

CMSC 433

Programming Language Technologies and Paradigms

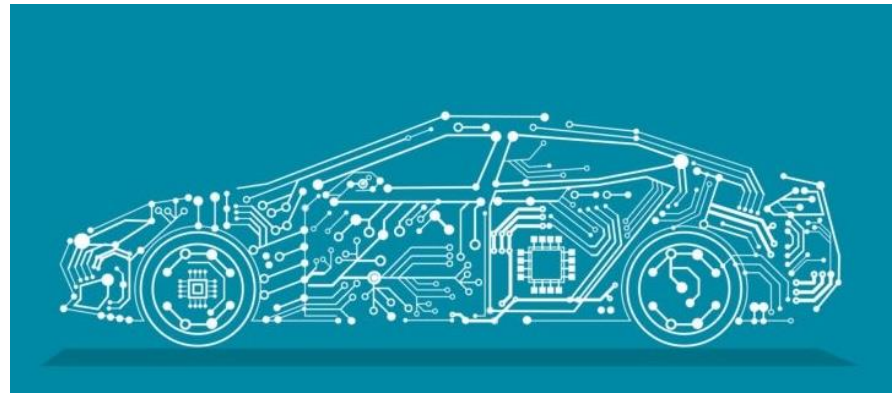
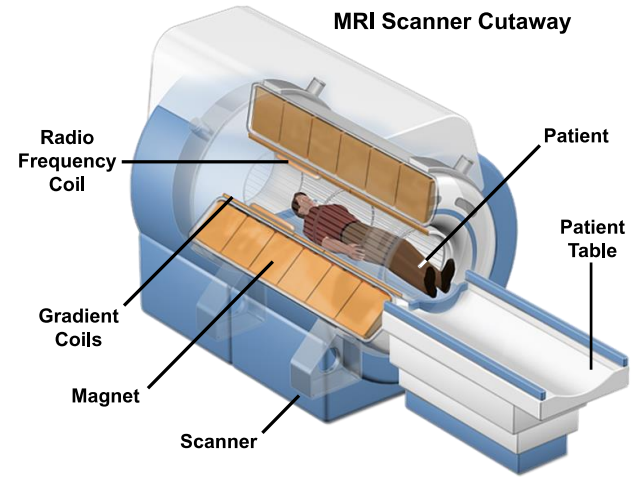
Introduction

CMSC433 History

CMSC433 used to be a study of Concurrent programming.

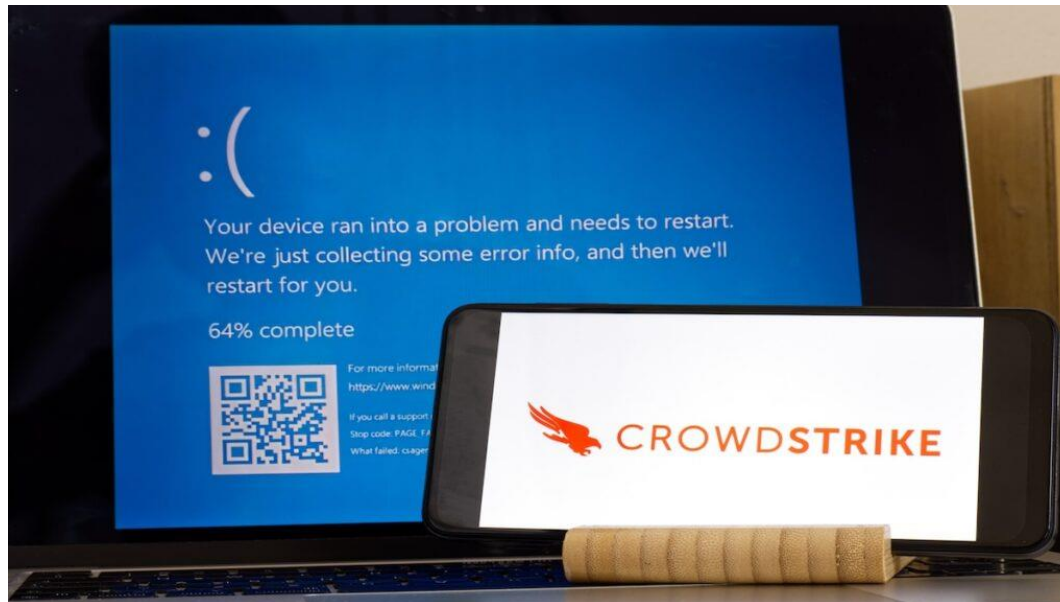
Now it is on program proofs and verification.

