

Final Exam

CMSC 433: Programming Language Technologies and Paradigms

December 14, 2014

Name _____

Instructions

- This exam has 15 pages (including this one); make sure you have them all.
- You have 120 minutes to complete the exam
- The exam is worth 110 points. Allocate your time wisely: some hard questions are worth only a few points, and some easy questions are worth a lot of points.
- If you have a question, please raise your hand and wait for the instructor.
- **Write neatly and clearly indicate your answers.**

	Question	Points	Score
	Definitions	20	
	Distributed Computing	20	
	Java Concurrency	30	
	Erlang Execution	20	
	Erlang Coding	10	
	Parallelization	10	
	Total	110	

Enjoy your break!

1. Definitions (20 points total, 2 points each).

Next to each term on the left, write the letter on the right that corresponds to the definition of the term, or is a true statement about it.

- | | |
|-----------------------------|---|
| ---- Marshaling | (a) A technique for ensuring the atomicity of multiple method calls |
| ---- Parallel GC | (b) Performing garbage collection using multiple parallel threads |
| ---- Sequential consistency | (c) Performing garbage collection in parallel with the mutator |
| ---- Thread-safe class | (d) A property of the Java Memory Model (JMM) |
| ---- Client-side locking | (e) A class that satisfies its specification when used by multiple threads |
| ---- State-dependent action | (f) An example of which is calling <code>BlockingQueue.take</code> |
| ---- Cast (as in Erlang) | (g) A lock that can be acquired by multiple threads |
| ---- Dynamic typing | (h) A property of executions in which the effects of separate threads' operations can be viewed as happening in a total order |
| ---- Reentrant lock | (i) Erlang uses this style of typing |
| ---- Amdahl's Law | (j) A lock that is particularly important when synchronized methods call other synchronized methods |
| | (k) Predicts a task's possible speedup based on the number of processors and the amount of strictly-serial work |
| | (l) An action involving an object implementing a finite state machine |
| | (m) A message sent to a server that expects a response |
| | (n) Predicts how many threads you should create based on the number of processors and the ratio of wait time to compute time |
| | (o) A message sent to a server that needs no response |
| | (p) A class with no data races |
| | (q) Java uses this style of typing |
| | (r) The process of preparing a remote method invocation |

2. (Distributed computing, 20 points)

- (a) (Scala Actors, 4 points) You can send a message to a Scala actor `a` using the code `a ! msg`. Scala also allows you to send messages via the syntax `a !! msg` and the syntax `a !? msg`. How will these two invocations differ from the first?

- (b) (Java RMI, 4 points) Consider the following two possible definitions of a class `MyString`:

Implementation (a):

```
1 public class MyString {
2     private final String contents;
3     public String getContents() { return contents; }
4 }
```

Implementation (b):

```
1 import java.rmi.Remote;
2 import java.rmi.RemoteException;
3 public interface RemoteString extends Remote {
4     public String getContents() throws RemoteException;
5 }
6 public class MyString implements RemoteString {
7     private final String contents;
8     public String getContents() throws RemoteException { return contents; }
9 }
```

Suppose `x` is a reference to a *remote* object of type `Foo`, which has a method `f` that takes a `MyString` argument, and `y` is a reference to a local `MyString` object. If we invoke `x.f(y)` what will be the difference in the case that we use implementation (a) vs. implementation (b) ?

(c) (MapReduce, 12 points) The following Hadoop code is based on the project 5 tutorial:

```
1  public static class TokenizerMapper extends
2      Mapper<Object, Text, IntWritable, IntWritable> {
3      private final static IntWritable one = new IntWritable(1);
4      private IntWritable count = new IntWritable(0);
5
6      public void map(Object file, Text value, Context context)
7          throws IOException, InterruptedException {
8          StringTokenizer itr = new StringTokenizer(value.toString());
9          while (itr.hasMoreTokens()) {
10             int length = itr.nextToken().length();
11             count.set(length);
12             context.write(count, one);
13         } } }
14
15 public static class IntSumReducer extends
16     Reducer<IntWritable, IntWritable, IntWritable, IntWritable> {
17     private IntWritable result = new IntWritable();
18
19     public void reduce(IntWritable key, Iterable<IntWritable> values,
20         Context context) throws IOException, InterruptedException {
21         int sum = 0;
22         for (IntWritable val : values) {
23             int v = val.get();
24             sum += v;
25         }
26         result.set(sum);
27         context.write(key, result);
28     } }
```

i. (6 points) What will the keys and values be in the final output file, assuming this code is run on a directory of text files (which consist of, say, the collected works of Shakespeare)?

ii. (6 points) Could `IntSumReducer` also be used as a combiner? Why or why not?

