Linear Data Structures

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Overview

- Linear data structures
  - General properties
- Implementations
  - Array
  - Linked list
- Restricted abstractions
  - Stack
  - Queue
**Linear Data Structures**

- **1-to-1 relationship between elements**
  - Each element has unique *predecessor* & *successor*
  - Results in total ordering over elements
  - For any two distinct elements $x$ and $y$, either $x$ comes before $y$ or $y$ comes before $x$
Linear Data Structures

Terminology

- Head (first element in list) \(\Rightarrow\) no predecessor
- Tail (last element in list) \(\Rightarrow\) no successor

Operations

- Add element
- Remove element
- Find element
Add & Remove Elements

Add an element

Where?

- At head (front) of list
- At tail (end) of list
- After a particular element

Remove an element

- Remove first element
- Remove last element
- Remove a particular element (e.g., String “Happy”)
  - What if “Happy” occurs more than once in list?
## Accessing Elements

### How do you find an element?
- At head (front) of list
- At tail (end) of list
- By position
  - Example: the 5th element
- By iterating through the list, and using relative position
  - Next element (successor)
  - Previous element (predecessor)
Two basic implementation techniques for lists

- **Store elements in an array**

- **Store as a linked list**
  - Place each element in a separate object (node)
  - Node contains reference to other node(s)
  - Link nodes together
Linked List

Properties
- Elements in linked list are ordered
- Element has successor

State of List
- Head
- Tail
- Cursor (current position)
Array Implementations

Advantages

- Can efficiently access element at any position
- Efficient use of space
  - Space to hold reference to each element

Disadvantages

- Expensive to grow / shrink array
  - Can amortize cost (grow / shrink in spurts)
- Expensive to insert / remove elements in middle
- Tricky to insert / remove elements at both ends
Linked Implementation

Advantages
- Can efficiently insert / remove elements anywhere

Disadvantages
- Cannot efficiently access element at any position
  - Need to traverse list to find element
- Less efficient use of space
  - 1-2 additional references per element
Efficiency of Operations

- **Array**
  - Insertion / deletion = $O(n)$
  - Indexing = $O(1)$

- **Linked list**
  - Insertion / deletion = $O(1)$
  - Indexing = $O(n)$
Linked List – Insert (After Cursor)

1. Original list & new element `temp`

![Diagram of original list and new element temp]

2. Modify `temp.next` → `cursor.next`

![Diagram showing modification after inserting temp]
3. Modify `cursor.next` \(\rightarrow\) `temp`

4. Modify `cursor` \(\rightarrow\) `temp`
Linked List – Delete (Cursor)

1. Find **before** such that \( \text{before.next} = \text{cursor} \)

2. Modify \( \text{before.next} \rightarrow \text{cursor.next} \)
Linked List – Delete (Cursor)

3. **Delete** cursor

![Diagram showing deletion of cursor]

4. **Modify** cursor → before.next

![Diagram showing modification of cursor]
Doubly Linked List

Linked list where

- **Element has predecessor & successor**

Issues

- Easy to find preceding / succeeding elements
- Extra work to maintain links (for insert / delete)
- More storage per node
Doubly Linked List – Insertion

Example

Must update references in both predecessor and successor nodes
Node Structures for Linked Lists

- **Linked list**
  
  ```java
  Class Node {
      Object data;
      Node next;
  }
  ```

- **Doubly linked list**
  
  ```java
  Class Node {
      Object data;
      Node next;
      Node previous;
  }
  ```
Restricted Abstractions

Restricting the operations an abstraction supports can be a good thing

- Efficiently supporting only a few operations efficiently is easier
- If limited abstraction is sufficient, easier to reason about limited abstraction than a more general one

Restricted list abstractions

- Stack (aka LIFO queue)
- Queue (aka FIFO queue)
- Dequeue (aka double ended queue)
Stack

Stack operations

- **Push**  = add element (to top)
- **Pop**    = remove element (from top)

Example

(a) A three-element stack

(b) After a `pop()` operation

(c) After a `push(W)` operation
Stack

Properties

- Elements removed in opposite order of insertion
- Last-in, First-out (LIFO)

A restricted list where

- Access only to elements at one end
- Can add / remove elements only at one end
Stack Applications

- Run-time procedure information

(a) Example of nested procedure calls

(b) Run-time stack while in procedure D

- Arithmetic computations
  - Postfix notation

- Simplified instruction set
  - Java bytecode
Stack Implementations

- **Linked list**
  - Add / remove from head of list

  (a) Logical view of the stack

<table>
<thead>
<tr>
<th>top</th>
<th>Z</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>Y</td>
</tr>
<tr>
<td></td>
<td>X</td>
</tr>
</tbody>
</table>

  (b) Its linked list implementation

  head → Z → Y → X

- **Array**
  - Increment / decrement Top pointer after push / pop

  ![Array Diagram]

  X   Y   Z   ...
Queue

Queue operations

- **Enqueue** = add element (to back)
- **Dequeue** = remove element (from front)

Example

(a) Three-element queue
(b) After deletion of X
(c) After insertion of W
Queue

Properties
- Elements removed in order of insertion
- First-in, First-out (FIFO)

A restricted list where
- Access only to elements at beginning / end of list
  - Add elements only to end of list
  - Remove elements only from front of list
- Alternatively, can add to front & remove from end
Queue Applications

**Examples**
- Songs to be played
- Jobs to be printed
- Customers to be served
- Citizens to cast votes

South Africa, 2004
Queue Implementations

- **Linked list**
  - Add to tail (back) of list
  - Remove from head (front) of list

![Linked list diagram]

- **Array**
- **Circular array**
Queue – Array

Store queue as elements in array

Problem

Queue contents move ("inchworm effect")

As result, can not add to back of queue, even though queue is not full
Queue – Circular Array

Circular array (ring)
- q[ 0 ] follows q[ MAX – 1 ]
- Index using q[ i % MAX ]

Problem
- Detecting difference between empty and nonempty queue
Queue – Circular Array

Approach 1
- Keep Front at first in
- Keep Back at last in

Problem
- Empty queue identical to queue with 1 element
Queue – Circular Array

Approach 2
- Keep Front at first in
- Keep Back at last in – 1

Problem
- Empty queue identical to full queue
Queue – Circular Array

Inherent problem for queue of size $N$
- Only $N$ possible (Front – Back) pointer locations
- $N+1$ possible queue configurations
  - Queue with 0, 1, … $N$ elements

Solutions
- Maintain additional state information
  - Use state to recognize empty / full queue
- Examples
  - Record Size
  - Record QueueEmpty flag
- Leave empty element in queue
- Store marker in queue