

CMSC 132:

OBJECT-ORIENTED PROGRAMMING II

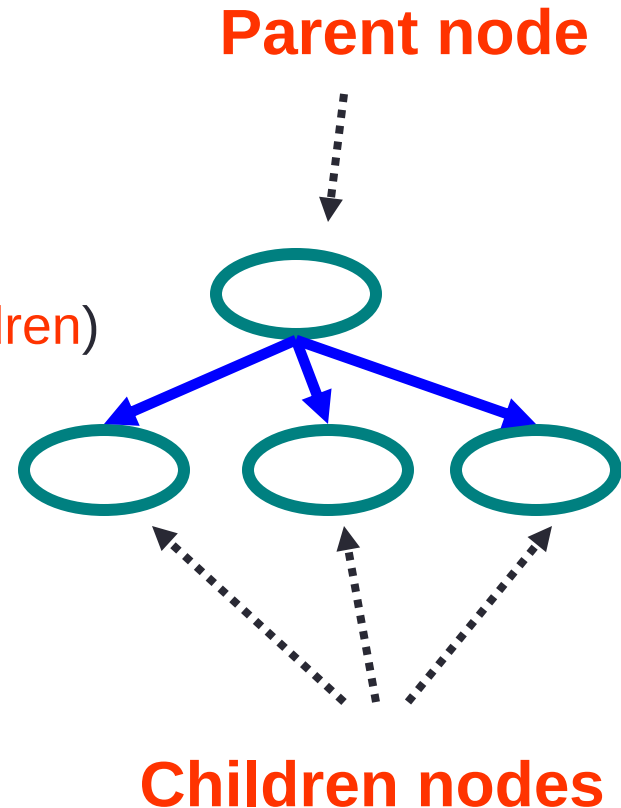


Trees & Binary Search Trees

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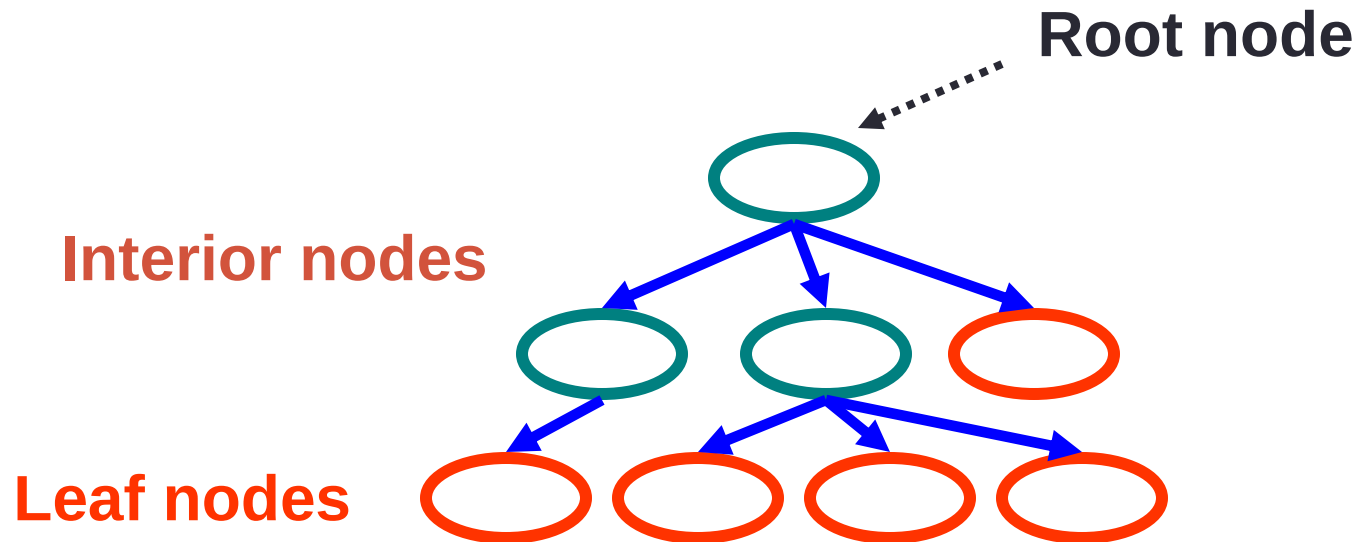
Trees

- Trees are hierarchical data structures
 - One-to-many relationship between elements
- Tree node / element
 - Contains data
 - Referred to by only 1 (**parent**) node
 - Contains links to any number of (**children**) nodes



