CMSC 330: Organization of Programming Languages

CMSC330 Fall 2017

OCaml Data Types

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OCaml Data

- So far, we've seen the following kinds of data
 - Basic types (int, float, char, string)
 - Lists
 - > One kind of data structure
 - > A list is either [] or h::t, deconstructed with pattern matching
 - Tuples and Records
 - Let you collect data together in fixed-size pieces
 - Functions
- How can we build other data structures?
 - Building everything from lists and tuples is awkward

User Defined Types

- type can be used to create new names for types
 - Useful for combinations of lists and tuples
- Examples
 - type my_type = int * (int list)
 - let (x:my_type) = (3, [1; 2])
 - type my_type2 = int*char*(int*float)
 - let (y:my_type2) = (3, 'a', (5, 3.0))

(User-Defined) Variants

type coin = Heads Tails let flip $x =$	In simplest form: Like a C enum
<pre>match x with Heads -> Tails Tails -> Heads</pre>	Basic pattern matching resembles C switch
<pre>let rec count_heads x = match x with [] -> 0</pre>	Combined list and variant patterns possible
(Heads:: x') -> 1 + count_he (_:: x') -> count_heads x'	ads x'

Constructing and Destructing Variants

• Syntax

- type t = C1 | ... | Cn
- the Ci are called constructors

> Must begin with a capital letter

- Evaluation
 - A constructor *Ci* is already a value
 - Destructing a value v of type t is done by pattern matching on v; the patterns are the constructors Ci
- Type Checking
 - Ci : t (for each Ci in t's definition)

Data Types: Variants with Data

- We can define variants that "carry data" too
 - Not just a constructor, but a constructor *plus values*

type shape =			
Rect of float * float	(*	width*length	*)
Circle of float	(*	radius *)	

- Rect and Circle are constructors
 - where a shape is either a Rect(w, 1)
 - > for any floats w and 1
 - or a Circle r
 - ▹ for any float r

Data Types (cont.)

```
let area s =
   match s with
        Rect (w, 1) -> w *. 1
        | Circle r -> r *. r *. 3.14
;;
area (Rect (3.0, 4.0));; (* 12.0 *)
area (Circle 3.0);; (* 9.42 *)
```

- Use pattern matching to deconstruct values
 - Can bind pattern values to data parts
- Data types are aka algebraic data types are aka tagged unions

Data Types (cont.)

```
type shape =
    Rect of float * float (* width*length *)
    | Circle of float (* radius *)
let lst = [Rect (3.0, 4.0) ; Circle 3.0]
```

- What's the type of lst?
 - shape list
- What's the type of lst's first element?
 - shape

Variation: Shapes in Java Compare this to OCaml

```
public interface Shape {
    public double area();
}
```

```
class Rect implements Shape {
   private double width, length;
   Rect (double w, double l) {
     this.width = w;
     this.length = l;
   }
   double area() {
     return width * length;
   }
}
```

```
class Circle implements Shape {
  private double rad;
  Circle (double r) {
    this.rad = r;
  }
  double area() {
   return rad * rad * 3.14159;
  }
}
```

Option Type

```
type optional_int =
   None
   Some of int
let divide x y =
   if y != 0 then Some (x/y)
   else None
let string_of_opt o =
   match o with
      Some i -> string_of_int i
   | None -> "nothing"
```

```
let p = divide 1 0;;
print_string
  (string_of_opt p);;
(* prints "nothing" *)
let q = divide 1 1;;
print_string
  (string_of_opt q);;
(* prints "1" *)
```

Comparing to Java: None is like null, while
 Some *i* is like an Integer (*i*) object

Polymorphic Option Type

- A Polymorphic version of option type can work with any kind of data
 - Polymorphic parameter: • As int option, char option, etc... like Option<T> in Java



In fact, this option type is built into OCaml

X:: Some let p = opthd [];; (* p = None *) let q = opthd [1;2];; (* q = Some 1 *) let r = opthd [``a''];; (* r = Some ``a'' *)

X

Recursive Data Types

• We can build up lists with recursive variant types

```
type 'a mylist =
   Nil
   | Cons of 'a * 'a mylist
let rec len = function
   Nil -> 0
   | Cons (_, t) -> 1 + (len t)
len (Cons (10, Cons (20, Cons (30, Nil))))
(* evaluates to 3 *)
```

