CMSC 330: Organization of Programming Languages

OCaml Imperative Programming

So Far, Only Functional Programming

- We haven't given you any way so far to change something in memory
  - All you can do is create new values from old
- This makes programming easier since it supports mathematical (i.e., functional) reasoning
  - Don't care whether data is shared in memory
    - Aliasing is irrelevant
  - Calling a function f with argument x always produces the same result
    - f x = f x for all x

Imperative OCaml

- Sometimes it is useful for values to change
  - Call a function that returns an *incremented* counter
  - Store aggregations in *efficient* hash tables
- OCaml variables are *immutable*, but
- OCaml has references, fields, and arrays that are actually *mutable*
  - I.e., they can change

References

- **`'a ref`**: Pointer to a mutable value of type **`'a`**
- There are three basic operations on references:

  `ref : 'a -> 'a ref`

  - Allocate a reference

  `! : 'a ref -> 'a`

  - Read the value stored in reference

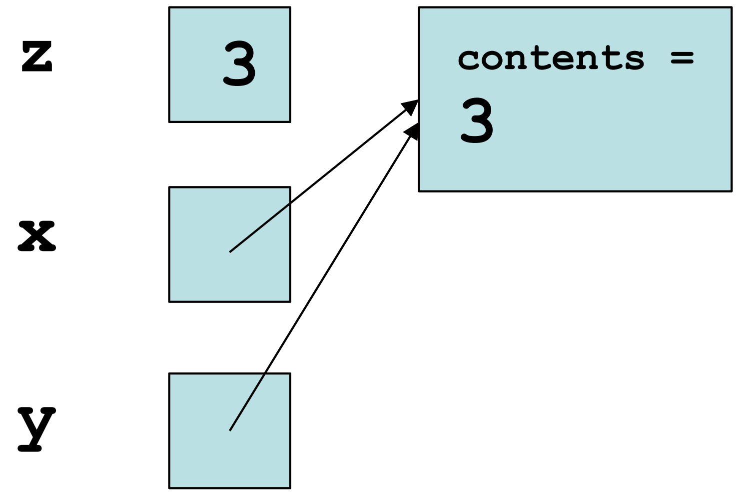
  `:= : 'a ref -> 'a -> unit`

  - Change the value stored in reference
- Binding variable **`x`** to a reference is immutable
  - The contents of the reference **`x`** points to may change

References Usage

Example:

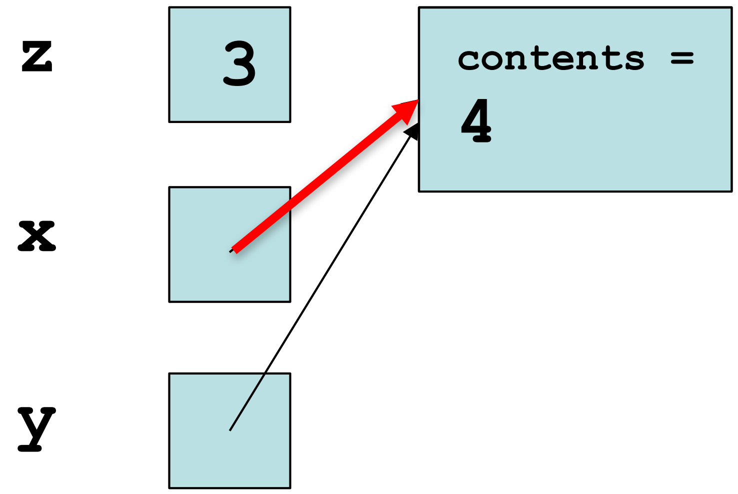
```
# let z = 3;;
val z : int = 3

# let x = ref z;;
val x : int ref = {contents = 3}

# let y = x;;
val y : int ref = {contents = 3}
```



References Usage

Example:

