

Reminder

- Homework 3 due Thursday
- Any questions about homework?

Onto Project 2...

- Hopefully everyone got all 50 points for public tests and passed the 8 release tests

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- Hopefully everyone got all 50 points for public tests and passed the 8 release tests
 - ...just kidding
- Project 2 builds on P1's Sys_Kill, and extends it to a more general concept of signal handling
- Due Friday, March 4, 11:59pm
 - **ELEVEN DAYS LEFT...**

Context Switching

- Does anyone know what context switching is good for?

Context Switching

- The OS can only run one process at a time
 - but this isn't very useful
- Simulate multiple processes with time-sharing
 - It gives each process a certain amount of time
 - When the time runs out, context-switch!

Context Switching

- Context switching does a few important things:
 - 1) Save the state of the current thread
 - 2) Load the state of the new thread
 - 3) Start the new thread
 - Look at `Switch_To_Thread()` in `geekos/lowlevel.asm`
- What is a state?
 - Kernel → registers, program counter, selectors
 - User → data (user variables, etc)

What To Do

- (1) Handle signals
 - What is a signal?
 - different than context switching
 - signals require jumping to an arbitrary function in user context
 - therefore you must save the old user context before going into the new user context

What To Do

- (2) Add a collection of syscalls to setup and call signals
- (3) Background processes are no longer orphans
- No p1 parts will be tested, but you will need background processes (and some version of your kill)
- P2 is overall not much code, but **VERY COMPLICATED** → **generally all or nothing**

Signals

- **SIGKILL**
 - Produces the same results as `Sys_Kill` in p1
 - Can't be handled by a process
- **SIGUSR1, SIGUSR2**
 - User-defined signals
- **SIGCHLD**
 - You trigger it when a child dies and its parent abandoned it
 - In p2, ALL processes have parents!

System Calls To Implement

- `Sys_WaitNoPID(int* status)`
 - Looks for a zombie and fills the exit status & reaps
- `Sys_Signal(ptr, sig_num)`
 - Handle one of the three signals allowed with a user-space function
 - `void handler(int sig_num)`
 - Store this pointer for lookup later

System Calls To Implement

- `Sys_RegDeliver(ptr)`
 - Register a trampoline given to you
 - a special function that calls `Sys_ReturnSignal` at the end of a handler function
- `Sys_Kill(pid, sig_num)`
 - Send a given signal to a PID. Asynchronous. It's not handled until kernel context switches to the pid.
- `Sys_ReturnSignal`
 - Signal handling is complete! Restore our old state stored on the user stack.

Example: shell

- Take shell.c
- You need to handle dead background processes
 - `Sys_Signal(some_cleaning_fn, SIGCHLD)`
- Your cleaning function will get called when there is a SIGCHLD
- But it can't do much on its own
 - Must call `Sys_WaitNoPID` to do the job

How it fits together

- When foo.exe starts:
 - Sys_RegDeliver(ptr) → store given "trampoline" function (not something you control)
- Sys_Signal(bar, SIGUSR1)
- Someone calls Sys_Kill(your pid, SIGUSR1)
 - Sets a flag so foo.exe knows it has a signal waiting
- Check_Pending_Signal (signal.c) is called to check for flags before context-switching
 - If there's a signal, call Setup_Frame
 - If not, just switch to the thread as normal

How it fits together

- Setup_Frame → save state & modify where in foo.exe we switch to
 - Look up the right handler (bar)
 - Push current kernel interrupt state onto user stack
 - We don't want to lose our registers
 - Push signal # for later
 - Push address of "trampoline" for later
 - Changes kernel stack so it has
 - Updated user stack ptr → free space
 - Saved program counter = signal handler (bar)
- Now, after context switch, bar gets called!

How it fits together

- Bar finishes. Say it computed the square root of three. (All this coding work for that lousy function?)
- Because of its position in the stack, the trampoline function gets called
 - Anyone know why?
- Trampoline function calls `Sys_ReturnSignal`
- `Sys_ReturnSignal` calls `Complete_Handler` (you write)
 - Restores the kernel stack to where it was before `Setup_Frame`
- Now the user process can be context switched to without knowing what hit it!

Final Words

- More details in the spec
 - Reading is **MANDATORY**
- Three rules to CS412
 - Start early
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 - Also, ask questions, which is easier if you start early
- Any questions?