3. (Java concurrency, 30 points) Each of the next few pages contains a small program with at least two classes, one of which is called `TestCase`. Consider what would happen when running `TestCase.main()`.

If the program would terminate normally and always print the same answer, indicate that answer. Otherwise, indicate *one or more* of the following things the program will do: (a) exhibit a data race, (b) exhibit an atomicity violation, (c) exhibit a deadlock, (d) run forever, (e) print different things on different runs.

Each problem is worth 6 points.

Be careful: I am only interested in what will happen for executions of the given `TestCase.main`, not hypothetical ways in which the classes could be used.

Also, note that in several cases we are ignoring that methods could throw `InterruptedException`, to keep the code shorter.

(a)

```
1 public class GlobalRef {
2     public int x = 0;
3     public Object y = new Object();
4     public void increment() { synchronized(y) { x = x + 1; } }
5 }
6
7 public class TestCase {
8     public static void main(String args[]) {
9         final GlobalRef r = new GlobalRef();
10        new Thread() { public void run() { r.increment(); } }.start ();
11        new Thread() { public void run() { r.increment(); } }.start ();
12        Thread.yield ();
13        System.out.println(r.x);
14    }
15 }
```

(b)

```
1 public class OnePlaceBuffer {
2     private String value = null;
3     public synchronized String get() {
4         while (value == null) wait (); // ignore IE
5         String ret = value;
6         value = null;
7         notify ();
8         return ret;
9     }
10    public synchronized void put(String x) {
11        while (value != null) wait (); // ignore IE
12        value = x;
13        notify ();
14    }
15 }
16
17 public class TestCase {
18     public static void main(String args[]) {
19         final OnePlaceBuffer b = new OnePlaceBuffer();
20         Runnable p = new Runnable() {
21             public void run() { b.put("Hello"); }
22         };
23         new Thread(p).start();
24         new Thread(p).start();
25         System.out.println(b.get ());
26     }
27 }
```

(c)

```
1 public class SillyFact {
2     private ExecutorService exec;
3     public SillyFact (ExecutorService e) { exec = e; }
4     public int nth(int n) {
5         if (n == 0) return 1;
6         else {
7             final int m = n-1;
8             final SillyFact me = this;
9             Callable<Integer> c = new Callable<Integer>() {
10                public Integer call () { return me.nth(m); }
11            };
12            Future<Integer> f = exec.submit(c);
13            return n + f.get(); // ignore IE
14        }
15    }
16 }
17
18 public class TestCase {
19     public static void main(String args[]) {
20         ExecutorService e = Executors.newFixedThreadPool(4);
21         SillyFact f = new SillyFact(e);
22         System.out.println(f.nth(6));
23         e.shutdown();
24     }
25 }
```

(d)

```
1 public class Pair {
2     public final int left , right ;
3     public Pair(int x, int y) { left = x; right = y; }
4 }
5 public class ConcPair {
6     private Pair p = new Pair(0,0);
7     public void updateLeft(int x) {
8         int y = p. right ;
9         p = new Pair(x,y);
10    }
11    public void updateRight(int x) {
12        int y = p. left ;
13        p = new Pair(y,x);
14    }
15    public int getLeft() { return p. left ; }
16    public int getRight() { return p. right ; }
17 }
18
19 public class TestCase {
20     public static void main(String args[]) throws InterruptedException {
21         final ConcPair m = new ConcPair ();
22         Thread t1 = new Thread() { public void run() { m.updateLeft(1); } };
23         Thread t2 = new Thread() { public void run() { m.updateRight(1); } };
24         t1 . start (); t2 . start ();
25         t1 . join (); t2 . join ();
26         System.out.println("(" + m.getLeft() + ", " + m.getRight() + ")");
27     }
28 }
```


- (e) Note: Here, OptWhiteList is a different implementation of ConcPair from the previous question; the Pair class is the same as before; the TestCase class is the same, but refers to OptWhiteList

```
1 public class OptWhiteList {
2     private AtomicReference<Pair> p;
3     public OptWhiteList() { p = new AtomicReference<Pair>(new Pair(0,0)); }
4     private void update(boolean doLeft, int x) {
5         while (true) {
6             Pair q, old = p.get();
7             if (doLeft) q = new Pair(x,old. right );
8             else      q = new Pair(old. left ,x);
9             if (p.compareAndSet(old,q)) break;
10        }
11    }
12    public void updateLeft(int x) { update(true,x); }
13    public void updateRight(int x) { update(false,x); }
14    public int getLeft() { return p.get(). left ; }
15    public int getRight() { return p.get(). right ; }
16 }
17
18 public class TestCase {
19     public static void main(String args[]) throws InterruptedException {
20         final OptWhiteList m = new OptWhiteList ();
21         Thread t1 = new Thread() { public void run() { m.updateLeft(1); } };
22         Thread t2 = new Thread() { public void run() { m.updateRight(1); } };
23         t1.start (); t2.start ();
24         t1.join (); t2.join ();
25         System.out.println("(" + m.getLeft() + ", " + m.getRight() + ")");
26     }
27 }
```