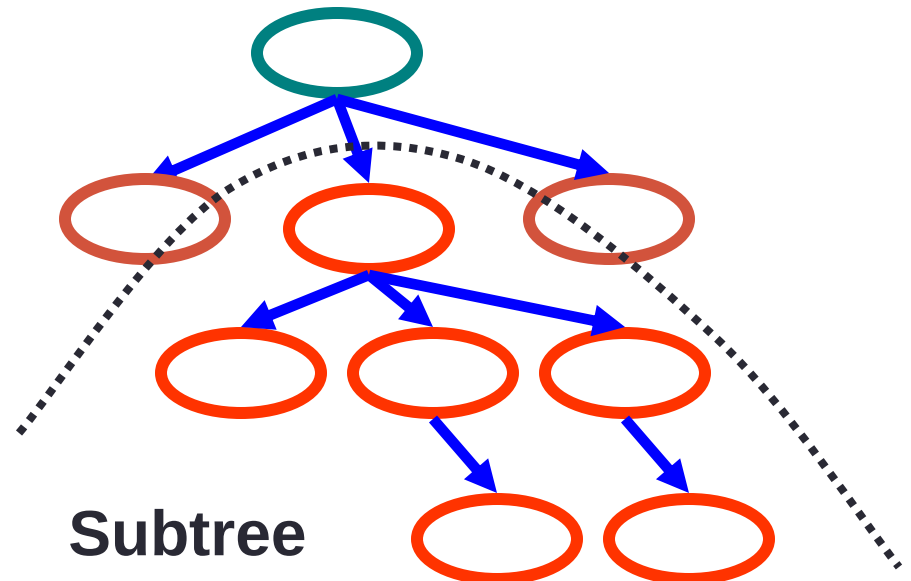
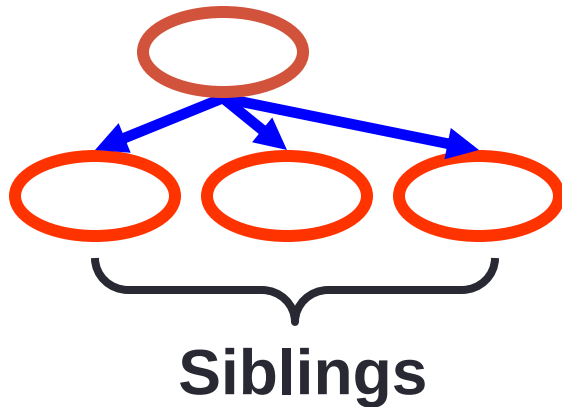
Trees

- Terminology
 - Root $\Rightarrow$ node with no parent
 - Leaf $\Rightarrow$ all nodes with no children
 - Interior $\Rightarrow$ all nodes with children



Trees

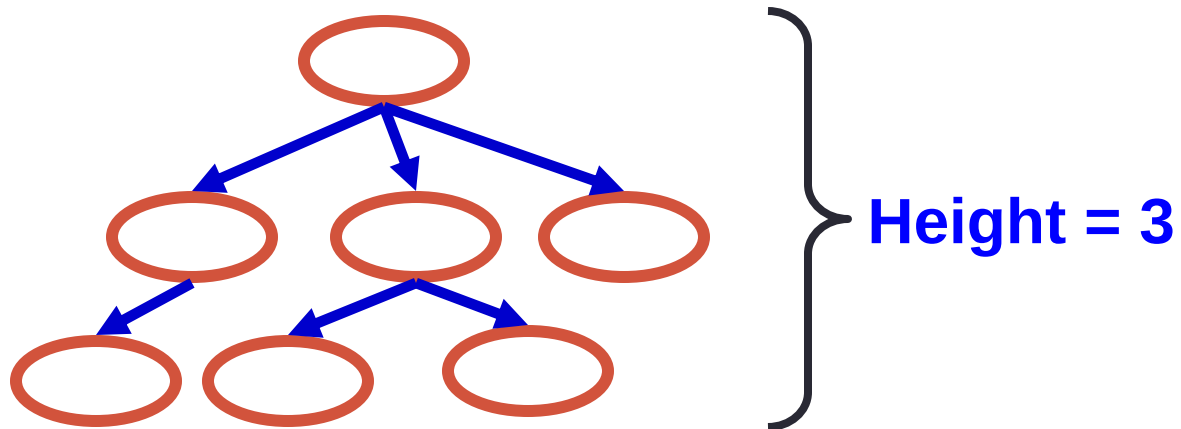
- Terminology
 - Sibling $\Rightarrow$ node with same parent
 - Descendent $\Rightarrow$ children nodes & their descendants
 - Subtree $\Rightarrow$ portion of tree that is a tree by itself
 $\Rightarrow$ a node and its descendants



