

CMSC 330: Organization of Programming Languages

OCaml 1 Functional Programming

Dialects of ML

- ▶ ML (Meta Language)
 - Univ. of Edinburgh, 1973
 - Part of a theorem proving system LCF
 - The Logic of Computable Functions
- ▶ SML/NJ (Standard ML of New Jersey)
 - Bell Labs and Princeton, 1990
 - Now Yale, AT&T Research, Univ. of Chicago, etc...
- ▶ OCaml (Objective CAML)
 - INRIA, 1996
 - French Nat'l Institute for Research in Computer Science

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2

Dialects of ML (cont.)

- ▶ Other dialects
 - MoscowML, ML Kit, Concurrent ML, etc...
 - SML/NJ and OCaml are most popular
- ▶ Languages all have the same core ideas
 - But small and annoying syntactic differences
 - So you should not buy a book with ML in the title
 - Because it probably won't cover OCaml

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3

Features of ML

- ▶ "Mostly functional"
 - Some assignments
- ▶ Higher-order functions
 - Functions can be parameters and return values
- ▶ Type inference
 - No need to write types in the source language
 - But the language is statically typed
 - Supports parametric polymorphism
 - Generics in Java, templates in C++

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4

Features of ML (cont.)

- ▶ Data types and pattern matching
 - Convenient for certain kinds of data structures
- ▶ Exceptions
- ▶ Garbage collection

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5

Functional Languages

- ▶ In a pure functional language
 - Every program is just an expression evaluation
- ```
let add1 x = x + 1;;

let rec add (x,y) = if x=0 then y else add(x-1, add1(y));;

add(2,3) = add(1,add1(3)) = add(0,add1(add1(3)))
 = add1(add1(3)) = add1(3+1) = 3+1+1
 = 5
```

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6

## Functional Languages (cont.)

- OCaml has similar basic behavior
  - Program = expression evaluation
- But has additional features
  - To ease the programming process
  - Features support
    - Less emphasis on data storage
    - More emphasis on function execution

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7

## A Small OCaml Program – Things to Notice

Use `let` to bind variables

Use `(*)` for comments (may nest)

No type declarations

Need to use correct print function (OCaml also has `printf`)

Line breaks, spacing ignored (like C, C++, Java, not like Ruby)

`::` ends a top-level expression

```
(* Small OCaml program *)
let x = 37;;
let y = x + 5;;
print_int y;;
print_string "\n";;
```

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8

## OCaml Interpreter

Expressions can be typed and evaluated at the top-level

```
3 + 4;;
- : int = 7
let x = 37;;
val x : int = 37
x;;
- : int = 37
let y = 5;;
val y : int = 5
let z = 5 + x;;
val z : int = 42
print_int z;;
42- : unit = ()
print_string "Colorless green ideas sleep furiously";;
Colorless green ideas sleep furiously- : unit = ()
print_int "Colorless green ideas sleep furiously";;
This expression has type string but is here used with type int
```

gives type and value of each expr

"-" = "the expression you just typed"

unit = "no interesting value" (like void)

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9

## OCaml Interpreter (cont.)

- Files can be loaded at top level

```
% ocaml
Objective Caml version 3.08.3

#use "ocaml1.ml";;
val x : int = 37
val y : int = 42
42- : unit = ()

- : unit = ()
x;;
- : int = 37
```

ocaml1.ml

```
(* Small OCaml program *)
let x = 37;;
let y = x + 5;;
print_int y;;
print_string "\n";;
```

#use loads in a file one line at a time

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10

## Run, OCaml, Run

- OCaml programs can be compiled using `ocamlc`
  - Produces `.cmo` ("compiled object") and `.cmi` ("compiled interface") files
    - We'll talk about interface files later
  - By default, also links to produce executable `a.out`
    - Use `-o` to set output file name
    - Use `-c` to compile only to `.cmo/.cmi` and not to link
    - You can use a `Makefile` if you need to compile your files
- Not needed for this course

