $\frac{n^t}{C^{(1+o(1))t^2}}$ mono K_t 's where $C \sim 2.18$. Note that Conlon's result is an improvement of Theorem 5.2.
- 4. Jacob Fox [4] looks at the problem for target graphs G other than K_t .
- 5. For more results (1) look at the bibliographies of the papers above, (2) Google *Ramsey Multiplicity*, and (3) watch this cool lecture by David Conlon:

https://www.ima.umn.edu/2014-2015/W9.8-12.14/21327

6 Open Problems

Here is a large class of open problems with the same theme as Theorem 4.1. First we need some notation.

Notation 6.1 Let $k \geq 1$.

- 1. C_k is the cycle graph on k vertices. Vertex i has an edge to vertices $i-1 \pmod{k}$ and $i+1 \pmod{k}$ but no other vertices.
- 2. P_k is the path graph on k vertices. If $i \leq k-1$ then vertex i has an edge to vertex i+1.
- 3. W_k is the wheel graph on k vertices. It is C_{k-1} with one more vertex that has an edge to all other vertices.
- 4. $K_{1,k}$ is the star graph.
- 5. If G is a graph on k vertices, R(G) is the least number n such that any 2-coloring of the edges of K_n has a mono G. By Ramsey's theorem for graphs, for all G, R(G) exists; however, for many graphs, R(G) is much lower than R(k).

6. The values of R(G) are known for $G \in \{C_k, P_k, W_k, K_{1,k}\}$. See Stanislaw Radziszowski's survey of Small Ramsey Numbers [7].

We now state some open problems.

- 1. Fix k. Find the function f such that, for all 2-colorings of K_n there are f(n) mono K_k 's. Try to make f(n) as large as possible. The k=3 case is solved. The k=4 case is not solved, but it is plausible that it will be. The k=5 case is harder than finding R(5) and hence is unlikely to be solved ... ever.
- 2. Replace K_k in the last item with C_k , P_k , W_k , $K_{1,k}$ or whatever your favorite parameterized set of graphs is. When R(G) is known there is hope of solving this problem. If R(G) is unknown, then note that the problem of finding f(n) is harder than finding R(G).
- 3. Let G be your favorite graph and L be your favorite number. Find the least n (or at least a non-obvious k) such that every 2-coloring of the edges of K_n yields L copies of a mono G.

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