

CMSC430 Spring 2014 Quiz 2

Name _____

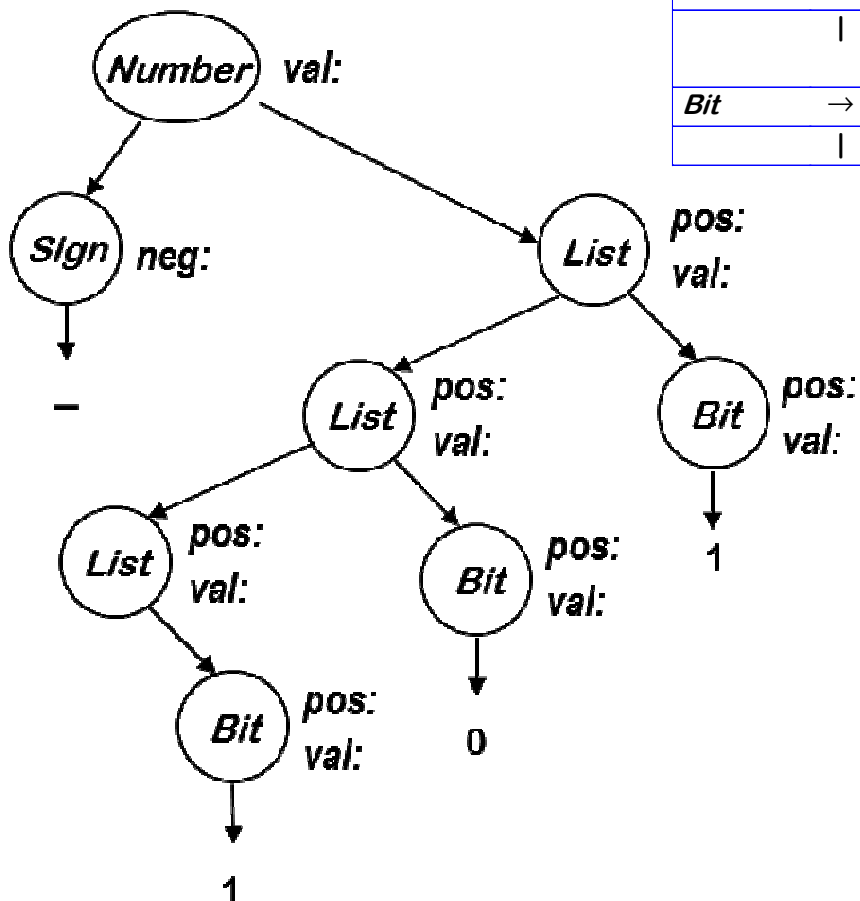
Instructions

- This quiz is worth 30 points. You have 20 minutes for to take this quiz.
- This is a closed book exam. No notes or other aids are allowed.
- Write neatly. Credit cannot be given for illegible answers.

1. (6 pts) Attribute grammars

Consider the following attribute grammar for calculating values of signed binary numbers. Show how the attribute grammar can calculate the value for “-101” by annotating its parse tree, showing how each attribute is calculated.

<i>Productions</i>	<i>Attribution Rules</i>
$Number \rightarrow Sign\ List$	$List.pos \leftarrow 0$ If $Sign.neg$ then $Number.val \leftarrow -List.val$ else $Number.val \leftarrow List.val$
$Sign \rightarrow \pm$	$Sign.neg \leftarrow false$
$Sign \rightarrow =$	$Sign.neg \leftarrow true$
$List_0 \rightarrow List_1\ Bit$	$List_1.pos \leftarrow List_0.pos + 1$ $Bit.pos \leftarrow List_0.pos$ $List_0.val \leftarrow List_1.val + Bit.val$
$List \rightarrow Bit$	$Bit.pos \leftarrow List.pos$ $List.val \leftarrow Bit.val$
$Bit \rightarrow 0$	$Bit.val \leftarrow 0$
$Bit \rightarrow 1$	$Bit.val \leftarrow 2^{Bit.pos}$



2. (8 pts) Syntax directed translation & type checking

Consider the following grammar fragment for an OCaml-style if-then-else expression:

$$exp \rightarrow IF\ exp\ THEN\ exp\ ELSE\ exp$$

Assume that the nonterminal *exp* has an attribute *type* that can be assigned a range of values, including *typeBoolean*, *typeInt*, *typeError*, etc. Add syntax-directed type checking rules to the grammar fragment to enforce the following:

- a. (2 pts) The guard of the IF expression must be *typeBoolean*. Otherwise call `parser.msg("Boolean required")`.
- b. (2 pts) The THEN and ELSE expressions must have the same type. Otherwise call `parser.msg("Else type mismatch")`.
- c. (4 pts) The type of the IF expression is assigned the type of the THEN expression if both (a) and (b) are true, else it is assigned *typeError*.

Answer:

$$exp_1 \rightarrow IF\ exp_2\ THEN\ exp_3\ ELSE\ exp_4$$

{:

:}

3. (6 pts) Symbol tables

Show the state of lexically nested symbol tables when the compiler reaches the point marked HERE in the following code fragment:

```
int a;  
void foo( ) { int b; { int c;} }  
void bar( ) { int c, d; /* HERE */ }
```

4. (10 pts) Intermediate Representations

Translate the following arithmetic expression into:

$$x = (a - b) - c$$

3-addr Instruction	Effect
load R1 x	$R1 \leftarrow x$
store x R1	$x \leftarrow R1$
add R1 R2 R3	$R1 \leftarrow R2 + R3$
sub R1 R2 R3	$R1 \leftarrow R2 - R3$
mult R1 R2 R3	$R1 \leftarrow R2 * R3$

Java Stack Code	Effect
nop	none
ldc_int c	push constant c onto stack
iload index(x)	push local variable X onto stack
istore index(x)	pop stack, store in local variable X
iadd	pop 2 elems off stack, add, push
isub	pop 2 elems, subtract top from 2nd, push
imult	pop 2 elems off stack, multiply, push

(2 pts) an abstract syntax tree (AST)	(4 pts) 3-address code	(4 pts) Java stack code