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11

## Run, OCaml, Run (cont.)

- Compiling & running the previous small program

```
ocaml1.ml (* Small OCaml program *)
let x = 37;;
let y = x + 5;;
print_int y;;
print_string "\n";;
```

```
% ocamlc ocaml1.ml
% ./a.out
42
%
```

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12

## This Lecture

- Basic OCaml types
- Defining variables & functions with `let`
- Lists
- Pattern matching with `match`

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13

## Basic Types in OCaml

- Read `e : t` as “expression `e` has type `t`”  
`42 : int`                      `true : bool`  
`"hello" : string`            `'c' : char`  
`3.14 : float`                `() : unit` (\* don't care value \*)
- OCaml has **static** types to help you avoid errors
  - Note: Sometimes the messages are a bit confusing  
`# 1 + true;;`  
This expression has type `bool` but is here used with type `int`
  - Watch for the underline as a hint to what went wrong
  - But not always reliable

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14

## The Let Construct

- `let` is often used for defining local variables
  - `let x = e1 in e2` means
    - > Evaluate `e1`
    - > Then evaluate `e2`, with `x` bound to result of evaluating `e1`
    - > `x` is *not* visible outside of `e2`

```
let pi = 3.14 in pi *. 3.0 *. 3.0;;
pi;;
```

error

bind pi in  
body of let

floating point  
multiplication

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15

## The Let Construct (cont.)

- Compare to similar usage in Java/C

|                                                                   |                                                                         |
|-------------------------------------------------------------------|-------------------------------------------------------------------------|
| <pre>let pi = 3.14 in<br/>  pi *. 3.0 *.<br/>3.0;;<br/>pi;;</pre> | <pre>{<br/>  float pi = 3.14;<br/>  pi * 3.0 * 3.0;<br/>}<br/>pi;</pre> |
|-------------------------------------------------------------------|-------------------------------------------------------------------------|

- In the top-level, omitting `in` means “from now on”  
`# let pi = 3.14;;`  
(\* `pi` is now bound in the rest of the top-level scope \*)

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16

## `;;` versus `;`

- `;;` ends an expression in the top-level of OCaml
  - Use it to say: “Give me the value of this expression”
  - Not used in the body of a function
  - Not needed after each function definition
    - > Though for now it won't hurt if used there
- `e1; e2` evaluates `e1` and then `e2`, and returns `e2`  
`let p (s,t) = print_int s; print_int t; "Done!"`
  - Notice `no ;` at end
  - `;` is a **separator**, not a **terminator**
  - Invoking  
    `p (1,2)`
    - > Prints            `"1 2"`
    - > Returns          `"Done!"`

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17

## Examples – Semicolon

- Definition
  - `e1 ; e2` (\* evaluate `e1`, evaluate `e2`, return `e2` \*)
- `1 ; 2 ;;`
  - (\* 2 – value of 2<sup>nd</sup> expression is returned \*)
- `(1 + 2) ; 4 ;;`
  - (\* 4 – value of 2<sup>nd</sup> expression is returned \*)
- `1 + (2 ; 4) ;;`
  - (\* 5 – value of 2<sup>nd</sup> expression is returned to 1 + \*)
- `1 + 2 ; 4 ;;`
  - (\* 4 – because + has higher precedence than ; \*)

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18

## Nested Let

- Uses of `let` can be nested

```
let pi = 3.14 in
let r = 3.0 in
 pi *. r *. r;;
(* pi, r no longer in scope *)
```

```
{
 float pi = 3.14;
 {
 float r = 3.0;
 pi * r * r;
 }
}
/* pi, r not in scope */
```

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19

## Defining Functions

Use `let` to  
define  
functions

List parameters  
after function name

```
let next x = x + 1;;
next 3;;
let plus (x, y) = x + y;;
plus (3, 4);;
```

No parentheses  
on function calls

No return  
statement

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20

## Local Variables

- You can use `let` inside of functions for locals

```
let area r =
 let pi = 3.14 in
 pi *. r *. r
```