Trees

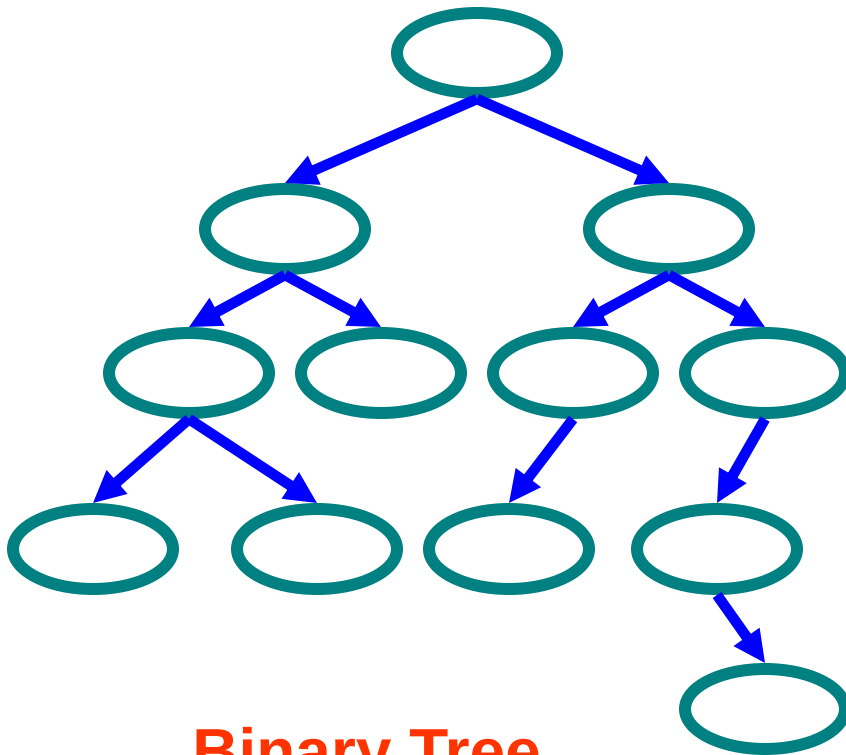
- Terminology

- Level $\Rightarrow$ is a measure of a node's distance from root
- Definition of level
 - If node is the root of the tree, its level is 1
 - Else, the node's level is 1 + its parent's level
- Height (depth) $\Rightarrow$ max level of any node in tree

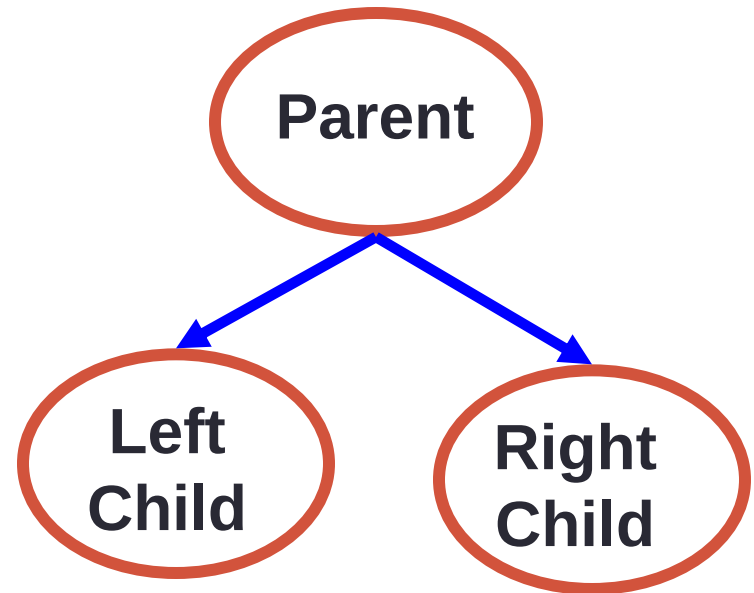


Binary Trees

- Binary tree
 - Tree with 0–2 children per node
 - Left & right child / subtree



Binary Tree



Tree Traversal

- Often we want to
 - Find all nodes in tree
 - Determine their relationship
- Can do this by
 - Walking through the tree in a prescribed order
 - Visiting the nodes as they are encountered
- Process is called **tree traversal**

Tree Traversal

- Goal
 - Visit every node in binary tree
- Approaches
 - **Depth first**
 - Preorder $\Rightarrow$ parent before children
 - Inorder $\Rightarrow$ left child, parent, right child
 - Postorder $\Rightarrow$ children before parent
 - **Breadth first** $\Rightarrow$ closer nodes first

Tree Traversal Methods

- Pre-order

1. Visit **node** // first
2. Recursively visit left subtree
3. Recursively visit right subtree

- In-order

1. Recursively visit left subtree
2. Visit **node** // second
3. Recursively right subtree

- Post-order

1. Recursively visit left subtree
2. Recursively visit right subtree
3. Visit **node** // last

Tree Traversal Methods

- Breadth-first

BFS(Node n) {

Queue Q = new Queue();

Q.enqueue(n); // insert node into Q

while (!Q.empty()) {

 n = Q.dequeue(); // remove next node

 if (!n.isEmpty()) {

 visit(n); // visit node

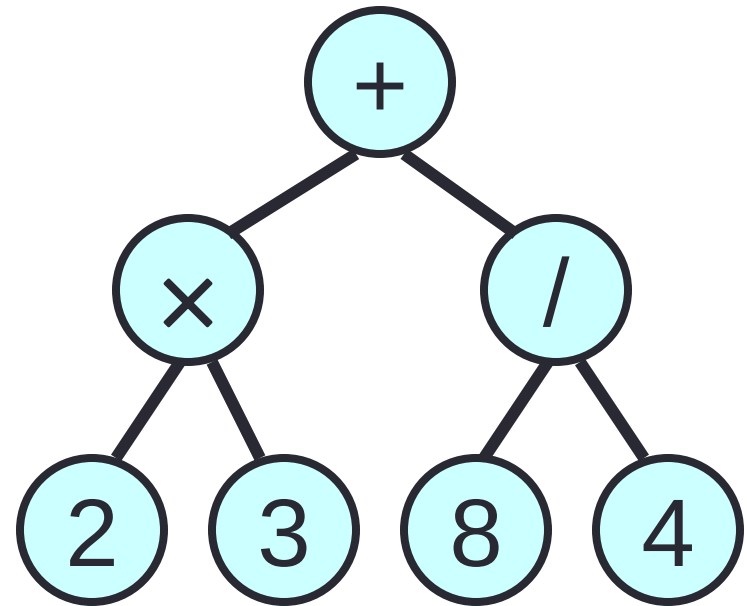
 Q.enqueue(n.Left()); // insert left subtree in Q