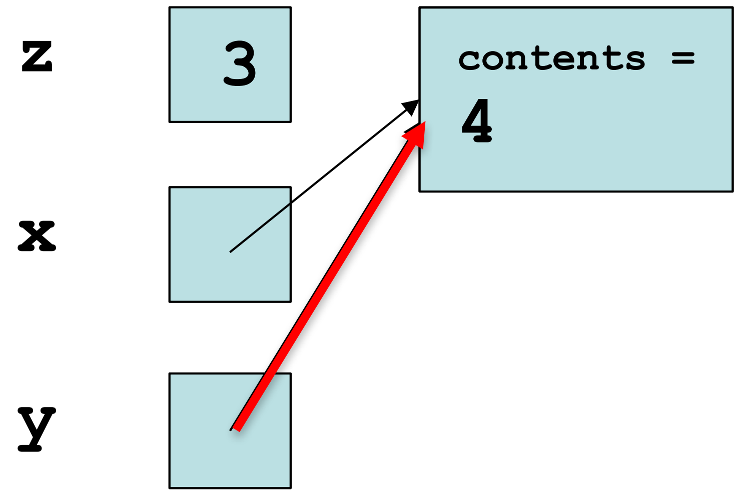
```
# let z = 3;;
val z : int = 3

# let x = ref z;;
val x : int ref = {contents = 3}

# let y = x;;
val y : int ref = {contents = 3}

# x := 4;;
- : unit = ()
```



References Usage

Example:

```
# let z = 3;;
val z : int = 3

# let x = ref z;;
val x : int ref = {contents = 3}

# let y = x;;
val y : int ref = {contents = 3}

# x := 4;;
- : unit = ()

# !y;;
- : int = 4
```



Aliasing

- Reconsider our example

```
let z = 3;;
let x = ref z;;
let y = x;;
x := 4;;
!y;;
```

- Here, variables **y** and **x** are aliases:
  - In **let y = x**, variable **x** evaluates to a location, and **y** is bound to the same location
  - So, changing the contents of that location will cause both **!x** and **!y** to change

Quiz 1: What is the value **`w`**?

```
let x = ref 42 in
let y = ref 42 in
let z = x in
let () = x := 43 in
let w = !y + !z in
w
```

A. 42
B. 84
C. 85
D. 86

Quiz 1: What is the value `w`?

```
let x = ref 42 in
let y = ref 42 in
let z = x in
let () = x := 43 in
let w = !y + !z in
w
```

A. 42
B. 84
C. 85
D. 86

Quiz 1a: What is the value **`w`**?

```
let x = ref 42 in
let y = ref 42 in
let z = !x in
let () = x := 43 in
let w = !y + z in
w
```

A. 42
B. 84
C. 85
D. Error

Quiz 1a: What is the value **w**?

```
let x = ref 42 in
let y = ref 42 in
let z = !x in
let () = x := 43 in
let w = !y + z in
w
```

A. 42
B. 84
C. 85
D. Error

Implement a Counter

```
# let counter = ref 0 ;;
val counter : int ref = { contents=0 }

# let next =
    fun () ->
     counter := !counter + 1; !counter ;;
val next : unit -> int = <fun>

# next ();;
- : int = 1

# next ();;
- : int = 2
```

Hide the Reference

```
# let next =
    let counter = ref 0 in
    fun () ->
      counter := !counter + 1; !counter ;;
val next : unit -> int = <fun>

# next ();;
- : int = 1

# next ();;
- : int = 2
```

Quiz 2: What is wrong with the counter?

```
let next =
      fun () ->
         let counter = ref 0 in
          counter := !counter + 1;
          !counter
```

A. Error, because **`counter`** isn't in scope in the final line
B. It returns a reference to an integer instead of an integer
C. It returns the same integer every time
D. Nothing is wrong

Quiz 2: What is wrong with the counter?

```
let next =
    fun () ->
      let counter = ref 0 in
       counter := !counter + 1;
       !counter
```

A. Error, because **`counter`** isn't in scope in the final line
B. It returns a reference to an integer instead of an integer
C. It returns the same integer every time
D. Nothing is wrong

