

# CMSC 330: Organization of Programming Languages

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## OCaml Expressions and Functions

# Lecture Presentation Style

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- Our focus: **semantics** and **idioms** for OCaml
  - *Semantics* is what the language does
  - *Idioms* are ways to use the language well
- We will also cover some useful **libraries**
- **Syntax** is what you type, not what you mean
  - In one lang: Different syntax for similar concepts
  - Across langs: Same syntax for different concepts
  - Syntax can be a source of fierce disagreement among language designers!

# Expressions

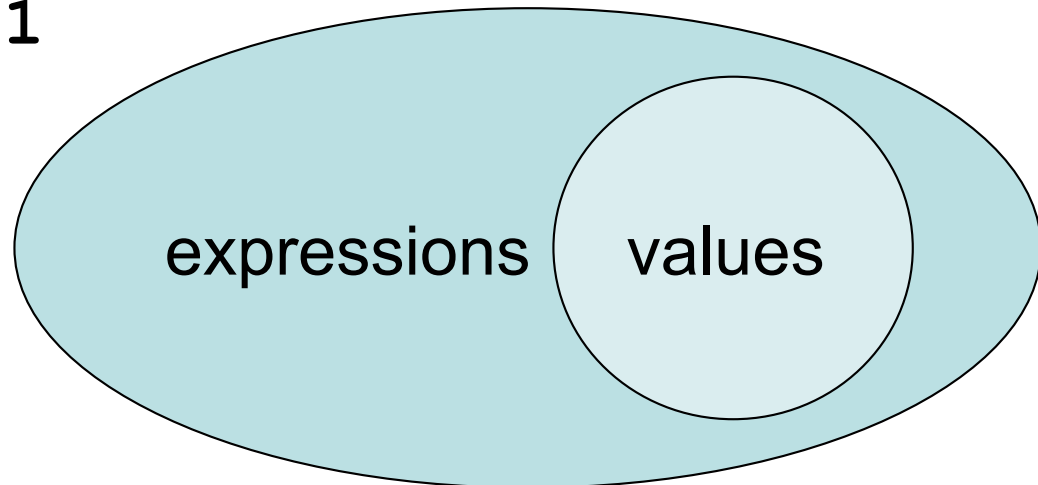
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- **Expressions** are our primary building block
  - Akin to *statements* in imperative languages
- Every kind of expression has
  - **Syntax**
    - We use metavariable  $e$  to designate an arbitrary expression
  - **Semantics**
    - **Type checking** rules (static semantics): produce a type or fail with an error message
    - **Evaluation** rules (dynamic semantics): produce a value
      - (or an exception or infinite loop)
      - Used *only* on expressions that type-check

# Values

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- A **value** is an expression that is final
  - **Evaluating** an expression means running it until it becomes a value
  - We use metavariable **v** to designate an arbitrary value
- **34** is a value, **true** is a value
- **34+17** is an *expression*, but *not* a value
  - It *evaluates* to **51**



# Types

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- **Types** classify expressions
  - Characterize the set of possible values an expression could evaluate to
  - We use metavariable  $t$  to designate an arbitrary type
    - Examples include `int`, `bool`, `string`, and more.
- Expression  $e$  has type  $t$  if  $e$  will (always) evaluate to a value of type  $t$ 
  - $\{ \dots, -1, 0, 1, \dots \}$  are values of type `int`
  - `34+17` is an expression of type `int`, since it evaluates to `51`, which has type `int`
  - Write  $e : t$  to say  $e$  has type  $t$

# If Expressions

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- Syntax

- `if e1 then e2 else e3`

- Evaluation

- If *e1* evaluates to `true`, and if *e2* evaluates to *v*, then `if e1 then e2 else e3` evaluates to *v*

- If *e1* evaluates to `false`, and if *e3* evaluates to *v*, then `if e1 then e2 else e3` evaluates to *v*

- Type checking

- If *e1* has type `bool` and *e2* has type *t* and *e3* has type *t* then `if e1 then e2 else e3` has type *t*

# If Expressions

---

- Syntax
  - `if e1 then e2 else e3`
- Evaluation
  - If `e1` evaluates to `true`, and if `e2` evaluates to `v`, then `if e1 then e2 else e3` evaluates to `v`
  - If `e1` evaluates to `false`, and if `e3` evaluates to `v`, then `if e1 then e2 else e3` evaluates to `v`
- Type checking
  - If `e1 : bool` and `e2 : t` and `e3 : t` then `if e1 then e2 else e3 : t`

# If Expressions

---

- Syntax
  - `if e1 then e2 else e3`
- Evaluation
  - If `e1` evaluates to `true`, and if `e2` evaluates to `v`, then `if e1 then e2 else e3` evaluates to `v`
  - If `e1` evaluates to `false`, and if `e3` evaluates to `v`, then `if e1 then e2 else e3` evaluates to `v`
- Type checking
  - If `e1 : bool` and `e2 : t` and `e3 : t` then `(if e1 then e2 else e3) : t`



# If Expressions: Examples

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```
# if 7 > 42 then "hello" else "goodbye";;
```

```
- : string = "goodbye"
```

```
# if true then 3 else 4;;
```

```
- : int = 3
```

```
# if false then 3 else 3.0;;
```

```
This expression has type float but is  
here used with type int
```

# Quiz 1

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To what value does this expression evaluate?