 Q.enqueue(n.Right()); // insert right subtree in Q

}}

Tree Traversal Examples

- Pre-order (prefix)
 - $+ \times 2 3 / 8 4$
- In-order (infix)
 - $2 \times 3 + 8 / 4$
- Post-order (postfix)
 - $2 3 \times 8 4 / +$
- Breadth-first
 - $+ \times / 2 3 8 4$



Expression tree

Binary Tree Implementation

- Using a class to represent a Node

```
Class Node {  
    KeyType key;  
    Node left, right;    // null if empty  
}
```

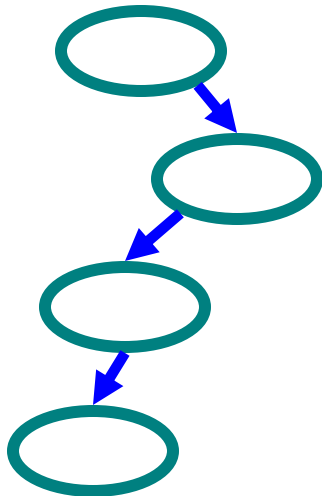
Node root = null; // Empty Tree

- Using a Polymorphic Binary Tree
 - We will talk about this implementation later on

Types of Binary Trees

- **Degenerate**

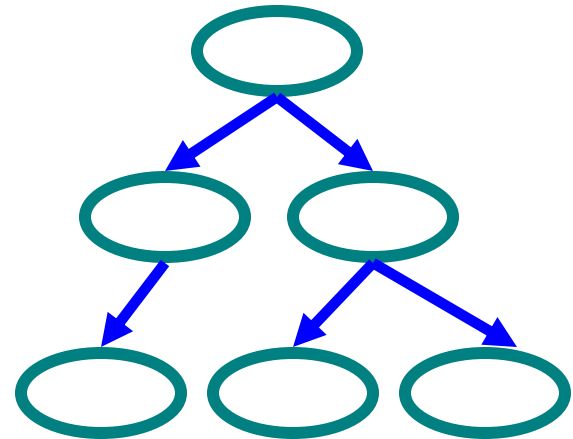
- Mostly 1 child / node
- Height = $O(n)$
- Similar to linear list



**Degenerate
binary tree**

- **Balanced**

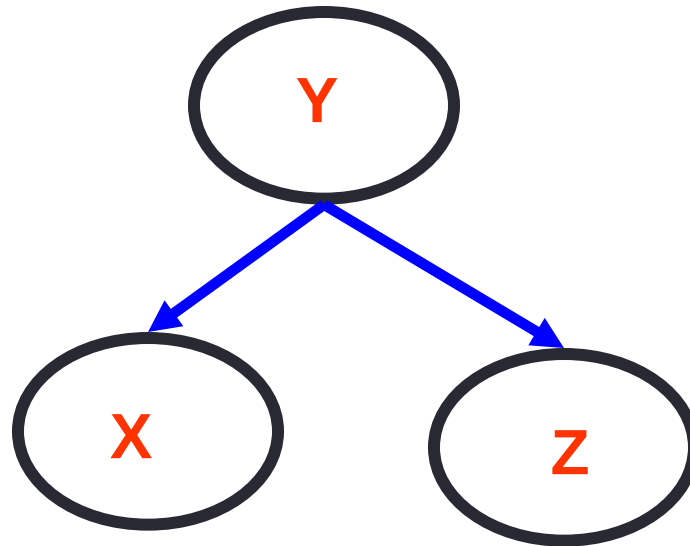
- Mostly 2 child / node
- Height = $O(\log(n))$
- $2^{\text{Height}} - 1 = n$ (# of nodes)
- Useful for searches



**Balanced
binary tree**

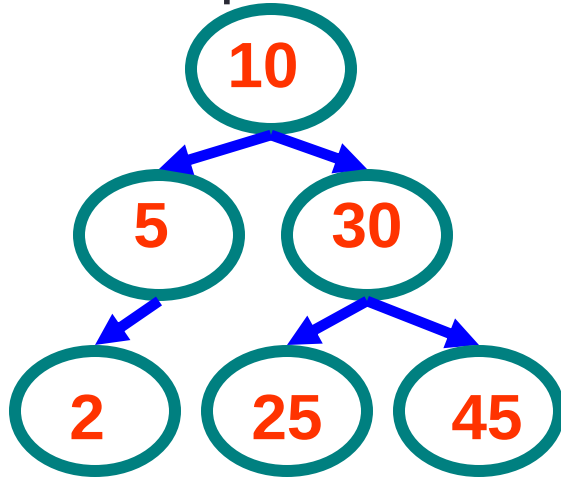
Binary Search Trees

- Key property
 - Value at node
 - Smaller values in left subtree
 - Larger values in right subtree
- Example
 - $Y > X$
 - $Y < Z$

