

# CMSC 330: Organization of Programming Languages

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## OCaml Expressions and Functions

# Lecture Presentation Style

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- Our focus: **semantics** and **idioms** for OCaml
  - *Semantics* is what the language does
  - *Idioms* are ways to use the language well
- We will also cover some useful **libraries**
- **Syntax** is what you type, not what you mean
  - In one lang: Different syntax for similar concepts
  - Across langs: Same syntax for different concepts
  - Syntax can be a source of fierce disagreement among language designers!

# Expressions

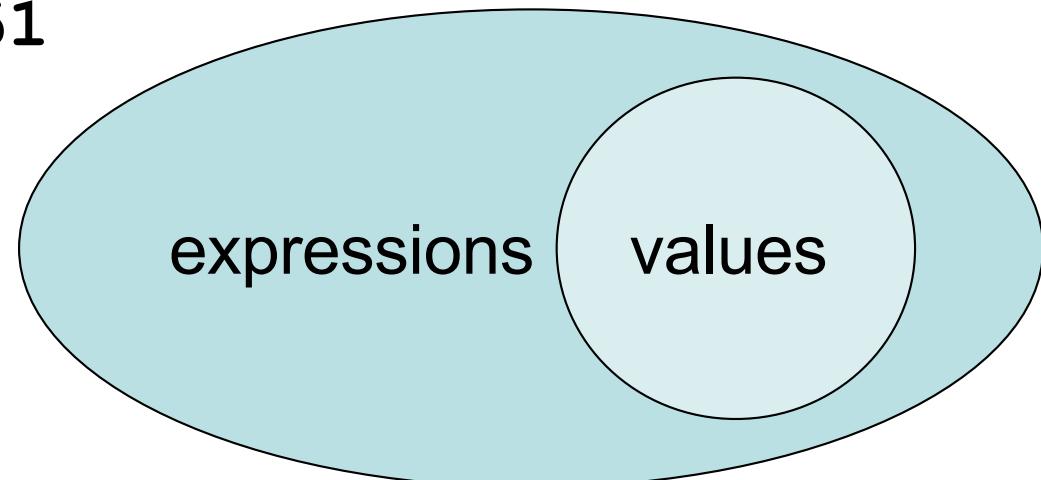
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- Expressions are our primary building block
  - Akin to *statements* in imperative languages
- Every kind of expression has
  - Syntax
    - We use metavariable  $e$  to designate an arbitrary expression
  - Semantics
    - Type checking rules (static semantics): produce a type or fail with an error message
    - Evaluation rules (dynamic semantics): produce a value
      - (or an exception or infinite loop)
      - Used *only* on expressions that type-check

# Values

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- A **value** is an expression that is final
  - Evaluating an expression means running it until it becomes a value
  - We use metavariable **v** to designate an arbitrary value
- **34** is a value, **true** is a value
- **34+17** is an *expression*, but *not* a value
  - It evaluates to 51



# Types

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- Types classify expressions
  - The set of values an expression could evaluate to
  - We use metavariable  $t$  to designate an arbitrary type
    - Examples include `int`, `bool`, `string`, and more.
- Expression  $e$  has type  $t$  if  $e$  will (always) evaluate to a value of type  $t$ 
  - $\{ \dots, -1, 0, 1, \dots \}$  are values of type `int`
  - $34+17$  is an expression of type `int`, since it evaluates to  $51$ , which has type `int`
  - Write  $e : t$  to say  $e$  has type  $t$
  - Determining that  $e$  has type  $t$  is called type checking (or simply, typing)

# If Expressions

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- Syntax
  - `if e1 then e2 else e3`
- Evaluation
  - If  $e1$  evaluates to `true`, and if  $e2$  evaluates to  $v$ , then `if e1 then e2 else e3` evaluates to  $v$
  - If  $e1$  evaluates to `false`, and if  $e3$  evaluates to  $v$ , then `if e1 then e2 else e3` evaluates to  $v$
- Type checking
  - If  $e1$  has type `bool` and  $e2$  has type  $t$  and  $e3$  has type  $t$  then `if e1 then e2 else e3` has type  $t$

