CMSC 313
Introduction to Computer Systems
Lecture 13
Assembly Language, cont.

Alan Sussman
als@cs.umd.edu

Administrivia

- Project 3 due tomorrow, 6PM
  - questions?
- Exam #1 questions?
- Quiz #3 on Wednesday
- Continue reading Bryant and O'Hallaron Section 4.1 (Y86 subset) and Chapter 3, for more info on IA-32 instruction set architecture
- AWC 400 level lecture series this Tuesday to Thursday, 4:45-6PM in CSIC 3117

Chapters 3 and 4.1, Bryant and O'Hallaron

ASSEMBLY LANGUAGE

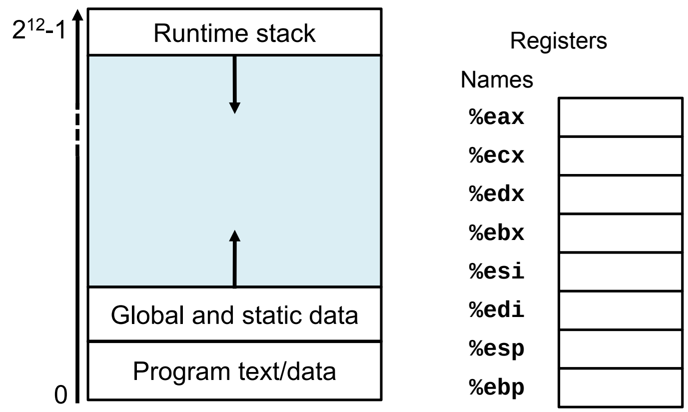
A full assembly

- This is the result of running the Y86 assembler on the example assembler source code:

```
0x000: 308000000000 |       irmovl $0,%eax    # sum = 0
0x006: 308101000000 |       irmovl $1,%ecx    # num = 1
0x00c: 6010         | Loop: addl   %ecx,%eax  # sum += num
0x00e: 308201000000 |       irmovl $1,%edx    # tmp = 1
0x014: 6021         |       addl   %edx,%ecx  # num++
0x016: 3082e8030000 |       irmovl $1000,%edx # lim = 1000
0x01c: 6112         |       subl   %ecx,%edx  # if lim - num >= 0
0x01e: 750c000000   |       jge    Loop       #   loop again
                    |
0x023: f308         |       wrint  %eax       # printf("%d", sum)
0x025: 30820a000000 |       irmovl $10,%edx   # ch = '\n'
0x02b: f128         |       wrch   %edx       # printf("%c", ch)
0x02d: 10           |       halt
```

Y86 program state

- 2^{12} bytes of memory
- You can set the stack to start somewhere other than 0x1000, but you have to explicitly set it



Working with Y86

- Source code is usually stored in **`*.ys`** files
- On the Grace systems, there are two programs available in **`~/313public/bin`** for working with Y86 programs
- **`yas`** is the Y86 assembler, which creates **`*.yo`** files
  - Run like this: **`yas prog.ys`**
- **`yis`** is the Y86 simulator, which operates on **`*.yo`** files
  - Run like this: **`yis prog.yo`**

Y86 data movement instructions

Instruction	Effect	Description
`irmovl V,R`	Reg[R] ← V	Immediate-to-register move
`rrmovl rA,rB`	Reg[rB] ← Reg[rA]	Register-to-register move
`rmmovl rA,D(rB)`	Mem[Reg[rB]+D] ← Reg[rA]	Register-to-memory move
`mrmovl D(rA),rB`	Reg[rB] ← Mem[Reg[rA]+D]	Memory-to-register move

- **`irmovl`** is used to place known numeric values (labels or numeric literals) into registers
- **`rrmovl`** copies a value between registers
- **`rmmovl`** stores a word in memory
- **`mrmovl`** loads a word from memory
- **`rmmovl`** and **`mrmovl`** are the only instructions that access memory - Y86 is a load/store architecture

Examples of data movement

```
irmovl $55,%edx      # d = 55
rrmovl %edx,%ebx     # b = d
irmovl Array,%eax    # a = Array
rmmovl %ebx,4(%eax)  # a[1] = 55
mrmovl 0(%eax),%ecx  # c = a[0]
halt
  .align 4
Array:
  .long 0x6f
  .long 0x84
```

Data movement example, cont.

- Assembler output:

```
0x000: 308237000000 | irmovl $55,%edx     # d = 55
0x006: 2023         | rrmovl %edx,%ebx    # b = d
0x008: 30801c000000 | irmovl Array,%eax   # a = Array
0x00e: 403004000000 | rmmovl %ebx,4(%eax) # a[1] = 55
0x014: 501000000000 | mrmovl 0(%eax),%ecx # c = a[0]
0x01a: 10           | halt
                    |
0x01c:              |   .align 4
0x01c:              | Array:
0x01c: 6f000000     |   .long 0x6f
0x020: 84000000     |   .long 0x84
```

- Simulator output:

```
Stopped in 6 steps at PC = 0x1b.  Exception 'HLT', CC Z=1 S=0 O=0
Changes to registers:
%eax:   0x00000000      0x0000001c
%ecx:   0x00000000      0x0000006f
%edx:   0x00000000      0x00000037
%ebx:   0x00000000      0x00000037

Changes to memory:
0x0020: 0x00000084      0x00000037
```