Software is everywhere



Software has bugs

A software bug is a defect in a computer program or system that causes an undesired result.



Build Better Software

- We test software to check whether software satisfies expectations. Testing can show errors but not their absence.
- Software errors in critical systems can cause major disasters.
- This course is an introduction to techniques to get certainty that your program does what it is supposed to do.

Course logistics

- ▶ 6 assignments
- ▶ Midterm 10/23 (Wednesday)
- ▶ Final exam: Monday, December 15
- ▶ Several surveys and quizzes (on ELMS)

Grading

Assignments	50%
Quizzes & surveys	5%
Attendance	5%
Midterm	15%
Final	25%

Key resources

- ▶ Class web page (syllabus, assignments, course notes)
 - <https://www.cs.umd.edu/class/fall2025/cmssc433/>
- ▶ ELMS (announcements, grades)
- ▶ Piazza (communication, discussion)
- ▶ Gradescope (assignments, exams)
- ▶ Office Hours

Course Structure

- ▶ Topics:
 - Dafny
 - SAT solving and its applications
 - Solver aided programming
 - Computer aided (interactive) theorem proving
 - Testing
 - Supplementary reading

Dafny

- ▶ Dafny is both a **programming language** and a **formal verification** tool designed to help developers write programs that are mathematically proven to be correct.
- ▶ In Dafny, you can write **preconditions, postconditions, and invariants**. Dafny then uses an automated theorem prover (based on SMT solvers like Z3) to check whether your program **meets its specification**.
- ▶ This means Dafny doesn't just test code on some inputs—it proves properties hold for **all possible executions**.

SAT/SMT Solvers

- ▶ SAT/SMT solvers are computer programs which aim to solve the **Boolean satisfiability problem**.
- ▶ They are used in program verification (like **Dafny, SPARK, Why3**), theorem proving, and software/hardware correctness.
- ▶ Tools like Dafny **rely on SMT** solvers (e.g., Z3) to mathematically prove that your code always meets its specification.
- ▶ Modern SAT/SMT solvers (like Z3, CVC5, MiniSat) can solve *huge, real-world problems surprisingly fast*.

Computer Aided (interactive) Theorem Proving

- ▶ A software tool to assist with the development of **formal proofs** by human–machine collaboration.
 - Human guide the proof construction.
 - Machine checks the low-level details
- ▶ Examples:
 - Rocq (Coq), Isabelle HOL, F*, Lean4

Building Reliable Software

Cost of Software Errors

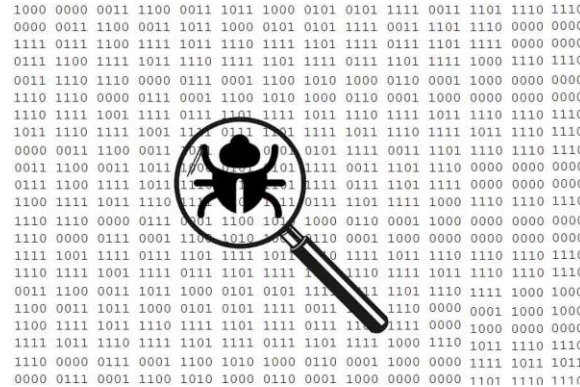
\$2.8 Trillion in 2020 alone

Source: Forbes

<https://www.forbes.com/councils/forbestechcouncil/2023/12/26/costly-code-the-price-of-software-errors/>

Cost of Software Errors

Estimated 50% of programmers time spent on finding and fixing bugs.



Software failure examples: 2024 CrowdStrike incident

On **July 19th 2024**, CrowdStrike distributed a faulty update to its Falcon Sensor security software that caused widespread problems with Microsoft Windows.

Roughly **8.5 million** systems crashed and were unable to properly restart.

The worldwide financial damage has been estimated to be at least US \$10 billion.



Nobody travels on 07/19



Software failure examples: Ariane flight V88

Ariane flight V88 (Ariane 5 rocket)
exploded right after launch in 1996.