```
if 22=0 then 1 else 2
```

- A. 0
- B. 1
- C. 2
- D. none of the above

# Quiz 1

---

To what value does this expression evaluate?

```
if 22=0 then 1 else 2
```

A. 0

B. 1

**C. 2**

D. none of the above

## Quiz 2

---

To what value does this expression evaluate?

```
if 22=0 then "bear" else 2
```

- A. 0
- B. 1
- C. 2
- D. none of the above

## Quiz 2

---

To what value does this expression evaluate?

```
if 22=0 then "bear" else 2
```

A. 0

B. 1

C. 2

**D. none of the above:** doesn't type check so never gets a chance to be evaluated

# Function Definitions

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- OCaml functions are like mathematical functions
  - Compute a result from provided arguments

```
(* requires n>=0 *)
(* returns: n! *)
let rec fact n =
  if n = 0 then
    1
  else
    n * fact (n-1)
```

function body

Use (\* \*) for comments  
(may nest)

Parameter  
(type inferred)

rec needed for recursion

Structural equality

Line breaks, spacing ignored  
(like C, C++, Java, not like Ruby)

# Function Types

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- In OCaml, `->` is the function type constructor
  - Type `t1 -> t` is a function with argument or *domain* type `t1` and return or *range* type `t`
  - Type `t1 -> t2 -> t` is a function that takes *two* inputs, of types `t1` and `t2`, and returns a value of type `t`. Etc.
- Examples
  - `let next x = x + 1 (* type int -> int *)`
  - `let fn x = (int_of_float x) * 3`  
`(* type float -> int *)`
  - `fact` `(* type int -> int *)`

# Type Checking Functions

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- Syntax `let rec f x1 ... xn = e`
- Type checking
  - Conclude that  $f : t1 \rightarrow \dots \rightarrow tn \rightarrow u$  if  $e : u$  under the following assumptions:
    - $x1 : t1, \dots, xn : tn$  (arguments with their types)
    - $f : t1 \rightarrow \dots \rightarrow tn \rightarrow u$  (for recursion)
- Example
  - Given  $n : \text{int}, \text{fact} : \text{int} \rightarrow \text{int}$
  - Does `if n = 0 then 1 ... : int` ?
    - It does!
  - Conclude  $\text{fact} : \text{int} \rightarrow \text{int}$

```
let rec fact n =  
  if n = 0 then  
    1  
  else  
    n * fact (n-1)
```



# Calling Functions

---

- Syntax  $f\ e1\ \dots\ en$ 
  - Parentheses not required around argument(s)
  - No commas; use spaces instead
- Type checking
  - If  $f : t1 \rightarrow \dots \rightarrow tn \rightarrow u$  and  $e1 : t1, \dots, en : tn$   
then  $f\ e1\ \dots\ en : u$
- Example:
  - `fact 1 : int`
  - since `fact : int -> int` and `1 : int`
- Function call *aka* function application

# Calling Functions

---

- Syntax  $f\ e1\ \dots\ en$
- Evaluation
  - Evaluate arguments  $e1\ \dots\ en$  to values  $v1\ \dots\ vn$ 
    - Order is actually right to left, not left to right
    - But this doesn't matter if  $e1\ \dots\ en$  don't have side effects
  - Find the definition of  $f$ 
    - `let rec f x1 ... xn = e`
  - Substitute  $vi$  for  $xi$  in  $e$ , yielding new expression  $e'$
  - Evaluate  $e'$  to value  $v$ , which is the final result

# Calling Functions

## Example evaluation

```
let rec fact n =  
  if n = 0 then  
    1  
  else  
    n * fact (n-1)
```

- fact 2

- if 2=0 then 1 else 2\*fact(2-1)

- 2 \* fact 1

- 2 \* (if 1=0 then 1 else 1\*fact(1-1))

- 2 \* 1 \* fact 0

- 2 \* 1 \* (if 0=0 then 1 else 0\*fact(0-1))

- 2 \* 1 \* 1

- 2

# Type Annotations

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- The syntax `(e : t)` asserts that “`e` has type `t`”
  - This can be added anywhere you like

```
let (x : int) = 3
let z = (x : int) + 5
```
- Define functions’ parameter and return types

```
let fn (x:int):float =
    (float_of_int x) *. 3.14
```

  - Note special position for return type
  - Thus `let g x:int = ...` means `g` returns `int`
    - *Not* that `x` has type `int`
- Checked by compiler: Very useful for debugging

## Quiz 3: What is the value of `foo 4 2`

---

```
let rec foo n m =  
  if n >= 9 || n < 0 then  
    m  
  else  
    n + m + 1
```

- Type Error
- 2
- 8
- 7

## Quiz 3: What is the value of `foo 4 2`

---

```
let rec foo n m =  
  if n >= 9 || n < 0 then  
    m  
  else  
    n + m + 1
```

- Type Error
- 2
- 8
- **7**

## Quiz 4: What is the value of `bar 4`

---

```
let rec bar(n:int):int =  
  if n = 0 || n = 1 then 1  
  else  
    bar (n-1) + bar (n-2)
```

- Syntax Error
- 4
- 5
- 8

## Quiz 4: What is the value of `bar 4`

---

```
let rec bar(n:int):int =  
  if n = 0 || n = 1 then 1  
  else  
    bar (n-1) + bar (n-2)
```

- Syntax Error
- 4
- **5**
- 8