

CMSC 132: Object-Oriented Programming II



Synchronization in Java

Department of Computer Science
University of Maryland, College Park

Multithreading Overview

■ Motivation & background

■ Threads

■ Creating Java threads

■ Thread states

■ Scheduling

■ Synchronization

■ Data races ←

■ Locks

■ Deadlock



Data Race

■ Definition

- Concurrent accesses to same shared variable, where at least one access is a write

■ Properties

- Order of accesses may change result of program
- May cause intermittent errors, very hard to debug

■ Example

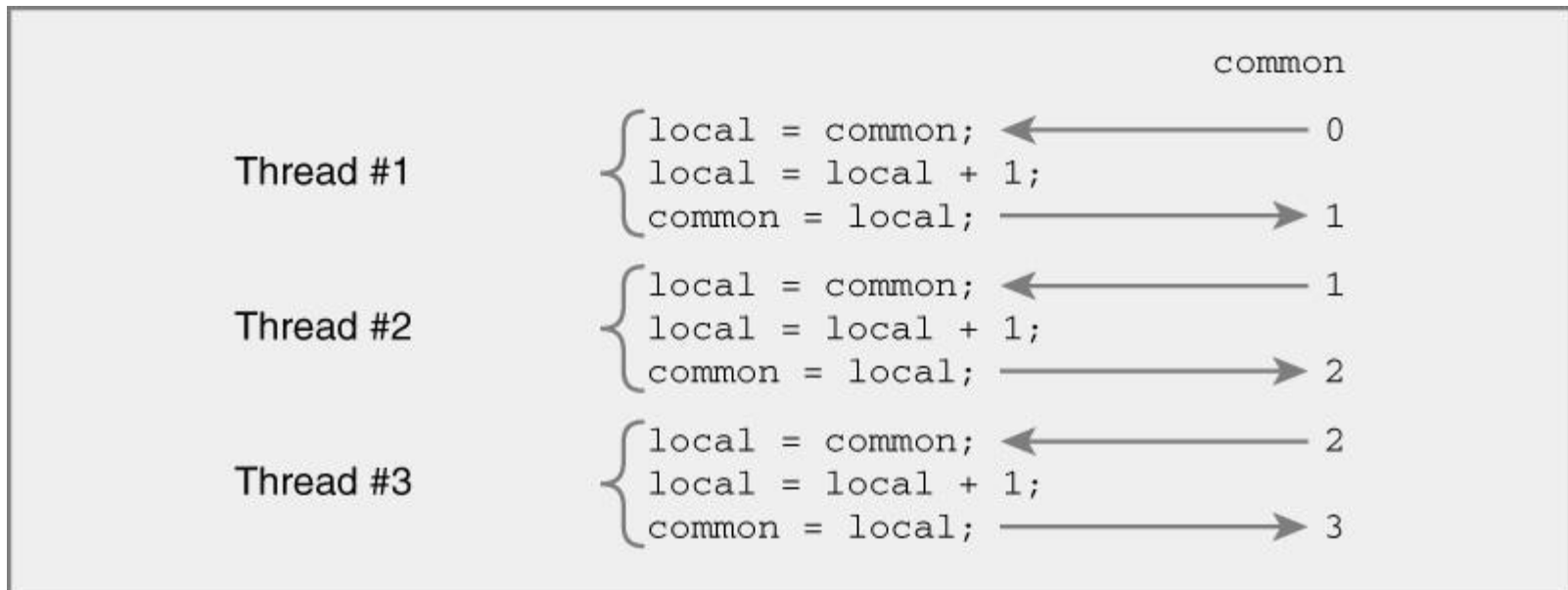
```
public class DataRace extends Thread {  
    static int x;    // shared variable x causing data race  
    public void run() { x = x + 1; } // access to x  
}
```

Data Race Example

```
public class DataRace extends Thread {  
    static int common = 0;  
    public void run() {  
        int local = common; // data race  
        local = local + 1;  
        common = local; // data race  
    }  
    public static void main(String[] args) throws InterruptedException {  
        int max = 3;  
        DataRace[] allThreads = new DataRace[max];  
        for (int i = 0; i < allThreads.length; i++)  
            allThreads[i] = new DataRace();  
        for (DataRace t : allThreads)  
            t.start();  
        for (DataRace t : allThreads)  
            t.join();  
        System.out.println(common); // may not be 3  
    }  
}
```

Data Race Example

■ Sequential execution output



Data Race Example

■ Concurrent execution output (possible case)

The diagram illustrates a concurrent execution scenario with three threads. On the right, a vertical box represents a shared memory or common variable, labeled 'common' at the top. It contains three '0' values at the top and three '1' values at the bottom. Arrows indicate the flow of data between the threads and this common state. The threads are listed on the left, with their actions and the corresponding state of the common variable shown to their right.

