

CMSC 313

Introduction to Computer Systems

Lecture 16

Process Control

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Administrivia

- Project 4 due Tuesday, Nov. 3
- Read Sections 8.2-8.5, on process control
- I'm out of town Monday and Tuesday
 - guest lecturer (Robert) on Monday
 - no office hours Tuesday

Sections 8.2-8.5, Bryant and O'Hallaron

PROCESS CONTROL

Process IDs

- Every active process has a unique process ID, which can be obtained via **getpid()**
- The process ID of the current process' parent (the process which created the current process) can be obtained via **getppid()**

- Function prototypes:
#include <unistd.h>
#include <sys/types.h>

```
pid_t getpid(void);  
pid_t getppid(void);
```

Creating new processes

- In UNIX, a new process is created by an existing process, making a parent-child relationship between the two processes
- The system call to do this is **fork()**; creates a new copy of the parent process
 - all variables (the whole address space) are copied
 - point of execution (PC) is copied
 - file table information is copied

The **fork()** function

- Prototype:
`#include <unistd.h>`
`#include <sys/types.h>`
`pid_t fork(void);`
- Returns **twice**, in BOTH parent and child
 - -1: error occurred
 - generally due to process table being full or resource limit reached
 - 0: returned to child process
 - > 0: returns pid of child process to parent

fork() example

```
/* #include statements omitted */

int main() {
    int var = 313;
    pid_t child_pid;
    if ((child_pid = fork()) < 0)
        err(EX_OSERR, "fork error");
    if (child_pid) { /* parent code */
        printf("Parent pid = %d; my child has pid = %d\n",
            getpid(), child_pid);
        var++;
        printf("Var in parent = %d\n", var);
    }
    else { /* child code */
        printf("Child pid = %d; my parent has pid = %d\n",
            getpid(), getppid());
        var--;
        printf("Var in child = %d\n", var);
    }
    return 0;
}
```

Execution order after a **fork()**

- The previous example's output on one machine was:
`Child pid = 18532; my parent has pid = 18531`
`Var in child = 312`
`Parent pid = 18531; my child has pid = 18532`
`Var in parent = 314`
- On Grace, it was:
`Parent pid = 726; my child has pid = 727`
`Var in parent = 314`
`Child pid = 727; my parent has pid = 726`
`Var in child = 312`
- It could even be:
`Parent pid = 23892; my child has pid = 23894`
`Child pid = 23894; my parent has pid = 23892`
`Var in child = 312`
`Var in parent = 314`
- Print order **within** a process is (usually) determinate
- Print order **between** processes is not

fork() semantics

- Some things inherited by a child from its parent process:
 - process credentials: user and group ID (UIDs and GIDs in UNIX terminology)
 - environment
 - a copy of the parent's memory contents, including program code, runtime stack, and heap
 - open file descriptors – **FILE** *'s from **fopen()**. The current file position is also shared between the parent and child, which can cause file consistency issues.
 - signal handling settings (a UNIX way of handling events external to the process, from the operating system or another program)

fork() semantics, con't.

- Some things inherited by a child from its parent process, con't:
 - current working directory (set with **cd**, viewed with **pwd**)
 - root directory
 - resource limits (that can be set and viewed with the **tcsh** **limit** command, or **ulimit** in **bash**)
 - the controlling terminal (the program that controls **stdin**, **stdout**, and **stderr** for the process, which is usually a shell), so the child reads input from and prints output to the same devices that the parent does
 - "nice" value (to determine process priority for scheduling by OS)

fork() semantics, con't.

- Some things that are unique to a child process:
 - its process ID
 - it has a different parent process ID (the parent, not the parent's parent)
 - it has its own copy of file descriptors and directory streams.
 - its process times are unique to it
 - its resource utilizations are initially set to 0
 - its pending signals are initialized to the empty set

The dangers of fork()

- The process table in the kernel can hold only a finite number of processes; what happens if you fill it up?

```
#include <unistd.h>
int main() {
    while (1) fork();
}
```
- That is a fork **bomb**, and it is a Very Bad Thing
- Fork bombs can be unintentional; a loop that doesn't terminate correctly can easily cause one
 - many students write them accidentally
- Often, it can require sysadmin intervention (e.g., reboot, killing all user processes)
- Lesson: Be careful when using **fork()** in a loop

Reaping child processes

- When a process exits, it is still tracked by the kernel (remember the termination process state?)
- Processes are released from the process table only when their parent *reaps* the terminated child; until this happens, the terminated process is called a zombie
- A parent can release its zombie children from the process table via either the **wait()** or **waitpid()** system calls
- If the parent terminates before the child, the child is orphaned, and then adopted by the **init** process (pid #1); **init** will reap children as soon as they terminate

wait() system calls

- Can be used to obtain the exit status of the reaped child (or not, if you don't care)
- **pid_t wait(int *status);**
 - requires **<sys/types.h>** and **<sys/wait.h>**
 - pass in a pointer to an int (or **NULL**) that will be populated by the status of the reaped process
 - will reap any single terminated child
 - **blocking** wait; does not return until a terminated child exists (if a child exists)
 - returns -1 on error (e.g., no unwaited-for children exist)
 - returns pid of reaped process on success

wait() system calls, cont.

