

CMSC 313

Introduction to Computer Systems

Lecture 21

Threads

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Chapter 13, Bryant and O'Hallaron

CONCURRENT PROGRAMMING

Administrivia

- Project 5 due tomorrow at 6PM
 - questions?
 - Do NOT use the submit server to test your projects; follow the instructions in the project specification that explain how to test on Grace
- Exam #2 returned next week
- Project 6 assigned next week
- Read Chapter 13 of Bryant and O'Hallaron

Threads example

```
#define THREAD_CT 2

void *print_stuff(void *ptr) {
    int i, id = * (int *) ptr;
    for (i = 0; i < 5; i++) {
        printf("Thread %d, loop %d\n", id, i);
        sleep(rand() % 2); /* sleep 0 or 1 seconds */
    }
    printf("Thread %d exiting\n", id);
    return NULL;
}

int main() {
    pthread_t tids[THREAD_CT];
    int i, ids[THREAD_CT];

    for (i = 0; i < THREAD_CT; i++) {
        ids[i] = i + 1;
        pthread_create(&tids[i], NULL, print_stuff, &ids[i]);
        printf("Thread 0 created thread %d\n", i + 1);
    }
    for (i = 0; i < THREAD_CT; i++) {
        pthread_join(tids[i], NULL);
        printf("Thread 0 reaped thread %d\n", i + 1);
    }
    return 0;
}
```

Execution of the example

```
$ ./simple
Thread 0 created thread 1
Thread 0 created thread 2
Thread 1, loop 0
Thread 2, loop 0
Thread 2, loop 1
Thread 1, loop 1
Thread 2, loop 2
Thread 1, loop 2
Thread 2, loop 3
Thread 2, loop 4
Thread 2 exiting
Thread 1, loop 3
Thread 1, loop 4
Thread 1 exiting
Thread 0 reaped thread 1
Thread 0 reaped thread 2
```

Return value example

```
#include <stdio.h>
#include <stdlib.h>
#include <pthread.h>

void *get_rand_num(void *args) {
    int *nump = malloc(sizeof(int));
    srand(pthread_self());
    *nump = rand();
    return nump;
}

int main() {
    pthread_t tid;
    void *ptr = NULL;

    pthread_create(&tid, NULL, get_rand_num, NULL);
    pthread_join(tid, &ptr);
    printf("Random number: %d\n", * (int *) ptr);
    return 0;
}
```

Bad return value example

```
#include <stdio.h>
#include <stdlib.h>
#include <pthread.h>

void *get_rand_num(void *args) {
    int num;
    srand(pthread_self());
    num = rand();
    return &num;
}

int main() {
    pthread_t tid;
    void *ptr = NULL;

    pthread_create(&tid, NULL, get_rand_num, NULL);
    pthread_join(tid, &ptr);
    printf("Random number: %d\n", * (int *) ptr);
    return 0;
}
```

Thread detachment

- Threads are by default joinable, meaning they can be reaped and killed by other threads, and a thread's memory resources stay until it is reaped
- We can detach threads so that they cannot be reaped or killed by other threads, and memory resources are automatically freed upon termination
- To avoid memory leaks, threads should either be reaped by another thread, or detached

Thread detachment, cont.

int pthread_detach(pthread_t tid);

- detaches the thread **tid**
- returns 0 on success, nonzero on error

- A thread can detach itself:
pthread_detach(pthread_self());
- We might use detachment if we have a constantly active process, like a mail daemon or web server, which shouldn't need to take time to reap terminated threads

Threads memory model

- Remember that threads run within the context of a single process
- Each thread has its own context, including a thread ID, stack, stack pointer, PC, CCs, and registers; everything else is shared
- It is possible for one thread to access another thread's stack if a pointer is made accessible

Threads memory model, cont.

- Global variables are always shared; there is only one of each global variable
- Automatic local variables are thread-local; each thread has its own copy of these in its stack
- Static local variables are also shared, just as globals
 - remember, static variables are just globals with restricted scope

Thread synchronization

- What will the following code output?

```
#define LOOPS 10000000
static int count = 0;

void *counter(void *args) {
    int i;
    for (i = 0; i < LOOPS; i++)
        count++;
    printf("Executed %d times\n", i);
    return NULL;
}

int main() {
    pthread_t tids[2];
    pthread_create(&tids[0], NULL, counter, NULL);
    pthread_create(&tids[1], NULL, counter, NULL);
    pthread_join(tids[0], NULL);
    pthread_join(tids[1], NULL);
    printf("Count: %d\n", count);
    return 0;
}
```

Thread synchronization, cont.

- The first two lines will be:
Executed 10000000 times
- And the last will be:
Count: 11398345
- Or maybe:
Count: 12398354
- Or even...
Count: 15892348
- But almost definitely not
"Count: 20000000" Why not?

Thread synchronization errors

- We might expect the increment to be an atomic operation, but it isn't
 - `i++` requires loading `i`'s value into a register, updating that value by adding 1 to it, then storing the result back into `i`'s location in memory
- Consider this schedule of events:
 - Thread A loads `i`'s value
 - Thread B loads `i`'s value
 - A updates register
 - B updates register
 - A stores register value in `i`
 - B stores register value in `i`
- What is the value stored in `i` now?
 - Only one more than it was before!

Thread synchronization errors, cont.

- In this example, we need to ensure that when one thread is in the middle of its load-update-store set of instructions, the other thread isn't in its own set
- For a given thread, these instructions constitute a *critical section* that should not have other threads' critical sections interleaved within them

Semaphores

- To properly implement a critical section, we can use semaphores
- These are special global variables, with nonnegative integer values, that can only be changed by two operations
 - in most literature, called *P* and *V*
 - we'll restrict ourselves to discussion of binary semaphores here, but semaphores can be more than just 0/1 values
- Used to ensure that only one thread is in a critical section at a time

Semaphores, cont.

- $P(s)$, or wait operation
 - if s is 0, waits until s is 1
 - if s is 1, sets s to 0 (decrements s)
- $V(s)$, or post operation
 - if s is 0, sets s to 1 (increments s), and restarts exactly one of the processes waiting for s to be 1
 - should not be called except after a $P(s)$ call
- All elements of these operations are indivisible (i.e., these operations are atomic)

Using semaphores

- Three functions, all in `<semaphore.h>`
 - Each return 0 on success, -1 on error
- ```
int sem_init(sem_t *s, 0,
 unsigned int value);
```
- initializes the semaphore pointed at by  $s$  to value
  - you must initialize semaphores before use
  - the second argument will always be 0 for our purposes
- ```
int sem_wait(sem_t *s);
```
- performs the $P(s)$ operation
- ```
int sem_post(sem_t *s);
```
- performs the  $V(s)$  operation

## Example using semaphores

```
#define LOOPS 10000000
static int count = 0;
static sem_t mutex;

void *counter(void *args) {
 int i;
 for (i = 0; i < LOOPS; i++) {
 sem_wait(&mutex);
 count++;
 sem_post(&mutex);
 }
 printf("Executed %d times\n", i);
 return NULL;
}

int main() {
 pthread_t tids[2];
 sem_init(&mutex, 0, 1);
 pthread_create(&tids[0], NULL, counter, NULL);
 pthread_create(&tids[1], NULL, counter, NULL);
 pthread_join(tids[0], NULL);
 pthread_join(tids[1], NULL);
 printf("Count: %d\n", count);
 return 0;
}
```