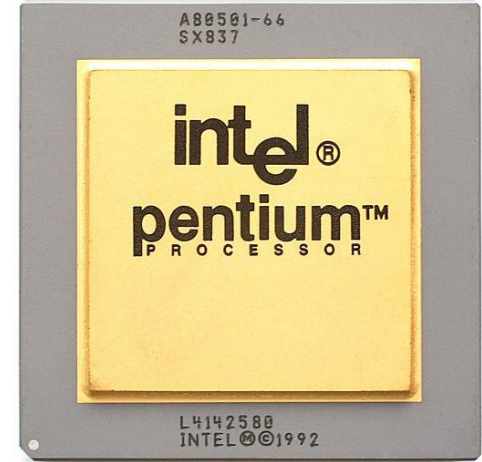
Conversion of 64-bit float to 16-bit
integer caused an exception (made it
crash)

European space agency spent **10**
years and \$7 billion to produce Ariane
5



Software failure examples: Pentium Floating Point (FDIV) Bug

A hardware bug affecting the floating-point unit (FPU) of the early **Intel** Pentium processors in 1994.



- Incorrect result through floating point division
- Rarely encountered in practice
- 1 in 9 billion floating point divides with random parameters would produce inaccurate results (Byte magazine)
- 475 million dollars, reputation of Intel.

Zune Leap Year Freeze

At midnight of December 31, 2008, all the millions of Zune 30 that Microsoft sold froze.

```
BOOL ConvertDays(UINT32 days, SYSTEMTIME* lpTime){
    ...
    year = ORIGINYEAR; /* = 1980 */
    while (days > 365){
        if (IsLeapYear(year)){
            if (days > 366){
                days -= 366;
                year += 1;
            }
        }else{
            days -= 365;
            year += 1;
        }
    }
    ...
}
```

Not just economic loss: **Toyota Unintended Acceleration**

- Bugs in electronic throttle control system (2009).
- Car kept accelerating on its own.
- May have caused up to 89 deaths in accidents.
- Recalls of 10 million vehicles.



Not just economic loss, Therac-25

- a computer-controlled radiation therapy machine (1985-1987)
- some patients were given massive overdoses of radiation.
- Killed four and left two others with lifelong injuries.



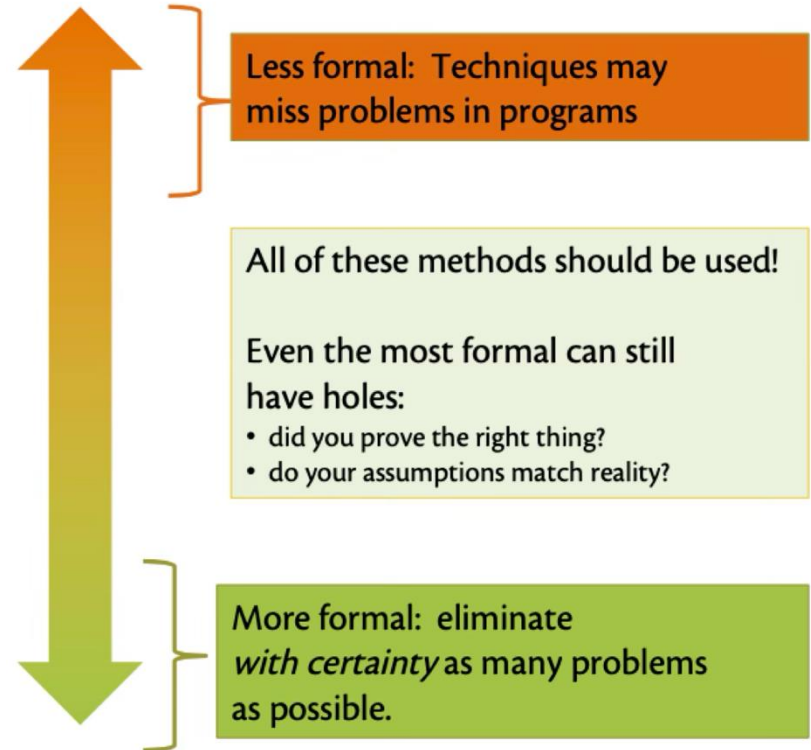
Problem Source

- **Requirements**: Incomplete, inconsistent, ...
- **Design**: Flaws in design
- **Implementation**: Programming errors,...
- **Tools**: Defects in support systems and tools used

How can you get some assurance that a program does what you want it to do?

