Internet Naming: DNS & DHCP

Slides from

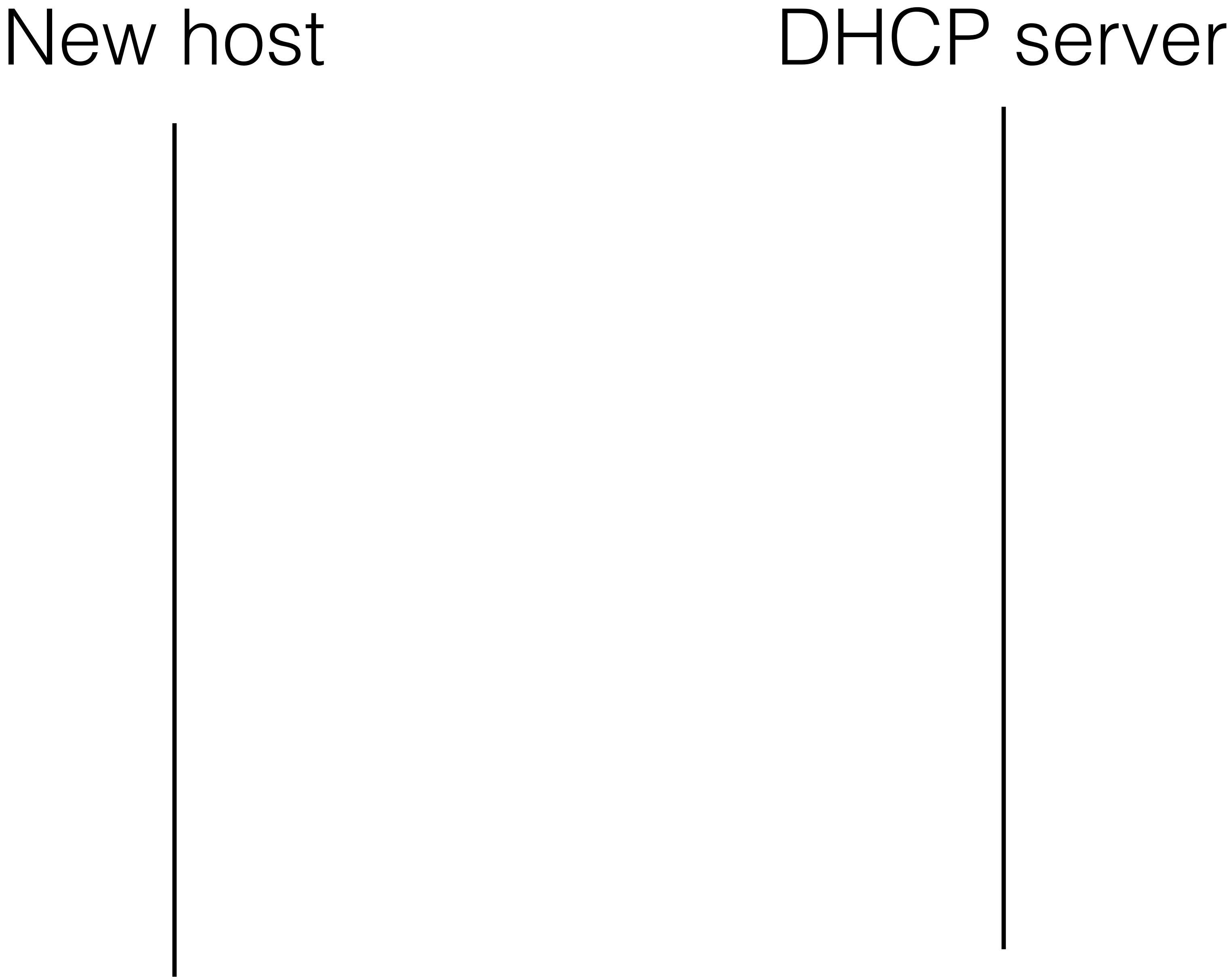
- Dave Levin 414-spring2016

Naming

- IP addresses allow global connectivity
- But they’re pretty useless for humans!
  - Can’t be expected to pick their own IP address
  - Can’t be expected to remember another’s IP address
- **DHCP** : Setting IP addresses
- **DNS** : Mapping a memorable name to a routable IP address

DHCP

Dynamic Host Configuration Protocol



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DHCP server

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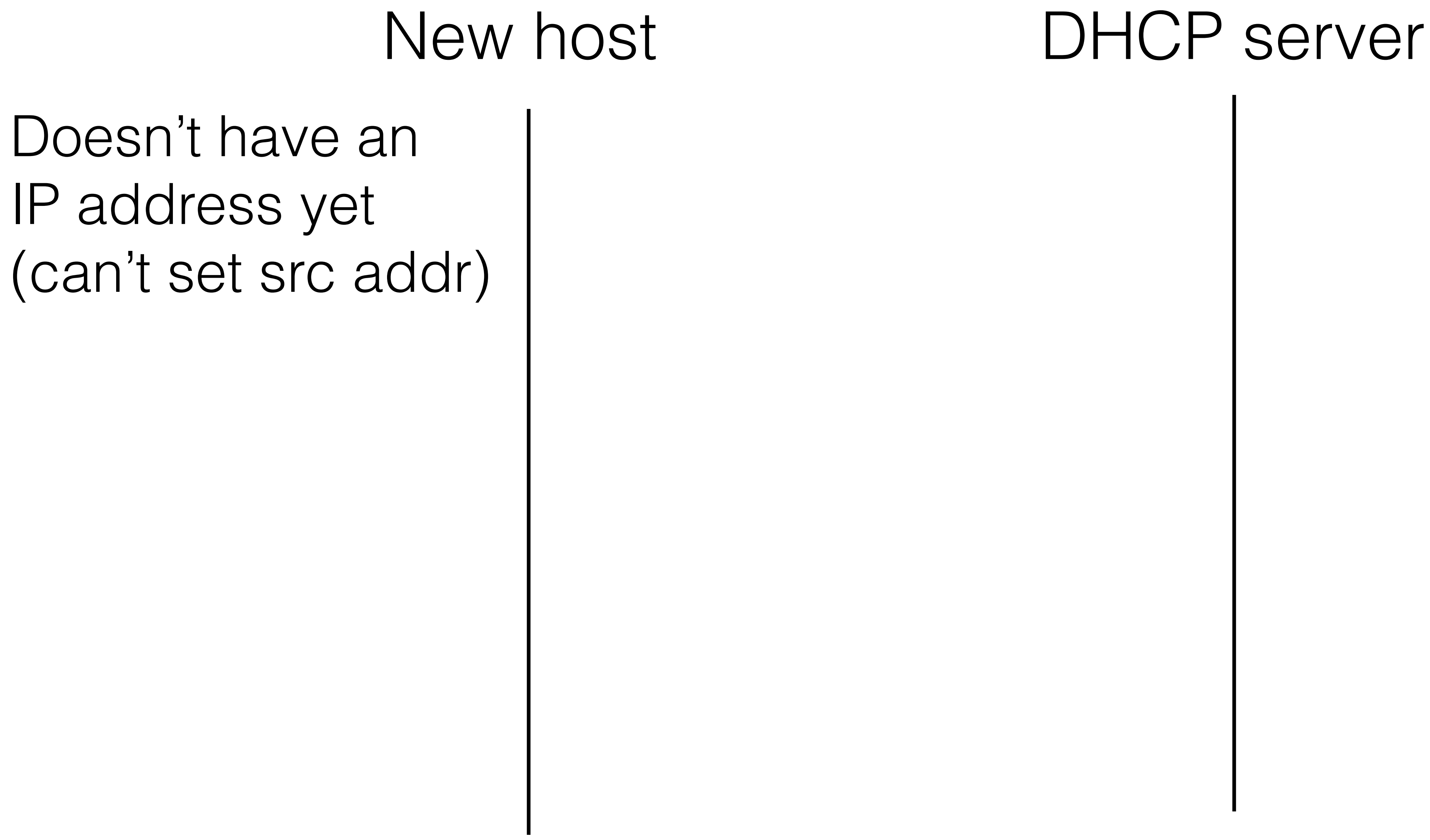

