

CMSC 330: Organization of Programming Languages

OCaml Expressions and Functions

Lecture Presentation Style

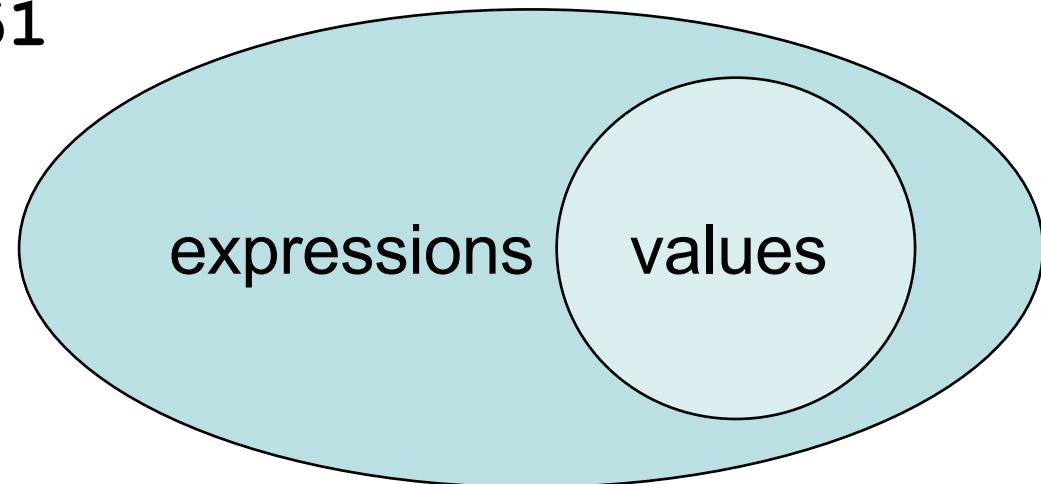
- Our focus: **semantics** and **idioms** for OCaml
 - *Semantics* is what the language does
 - *Idioms* are ways to use the language well
- We will also cover some useful **libraries**
- **Syntax** is what you type, not what you mean
 - In one lang: Different syntax for similar concepts
 - Across langs: Same syntax for different concepts
 - Syntax can be a source of fierce disagreement among language designers!

Expressions

- Expressions are our primary building block
 - Akin to *statements* in imperative languages
- Every kind of expression has
 - Syntax
 - We use metavariable e to designate an arbitrary expression
 - Semantics
 - Type checking rules (static semantics): produce a type or fail with an error message
 - Evaluation rules (dynamic semantics): produce a value
 - (or an exception or infinite loop)
 - Used *only* on expressions that type-check

Values

- A **value** is an expression that is final
 - Evaluating an expression means running it until it becomes a value
 - We use metavariable **v** to designate an arbitrary value
- **34** is a value, **true** is a value
- **34+17** is an *expression*, but *not* a value
 - It evaluates to 51



Types

- Types classify expressions
 - Characterize the set of possible values an expression could evaluate to
 - We use metavariable t to designate an arbitrary type
 - Examples include `int`, `bool`, `string`, and more.
- Expression e has type t if e will (always) evaluate to a value of type t
 - $\{ \dots, -1, 0, 1, \dots \}$ are values of type `int`
 - $34+17$ is an expression of type `int`, since it evaluates to 51 , which has type `int`
 - Write $e : t$ to say e has type t

If Expressions

- Syntax
 - `if e1 then e2 else e3`
- Evaluation
 - If $e1$ evaluates to `true`, and if $e2$ evaluates to v , then `if e1 then e2 else e3` evaluates to v
 - If $e1$ evaluates to `false`, and if $e3$ evaluates to v , then `if e1 then e2 else e3` evaluates to v
- Type checking
 - If $e1$ has type `bool` and $e2$ has type t and $e3$ has type t then `if e1 then e2 else e3` has type t

If Expressions

- Syntax
 - `if e1 then e2 else e3`
- Evaluation
 - If `e1` evaluates to `true`, and if `e2` evaluates to `v`, then `if e1 then e2 else e3` evaluates to `v`
 - If `e1` evaluates to `false`, and if `e3` evaluates to `v`, then `if e1 then e2 else e3` evaluates to `v`
- Type checking
 - If `e1 : bool` and `e2 : t` and `e3 : t` then
`if e1 then e2 else e3 : t`

If Expressions

- Syntax
 - `if e1 then e2 else e3`
- Evaluation
 - If `e1` evaluates to `true`, and if `e2` evaluates to `v`, then `if e1 then e2 else e3` evaluates to `v`
 - If `e1` evaluates to `false`, and if `e3` evaluates to `v`, then `if e1 then e2 else e3` evaluates to `v`
- Type checking
 - If `e1 : bool` and `e2 : t` and `e3 : t` then
`(if e1 then e2 else e3) : t`

If Expressions: Examples

```
# if 7 > 42 then "hello" else "goodbye";;
- : string = "goodbye"

# if true then 3 else 4;;
- : int = 3

# if false then 3 else 3.0;;
```

This expression has type float but is
here used with type int

Quiz 1

To what value does this expression evaluate?

if 22=0 then 1 else 2

- A. 0
- B. 1
- C. 2
- D. none of the above

Quiz 1

To what value does this expression evaluate?