Thread	Action	Common Value
Thread #1	local = common;	0
Thread #2	local = common;	0
Thread #3	local = common;	0
Thread #1	local = local + 1;	
Thread #2	local = local + 1;	
Thread #3	local = local + 1;	
Thread #1	common = local;	1
Thread #2	common = local;	1
Thread #3	common = local;	1

■ Result depends on thread execution order!

Synchronization

■ Definition

- Coordination of events with respect to time

■ Properties

- May be needed in multithreaded programs to eliminate **data races**
- Incurs runtime overhead
- Excessive use can reduce performance

Lock

■ Definition

- Entity can be held by only one thread at a time

■ Properties

- A type of synchronization
- Used to enforce **mutual exclusion**
 - Thread can acquire / release locks
 - Only 1 thread can acquire lock at a time
- Thread will wait to acquire lock (stop execution)
 - If lock held by another thread
- Used to implement **monitors**
 - Only 1 thread can execute (locked) code at a time

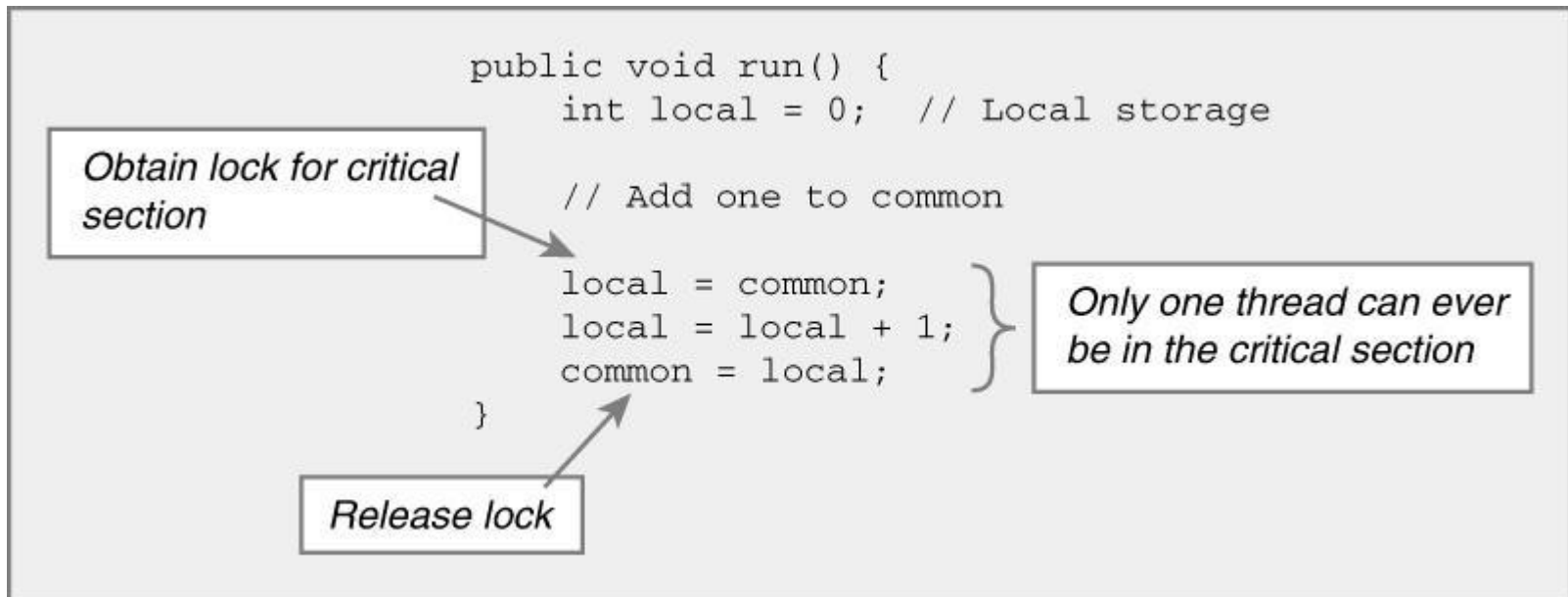


Synchronized Objects in Java

- Java objects provide locks
- To acquire the lock use the **synchronized** keyword
 - Apply **synchronized** keyword to object
 - Will acquire / release lock associated with object
 - Mutual exclusion for code in synchronization **block** (block defined by **synchronized() { }**)
- Example

```
Object x = new Object();  
  
block { synchronized(x) { // acquire lock on x on entry  
    ...                // hold lock on x in block  
}                      // release lock on x on exit
```

Fixing Data Race In Our Example



Lock Example

```
public class DataRace extends Thread {  
    static int common = 0;  
    static Object o;                // all threads use o's lock  
    public void run() {  
        synchronized(o) {         // single thread at once  
            int local = common;    // data race eliminated  
            local = local + 1;  
            common = local;  
        }  
    }  
    public static void main(String[] args) {  
        o = new Object();  
        ...  
    }  
}
```

- Keep in mind that lock objects do not need to be static (static is used in the above example to share the lock among all threads)
- How would you solve the data race without using a static lock object?

