Goal: Implement unreliable rate-controlled message delivery to neighbors.

1 Deadline

DUE: Oct 27. late deadline will be Monday, Oct 30; you have only one late deadline to play with.

2 Context

This project extends the previous project by providing rate-controlled message transmission. We’ll add a command called sendmsg that will inject a message into the network when doing so will not exceed a rate limit. The rate-control requirement ensures that messages from a host do not swamp its neighbors or the network. Because the service runs over UDP (over IP), without any retransmission, it remains unreliable.

The rate control is the so-called “leaky bucket” algorithm. Water (packets) pours into a leaky bucket (queue). When the bucket has water, it leaks steadily (sends at fixed rate). If the incoming water flow is so fast that it fills the bucket (queue is filled), water spills out (packets are dropped).

3 Service specification

The sendmsg service:

• Takes dst-address, the neighbor’s logical address (the 32 bit source address in the packet header), and msg, a message to send, as parameters. Consider: sscanf("sendmsg %u %s"), which might work.

• Limits the rate of sending to 10 packets (of maximum size 1000 bytes each) per second to each neighbor.

• Uses a per-neighbor output queue of maximum size 10 packets (irrespective of packet size) to limit the sending rate.

When sendmsg is invoked, it does the following in sequence:

• Look up the IP address and UDP port of the neighbor’s unicast socket in the neighbor table. (Don’t send messages to the multicast address!)

• Check if last packet was sent 0.1 seconds ago or earlier:
  – If so, send the message immediately using UDP sendto
  – If not, enqueue the message in the output queue subject to maximum queue size (10).
  – If the queue is full, drop packet

• Ensures that the queue, if not empty, is drained by one packet every 0.1 seconds.

1 Think of sendmsg as a wrapper over UDP sendto. This is an example of how services are composed over layers, i.e., a higher protocol layer uses a lower layer’s service specification to provide something more.
4 Implementation requirements

In addition to the requirements of P2, P3 should:

- Implement a command from STDIN with signature `sendmsg dst-address message` providing the semantics of the `sendmsg` service.
- Handle queue overflow by printing “ERROR:NOBUFF”. Do not block the `sendmsg` command waiting for the queue to drain.
- Handle unknown or dead neighbor by printing “ERROR:NOROUTE”
- Not leak memory
- Print received (version 1, protocol 1) messages on the unicast socket with `printf("U: %.s",...).

5 Hints

- Figure out what events will be required to drain the buffer (queue) periodically and use your event queue appropriately.
- Be careful when `free`ing packets; you want to free them, but only when you actually can.
- You should be able to test by sending to yourself.
- `sendmsg` is unrelated to the system call of the same name. You may use that one if you feel so inclined, but there’s no reason to.
- Multicast transmissions may also be rate-limited if you like, but the only transmission that will remain in multicast is the periodic hello, which generated on
- Receiving a unicast message from a neighbor may refresh the neighbor’s entry in the neighbor table. It should not create a new entry, because eventually, we’ll have forwarding, and the source address of a forwarded message is not a neighbor.

6 Objectives

At the end of this assignment, you should understand:

- maintaining queues of packets without coredumps :)
- scheduling events to manipulate queues of packets.
- simple rate control.

7 Free reign


To process packets in ruby requires using pack and unpack functions that can convert between packets as strings and arrays of constituent header values.

Later assignments may also be completed in C or Ruby.