

# CMSC 132:

# OBJECT-ORIENTED PROGRAMMING II



## Graph Implementation

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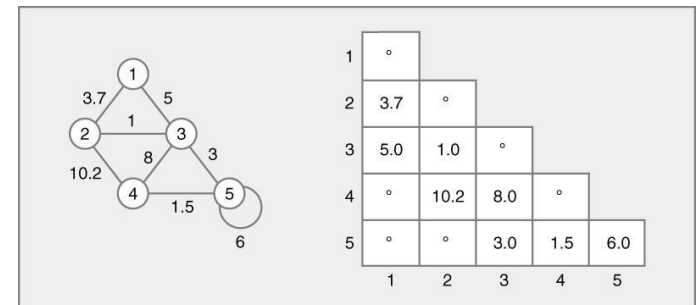
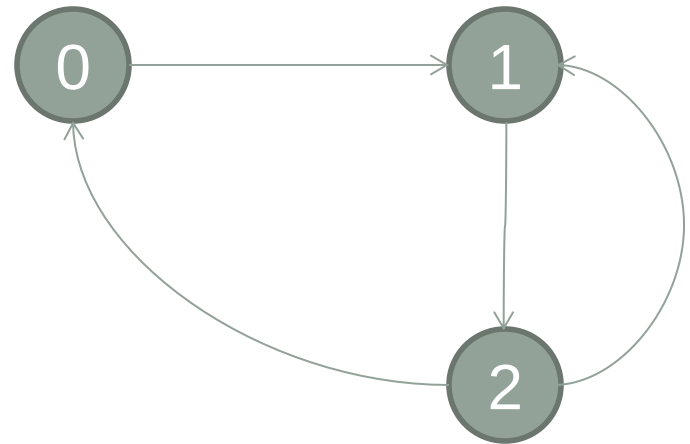
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# Graph Implementation

- How do we represent nodes/edges?
  - Adjacency matrix
    - 2D array of neighbors
  - Adjacency list/set/map
    - List/set/map of neighbors
- Important for very large graphs
  - Affects efficiency / storage

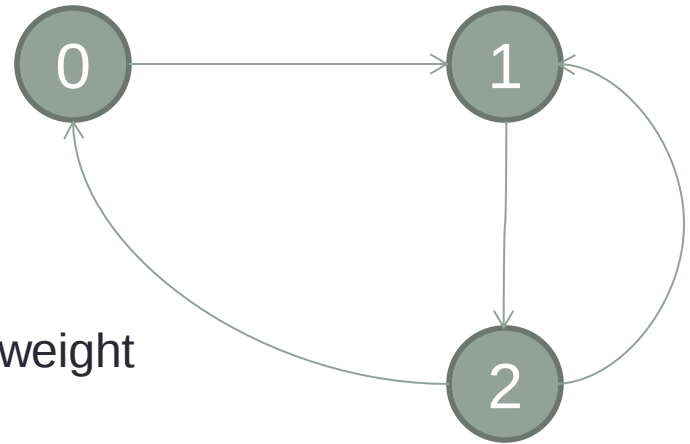
# Adjacency Matrix

- Single **two-dimensional array** for entire graph
- **Directed Graph**
  - **Unweighted graph**
    - Matrix elements  $\Rightarrow$  boolean
  - **Weighted graph**
    - Matrix elements  $\Rightarrow$  values
  - Let's see an example of each
- **Undirected Graph**
  - Let's see an example for unweighted graph
  - Let's see an example for weighted graph
- **For Undirected graph**
  - Only upper / lower triangle matrix needed
  - Since  $n_j, n_k$  implies  $n_k, n_j$



# Adjacency List/Set/Map

- For each node, **store neighbor information in a list, set, or map**
- The main structure can be a list, set, or map
- **Directed Graph**
  - **Unweighted graph**
    - List or set of neighbors
  - **Weighted graph**
    - Each entry keeps track of neighbor and weight
    - Easy to implement with maps
      - Maps of Maps (using HashMaps for efficiency)
  - Let's see an example of each
- **Undirected Graph**
  - Let's see an example for unweighted graph
  - Let's see an example for weighted graph



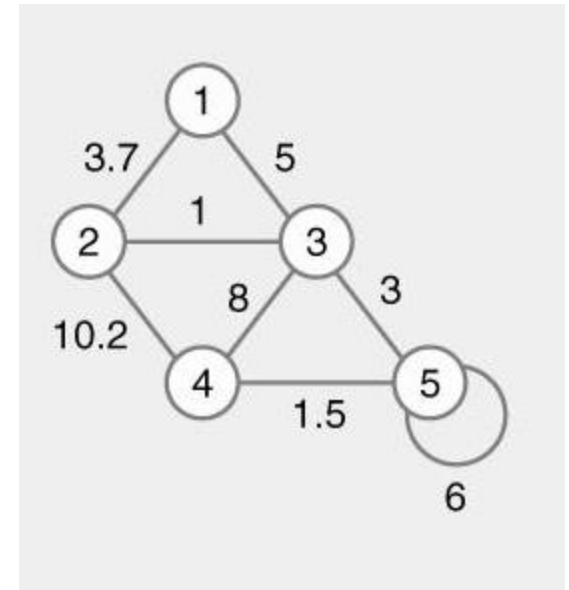
# Additional Examples

- Examples
  - **Unweighted graph**

node 1: {2, 3}  
node 2: {1, 3, 4}  
node 3: {1, 2, 4, 5}  
node 4: {2, 3, 5}  
node 5: {3, 4, 5}

- **Weighted graph**

node 1: {2=3.7, 3=5}  
node 2: {1=3.7, 3=1, 4=10.2}  
node 3: {1=5, 2=1, 4=8, 5=3}  
node 4: {2=10.2, 3=8, 5=1.5}  
node 5: {3=3, 4=1.5, 5=6}



# Graph Properties

- **Graph density**
  - Ratio edges to nodes (dense vs. sparse)
  - For adjacency matrix many empty entries for large, sparse graph
- **Adjacency matrix**
  - Can find individual edge (a,b) quickly
  - Examine entry in array `Edge[a,b]`
    - Constant time operation
- **Adjacency list / set / map**
  - Can find all edges for node (a) quickly
  - Iterate through collection of edges for a
    - On average  $E / N$  edges per node

# Complexity

- Average Complexity of operations
  - For graph with  $N$  nodes,  $E$  edges

| Operation                | Adj Matrix | Adj List | Adj Set/Map |
|--------------------------|------------|----------|-------------|
| Find edge                | $O(1)$     | $O(E/N)$ | $O(1)$      |
| Insert edge              | $O(1)$     | $O(E/N)$ | $O(1)$      |
| Delete edge              | $O(1)$     | $O(E/N)$ | $O(1)$      |
| Enumerate edges for node | $O(N)$     | $O(E/N)$ | $O(E/N)$    |

# Choosing Graph Implementations

- **Factors to consider**

- **Graph density**
- **Graph algorithm**

- Neighbor based

For each node X in graph

For each neighbor Y of X // adj list faster if sparse

doWork()

- Connection based

For each node X in ...

For each node Y in ...

if (X,Y) is an edge // adj matrix faster if dense

doWork()