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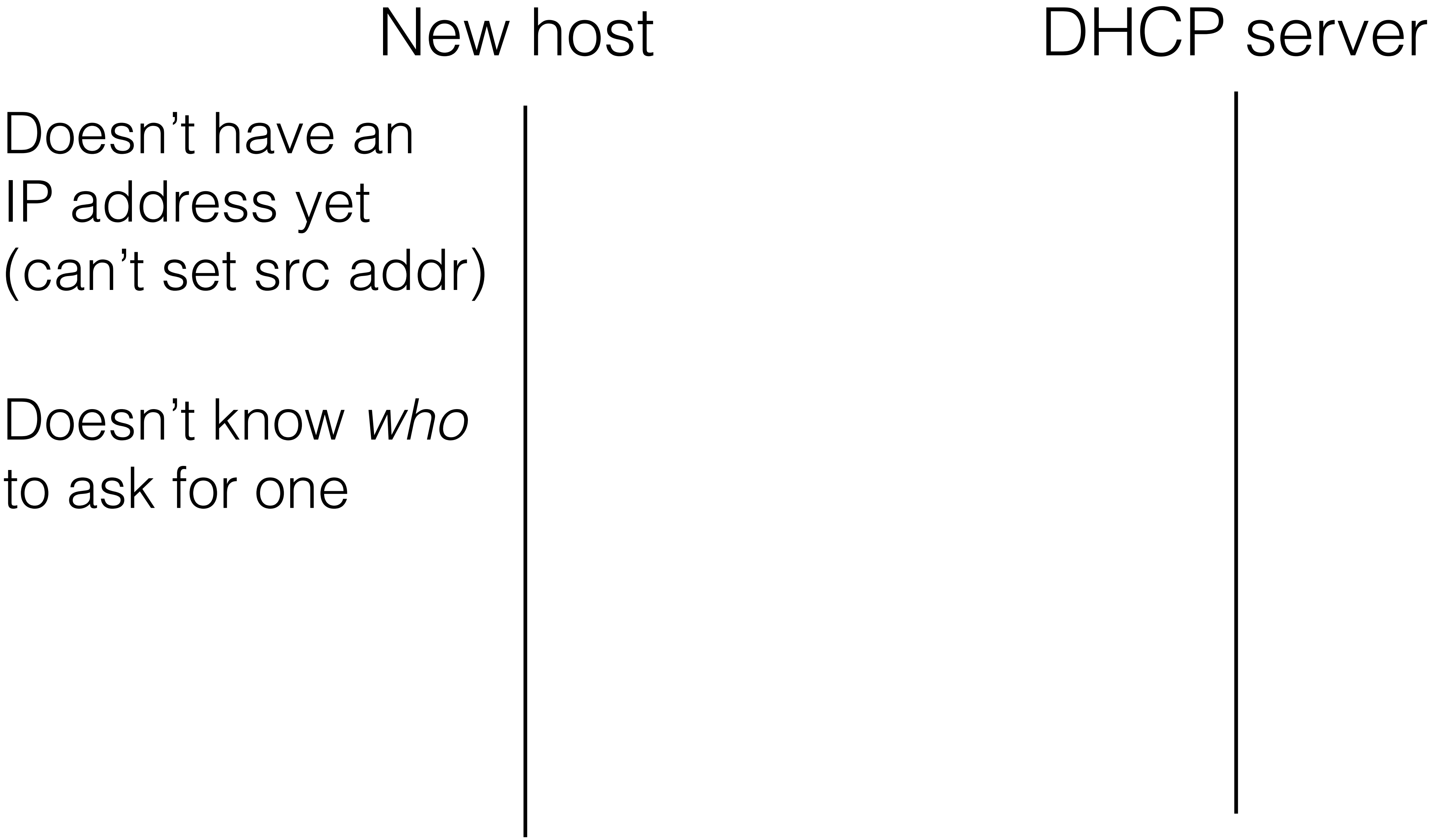

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Solution: Discover one on the local subnet

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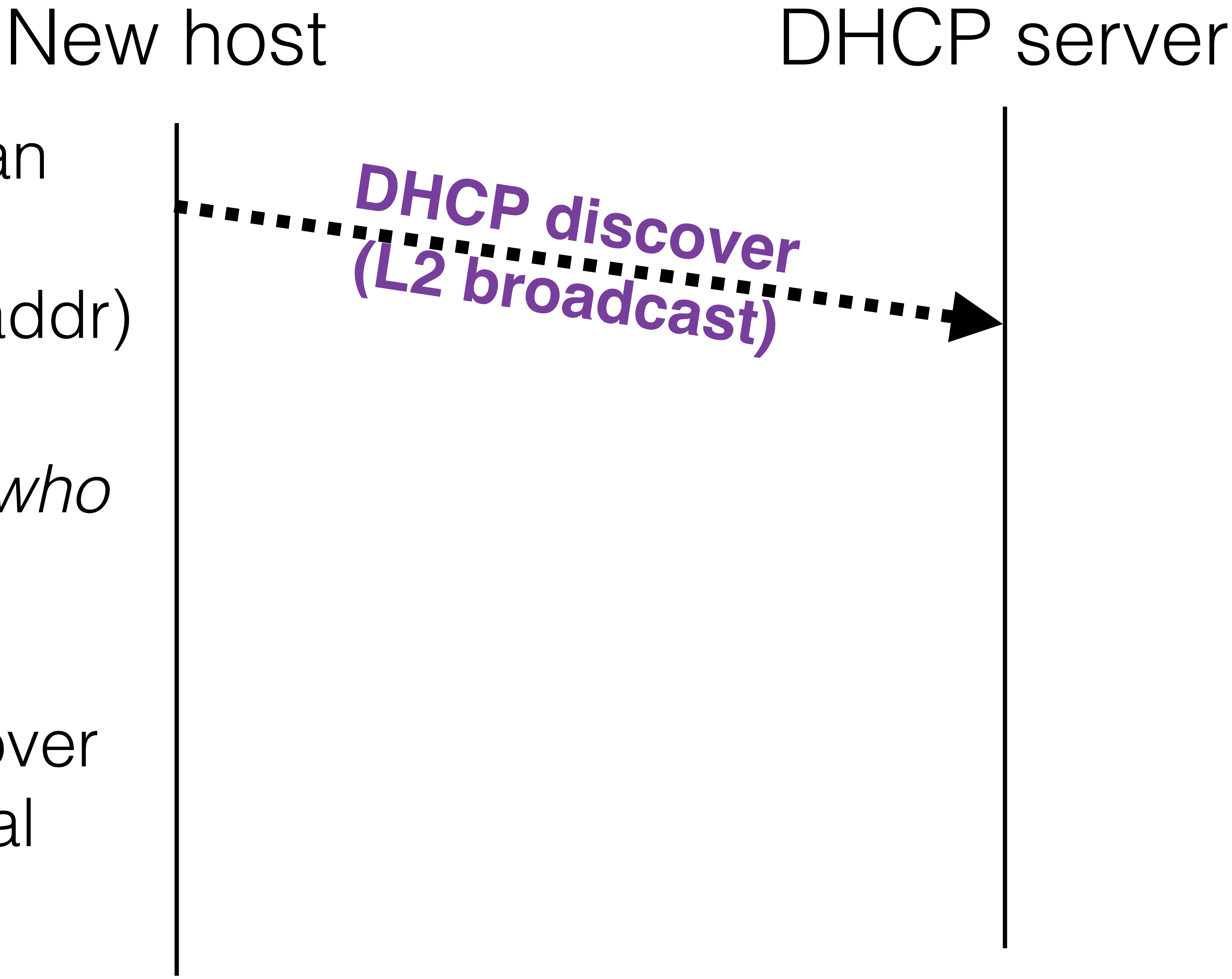