Locks in Java

■ Properties

- A lock can be held by only one thread at a time
- Locked block of code $\Rightarrow$ **critical section**
- Note: critical section should not be confused with the term critical section use for algorithmic complexity analysis
- You must protect the critical section wherever it appears in your code, otherwise several threads may access the critical section simultaneously
 - In other words, protecting the critical section with a lock, in a section(s) of your code, will not automatically protect the critical section everywhere it appears in your code

■ Lock is released when block terminates

- End of block reached
- Exit block due to return, continue, break
- Exception thrown

Synchronized Methods In Java

- Apply synchronized keyword to method
 - Mutual exclusion for entire body of method
 - Synchronizes on current object

- Example

block


```
synchronized foo() { ...code... }
```

// shorthand notation for

```
foo() {  
    ↓  
    synchronized (this) { ...code... }  
}
```

Synchronized Methods In Java

```
public synchronized void enqueue( Object item ) {  
    // Body of method goes here  
}
```



Shorthand notation for

```
public void enqueue( Object item ) {  
    synchronized ( this ) {  
        // Body of method goes here  
    }  
}
```

Examples

- **Code distribution provides an example of an account shared by two kinds of buyers (Normal buyer and Excessive buyer)**
- **Four versions**
 - **Version1 → Buyers synchronize using the account object as lock object**
 - **Version 2 → deposit/withdrawal/getBalance methods synchronize using a lock object (defined as instance variable of Account object)**
 - **Version 3 → As in Version 2 but lock object is the current object**
 - **Version 4 → deposit/withdrawal/getBalance methods defined as synchronized methods**

Synchronization Issues

- 1. Use same lock to provide mutual exclusion**
- 2. Ensure atomic transactions**
- 3. Avoiding deadlock**

Issue 1) Using Same Lock

■ Potential problem

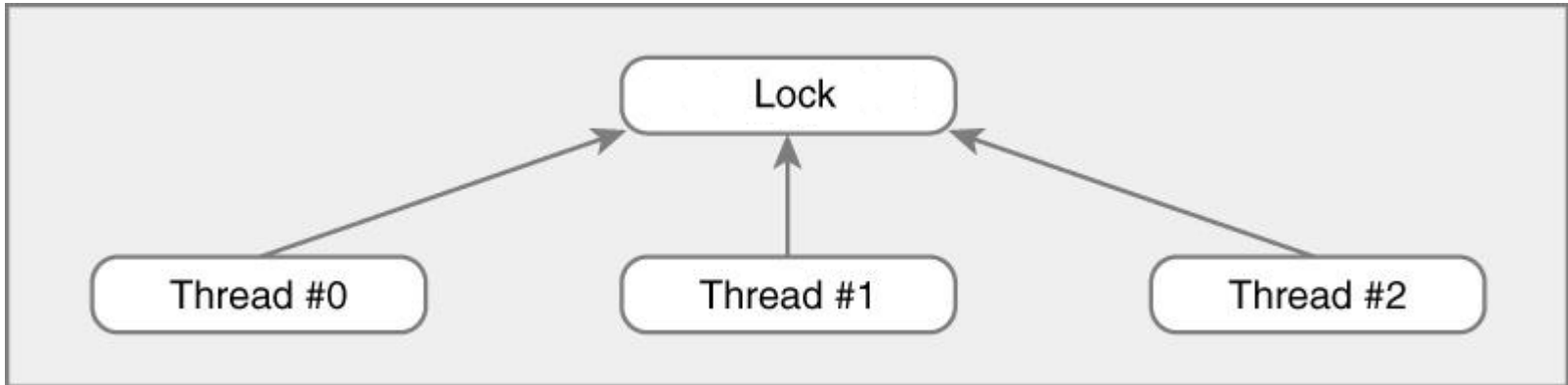
- Mutual exclusion depends on threads acquiring same lock
- No synchronization if threads have different locks