Hide the Reference, Visualized

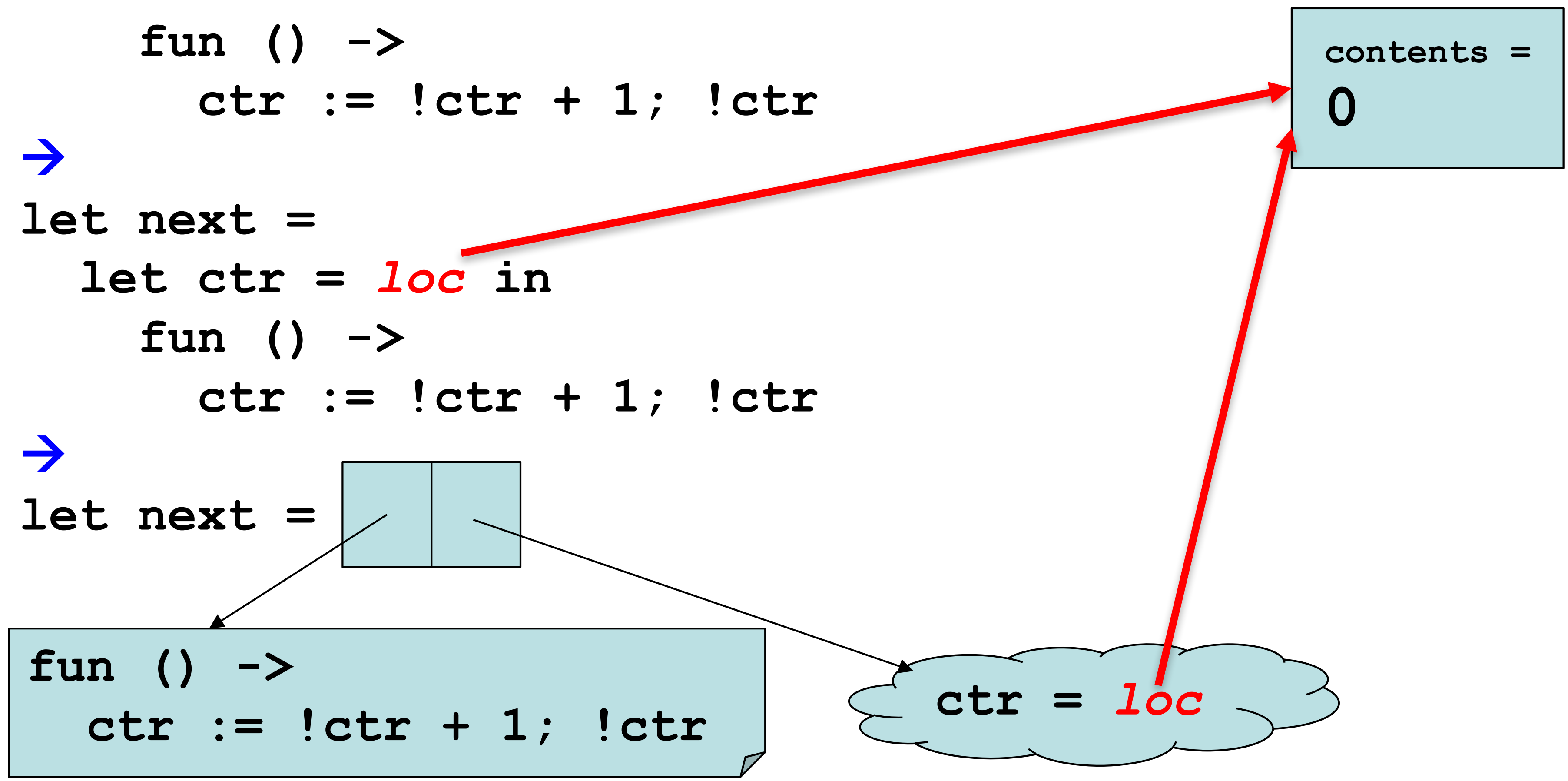
```
let next =
  let ctr = ref 0 in
    fun () ->
      ctr := !ctr + 1; !ctr
```

→

```
let next =
  let ctr = loc in
    fun () ->
      ctr := !ctr + 1; !ctr
```

→

let next =



References: Syntax and Semantics

- Syntax: **`ref`** ***`e`***
- Evaluation
  - Evaluate ***`e`*** to a value ***`v`***
  - Allocate a new location ***`loc`*** in memory to hold ***`v`***
  - Store ***`v`*** in contents of memory at ***`loc`***
  - Return ***`loc`***
    - Note: locations are first-class values
- Type checking
  - (**`ref`** ***`e`***) **`:`** ***`t`*** **`ref`**
    - if ***`e`*** **`:`** ***`t`***