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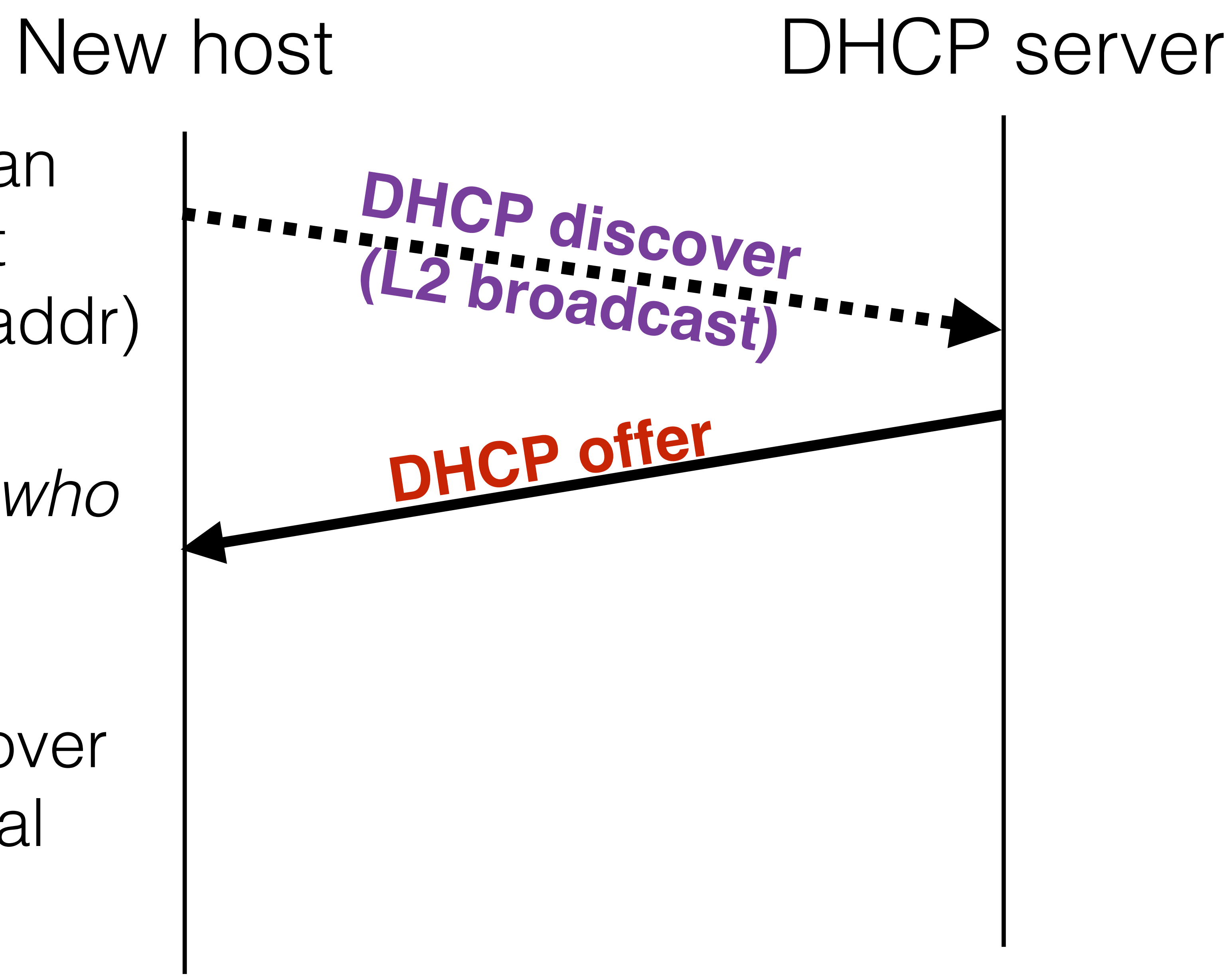

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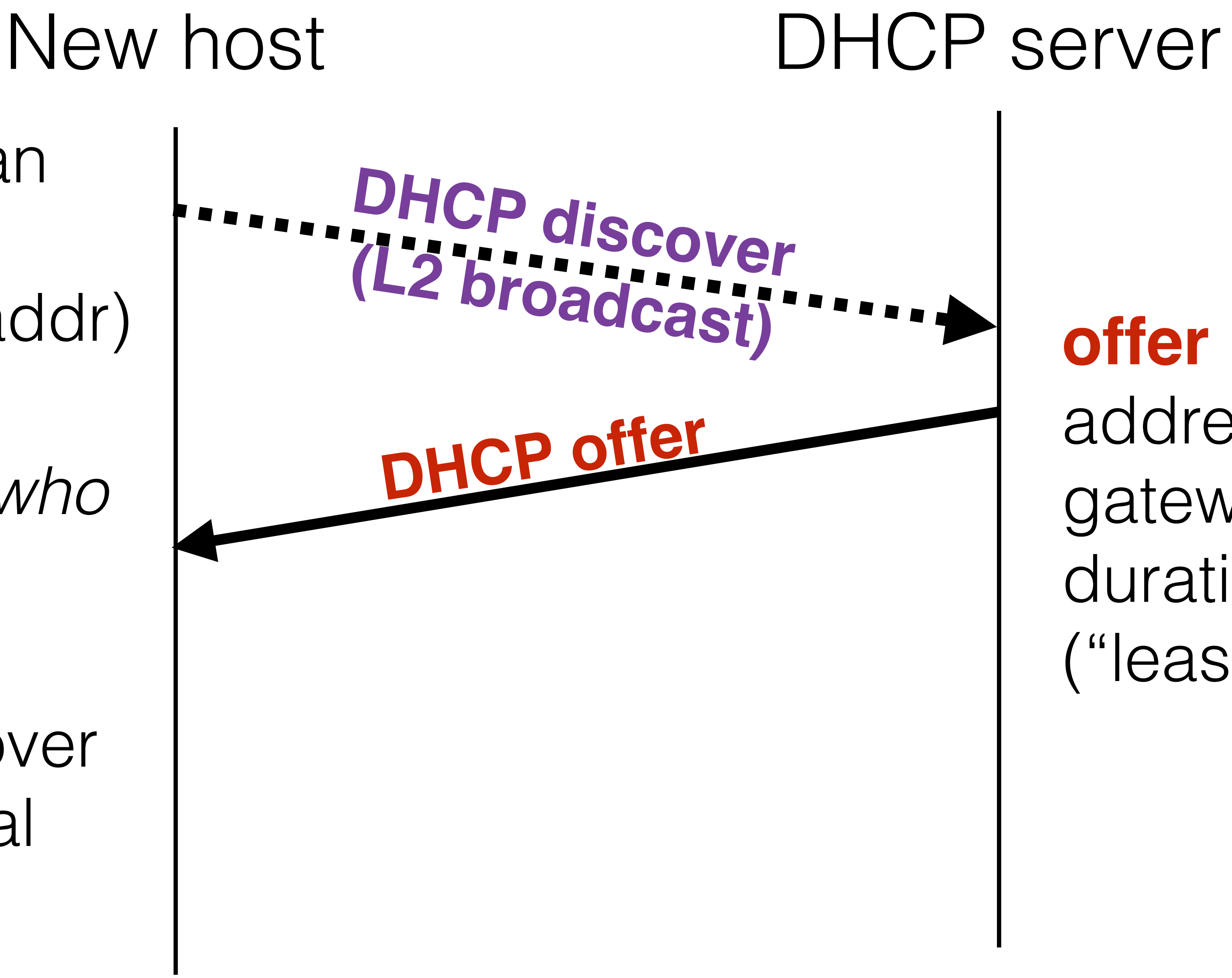