• Won't have nice [1; 2; 3] syntax for this kind of list

Constructing and Destructing Variants

- Syntax
 - type t = C1 [of t1] | ... | Cn [of tn]
 - the Ci are called constructors
 - Must begin with a capital letter; may include associated data notated with brackets [] to indicate it's optional
- Evaluation
 - A constructor *Ci* is a value if it has no assoc. data
 Ci vi is a value if it does
 - Destructing a value of type *t* is by pattern matching
 patterns are constructors *Ci* with data components, if any
- Type Checking
 - Ci [vi] : t [if vi has type ti]

Data Type Representations

- · Values in a data type are stored
 - 1. Directly as integers
 - 2. As pointers to blocks in the heap

```
type t =
   A of int
| B
| C of int * int
| D
```

Exercise: A Binary Tree Data Type

- Write type bin_tree for binary trees over int
 - Trees should be ordered (binary search tree)
- Implement the following empty : bin tree is empty : bin tree -> bool member : int -> bin tree -> bool insert : int -> bin tree -> bin tree remove: int -> bin tree -> bin tree equal : bin tree -> bin tree -> bool fold : (int -> 'a -> 'a) -> bin tree -> 'a -> 'a

type foo = (int * (string list)) list

Which one of the following could match foo?

- A. [(3, "foo", "bar")]
- в. [(5, ["foo", "bar"])]
- c. [(7, ["foo"; "bar"])]
- D. [(9, [("foo", "bar")])]

type foo = (int * (string list)) list

Which one of the following could match foo?

- A. [(3, "foo", "bar")]
- B. [(5, ["foo", "bar"])]
- c. [(7, ["foo"; "bar"])]
- D. [(9, [("foo", "bar")])]

Quiz 2: What does this evaluate to?

```
type num = Int of int | Float of float;;
let plus a b =
  match a, b with
  | Int i, Int j -> Int (i+j)
  | Float i, Float j -> Float (i +. j)
  | Float i, Int j -> Float (i +. float_of_int j)
;;
plus (Float 3.0) (Int 2);;
```

- A. float = 5.
- B. num = Int 5
- c. Type Error
- D. num = Float 5.

Quiz 2: What does this evaluate to?

```
type num = Int of int | Float of float;;
let plus a b =
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  | Int i, Int j -> Int (i+j)
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  | Float i, Int j -> Float (i +. float_of_int j)
;;
plus (Float 3.0) (Int 2);;
```

- A. float = 5.
- B. num = Int 5
- c. Type Error
- D. num = Float 5.

Quiz 3: What does this evaluate to?

```
let foo f = match f with
    None -> 42.0
    | Some n -> n +. 42.0
;;
foo 3.3;;
```

- A. float = 45.3
- B. Error
- c. float = 42.0
- D. No output

Quiz 3: What does this evaluate to?

- A. float = 45.3
- B. Error
- c. float = 42.0
- D. No output

OCaml Exceptions

```
exception My exception of int
let f n =
  if n > 0 then
    raise (My exception n)
  else
    raise (Failure "foo")
let bar n =
  try
    f n
  with My exception n \rightarrow
      Printf.printf "Caught %d\n" n
    | Failure s ->
      Printf.printf "Caught %s\n" s
```

Exceptions (cont.)

- Exceptions are declared with exception
 - They may appear in the signature as well
- Exceptions may take arguments
 - Just like type constructors
 - May also have no arguments
- Catch exceptions with try...with...
 - Pattern-matching can be used in with
 - If an exception is uncaught
 - Current function exits immediately
 - Control transfers up the call chain
 - > Until the exception is caught, or until it reaches the top level

OCaml Exceptions (cont.)

• Exceptions may be thrown by I/O statements

- Common way to detect end of file
- Need to decide how to handle exception
- Example

```
try
  (input_char stdin) (* reads 1 char *)
with End_of_file -> 0 (* return 0? *)

try
  read_line () (* reads 1 line *)
with End_of_file -> "" (* return ""? *)
```

OCaml Exceptions (cont.)

- **failwith**: Raise exception Failure with the given string.
- invalid_arg: Raise exception Invalid_argument with the given string
- Not_found: Raised if the object does not exist

```
let div x y =
  if y = 0 failwith "divide by zero" else x/y;;
let lst =[(1,"alice");(2,"bob");(3,"cat")];;
let lookup key lst =
  try
   List.assoc key lst
  with
   Not_found -> "key does not exist"
```