References: Syntax and Semantics

- Syntax: `e1 := e2`
- Evaluation
  - Evaluate `e2` to a value `v2`
  - Evaluate `e1` to a location `loc`
  - Store `v2` in contents of memory at `loc`
  - Return `()`
- Type checking
  - (`e1 := e2`) `: unit`
    - if `e1 : t ref` and `e2 : t`

References: Syntax and Semantics

- Syntax: `!`***`e`***
  - *This is not negation*
- Evaluation
  - Evaluate ***`e`*** to a location ***`loc`***
  - Return contents ***`v`*** of memory at ***`loc`***
- Type checking
  - `!`***`e`*** `:` ***`t`***
    - if ***`e`*** `:` ***`t`*** **`ref`**

Sequences: Syntax and Semantics

- Syntax: ***e1*** `;` ***e2***
  - ***e1*** `;` ***e2*** is the same as `let () =` ***e1*** `in` ***e2***
- Evaluation
  - Evaluate ***e1*** to a value ***v1***
  - Evaluate ***e2*** to a value ***v2***
  - Return ***v2***
    - We throw away ***v1*** – so ***e1*** is useful only if it has *effects*, e.g., if it changes a reference's contents or accesses a file
- Type checking
  - ***e1*** `;` ***e2*** `:` ***t***
    - if ***e1*** `: unit` and ***e2*** `:` ***t***

;; versus ;

- ;; ends an expression in the top-level of OCaml
  - Use it to say: "Give me the value of this expression"
  - Not used in the body of a function
  - Not needed after each function definition
    - Though for now it won't hurt if used there
- ***e1; e2*** evaluates ***e1*** and then ***e2***, and returns ***e2***

```
let print_both (s, t) = print_string s; print_string t;
                        "Printed s and t"
```

- notice no ; at end – it's a separator, not a terminator

```
print_both (”Colorless green ", ”ideas sleep")
```

Prints `”Colorless green ideas sleep"`, and returns `"Printed s and t"`

Grouping Sequences

- If you're not sure about the scoping rules, use begin...end, or *parentheses*, to group together statements with semicolons

```
let x = ref 0
let f () =
  begin
    print_string "hello";
    x := !x + 1
  end
```

```
let x = ref 0
let f () =
  (
    print_string "hello";
    x := !x + 1
  )
```

The Trade-Off Of Side Effects

- Side effects are absolutely necessary
  - That's usually why we run software! We want something to happen that we can observe
- They also make reasoning harder
  - Order of evaluation now matters
  - No referential transparency
    - Calling the same function with the same arguments may produce different results
  - Aliasing may result in hard-to-understand bugs
    - If we call a function with refs r1 and r2, it might do strange things if r1 and r2 are aliased

Quiz 3: What is the value **`w`**?

```
let f _ z = z+1 in
let y = ref 1 in
let w = f (y:=2) !y in
w
```

A. 3
B. 2
C. Type Error
D. ()

Quiz 3: What is the value **`w`**?

```
let f _ z = z+1 in
let y = ref 1 in
let w = f (y:=2) !y in
w
```

A. 3
B. 2
C. Type Error
D. ()

Quiz 4: What is the value **`w`**?

```
let f z _ = z+1 in
let y = ref 1 in
let w = f !y (y:=2) in
w
```

A. 3
B. 2
C. Type Error
D. ()

Quiz 4: What is the value **w**?

```
let f z _ = z+1 in
let y = ref 1 in
let w = f !y (y:=2) in
w
```