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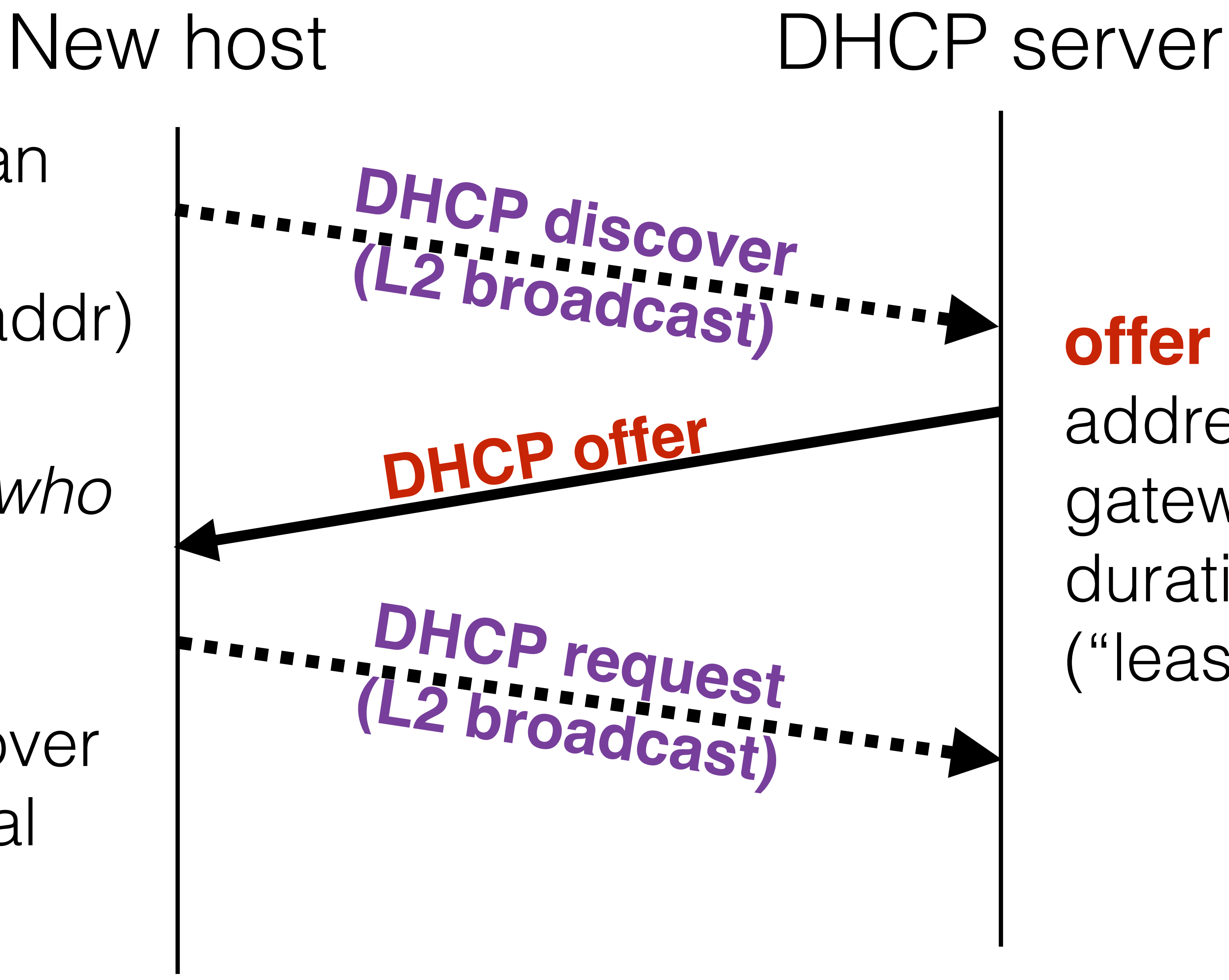

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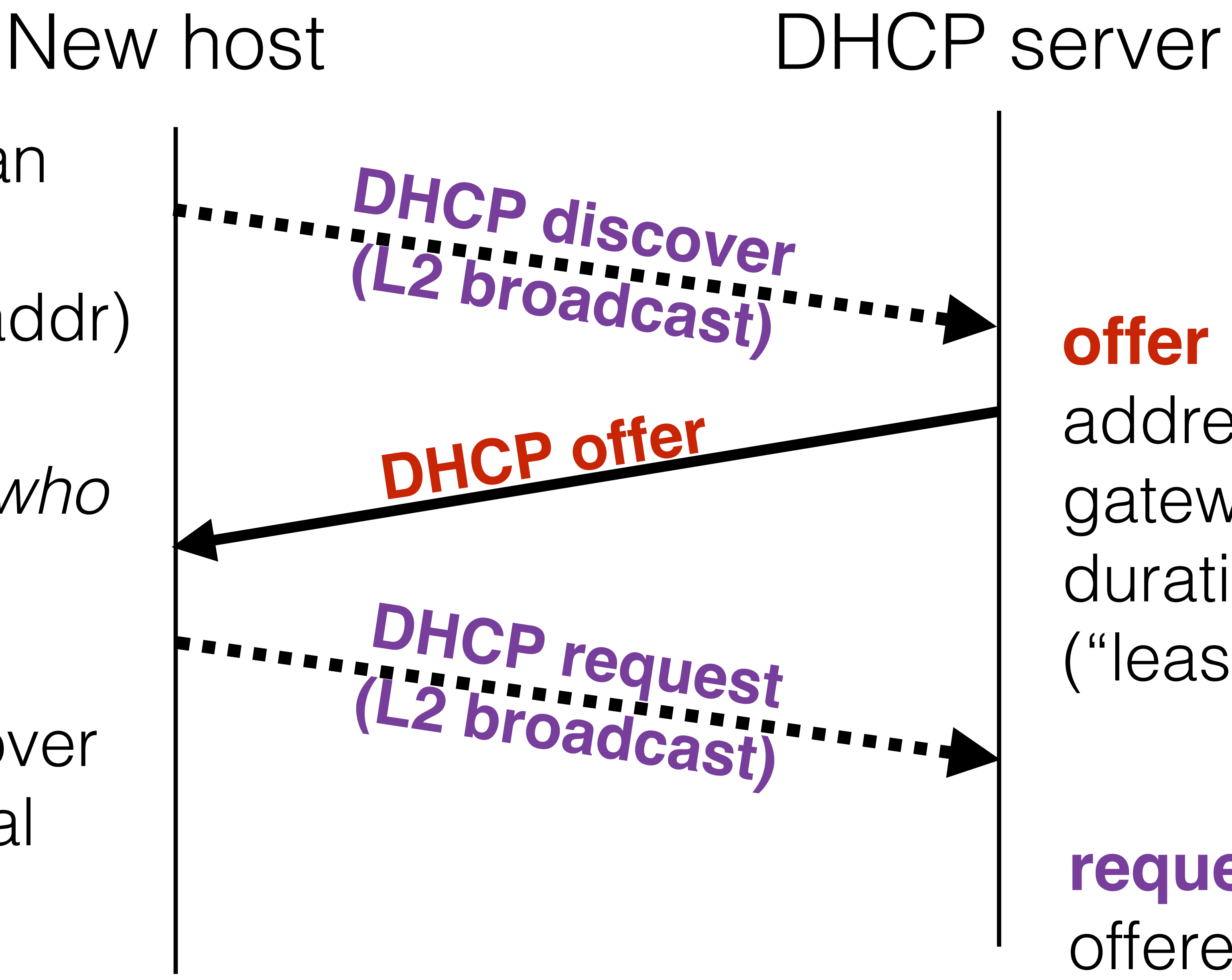

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DHCP server

DHCP discover (L2 broadcast)

DHCP offer

DHCP request (L2 broadcast)

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request asks for the offered IP address

