Announcements

- Midterm is next Thursday (10/13/11)
  - Covers up through deadlock (this lecture)

- Project #2 is due Friday at 6:00 PM
Deadlocks

- **System contains finite set of resources**
  - memory space
  - printer
  - tape
  - file
  - access to non-reentrant code

- **Process requests resource before using it, must release resource after use**

- **Process is in a deadlock state when every process in the set is waiting for an event that can be caused only by another process in the set**
Formal Deadlocks

- 4 \textit{necessary} deadlock conditions:
  - Mutual exclusion - at least one resource must be held in a non-sharable mode, that is, only a single process at a time can use the resource. If another process requests that resource, the requesting process must be delayed until the resource is released.
  - Hold and wait - There must exist a process that is holding at least one resource and is waiting to acquire additional resources that are currently held by other processors.
Formal Deadlocks

– No preemption: Resources cannot be preempted; a resource can be released only voluntarily by the process holding it, after that process has completed its task

– Circular wait: There must exist a set \{P_0,\ldots,P_n\} of waiting processes such that \(P_0\) is waiting for a resource that is held by \(P_1\), \(P_1\) is waiting for a resource held by \(P_2\) etc.

• Note that these are not sufficient conditions
Detecting Deadlock

Work is a vector of length m (resources)
Finish is a vector of length n (processes)
- Allocation is an n x m matrix indicating the number of each resource type held by each process
- Request is an m x n matrix indicating the number of additional resources requested by each process

1. Work = Available;
   if Allocation[i] != 0 Finish = false else Finish = true;
2. Find an \( i \) such that Finish[i] = false and Request\( i \) <= Work if no such \( i \), go to 4
3. Work += Allocation; Finish[i] = true; goto step 2
4. If Finish[i] = false for some \( i \), system is in deadlock

**Note:** this requires \( m \times n^2 \) steps
Recovery from deadlock

- **Must free up resources by some means**
- **Process termination**
  - kill all deadlocked processes
  - select one process and kill it
    - must re-run deadlock detection algorithm again to see if it is freed.
- **Resource Preemption**
  - select a process, resource and de-allocate it
  - rollback the process
    - needs to be reset the process to a safe state
    - this requires additional state
  - starvation
    - what prevents a process from never finishing?
Deadlock Prevention

- Ensure that one (or more) of the necessary conditions for deadlock do not hold
- Hold and wait
  - guarantee that when a process requests a resource, it does not hold any other resources
  - Each process could be allocated all needed resources before beginning execution
  - Alternately, process might only be allowed to wait for a new resource when it is not currently holding any resource
Deadlock Prevention

● **Mutual exclusion**
  – Sharable resources do not require mutually exclusive access and cannot be involved in a deadlock.

● **Circular wait**
  – Impose a total ordering on all resource types and make sure that each process claims all resources in increasing order of resource type enumeration

● **No Premption**
  – Virtualize resources and permit them to be preempted. For example, CPU can be preempted.
Deadlock Avoidance

- Require additional information about how resources are to be requested - decide to approve or disapprove requests on the fly
- Assume that each process lets us know its maximum resource request
- Safe state:
  - system can allocate resources to each process (up to its maximum) in some order and still avoid a deadlock
  - A system is in a safe state if there exists a safe sequence
Safe Sequence

- Sequence of processes \(<P_1, \ldots, P_n>\) is a safe sequence if for each \(P_i\), the resources that \(P_i\) can request can be satisfied by the currently available resources plus the resources held by all \(P_j, j<i\).
- If the necessary resources are not immediately available, \(P_i\) can always wait until all \(P_j, j<i\) have completed.
Banker’s Algorithm

- Each process must declare the maximum number of instances of each resource type it may need
- Maximum can’t exceed resources available to system
- Variables:
  - \( n \) is the number of processes
  - \( m \) is the number of resource types
  - Available - vector of length \( m \) indicating the number of available resources of each type
  - Max - \( n \) by \( m \) matrix defining the maximum demand of each process
  - Allocation - \( n \) by \( m \) matrix defining number of resources of each type currently allocated to each process
  - Need: \( n \) by \( m \) matrix indicating remaining resource needs of each process
- Work is a vector of length $m$ (resources)
- Finish is a vector of length $n$ (processes)

1. Work = Available; Finish = false
2. Find an $i$ such that Finish[$i$] = false and Need[$i$] $\leq$ Work if no such $i$, go to 4
3. Work += Allocation[$i$]; Finish[$i$] = true; goto step 2
4. If Finish[$i$] = true for all $i$, system is in a safe state

Note this requires $m \times n^2$ steps
Banker’s Algorithm - Example

Three resources: A, B, C (10, 5, 7 instances each)

Consider the snapshot of the system at this time

<table>
<thead>
<tr>
<th>Alloc</th>
<th>Max</th>
<th>Avail</th>
<th>Need</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>A B C</td>
<td>A B C</td>
<td>A B C</td>
</tr>
<tr>
<td>P0</td>
<td>0 1 0</td>
<td>7 5 3</td>
<td>3 3 2</td>
</tr>
<tr>
<td>P1</td>
<td>2 0 0</td>
<td>3 2 2</td>
<td></td>
</tr>
<tr>
<td>P2</td>
<td>3 0 2</td>
<td>9 0 2</td>
<td></td>
</tr>
<tr>
<td>P3</td>
<td>2 1 1</td>
<td>2 2 2</td>
<td></td>
</tr>
<tr>
<td>P4</td>
<td>0 0 2</td>
<td>4 3 3</td>
<td></td>
</tr>
</tbody>
</table>

System is in a safe state, since the sequence <P1, P3, P4, P2, P0> satisfy the safety criteria.
Resource Request Algorithm

(1) If $\text{Request}_i \leq \text{Need}_i$ then goto 3
   - otherwise - the process has exceeded its maximum claim

(2) If $\text{Request}_i \leq \text{Available}$ then goto 3
   - otherwise process must wait since resources are not available

(3) Check request by having the system pretend that it has
allocated the resources by modifying the state as follows:
   - $\text{Available} = \text{Available} - \text{Request}_i$
   - $\text{Allocation} = \text{Allocation} + \text{Request}_i$
   - $\text{Need}_i = \text{Need}_i - \text{Request}_i$

* Find out if resulting resource allocation state is safe, otherwise the request must wait.