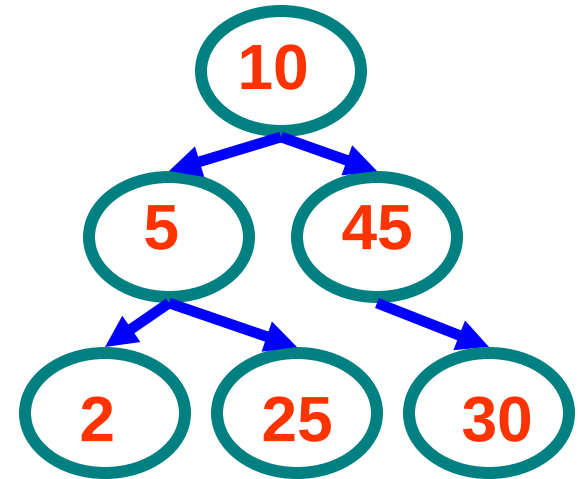
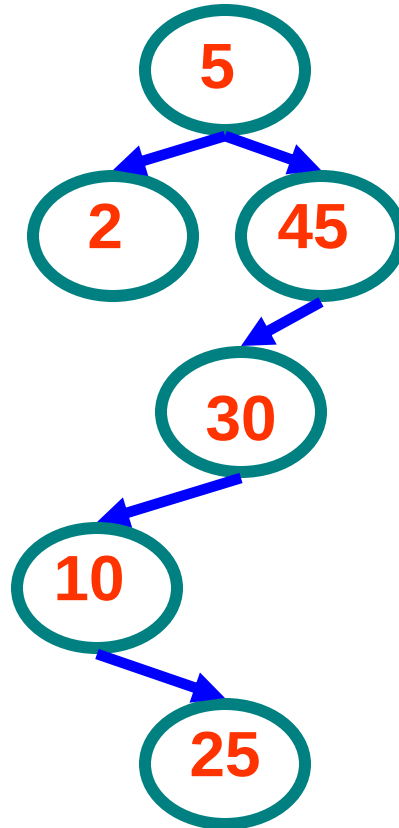


Binary Search Trees

- Examples



**Binary
search trees**



**Non-binary
search tree**

Tree Traversal Examples

- Pre-order

- 44, 17, 32, 78, 50, 48, 62, 88

Sorted order!