- And you can use as many `lets` as you want

```
let area d =
 let pi = 3.14 in
 let r = d /. 2.0 in
 pi *. r *. r
```

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21

## Function Types

- In OCaml, `->` is the function type constructor

- The type `t1 -> t2` is a function with argument or domain type `t1` and return or range type `t2`

- Examples

```
• let next x = x + 1 (* type int -> int *)
• let fn x = (float_of_int x) *. 3.14
 (* type int -> float *)
• print_string (* type string -> unit *)
```

- Type a function name at top level to get its type

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22

## Lists in OCaml

- The basic data structure in OCaml is the `list`

- Lists are written as `[e1; e2; ...; en]`

```
[1;2;3]
```

```
- : int list = [1;2;3]
```

- Notice type of list is `int list`

> Lists must be homogeneous

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23

## Lists in OCaml (cont.)

- More on OCaml lists

- The empty list is `[]`

```
[]
```

```
- : 'a list
```

- The `'a` means "a list containing anything"

> We'll find out more about this later

- List elements are separated using semicolons

- Warning: Don't use a comma instead of a semicolon

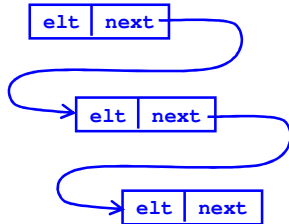
> Means something different (we'll see in a bit)

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24

## Consider a Linked List in C

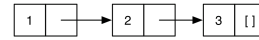
```
struct list {
 int elt;
 struct list *next;
};
...
struct list *l;
i = 0;
while (l != NULL) {
 i++;
 l = l->next;
}
```



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25

## Lists in OCaml are Linked

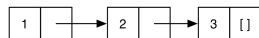


- `[1;2;3]` is represented above
  - A nonempty list is a pair (element, rest of list)
  - The element is the **head** of the list
  - The pointer is the **tail** or **rest** of the list
    - ...which is itself a list!
- Thus in math a list is either
  - The empty list `[]`
  - Or a pair consisting of an element and a list
    - This recursive structure will come in handy shortly

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26

## Lists are Linked (cont.)



- `::` prepends an element to a list
  - `h::t` is the list with `h` as the element at the beginning and `t` as the "rest"
  - `::` is called a **constructor**, because it builds a list
  - Although not emphasized, `::` does allocate memory
- Examples
  - `3::[]` (\* The list `[3]` \*)
  - `2::(3::[])` (\* The list `[2; 3]` \*)
  - `1::(2::(3::[]))` (\* The list `[1; 2; 3]` \*)

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27

## More Examples

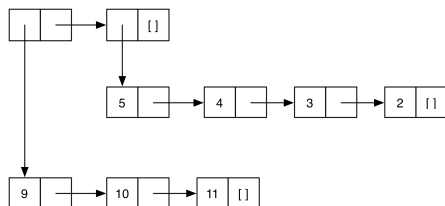
```
let y = [1;2;3] ;;
val y : int list = [1; 2; 3]
let x = 4::y ;;
val x : int list = [4; 1; 2; 3]
let z = 5::y ;;
val z : int list = [5; 1; 2; 3]
> not modifying existing lists, just creating new lists
let w = [1;2]::y ;;
This expression has type int list but is here
used with type int list list
> The left argument of :: is an element
> Can you construct a list y such that [1;2]::y makes sense?
```

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28

## Lists of Lists

- Lists can be nested arbitrarily
  - Example: `[ [9; 10; 11]; [5; 4; 3; 2] ]`
    - Type = `int list list`



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29

## Practice

- What is the type of
  - `[1;2;3]` `int list`
  - `[[]; []; [1.3;2.4]]` `float list list list`
  - `let func x = x::(0::[])` `int -> int list`