# If Expressions

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- Syntax
  - `if e1 then e2 else e3`
- Evaluation
  - If `e1` evaluates to `true`, and if `e2` evaluates to `v`, then `if e1 then e2 else e3` evaluates to `v`
  - If `e1` evaluates to `false`, and if `e3` evaluates to `v`, then `if e1 then e2 else e3` evaluates to `v`
- Type checking
  - If `e1 : bool` and `e2 : t` and `e3 : t` then  
`if e1 then e2 else e3 : t`

# If Expressions

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- Syntax
  - `if e1 then e2 else e3`
- Evaluation
  - If `e1` evaluates to `true`, and if `e2` evaluates to `v`, then `if e1 then e2 else e3` evaluates to `v`
  - If `e1` evaluates to `false`, and if `e3` evaluates to `v`, then `if e1 then e2 else e3` evaluates to `v`
- Type checking
  - If `e1 : bool` and `e2 : t` and `e3 : t` then  
`(if e1 then e2 else e3) : t`

# If Expressions: Examples

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```
# if 7 > 42 then "hello" else "goodbye";;  
- : string = "goodbye"
```

```
# if true then 3 else 4;;  
- : int = 3
```

```
# if false then 3 else 3.0;;
```

Error: This expression has type float but  
an expression was expected of type  
**int**

# Quiz 1

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To what value does this expression evaluate?

**if 22<0 then 2 else 1**

- A. 2
- B. 1
- C. 0
- D. none of the above

# Quiz 1

---

To what value does this expression evaluate?

`if 22<0 then 2 else 1`

- A. 2
- B. 1
- C. 0
- D. none of the above

## Quiz 2

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To what value does this expression evaluate?

**if 22<0 then "bear" else 2**

- A. 2
- B. 1
- C. 0
- D. none of the above

## Quiz 2

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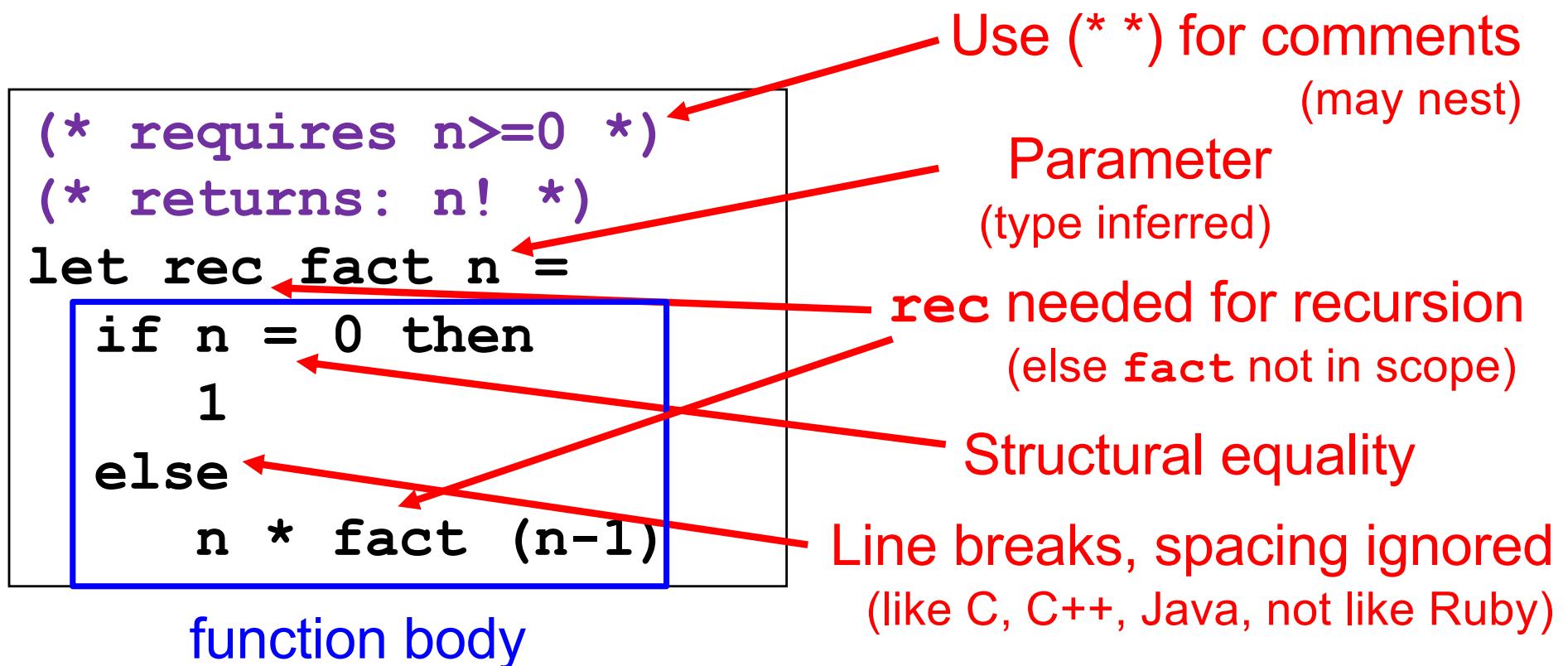
To what value does this expression evaluate?

`if 22<0 then "bear" else 2`

- A. 2
- B. 1
- C. 0
- D. **none of the above:** doesn't type check so never gets a chance to be evaluated

# Function Definitions

- OCaml functions are like mathematical functions
  - Compute a result from provided arguments



The diagram shows an OCaml function definition with various annotations:

```
(* requires n>=0 *)
(* returns: n! *)
let rec fact n =
  if n = 0 then
    1
  else
    n * fact (n-1)
```

- Annotations:
  - Use (\* \*) for comments (may nest)
  - Parameter (type inferred)
  - rec needed for recursion (else fact not in scope)
  - Structural equality
  - Line breaks, spacing ignored (like C, C++, Java, not like Ruby)
- A blue box encloses the body of the function definition.
- A label "function body" is positioned below the blue box.

# Type Inference

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- As we just saw, a declared variable need not be annotated with its type
  - The type can be inferred

```
(* requires n>=0 *)
(* returns: n! *)
let rec fact n =
  if n = 0 then
    1
  else
    n * fact (n-1)
```

n's type is **int**. Why?

= is an infix function that takes two **ints** and returns a **bool**; so **n** must be an **int** for **n = 0** to type check

- Type inference happens as a part of type checking
  - Determines a type that satisfies code's constraints

# Function Types

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- In OCaml, `->` is the function type constructor
  - Type  $t_1 \rightarrow t$  is a function with argument or *domain* type  $t_1$  and return or *range* type  $t$
  - Type  $t_1 \rightarrow t_2 \rightarrow t$  is a function that takes *two* inputs, of types  $t_1$  and  $t_2$ , and returns a value of type  $t$ . Etc.
- Examples
  - `let next x = x + 1 (* type int -> int *)`
  - `let fn x = (int_of_float x) * 3`  
 $\qquad\qquad\qquad (* type float -> int *)$
  - `fact`  
 $\qquad\qquad\qquad (* type int -> int *)$

# Function Types

# Considering inference

- `+` has type `int -> int -> int`.
    - Therefore, `x + 1` forces `x` to be an `int`.
  - `int_of_float` has type `float -> int`.
    - Therefore `(int_of_float x)` forces `x` to be a `float`

## • Examples