- In-order

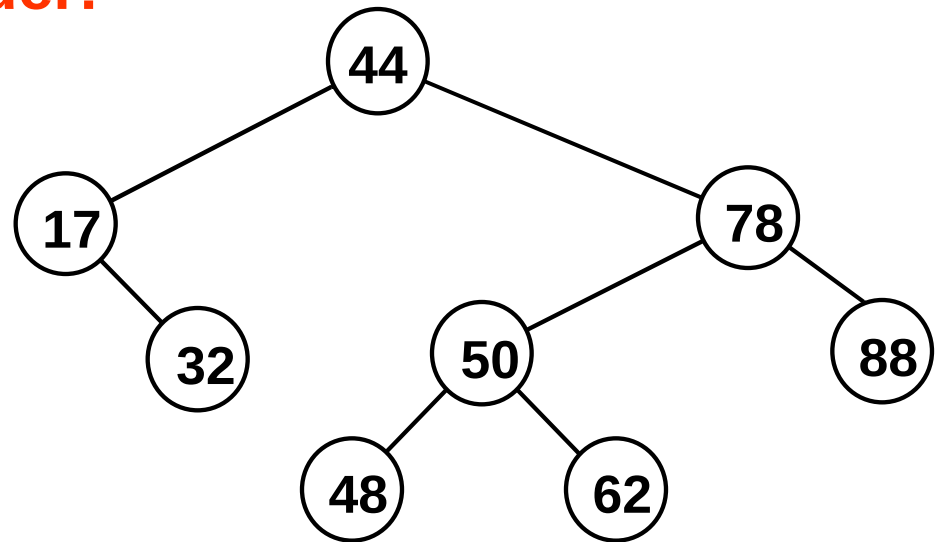
- 17, 32, 44, 48, 50, 62, 78, 88

Post-order

32, 17, 48, 62, 50, 88, 78, 44

Breadth-first

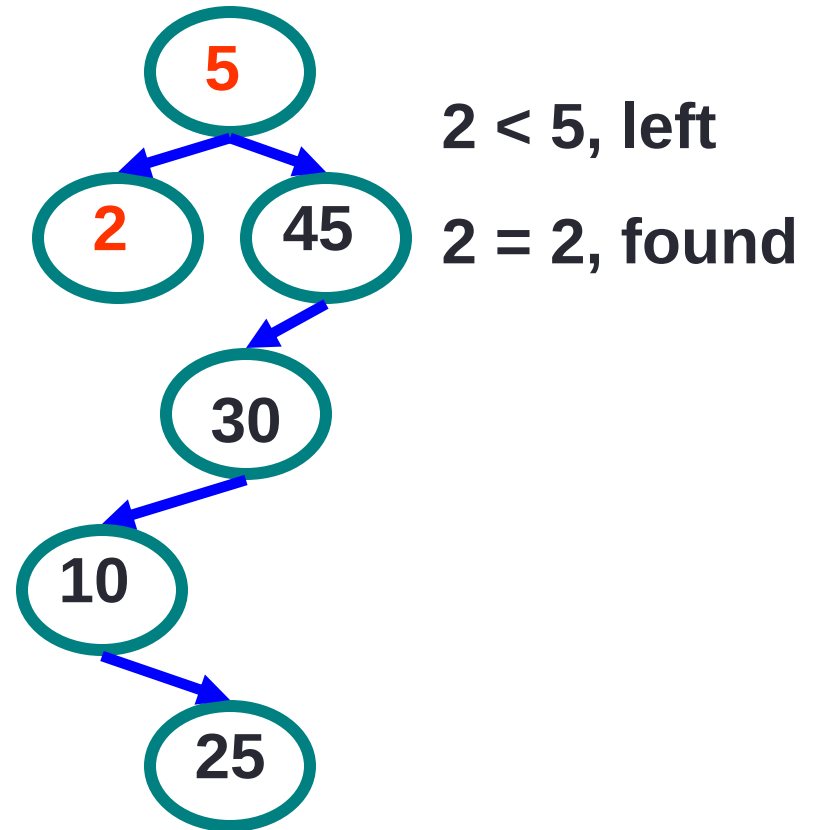
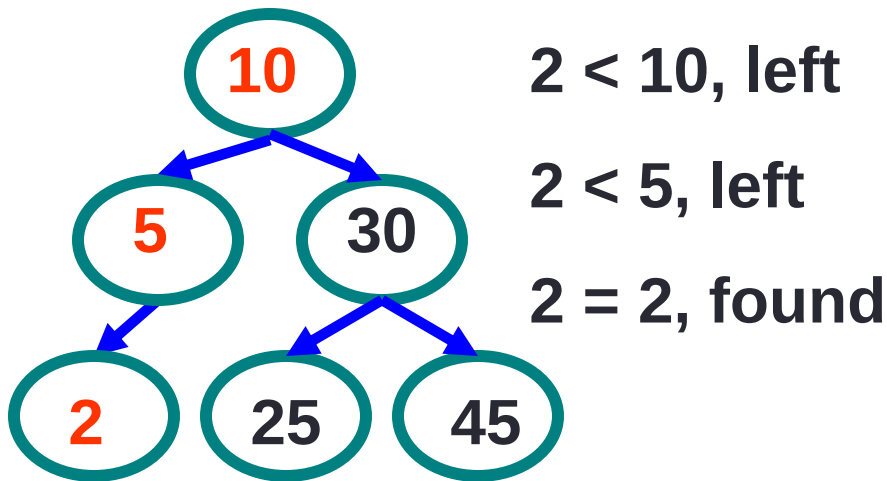
44, 17, 78, 32, 50, 88, 48, 62



Binary search tree

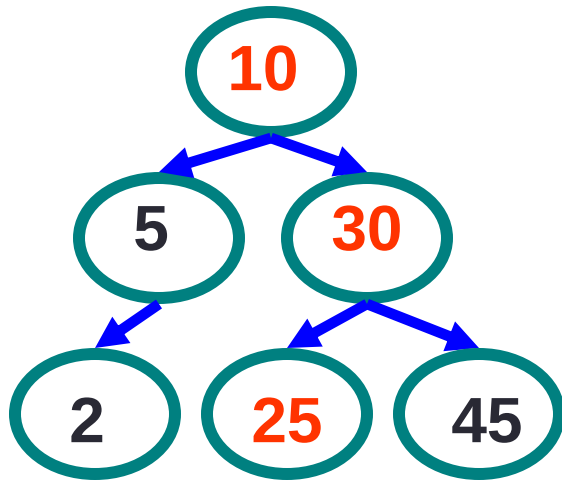
Example Binary Searches

- Find (2)



Example Binary Searches

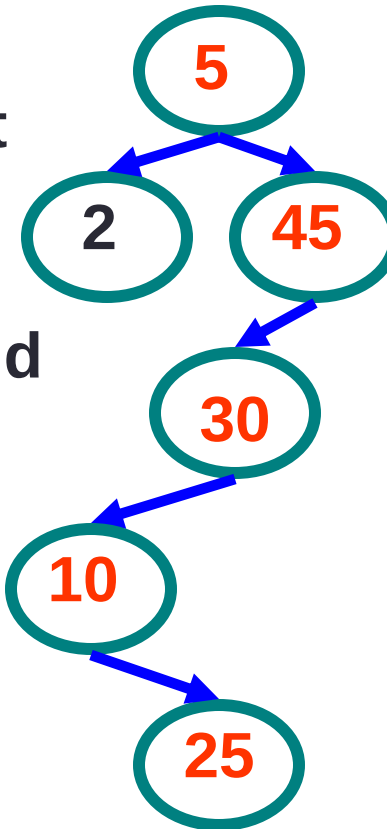
- Find (25)



$25 > 10$, right

$25 < 30$, left

$25 = 25$, found



$25 > 5$, right

$25 < 45$, left

$25 < 30$, left

$25 > 10$, right

$25 = 25$, found

Binary Search Properties

- Time of search
 - Proportional to height of tree
 - Balanced binary tree
 - $O(\log(n))$ time
 - Degenerate tree
 - $O(n)$ time
 - Like searching linked list / unsorted array
- Requires
 - Ability to **compare** key values