■ Example

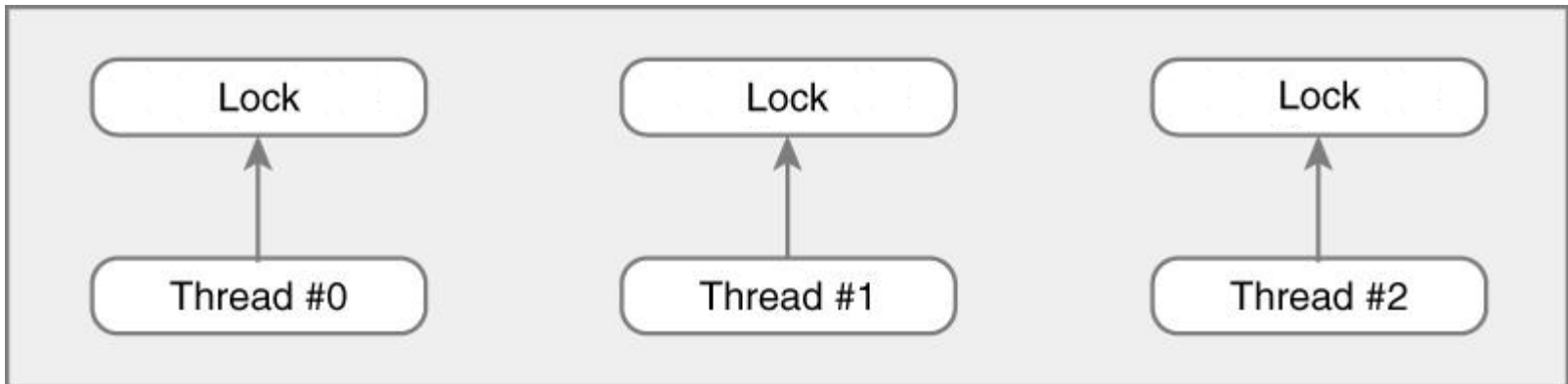
```
foo() {  
    Object o = new Object(); // different o per thread  
    synchronized(o) {  
        ... // potential data race  
    }  
}
```

Locks in Java

- **Single lock for all threads (mutual exclusion)**



- **Separate locks for each thread (no synchronization)**



Lock Example – Incorrect Version

```
public class DataRace extends Thread {  
    static int common = 0;  
    public void run() {  
        Object o = new Object(); // different o per thread  
        synchronized(o) {  
            int local = common;    // data race  
            local = local + 1;  
            common = local;        // data race  
        }  
    }  
    public static void main(String[] args) {  
        ...  
    }  
}
```

Issue 2) Atomic Transactions

■ Potential problem

- Sequence of actions must be performed as single **atomic transaction** to avoid data race
- Ensure lock is held for duration of transaction

■ Example

```
synchronized(o) {  
    int local = common;  
    local = local + 1;  
    common = local;  
}
```

} // all 3 statements must
// be executed together
// by single thread

Lock Example – Incorrect Version

```
public class DataRace extends Thread {  
    static int common = 0;  
    static Object o;           // all threads use o's lock  
    public void run() {  
        int local;  
        synchronized(o) {  
            local = common;  
        }  
        synchronized(o) {  
            local = local + 1;  
            common = local;  
        }  
    }  
}
```

} // transaction not atomic
// data race may occur
// even using locks

Issue 3) Avoiding Deadlock

■ Potential problem

- Threads holding lock may be unable to obtain lock held by other thread, and vice versa
- Thread holding lock may be waiting for action performed by other thread waiting for lock
- Program is unable to continue execution (**deadlock**)

Deadlock Example 1

Object **a** = new Object()

Object **b** = new Object()

```
Thread1() {  
    synchronized(a) {  
        synchronized(b) {  
            ...  
        }  
    }  
}
```

```
Thread2() {  
    synchronized(b) {  
        synchronized(a) {  
            ...  
        }  
    }  
}
```

// Thread1 holds lock for a, waits for b

// Thread2 holds lock for b, waits for a

Deadlock Example 2

```
void swap(Object a, Object b) {  
    Object local;  
    synchronized(a) {  
        synchronized(b) {  
            local = a; a = b; b = local;  
        }  
    }  
}
```

```
Thread1() { swap(a, b); } // holds lock for a, waits for b  
Thread2() { swap(b, a); } // holds lock for b, waits for a
```

Deadlock

■ Avoiding deadlock

- In general, avoid holding lock for a long time
- Especially avoid trying to hold two locks
 - May wait a long time trying to get 2nd lock

Thread-safe

- **Thread-safe** – Code is considered thread-safe if it works correctly when executed by multiple threads simultaneously.
- **Example: ArrayList is not thread-safe**

From Java API: “Note that this implementation is not synchronized. If multiple threads access an ArrayList instance concurrently, and at least one of the threads modifies the list structurally, it *must* be synchronized externally.”

Synchronization Summary

- Needed in multithreaded programs
- Can prevents data races
- Java objects support synchronization
- Many other tricky issues
 - To be discussed in future courses