Approaches to Validation

- Social
 - Code reviews
 - Extreme/pair programming
- Methodological
 - Design patterns
 - Test-driven development
 - Version control
 - Bug Tracking
- Technological
 - Static analysis
 - Fuzzers
- Mathematical
 - Sound Type Systems
 - Formal verification



Testing

- ▶ Evaluating software by observing its execution
- ▶ Execute program with the intent of finding failures (try out inputs, see if outputs are correct)

How do we test?

- ▶ Test: try out inputs, see if outputs are correct
- ▶ Testing means to execute a program with the intent of detecting failure
- ▶ terminology, testing levels, unit testing, black box vs white box, principles of test-set construction/coverage, automated and repeatable testing (JUnit)

Formal verification

- ▶ Determine whether a piece of software fulfils a set of **formal** requirements in **every** execution
 - Formally prove method correct (find evidence of absence of failure)

Formal verification

- ▶ Scaled to 10s of lines of code in 1970s
- ▶ Now, research projects scale to real software:
 - CompCert: A verified C compiler
 - seL4: verified microkernel OS
 - Ynot: verified DBMS, web services
- ▶ In another 50 years?

Some failures are obvious

- ▶ obviously wrong output/behavior
 - non-termination
 - crash
 - freeze

- ▶ . . . but most are not!

In general, what constitutes a failure, is defined by: a **specification!**

Specification

- ▶ **Specification**: An **unambiguous** description of what a program should do.
- ▶ **Bug**: Failure to meet specification.
- ▶ Unclear Specification leads to failure



Specification: Example

`Sort(src: Integer Array) -> Integer Array`

- Specification:
 - **Requires**: `src` is an array of integer
 - **Ensures**: returns a sorted array

Is this a good specification?

Specification: Example

Sort(src: Integer Array) -> Integer Array

► Specification:

- **Requires:** src is an array of integer
- **Ensures:** returns a sorted array

`Sort([3,1,4,5]) == []` ✗

`Sort([3,1,4,5]) == [1,2,3]` ✗

Specification: Example

Sort(src: Integer Array) -> Integer Array

► Specification:

- **Requires:** **src** is an array of integer
- **Ensures:** returns a sorted array with only elements from the input

`Sort([3,1,4,5]) == [1,1,4]` ✗

`Sort([3,1,4,5]) == [1,3,3,5]` ✗

Specification: Example

Sort(src: Integer Array) -> Integer Array

► Specification:

- **Requires:** src is an array of integer
- **Ensures:** returns a permutation of src that is sorted

Sort([]) == ?

Sort(null) == ?

Permutation?

Specification: Example

Sort(src: Integer Array) -> Integer Array

▸ Specification:

- **Requires:** src is a **non-null** array of integer
- **Ensures:** returns a permutation of src that is sorted

Specification of a method

method **m()**

Requires: Precondition

Ensures: Postcondition

Means:

- If a caller of **m()** fulfills the required Precondition, then the callee **m()** ensures that the Postcondition holds after **m()** finishes.
- Garbage in, garbage out

Failure vs Correctness

- ▶ What constitutes a **failure**

A method fails when it is called in a state fulfilling the required precondition of its contract and it does not terminate in a state fulfilling the postcondition to be ensured.

Failure vs Correctness

- A method is **correct** means:
 - whenever it is started in a state fulfilling the required precondition, then it **terminates** in a state fulfilling the postcondition to be ensured.
- Correctness amounts to proving **absence of failures**! A correct method cannot fail!

Verification

- ▶ Testing cannot guarantee correctness, i.e., absence of failures
- ▶ Verification: Mathematically prove method correct
 - Goal: find evidence for absence of failures
- ▶ This course: Formal verification (logics, tool support)

Summary

- ▶ CS433 introduces techniques for **ensuring program correctness**—that is, proving that a program does what it is intended to do. The course provides a mathematical foundation for the rigorous analysis of real-world software systems.
- ▶ The skills developed in this course are in high demand, as nearly all human interactions are now mediated by software—and it is especially critical to guarantee the correctness of AI-generated programs.