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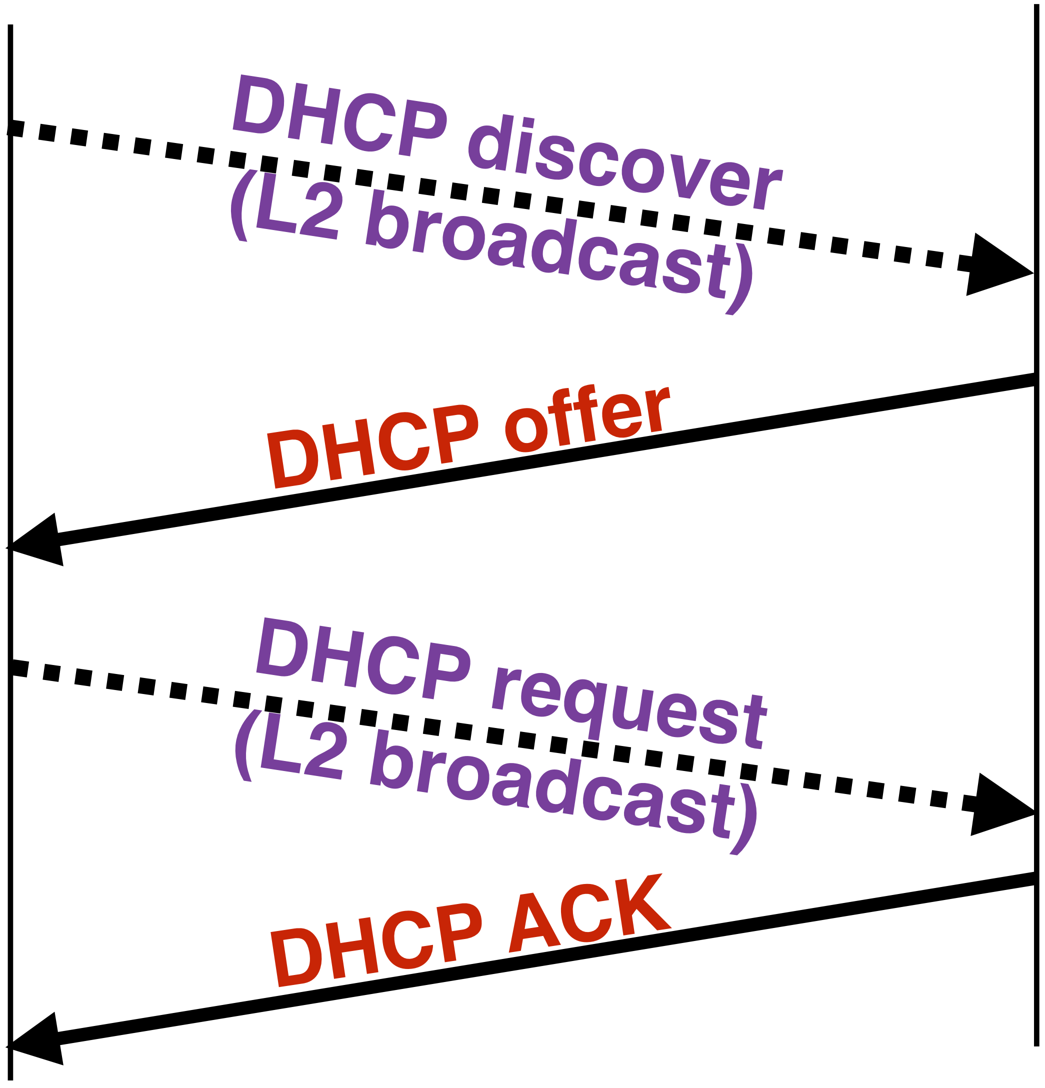


DHCP discover (L2 broadcast)

DHCP offer

DHCP request (L2 broadcast)

DHCP ACK

offer includes: IP address, DNS server, gateway router, and duration of this offer (“lease” time)

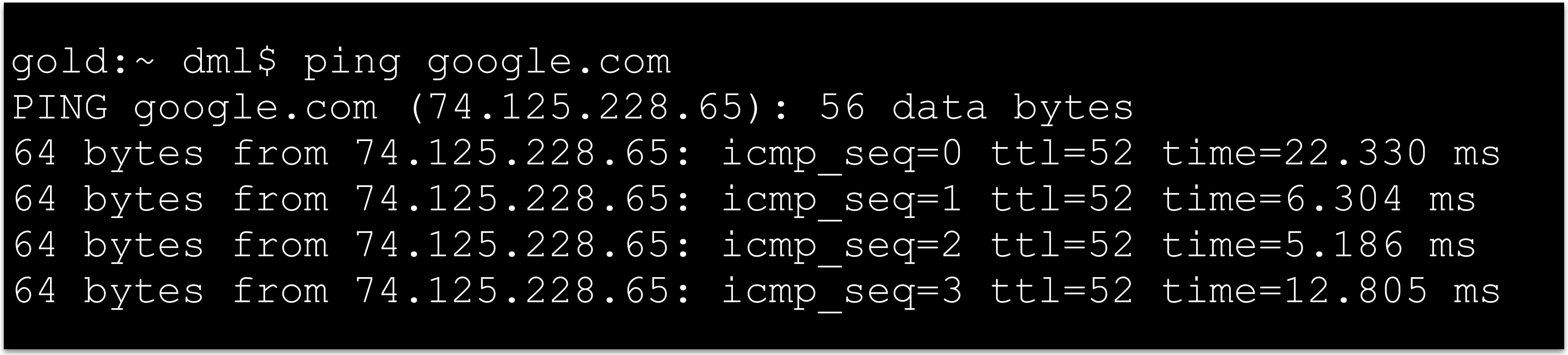
request asks for the offered IP address

DHCP attacks

- Requests are broadcast: attackers on the same subnet can hear new host’s request

- Race the *actual* DHCP server to replace:
  - DNS server
    - Redirect any of a host’s lookups (“what IP address should I use when trying to connect to google.com?”) to a machine of the attacker’s choice
  - Gateway
    - The gateway is where the host sends all of its outgoing traffic (so that the host doesn’t have to figure out routes himself)
    - Modify the gateway to intercept all of a user’s traffic
    - Then relay it to the gateway (MITM)
    - How could the user detect this?

Hostnames & IP addresses



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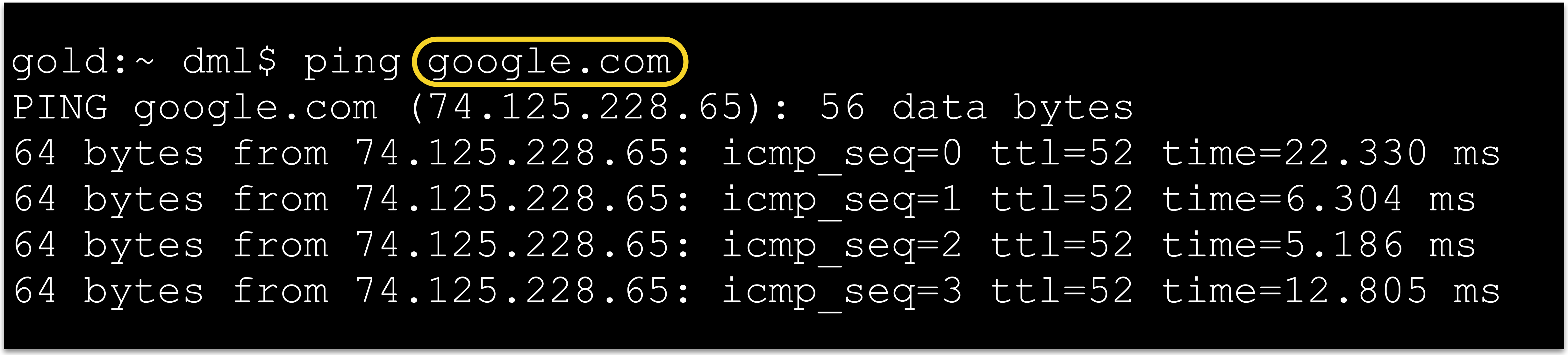

