Overview

- Binary trees
  - Complete

- Heaps
  - Insert
  - getSmallest

- Heap applications
  - Heapsort
  - Priority queues
Complete Binary Trees

An binary tree (height h) where

- Perfect tree to level \( h-1 \)
- Leaves at level \( h \) are as far left as possible
Complete Binary Trees

Basic complete tree shape

Not Allowed
Heaps

Two key properties
- Complete binary tree
- Value at node
  - Smaller than or equal to values in subtrees

Example heap
- \(X \leq Y\)
- \(X \leq Z\)
Heap & Non-heap Examples

Heaps

Non-heaps
Heap Properties

- Heaps are balanced trees
  - Height = \( \log_2(n) = O(\log(n)) \)

- Can find smallest element easily
  - Always at top of heap!

- Can organize heap to find maximum value
  - Value at node larger than values in subtrees
  - Heap can track either min or max, but not both
Heap

Key operations
- Insert (X)
- getSmallest()

Key applications
- Heapsort
- Priority queue
Heap Operations – Insert( X )

Algorithm
1. Add X to end of tree
2. While (X < parent)
   Swap X with parent // X bubbles up tree

Complexity
- # of swaps proportional to height of tree
- O( log(n) )
Heap Insert Example

1) Insert to end of tree
2) Compare to parent, swap if parent key larger
3) Insert complete

Insert (20)
Heap Insert Example

Insert (8)

1) Insert to end of tree
2) Compare to parent, swap if parent key larger
3) Insert complete
Heap Operation – getSmallest()

Algorithm

1. Get smallest node at root
2. Replace root with X at end of tree
3. While ( X > child )
   - Swap X with smallest child  // X drops down tree
4. Return smallest node

Complexity

- # swaps proportional to height of tree
- $O(\log(n))$
Heap GetSmallest Example

getSmallest ()

1) Replace root with end of tree
2) Compare node to children, if larger swap with smallest child
3) Repeat swap if needed
Heap GetSmallest Example

getSmallest ()

1) Replace root with end of tree
2) Compare node to children, if larger swap with smallest child
3) Repeat swap if needed
Heap Implementation

Can implement heap as array
- Store nodes in array elements
- Assign location (index) for elements using formula

(a) Heap represented as a tree
(b) Heap represented as an array
Observations

- Compact representation
- Edges are implicit (no storage required)
- Works well for complete trees (no wasted space)
Heap Implementation

Calculating node locations

- Array index $i$ starts at 0
- Parent$(i) = \left\lfloor \frac{(i - 1)}{2} \right\rfloor$
- LeftChild$(i) = 2 \times i + 1$
- RightChild$(i) = 2 \times i + 2$
Heap Implementation

Example

- Parent(1) = \( \lfloor \frac{1 - 1}{2} \rfloor = \lfloor 0/2 \rfloor = 0 \)
- Parent(2) = \( \lfloor \frac{2 - 1}{2} \rfloor = \lfloor 1/2 \rfloor = 0 \)
- Parent(3) = \( \lfloor \frac{3 - 1}{2} \rfloor = \lfloor 2/2 \rfloor = 1 \)
- Parent(4) = \( \lfloor \frac{4 - 1}{2} \rfloor = \lfloor 3/2 \rfloor = 1 \)
- Parent(5) = \( \lfloor \frac{5 - 1}{2} \rfloor = \lfloor 4/2 \rfloor = 2 \)
Heap Implementation

Example

- LeftChild(0) = $2 \times 0 + 1 = 1$
- LeftChild(1) = $2 \times 1 + 1 = 3$
- LeftChild(2) = $2 \times 2 + 1 = 5$

Diagram:

```
    X0
   /   \
X1    X2
  /   /   \
X3 X4 X5
```

Array $X$:

```
<table>
<thead>
<tr>
<th>X0</th>
<th>X1</th>
<th>X2</th>
<th>X3</th>
<th>X4</th>
<th>X5</th>
</tr>
</thead>
<tbody>
<tr>
<td>index: 0</td>
<td>1</td>
<td>2</td>
<td>3</td>
<td>4</td>
<td>5</td>
</tr>
</tbody>
</table>
```
Heap Implementation

Example

- RightChild(0) = 2 × 0 + 2 = 2
- RightChild(1) = 2 × 1 + 2 = 4
Heap Application – Heapsort

- Use heaps to sort values
  - Heap keeps track of smallest element in heap

Algorithm
1. Create heap
2. Insert values in heap
3. Remove values from heap (in ascending order)

Complexity
- $O(n \log(n))$
Heapsort Example

Input
- 11, 5, 13, 6, 1

View heap during insert, removal
- As tree
- As array
Heapsort – Remove Values

(a) Print root = 1

(b) Rebuild heap

(c) Print root = 5

(d) Rebuild heap

(e) Print root = 6

(f) Rebuild heap

(g) Print root = 11

(h) Rebuild heap

(f) Print root = 13

Done
Heapsort – Insert in to Array 1

Input

11, 5, 13, 6, 1

Index = 0 1 2 3 4

Insert 11

11
Heapsort – Insert in to Array 2

Input
- 11, 5, 13, 6, 1

Index = 0 1 2 3 4

Insert 5
- 11 5

Swap
- 5 11
Heapsort – Insert in to Array 3

Input

11, 5, 13, 6, 1

Index = 0 1 2 3 4

Insert 13

<table>
<thead>
<tr>
<th></th>
<th>5</th>
<th>11</th>
<th>13</th>
<th></th>
<th></th>
</tr>
</thead>
</table>
Heapsort – Insert in to Array 4

Input

11, 5, 13, 6, 1

Index = 0 1 2 3 4

Insert 6

Insert 6

Swap

...
Heapsort – Remove from Array 1

Input
11, 5, 13, 6, 1

Index = 0 1 2 3 4

Remove root
1 5 13 11 6

Replace
6 5 13 11

Swap w/ child
5 6 13 11
Heapsort – Remove from Array 2

Input
- 11, 5, 13, 6, 1

Index =
<table>
<thead>
<tr>
<th>0</th>
<th>1</th>
<th>2</th>
<th>3</th>
<th>4</th>
</tr>
</thead>
<tbody>
<tr>
<td>11</td>
<td>5</td>
<td>6</td>
<td>13</td>
<td>11</td>
</tr>
</tbody>
</table>

Remove root

Replace

Swap w/ child
Queue

- Linear data structure
- First-in First-out (FIFO)
- Implement as array / linked list
Heap Application – Priority Queue

Priority queue
- Elements are assigned priority value
- Higher priority elements are taken out first
- Implement as heap
  - Enqueue ⇒ insert( )
  - Dequeue ⇒ getSmallest( )

Enqueue
Dequeue
Priority Queue

Properties
- Lower value = higher priority
- Heap keeps highest priority items in front

Complexity
- Enqueue ⇒ insert( ) = O( log(n) )
- Dequeue ⇒ getSmallest( ) = O( log(n) )
- For any heap
Heap vs. Binary Search Tree

**Binary search tree**
- Keeps values in sorted order
- Find any value
  - $O(\log(n))$ for balanced tree
  - $O(n)$ for degenerate tree (worst case)

**Heap**
- Keeps smaller values in front
- Find minimum value
  - $O(\log(n))$ for any heap