Binary Search Tree Construction

- How to build & maintain binary trees?
 - Insertion
 - Deletion
- Maintain key property (invariant)
 - Smaller values in left subtree
 - Larger values in right subtree

Binary Search Tree – Insertion

- Algorithm

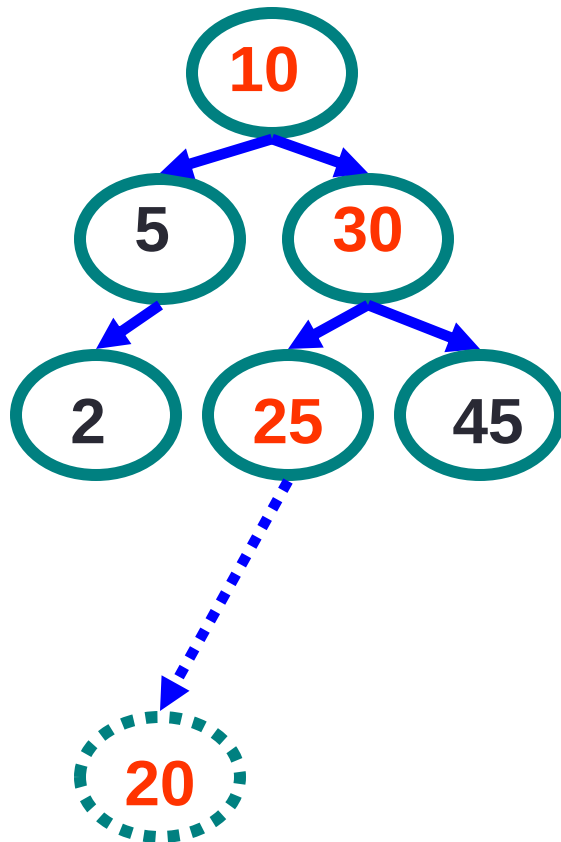
1. Perform search for value X
2. Search will end at node Y (if X not in tree)
3. If $X < Y$, insert new leaf X as new left subtree for Y
4. If $X > Y$, insert new leaf X as new right subtree for Y

- Observations

- $O(\log(n))$ operation for balanced tree
- Insertions may unbalance tree

Example Insertion

- Insert (20)



$20 > 10$, right

$20 < 30$, left

$20 < 25$, left

Insert 20 on left

Binary Search Tree – Deletion

- Algorithm

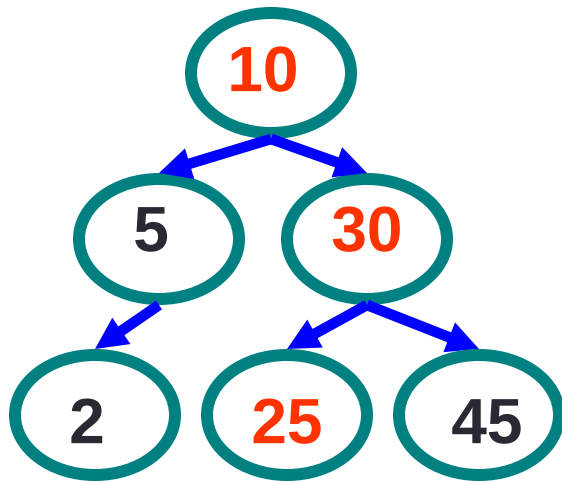
1. Perform search for value X
2. If X is a leaf, delete X
3. Else // must delete internal node
 - a) Replace with largest value Y on left subtree
OR smallest value Z on right subtree
 - b) Delete replacement value (Y or Z) from subtree

- Observation

- $O(\log(n))$ operation for balanced tree
- Deletions may unbalance tree

Example Deletion (Leaf)

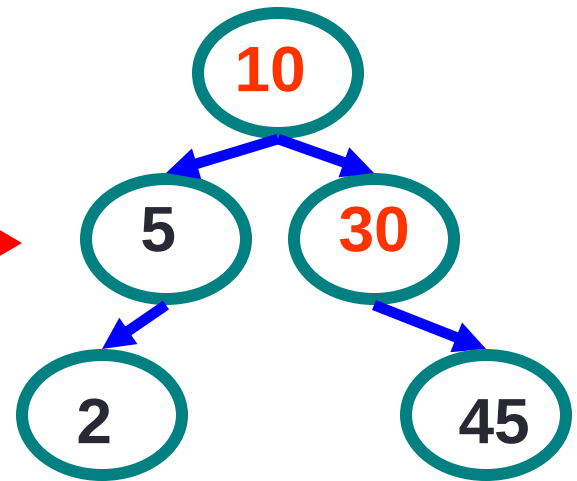
- Delete (25)