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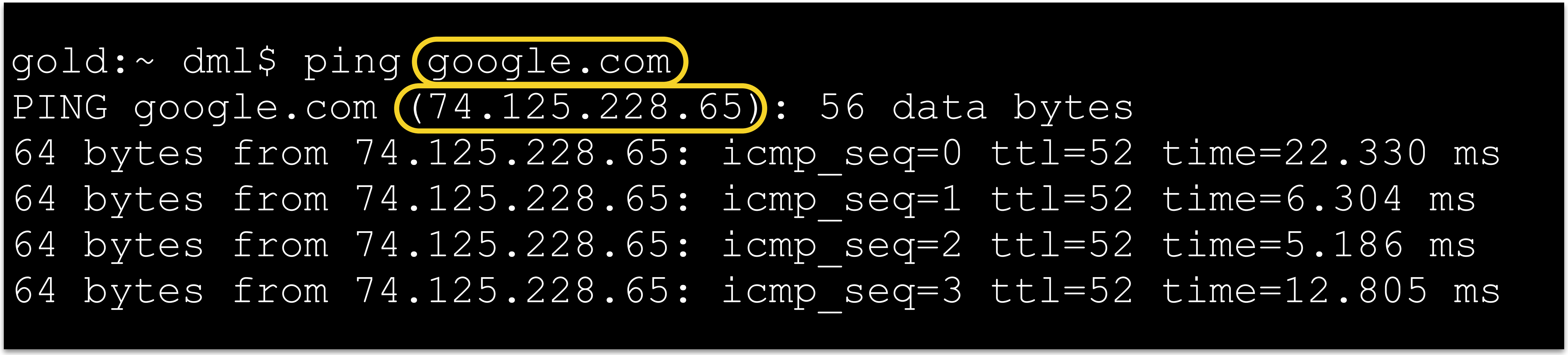

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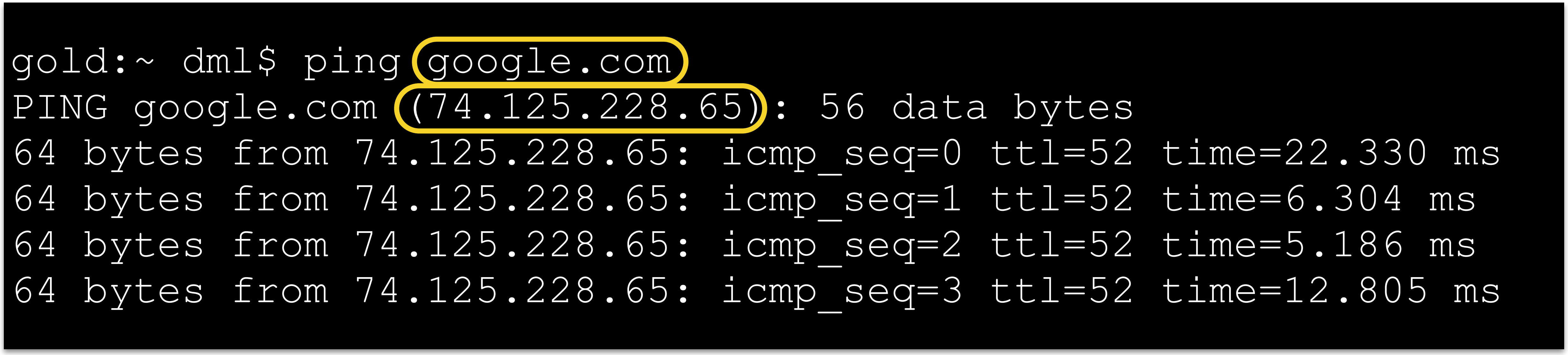

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google.com is easy to remember, but not routable

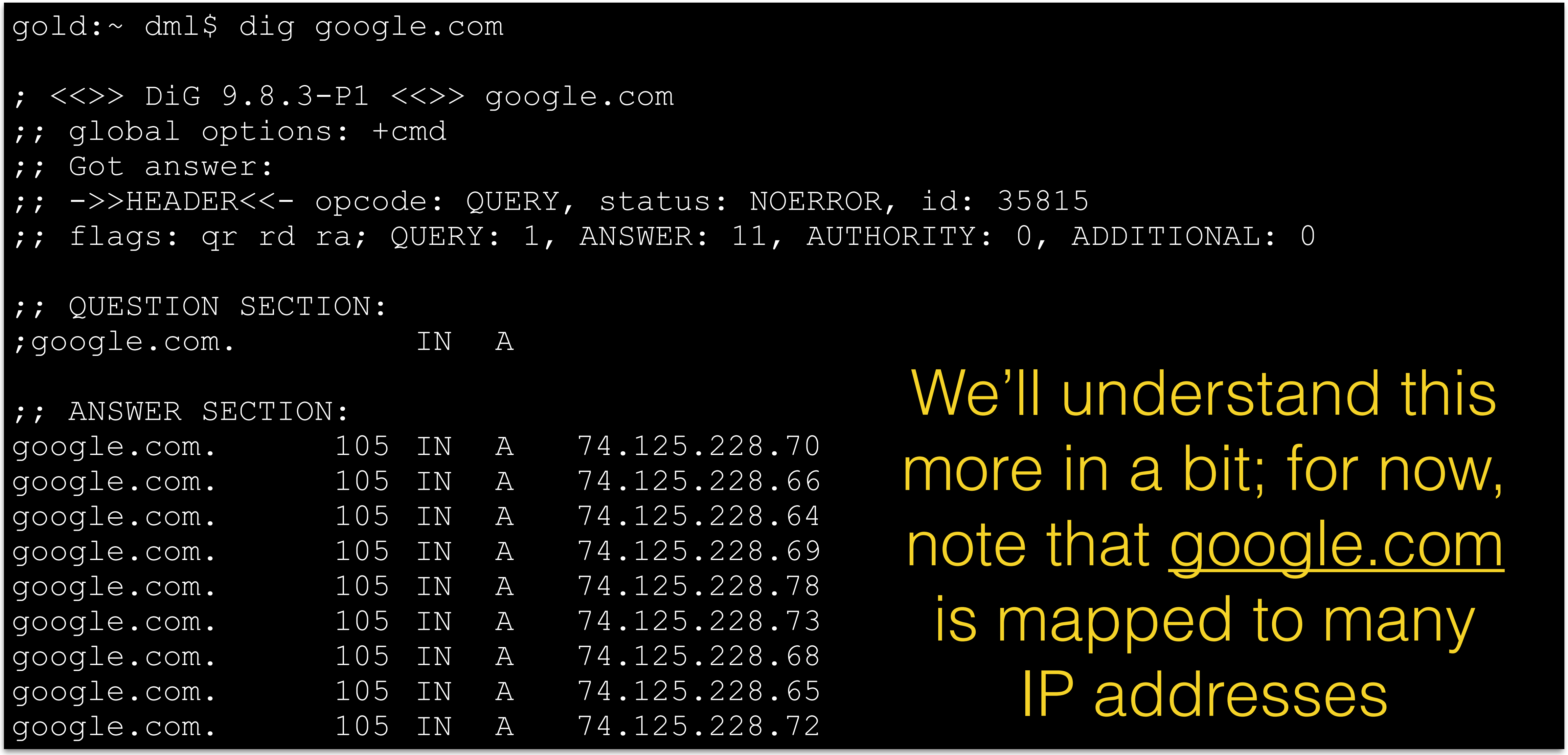
74.125.228.65 is routable

Name resolution:
The process of mapping from one to the other

Terminology

- www.cs.umd.edu = “**domain name**”
  - www.cs.umd.edu is a “subdomain” of cs.umd.edu

- Domain names can map to a set of IP addresses



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;; global options: +cmd
;; Got answer:
;; ->>HEADER<<- opcode: QUERY, status: NOERROR, id: 35815
;; flags: qr rd ra; QUERY: 1, ANSWER: 11, AUTHORITY: 0, ADDITIONAL: 0

;; QUESTION SECTION:
;google.com.          IN   A

;; ANSWER SECTION:
google.com.      105  IN   A    74.125.228.70
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We’ll understand this more in a bit; for now, note that google.com is mapped to many IP addresses