A. 3
B. 2
C. Type Error
D. ()

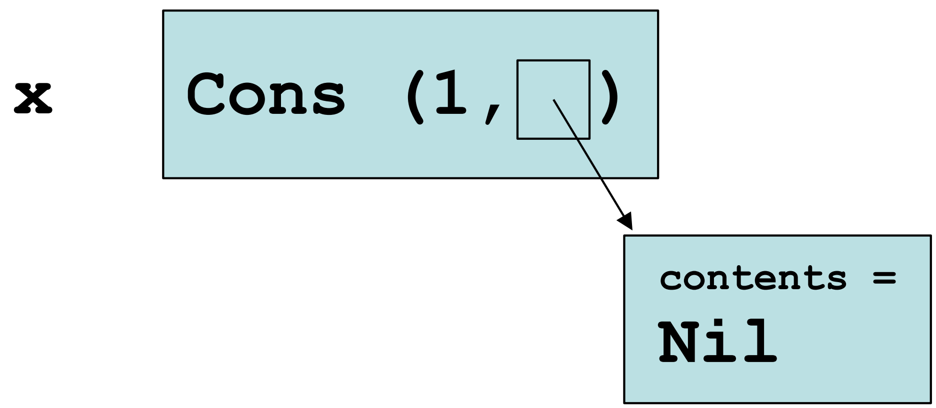
Structural vs. Physical Equality

- In OCaml, the = operator compares objects structurally
  - [1;2;3] = [1;2;3]   (* true *)
  - (1,2) = (1,2)   (* true *)
  - The = operator is used for pattern matching
- The == operator compares objects physically
  - [1;2;3] == [1;2;3]   (* false *)
- Mostly you want to use the first one
  - But it's a problem with cyclic data structures

Cyclic Data Structures Possible With Ref

```
type 'a rlist =
  Nil | Cons of 'a * ('a rlist ref);;

let newcell x y = Cons(x,ref y);;

let updnext (Cons (_,r)) y = r := y;;
```

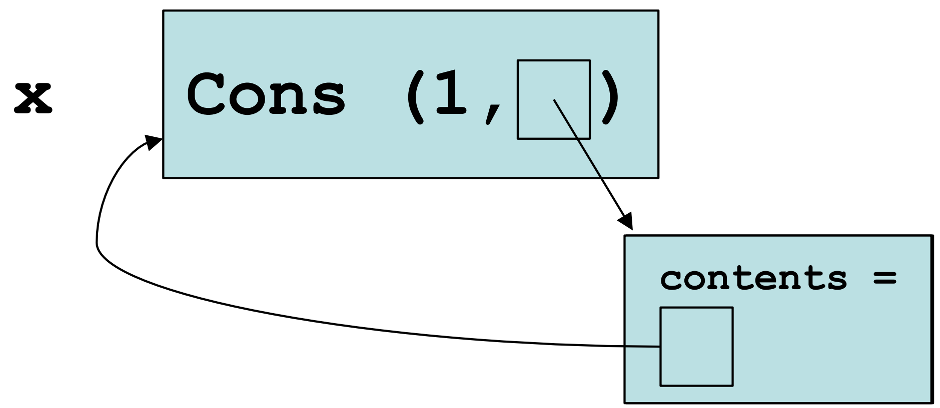
```
# let x = newcell 1 Nil;;
val x : int reflist = Cons (1, {contents = Nil})
```



Cyclic Data Structures Possible With Ref

```
type 'a rlist =
  Nil | Cons of 'a * ('a rlist ref);;

let newcell x y = Cons(x,ref y);;

let updnext (Cons (_,r)) y = r := y;;
```

```
# let x = newcell 1 Nil;;
val x : int reflist = Cons (1, {contents = Nil})

# updnext x x;;
- : unit = ()

# x == x;;
- : bool = true

# x = x;; (* hangs *)
```



Mutable fields

- Fields of a record type can be declared as mutable:

```
# type point = {x:int; y:int; mutable c:string};;
type point = { x : int; y : int; mutable c : string; }

# let p = {x=0; y=0; c="red"};;
val p : point = {x = 0; y = 0; c = "red"}

# p.c <- “white”;;
- : unit = ()

# p;;
val p : point = {x = 0; y = 0; c = ”white"}

# p.x <- 3;;
Error: The record field x is not mutable
```

Implementing Refs

- Ref cells are essentially syntactic sugar:

```
type 'a ref = { mutable contents: 'a }
let ref x = { contents = x }
let (!) r = r.contents
let (:=) r newval = r.contents <- newval
```

- ref type is declared in **`Pervasives`**
- ref functions are compiled to equivalents of above

Arrays

- **Arrays** generalize ref cells from a single mutable value to a sequence of mutable values

```
# let v = [|0.; 1.|];;
val v : float array = [|0.; 1.|]

# v.(0) <- 5.;;
- : unit = ()

# v;;
- : float array = [|5.; 1.|]
```