$25 > 10$, right

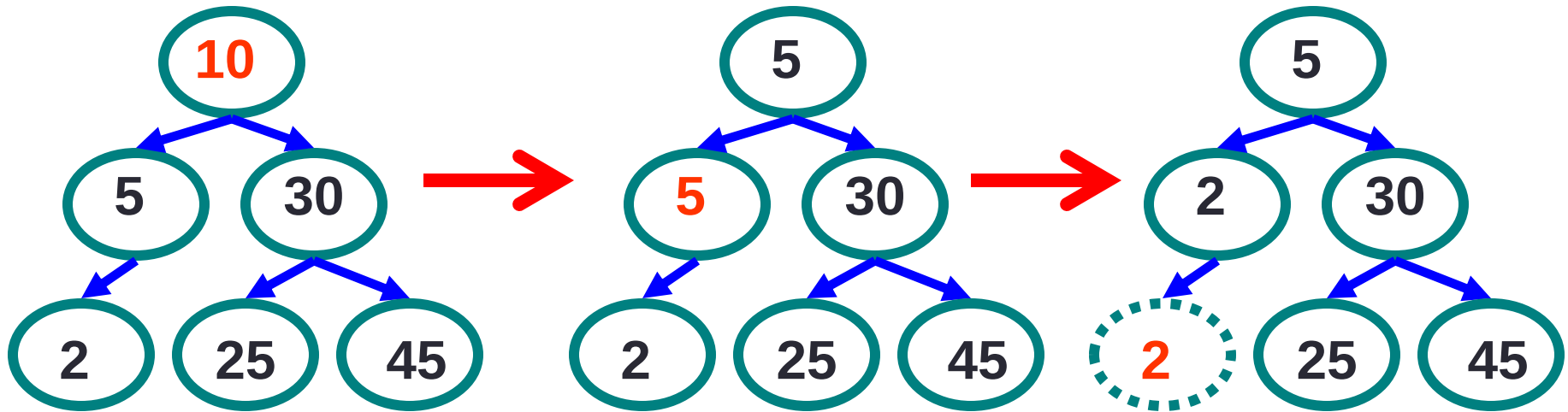
$25 < 30$, left

$25 = 25$, delete



Example Deletion (Internal Node)

- Delete (10)



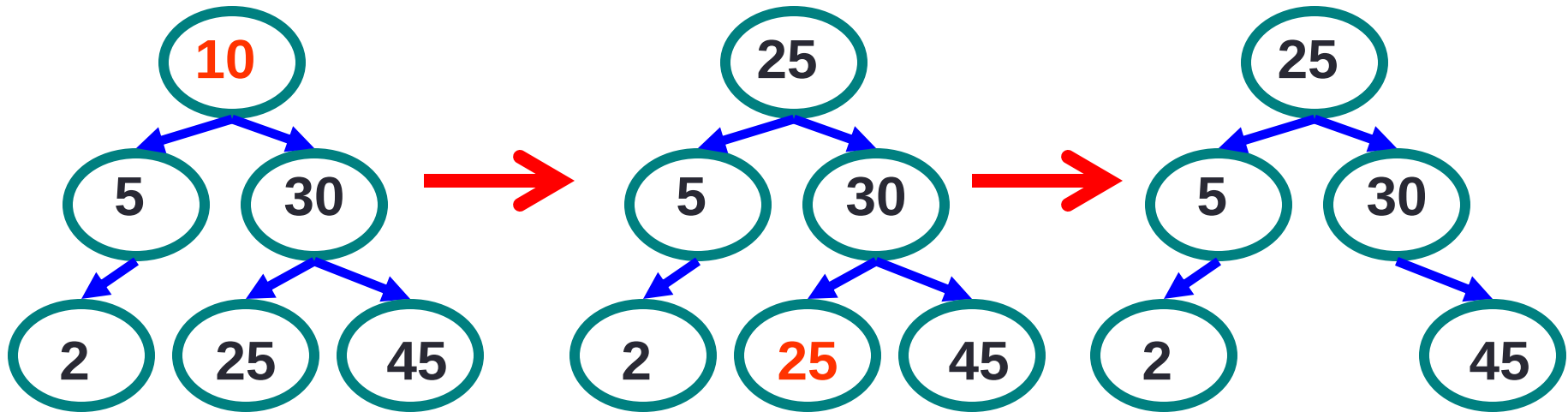
Replacing 10
with **largest**
value in left
subtree

Replacing 5
with **largest**
value in left
subtree

Deleting leaf

Example Deletion (Internal Node)

- Delete (10)



Replacing 10
with **smallest**
value in right
subtree

Deleting leaf

Resulting tree

Building Maps w/ Search Trees

- Binary Search trees often used to implement **maps**
 - Each non-empty node contains
 - Key
 - Value
 - Left and right child
- Need to be able to compare keys
 - Generic type `<K extends Comparable<K>>`
 - Denotes any type K that can be compared to K's

BST (Binary Search Tree) Implementation

- Implementing Tree using traditional approach
- Based on the BST definition below let's see how to implement typical BST Operations ([constructor](#), [add](#), [print](#), [find](#), [isEmpty](#), [isFull](#), [size](#), [height](#), etc.)

```
public class BinarySearchTree <K extends Comparable<K>, V> {  
    private class Node {  
        private K key;  
        private V data;  
        private Node left, right;  
        public Node(K key, V data) {  
            this.key = key;  
            this.data = data;  
        }  
    }  
    private Node root;  
}
```

- See code distribution [BinaryTreeCode.zip](#)

BST Testing

- How can we test the correctness of BST Methods?
- What is the best approach?