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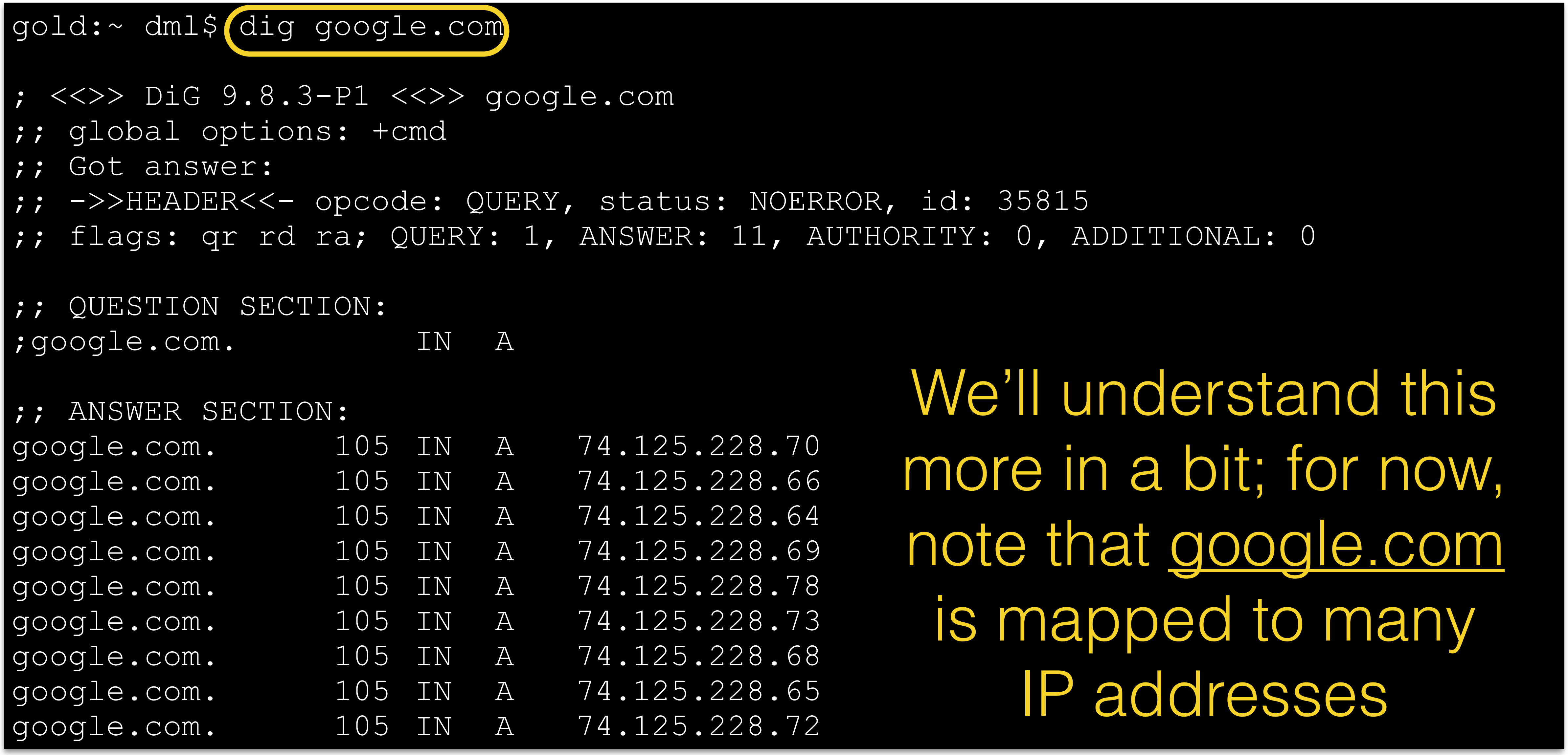

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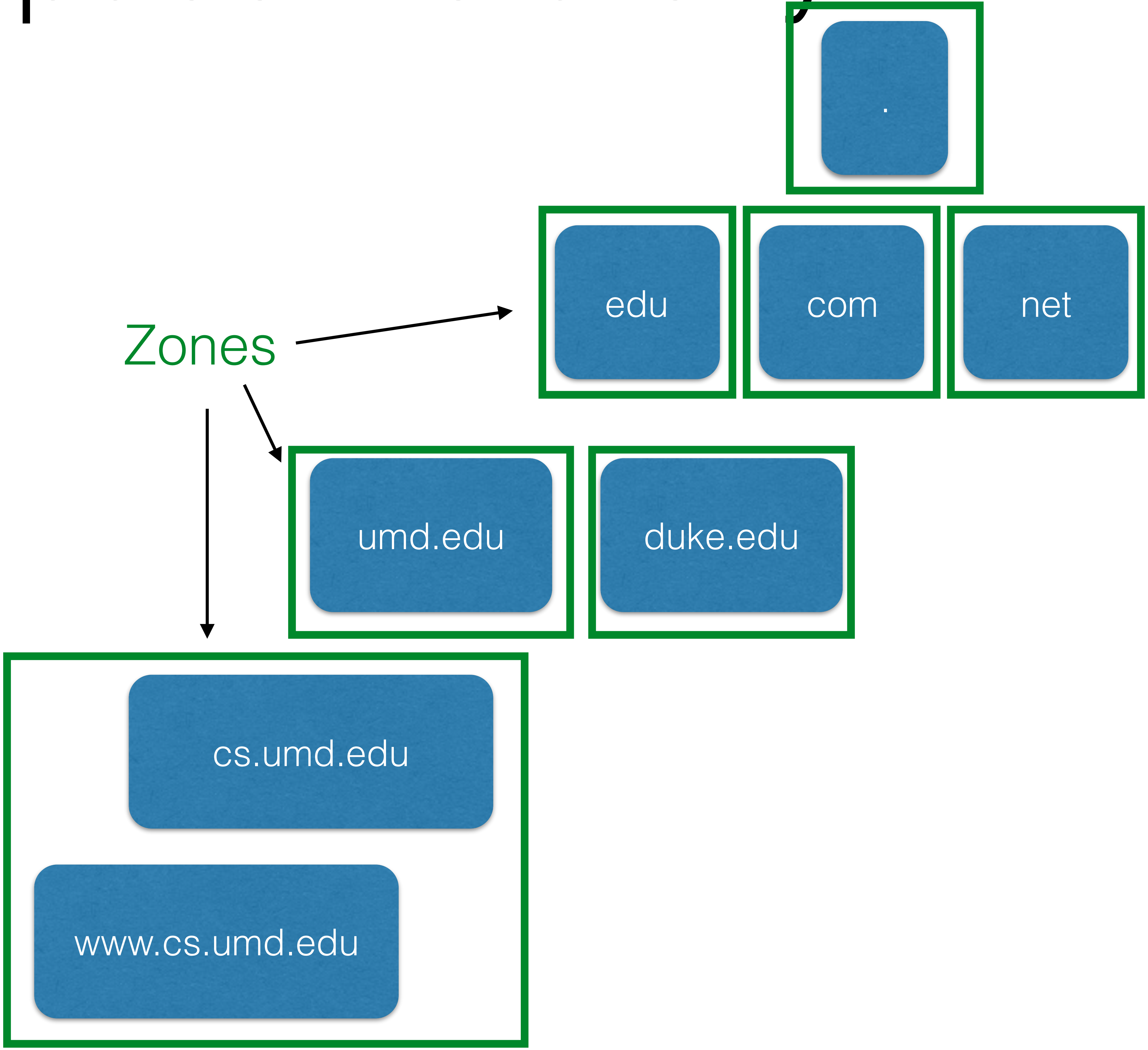
Terminology

- "**zone**" = a portion of the DNS namespace, divided up for administrative reasons
    - Think of it like a collection of hostname/IP address pairs that happen to be lumped together
        - www.google.com, mail.google.com, dev.google.com, …

- Subdomains do not need to be in the same zone
    - Allows the owner of one zone (umd.edu) to delegate responsibility to another (cs.umd.edu)

Namespace hierarchy



Terminology

- "**Nameserver**" = A piece of code that answers queries of the form "What is the IP address for foo.bar.com?"
  - Every zone must run ≥2 nameservers
  - Several very common nameserver implementations: BIND, PowerDNS (more popular in Europe)

- "**Authoritative nameserver**":
  - Every zone has to maintain a file that maps IP addresses and hostnames ("www.cs.umd.edu is 128.8.127.3")
  - One of the name servers in the zone has the *master* copy of this file. It is the authority on the mapping.