Y86 input/output instructions

Instruction	Effect	Description
`rdch R`	`scanf("%c", &Reg[R])`	Read character
`rdint R`	`scanf("%d", &Reg[R])`	Read integer
`wrch R`	`printf("%c", Reg[R])`	Write character
`wrint R`	`printf("%d", Reg[R])`	Write integer

- All these instructions are extensions to Y86 we've added to the ones in the book
- These are what allow you to interact with the simulator and write more interesting programs

I/O example

- Assembler output:

```
0x000: f208         | rdint  %eax     # a = 65 (via scanf())
0x002: f038         | rdch   %ebx     # b = 'B' (via scanf())
0x004: f308         | wrint  %eax     # printf("%d", a)
0x006: f108         | wrch   %eax     # printf("%c", a)
0x008: f338         | wrint  %ebx     # printf("%d", b)
0x00a: f138         | wrch   %ebx     # printf("%c", b)
0x00c: 30810a000000 | irmovl $10,%ecx # c = 10
0x012: f118         | wrch   %ecx     # printf("%c", c)
0x014: 10           | halt
```

- Simulator run:

```
$ echo 65B | yis io.yo
65A66B
Stopped in 9 steps at PC = 0x15.  Exception 'HLT', CC Z=1 S=0 O=0
Changes to registers:
%eax:   0x00000000      0x00000041
%ecx:   0x00000000      0x0000000a
%ebx:   0x00000000      0x00000042

Changes to memory:
```


Y86 integer instructions

Instruction	Effect	Description
`addl S,D`	Reg[D] ← Reg[D] + Reg[S]	Addition
`subl S,D`	Reg[D] ← Reg[D] - Reg[S]	Subtract
`andl S,D`	Reg[D] ← Reg[D] & Reg[S]	Bitwise AND
`xorl S,D`	Reg[D] ← Reg[D] ^ Reg[S]	Bitwise XOR
`multl S,D`	Reg[D] ← Reg[D] * Reg[S]	Multiplication*
`divl S,D`	Reg[D] ← Reg[D] / Reg[S]	Integer division*
`modl S,D`	Reg[D] ← Reg[D] % Reg[S]	Remainder*

- All these instructions operate on two integers, and set the condition code flags appropriately
- Instructions marked with an asterisk (*) are extensions to Y86 we've added to the ones in the book

Integer instruction example

- Assembler output:

```
0x000: 308003000000 | irmovl $3,%eax   # a = 3
0x006: 308305000000 | irmovl $5,%ebx   # b = 5
0x00c: 6003         | addl   %eax,%ebx # b = a + b
0x00e: f308         | wrint  %eax
0x010: 308620000000 | irmovl $32,%esi  # 32 == ' '
0x016: f168         | wrch   %esi
0x018: f338         | wrint  %ebx
0x01a: 30860a000000 | irmovl $10,%esi  # 10 == '\n'
0x020: f168         | wrch   %esi
0x022: 10           | halt
```

- Simulator run:

```
3 8
...
```

- Notice these instructions are destructive; they overwrite the second operand
  – Need to make copies if you need old values

Condition codes

- Performing integer operations causes various flags to be set, describing the attributes of the result of the operation
- These are used by other, subsequent instructions to perform conditional branching
- The three we are concerned with are:
  – OF: overflow flag; did the operation overflow?
  – SF: sign flag; is the result negative?
  – ZF: zero flag; is the result zero?

Branch instructions

- These are used to perform the effect of if statements, loops, and switches
- When encountered, if a certain condition is true, control flow will then go to the address specified, rather than advancing to the next instruction
  – The address of the next instruction to be executed is held in the program counter; in many architectures, this is held in an accessible register (not so with Y86).

Y86 branch instructions

Instruction	Branch if...	Description
`jmp Label`	1	Unconditional jump
`jle Label`	(SF ^ OF) | ZF	Jump if less or equal
`jl  Label`	SF ^ OF	Jump if less
`je  Label`	ZF	Jump if equal
`jne Label`	~ZF	Jump if not equal
`jge Label`	~(SF ^ OF)	Jump if greater or equal
`jg  Label`	~(SF ^ OF) & ~ZF	Jump if greater

- Each instruction relies on the condition codes set by the most recent integer instruction

Branch example 1

- Assembler output:

```
0x000: f208           |          rdint  %eax
0x002: 308700000000   |          irmovl $0,%edi    # consistent zero
0x008: 308600000000   |          irmovl $0,%esi    # sum = 0
0x00e: 6070           |          addl   %edi,%eax
0x010: 7320000000     |          je     EndLoop
                      |
0x015: 6006           | Loop:    addl   %eax,%esi  # sum += n
0x017: f208           |          rdint  %eax
0x019: 6070           |          addl   %edi,%eax
0x01b: 7415000000     |          jne    Loop
                      |
0x020: f368           | EndLoop: wrint  %esi
0x022: 30830a000000   |          irmovl $10,%ebx
0x028: f138           |          wrch   %ebx
0x02a: 10             |          halt
```

- Simulator output:

```
$ echo 1 4 9 16 25 0 | yis io.yo
55
Stopped in 29 steps at PC = 0x2b.  Exception 'HLT', CC Z=1 S=0 O=0
...
```