- ```
- let next x = x + 1 (* type int -> int *)
- let fn x = (int_of_float x) * 3
              (* type float -> int *)
- fact          (* type int -> int *)
```

# Type Checking Functions

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- Syntax `let rec f x1 ... xn = e`
- Type checking
  - Conclude that  $f : t_1 \rightarrow \dots \rightarrow t_n \rightarrow u$  if  $e : u$  under the following assumptions:
    - $x_1 : t_1, \dots, x_n : t_n$  (arguments with their types)
    - $f : t_1 \rightarrow \dots \rightarrow t_n \rightarrow u$  (for recursion)
- Example
  - Given `n : int, fact : int -> int`
  - Does `if n = 0 then 1 ... : int` ?
    - It does!
  - Conclude `fact : int -> int`

```
let rec fact n =
  if n = 0 then
    1
  else
    n * fact (n-1)
```

# Calling Functions

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- Syntax  $f \ e_1 \dots \ e_n$ 
  - Parentheses not required around argument(s)
  - No commas; use spaces instead
- Type checking
  - If  $f : t_1 \rightarrow \dots \rightarrow t_n \rightarrow u$  and  $e_1 : t_1, \dots, e_n : t_n$   
then  $f \ e_1 \dots \ e_n : u$
- Example:
  - `fact 1 : int`
  - since `fact : int -> int` and `1 : int`
- Function call aka function application

# Calling Functions

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- Syntax  $f \ e_1 \dots \ e_n$
- Evaluation
  - Evaluate arguments  $e_1 \dots \ e_n$  to values  $v_1 \dots \ v_n$ 
    - Order is actually right to left, not left to right
    - But this doesn't matter if  $e_1 \dots \ e_n$  don't have side effects
  - Find the definition of  $f$ 
    - `let rec f x1 ... xn = e`
  - Substitute  $v_i$  for  $x_i$  in  $e$ , yielding new expression  $e'$
  - Evaluate  $e'$  to value  $v$ , which is the final result

# Calling Functions

## Example evaluation

- fact 2
  - if 2=0 then 1 else 2\*fact(2-1)
  - 2 \* fact 1
  - 2 \* (if 1=0 then 1 else 1\*fact(1-1))
  - 2 \* 1 \* fact 0
  - 2 \* 1 \* (if 0=0 then 1 else 0\*fact(0-1))
  - 2 \* 1 \* 1
  - 2

```
let rec fact n =  
  if n = 0 then  
    1  
  else  
    n * fact (n-1)
```

# Type Annotations

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- The syntax  $(e : t)$  asserts that “ $e$  has type  $t$ ”
  - This can be added (almost) anywhere you like

```
let (x : int) = 3
let z = (x : int) + 5
```
- Define functions’ parameter and return types

```
let fn (x:int):float =
    (float_of_int x) *. 3.14
```

  - Note special position for return type
  - Thus `let g x:int = ...` means `g` returns `int`
    - *Not* that `x` has type `int`
- Checked by compiler: Very useful for debugging

# Quiz 3: What is the type of `foo 4 2`

---

```
let rec foo n m =  
  if n >= 9 || n<0 then  
    m  
  else  
    n + m + 1
```

- a) Type Error
- b) int
- c) float
- d) int -> int -> int

# Quiz 3: What is the type of `foo 4 2`

---

```
let rec foo n m =  
  if n >= 9 || n<0 then  
    m  
  else  
    n + m + 1
```

- a) Type Error
- b) int
- c) float
- d) int -> int -> int

# Quiz 4: What is the value of **bar 4**

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```
let rec bar(n:int):int =  
  if n = 0 || n = 1 then 1  
  else  
    bar (n-1) + bar (n-2)
```

- a) Syntax Error
- b) 4
- c) 5
- d) 8

# Quiz 4: What is the value of **bar 4**

---

```
let rec bar(n:int):int =  
    if n = 0 || n = 1 then 1  
    else  
        bar (n-1) + bar (n-2)
```

- a) Syntax Error
- b) 4
- c) 5
- d) 8