Arrays

- Syntax: `[|e1; ...; en|]`
- Evaluation
  - Evaluates to an n-element array, whose elements are initialized to `v1` … `vn`, where `e1` evaluates to `v1`, ..., `en` evaluates to `vn`
    - Evaluates them *right to left*
- Type checking
  - `[|e1; …; en|] : t array`
    - If for all *i*, each `ei : t`

Arrays

- Syntax: `e1.(e2)`
- Evaluation
  - Evaluate `e2` to integer value `v2`
  - Evaluate `e1` to array value `v1`
  - If $0 \leq$ `v2` $< n$, where n is the length of array `v1`, then return element at offset `v2` of `v1`
  - Else raise `Invalid_argument` exception
- Type checking: `e1.(e2) : t`
  - if `e1 : t array` and `e2 : int`

Arrays

- Syntax: `e1.(e2) <- e3`
- Evaluation
  - Evaluate *e3* to *v3*
  - Evaluate *e2* to integer value *v2*
  - Evaluate *e1* to array value *v1*
  - If $0 \leq v2 < n$, where *n* is the length of array *v1*, then update element at offset *v2* of *v1* to *v3*
    - Else raise `Invalid_argument` exception
  - Return ()
- Type checking: `e1.(e2) <- e3 : unit`
  - if `e1 : t array` and `e2 : int` and `e3 : t`

Quiz 5: What is the value **`w`**?

```
let x = [| 0; 1 |] in
let w = x in
x.(0) <- 1;
w
```

A. 1
B. [| 0; 1 |]
C. Type Error
D. [| 1; 1 |]

Quiz 5: What is the value **`w`**?

```
let x = [| 0; 1 |] in
let w = x in
x.(0) <- 1;
w
```

A. 1
B. [| 0; 1 |]
C. Type Error
D. [| 1; 1 |]

Control structures

- Traditional loop structures are useful with imperative features:

```
while e1 do e2 done
for x=e1 to e2 do e3 done
for x=e1 downto e2 do e3 done
```

Comparison To OCaml

C

```
int x;
int y;

x = 3;

y = x;

3 = x;
```

OCaml

```
let x = ref 0;;
let y = ref 0;;

x := 3;;   (* x : int ref *)

y := (!x);;

3 := x;;   (* 3 : int; error *)
```

- In OCaml, an updatable location and the contents of the location have different types
  - The location has a **ref** type

OCaml Language Choices

- Implicit or explicit declarations?
  - Explicit – variables must be introduced with let before use
  - But you don't need to specify types

- Static or dynamic types?
  - Static – but you don't need to state types
  - OCaml does type inference to figure out types for you
  - Good: less work to write programs
  - Bad: easier to make mistakes, harder to find errors

OCaml Programming Tips

- Compile your program often, after small changes
  - The OCaml parser often produces inscrutable error messages
  - It's easier to figure out what's wrong if you've only changed a few things since the last compile

- If you're getting strange type error messages, add in type declarations
  - Try writing down types of arguments
  - For any expression e, can write (e:t) to assert e has type t

OCaml Programming Tips (cont.)

- Watch out for precedence and function application

```
let mult x y = x*y

mult 2 2+3    (* returns 7 *)
              (* parsed as (mult 2 2)+3 *)

mult 2 (2+3)  (* returns 10 *)
```

OCaml Programming Tips (cont.)

- All branches of a pattern match must return the same type

```
match x with
... -> -1    (* branch returns int *)
| ... -> ()   (* uh-oh, branch returns unit *)
| ... -> print_string “foo”
              (* also returns unit *)
```

OCaml Programming Tips (cont.)

- You cannot assign to ordinary variables!

```
# let x = 42;;
val x : int = 42
# x = x + 1;;          (* this is a comparison *)
-: bool = false
# x := 3;;
Error: This expression has type int but is here
used with type 'a ref
```

OCaml Programming Tips (cont.)

- Again: You cannot assign to ordinary variables!

```
# let x = 42;;
val x : int = 42
# let f y = y + x;;     (* captures x = 42*)
val f : int -> int = <fun>
# let x = 0;;           (* shadows binding of x *)
val x : int = 0
# f 10;;                (* but f still refers to x=42 *)
- : int = 52
```