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30

## Pattern Matching

- ▶ To pull lists apart, use the `match` construct

```
match e with p1 -> e1 | ... | pn -> en
```
- ▶ `p1...pn` are **patterns** made up of
  - `[]`, `::`, and **pattern variables**
- ▶ `match` finds the first `pk` that matches shape of `e`
  - Then `ek` is evaluated and returned
  - During evaluation of `pk`, pattern variables in `pk` are bound to the corresponding parts of `e`

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31

## Pattern Matching Example

- ▶ Match syntax
  - `match e with p1 -> e1 | ... | pn -> en`
- ▶ Code 1
  - `let is_empty l = match l with`

```
 [] -> true
 | (h::t) -> false
```
- ▶ Outputs
  - `is_empty []` (\* evaluates to true \*)
  - `is_empty [1]` (\* evaluates to false \*)
  - `is_empty [1;2]` (\* evaluates to false \*)

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32

## Pattern Matching Example (cont.)

- ▶ Code 2
  - `let hd l = match l with (h::t) -> h`
- ▶ Outputs
  - `hd [1;2;3]` (\* evaluates to 1 \*)
  - `hd [1;2]` (\* evaluates to 1 \*)
  - `hd [1]` (\* evaluates to 1 \*)
  - `hd []` (\* Exception: Match failure \*)

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33

## Pattern Matching Example (cont.)

- ▶ Code 3
  - `let tl l = match l with (h::t) -> t`
- ▶ Outputs
  - `tl [1;2;3]` (\* evaluates to [2;3] \*)
  - `tl [1;2]` (\* evaluates to [2] \*)
  - `tl [1]` (\* evaluates to [] \*)
  - `tl []` (\* Exception: Match failure \*)

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34

## Pattern Matching – Wildcards

- ▶ An underscore `_` is a wildcard pattern
  - Matches anything
  - Doesn't add any bindings
  - Useful when you want to know something matches
    - But don't care what its value is
- ▶ In previous examples
  - Many values of `h` or `t` ignored
  - Can replace with wildcard `_`
  - Code behavior is identical

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35

## Pattern Matching – Wildcards (cont.)

- ▶ Code using `_`
  - `let is_empty l = match l with`

```
 [] -> true | (_::_) -> false
```
  - `let hd l = match l with (h::_) -> h`
  - `let tl l = match l with (_::t) -> t`
- ▶ Outputs
  - `is_empty [1]` (\* evaluates to false \*)
  - `is_empty []` (\* evaluates to true \*)
  - `hd [1;2;3]` (\* evaluates to 1 \*)
  - `tl [1;2;3]` (\* evaluates to [2;3] \*)
  - `hd [1]` (\* evaluates to 1 \*)
  - `tl [1]` (\* evaluates to [] \*)

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36

## Pattern Matching – Missing Cases

- ▶ When pattern is defined
  - OCaml will warn you about non-exhaustive matches
- ▶ When pattern is used
  - Exceptions for inputs that don't match any pattern
- ▶ Example

```
let hd l = match l with (h::_) -> h;;
Warning: this pattern-matching is not exhaustive.
Here is an example of a value that is not matched:
[]
hd [];;
Exception: Match_failure ("", 1, 11).
```

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37

## Pattern Matching – An Abbreviation

- ▶ `let f p = e`, where `p` is a pattern
  - is shorthand for `let f x = match x with p -> e`
- ▶ Examples
  - `let hd (h::_) = h`
  - `let tl (_,t) = t`
  - `let f (x::y::_) = x + y`
  - `let g [x; y] = x + y`
- ▶ Useful if there's only one acceptable input

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38

## Pattern Matching – Lists of Lists

- ▶ Can pattern match on lists of lists as well
- ▶ Examples
  - `let addFirsts`  
`((x::_) :: (y::_) :: _) = x + y`  
`addFirsts [[1;2]; [4;5]; [7;8;9]] = 5`
  - `let addFirstSecond`  
`(x::_) :: (_,y::_) :: _ = x + y`  
`addFirstSecond [[1;2]; [4;5]; [7;8;9]] = 6`
- ▶ Note – you probably won't do this much or at all
  - You'll mostly write recursive functions over lists instead

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39