- **pid_t waitpid(pid_t pid, int *status, int options);**
 - will wait on one specified process
 - **pid** is pid of the child process
 - **options** is a number formed from the bitwise OR of several flags (or just 0); **WNOHANG** is the most useful of these flags (doesn't block)

wait() example

```
/* #include statements omitted */
int main() {
    pid_t child_pid;
    if ((child_pid = fork()) < 0)
        err(EX_OSERR, "fork error");
    if (child_pid) { /* parent code */
        int status;
        wait(&status); /* nothing happens until child exits */
        printf("Parent pid = %d; my child had pid = %d\n",
               getpid(), child_pid);
        printf("Child exited with status %d\n", status);
    }
    else { /* child code */
        printf("Child pid = %d; my parent has pid = %d\n",
               getpid(), getppid());
    }
    return 0;
}
```

Exit status

- The status argument points to an **int**; the **int** value is actually more than just the exit code
- We can use macros defined in **<sys/wait.h>** to learn information about the reaped child
 - **WIFEXITED(status)**: true if child terminated normally (via exit/return)
 - **WEXITSTATUS(status)**: the exit status of the normally terminated child
 - **WTERMSIG(status)**: the signal that caused the child to terminate

Environment variables

- Examples:
 - **PATH** (where does the shell look for a program?)
 - **PAGER** (what program do I want to use to view files one page at a time?) (hint: less)
- Are not shell variables
 - shell vars. only affect current shell, env. vars are copied to all child processes run by shell
- Shell commands to set shell variables
 - tcsh: **set var=value**
 - bash: **var=value**
- Shell commands to set environment variables
 - tcsh: **setenv VAR value**
 - bash: **export VAR=value**

Environment variables, cont.

- In the shell, accessed using **\$**
 - try "**echo \$PATH**"
- In C programs, can access with 3-param **main()**:

```
int main(int argc, char *argv[], char *envp[]) { ... }
```

 - **envp** is an array of strings of the form **NAME=VALUE**
- Can use **getenv()** to get value of these variables even without using the modified **main()**
- The **extern char **environ** (declared in **<unistd.h>**) also holds the current environment in the same form as **envp**

Loading a new program

- Since the creation of a new process with **fork()** is just a clone of the original, we need a way to change processes to run other programs
- The **execve()** system call can be used to load a program in the context of the current process (so it is the same process, but different program)
- Prototype:

```
#include <unistd.h>
int execve(const char *filename,
           char * const argv[],
           char * const envp[]);
```
- Returns -1 on error; doesn't return on success
 - doesn't return?!

Using `execve()`

- **filename** has to be the absolute path of the executable
- Both the **argv** and **envp** arrays are arrays of strings, with a **NULL** pointer as the final element
 - Not-so-coincidentally, just like the **argv** and **envp** arrays in the 2- and 3-param forms of **main()**
- **argv**: argument vector for the new program
- **envp**: list of environment variable strings, each in the form "**NAME=VALUE**"
 - can just use **environ** to pass along the environment

`execve()` example

```
/* #include statements omitted */
extern char **environ;
int main() {
    char *args[] = { "ls", "-l", NULL };
    pid_t child_pid;
    if ((child_pid = fork()) < 0)
        err(EX_OSERR, "fork error");
    if (child_pid) { /* parent code */
        wait(NULL);
        printf("Parent pid = %d; my child had pid = %d\n",
              getpid(), child_pid);
    }
    else { /* child code */
        printf("PID %d replacing myself\n", getpid());
        execve("/bin/ls", args, environ);
        err(EX_OSERR, "exec error"); /* why no if statement? */
    }
    return 0;
}
```

Other exec functions

- There are 5 other functions that perform exec operations
 - **l**: use an argument list, terminated by **NULL**
 - **v**: use an argument vector
 - **p**: search the **PATH** list of directories
 - **e**: can specify environment
- `execle(const char *filename, ..., NULL, char * const envp[]);`
- `execv(const char *filename, char * const argv[]);`
- `execl(const char *filename, ..., NULL);`
- `execvp(const char *progname, char * const argv[]);`
- `execlp(const char *progname, ..., NULL);`