4. (Erlang Execution, 20 points)

Look at each of the Erlang programs below. Along with each program, we will provide a call; say what the call will do. If it returns, indicate the value returned; if it fails with some exception, indicate the exception; or if it does not return, say so. Explain your reasoning if you hope for partial credit. You may assume all of the programs compile. Each question is worth 5 points.

(a) What will `go([1,2,3])` do?

```
go([]) -> 0;
go([H|T]) -> 1+go(T).
```

(b) What will `go(grill ,burger)` do?

```
machine_map() ->
    dict : from_list ([{ grill , {burger, 500}},
                       { frier , {fries , 250}},
                       {soda_fountain, {coke, 100}}]).
go(Machine,Food) ->
    Dict = machine_map(),
    {Food,Time} = dict:fetch(Machine,Dict),
    Time.
```

(c) What will go(goo) do?

```
init () -> register(s1, spawn(fun () -> receive {msg,Pid} -> Pid ! ok end end)).  
go(Msg) ->  
  init (),  
  s1 ! {Msg,self()},  
  receive  
    ok -> ok  
  end.
```

(d) What will go() do?

```
pmap(F,L) ->  
  Me = self(),  
  Pids = lists :map(fun (M) ->  
    spawn(fun () ->  
      Pid = self (),  
      R = F(M),  
      Me ! {Pid,R} end) end,  
    L),  
  lists :map(fun (Pid) ->  
    receive  
      {Pid,R} -> R  
    end  
  end, Pids).
```

```
go() ->  
  pmap(fun ([H|T]) -> H end, [[1,2,3],[3,2,1],[3,2],[1,7,5,4]]).
```

5. (Erlang coding, 10 points) The following function `go()` calls a function `f` for which we have not provided a definition. This function will return a process ID to which `go` subsequently sends messages. Give an implementation of `f` so that `go()` returns correctly with the atom `ok`. (You may define other functions too, if you need them.)

```
go() ->
  P = f (),
  P ! [1, 3, 7],
  receive
    {P,3} -> ok
  end,
  P ! [4, 5],
  receive
    {P,2} -> ok
  end,
  P ! done,
  receive
    finished -> ok
  end,
  ok.
```

Give your definition of `f` here:

6. (Java parallelization, 10 points) A `Data` object has a `Grid` as a private field. `Data` also defines a method `heatMap` to compute this grid's "heat map", which is itself a `Grid`. Assume that `ArrayGrid` is an implementation of the `Grid` interface backed by an array.

On the next page, provide the implementation of a new `Data` class method `parHeatMap` that produces the same result as `heatMap`, but does it in parallel.

Extra credit: Assume that the time to create and run a new task is roughly the same as the time to compute the heat map of 100 grid squares. For extra credit, incorporate this information in your solution.

```
1 public interface Grid {
2     public int getWidth(); // x axis length
3     public int getHeight(); // y axis length
4     public int get(int x, int y); // (0,0) is in the upper left
5     public void set(int x, int y, int v); // sets grid(x,y) to v
6 }
7 public class Data {
8     final private Grid grid;
9     public Data(Grid grid) {
10        this.grid = grid;
11    }
12    private int heat(int x, int y) {
13        int v = grid.get(x,y);
14        for (int i = x-1; i<=x+1; i++) {
15            for (int j = y-1; j<=y+1; j++) {
16                if (i >= 0 && i < grid.getWidth() &&
17                    j >= 0 && j < grid.getHeight() &&
18                    !(i == x && j == y)) {
19                    v = v + (grid.get(i, j) / 2);
20                }
21            }
22        }
23        return v;
24    }
25    public Grid heatMap() {
26        Grid results = new ArrayGrid(grid.getWidth(),grid.getHeight());
27        for (int i = 0; i<grid.getWidth(); i++) {
28            for (int j = 0; j<grid.getHeight(); j++) {
29                results.set(i, j, heat(i, j));
30            }
31        }
32        return results;
33    } }
```

You may find the following API calls useful:

`Runtime.getRuntime().availableProcessors()` returns the number of available processors

`e.awaitTermination(Long.MAX_VALUE,
TimeUnit.SECONDS)` when `e` is an `ExecutorService`, this call waits until all its tasks terminate.

Also see question 3(c) for some useful API calls.

(This page intentionally left blank; feel free to use as scratch paper.)