`if 22=0 then 1 else 2`

- A. 0
- B. 1
- C. 2
- D. none of the above

Quiz 2

To what value does this expression evaluate?

if 22=0 then "bear" else 2

- A. 0
- B. 1
- C. 2
- D. none of the above

Quiz 2

To what value does this expression evaluate?

`if 22=0 then "bear" else 2`

- A. 0
- B. 1
- C. 2
- D. **none of the above:** doesn't type check so never gets a chance to be evaluated

Function Definitions

- OCaml functions are like mathematical functions
 - Compute a result from provided arguments

```
(* requires n>=0 *)
(* returns: n! *)
let rec fact n =
  if n = 0 then
    1
  else
    n * fact (n-1)
```

function body

Use (* *) for comments
(may nest)

Parameter
(type inferred)

rec needed for recursion

Structural equality

Line breaks, spacing ignored
(like C, C++, Java, not like Ruby)

Function Types

- In OCaml, `->` is the function type constructor
 - Type $t_1 \rightarrow t$ is a function with argument or *domain* type t_1 and return or *range* type t
 - Type $t_1 \rightarrow t_2 \rightarrow t$ is a function that takes *two* inputs, of types t_1 and t_2 , and returns a value of type t . Etc.
- Examples
 - `let next x = x + 1 (* type int -> int *)`
 - `let fn x = (int_of_float x) * 3`
$$(* \text{ type float} \rightarrow \text{int} *)$$
 - `fact`
$$(* \text{ type int} \rightarrow \text{int} *)$$

Type Checking Functions

- Syntax `let rec f x1 ... xn = e`
- Type checking
 - Conclude that $f : t_1 \rightarrow \dots \rightarrow t_n \rightarrow u$ if $e : u$ under the following assumptions:
 - $x_1 : t_1, \dots, x_n : t_n$ (arguments with their types)
 - $f : t_1 \rightarrow \dots \rightarrow t_n \rightarrow u$ (for recursion)
- Example
 - Given `n : int, fact : int -> int`
 - Does `if n = 0 then 1 ... : int` ?
 - It does!
 - Conclude `fact : int -> int`

```
let rec fact n =
  if n = 0 then
    1
  else
    n * fact (n-1)
```

Calling Functions

- Syntax $f \ e_1 \dots \ e_n$
 - Parentheses not required around argument(s)
 - No commas; use spaces instead
- Type checking
 - If $f : t_1 \rightarrow \dots \rightarrow t_n \rightarrow u$ and $e_1 : t_1, \dots, e_n : t_n$
then $f \ e_1 \dots \ e_n : u$
- Example:
 - `fact 1 : int`
 - since `fact : int → int` and `1 : int`
- Function call aka function application

Calling Functions

- Syntax $f \ e_1 \dots \ e_n$
- Evaluation
 - Evaluate arguments $e_1 \dots \ e_n$ to values $v_1 \dots \ v_n$
 - Order is actually right to left, not left to right
 - But this doesn't matter if $e_1 \dots \ e_n$ don't have side effects
 - Find the definition of f
 - `let rec f x1 ... xn = e`
 - Substitute v_i for x_i in e , yielding new expression e'
 - Evaluate e' to value v , which is the final result

Calling Functions

Example evaluation

- fact 2
 - if 2=0 then 1 else 2*fact(2-1)
 - 2 * fact 1
 - 2 * (if 1=0 then 1 else 1*fact(1-1))
 - 2 * 1 * fact 0
 - 2 * 1 * (if 0=0 then 1 else 0*fact(0-1))
 - 2 * 1 * 1
 - 2

```
let rec fact n =  
  if n = 0 then  
    1  
  else  
    n * fact (n-1)
```

Type Annotations

- The syntax $(e : t)$ asserts that “ e has type t ”
 - This can be added (almost) anywhere you like

```
let (x : int) = 3
let z = (x : int) + 5
```
- Define functions’ parameter and return types

```
let fn (x:int):float =
    (float_of_int x) *. 3.14
```

 - Note special position for return type
 - Thus `let g x:int = ...` means `g` returns `int`
 - Not that `x` has type `int`
- Checked by compiler: Very useful for debugging

Quiz 3: What is the value of `foo 4 2`

```
let rec foo n m =  
  if n >= 9 || n<0 then  
    m  
  else  
    n + m + 1
```

- Type Error
- 2
- 8
- 7

Quiz 3: What is the value of `foo 4 2`

```
let rec foo n m =  
  if n >= 9 || n<0 then  
    m  
  else  
    n + m + 1
```

- Type Error
- 2
- 8
- 7

Quiz 4: What is the value of **bar 4**

```
let rec bar(n:int):int =  
    if n = 0 || n = 1 then 1  
    else  
        bar (n-1) + bar (n-2)
```

- Syntax Error
- 4
- 5
- 8

Quiz 4: What is the value of **bar 4**

```
let rec bar(n:int):int =  
  if n = 0 || n = 1 then 1  
  else  
    bar (n-1) + bar (n-2)
```

- Syntax Error
- 4
- 5
- 8