Terminology

- “**Resolver**” - while name servers *answer* queries, resolvers *ask* queries.

- Every OS has a resolver. Typically small and pretty dumb. All it typically does it forward the query to a local…

- “**Recursive nameserver**” - a nameserver which will do the heavy lifting, issuing queries on behalf of the client resolver until an authoritative answer returns.

- Prevalence
    - There is almost always a *local* (private) recursive name server
    - But very rare for name servers to support recursive queries otherwise

Terminology

- "**Record**" (or "resource record") = usually think of it as a mapping between hostname and IP address

- But more generally, it can map virtually anything to virtually anything

- Many record types:
  - (**A**)ddress records (IP <-> hostname)
  - Mail server (**MX**, mail exchanger)
  - SOA (start of authority, to delineate different zones)
  - Others for DNSSEC to be able to share keys

- Records are the unit of information

Terminology

Nameservers within a zone must be able to give:

- **Authoritative answers (A)** for hostnames in that zone
  - The umd.edu zone's nameservers must be able to tell us what the IP address for umd.edu is

"A" record: umd.edu = 54.84.241.99

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- Pointers to **name servers (NS)** who host zones in its subdomains
  - The umd.edu zone's nameservers must be able to tell us what the name and IP address of the cs.umd.edu zone's nameservers

"NS" record: cs.umd.edu = ipa01.cs.umd.edu.

Ask ipa01.cs.umd.edu for all cs.umd.edu subdomains

DNS

Domain Name Service at a very high level

Requesting host

What is an IP address
for cs.umd.edu?

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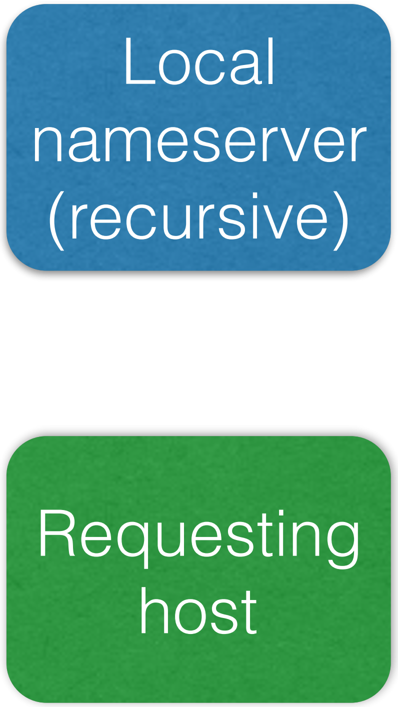

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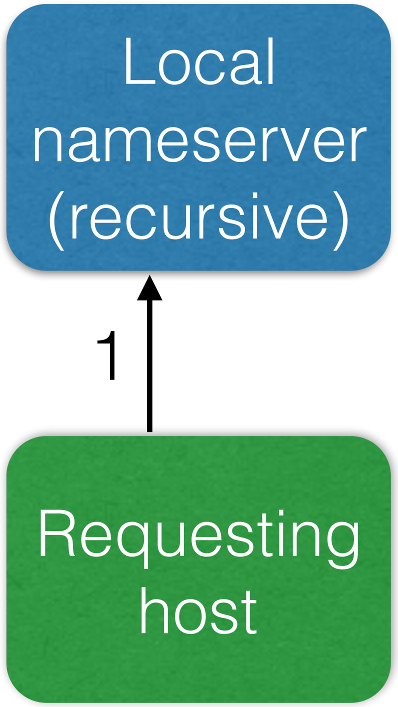

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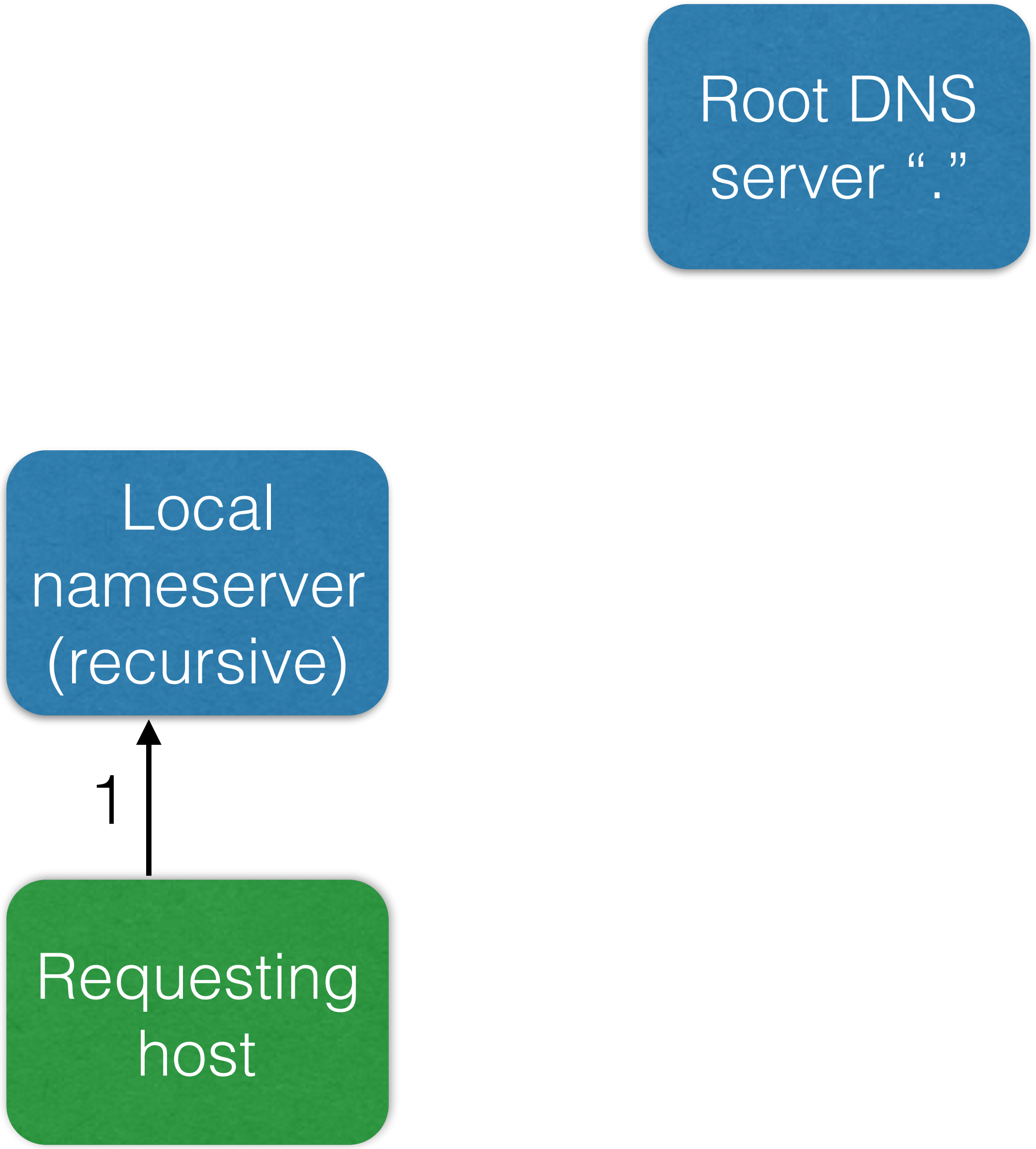

Root DNS server "."

Local nameserver (recursive)

1

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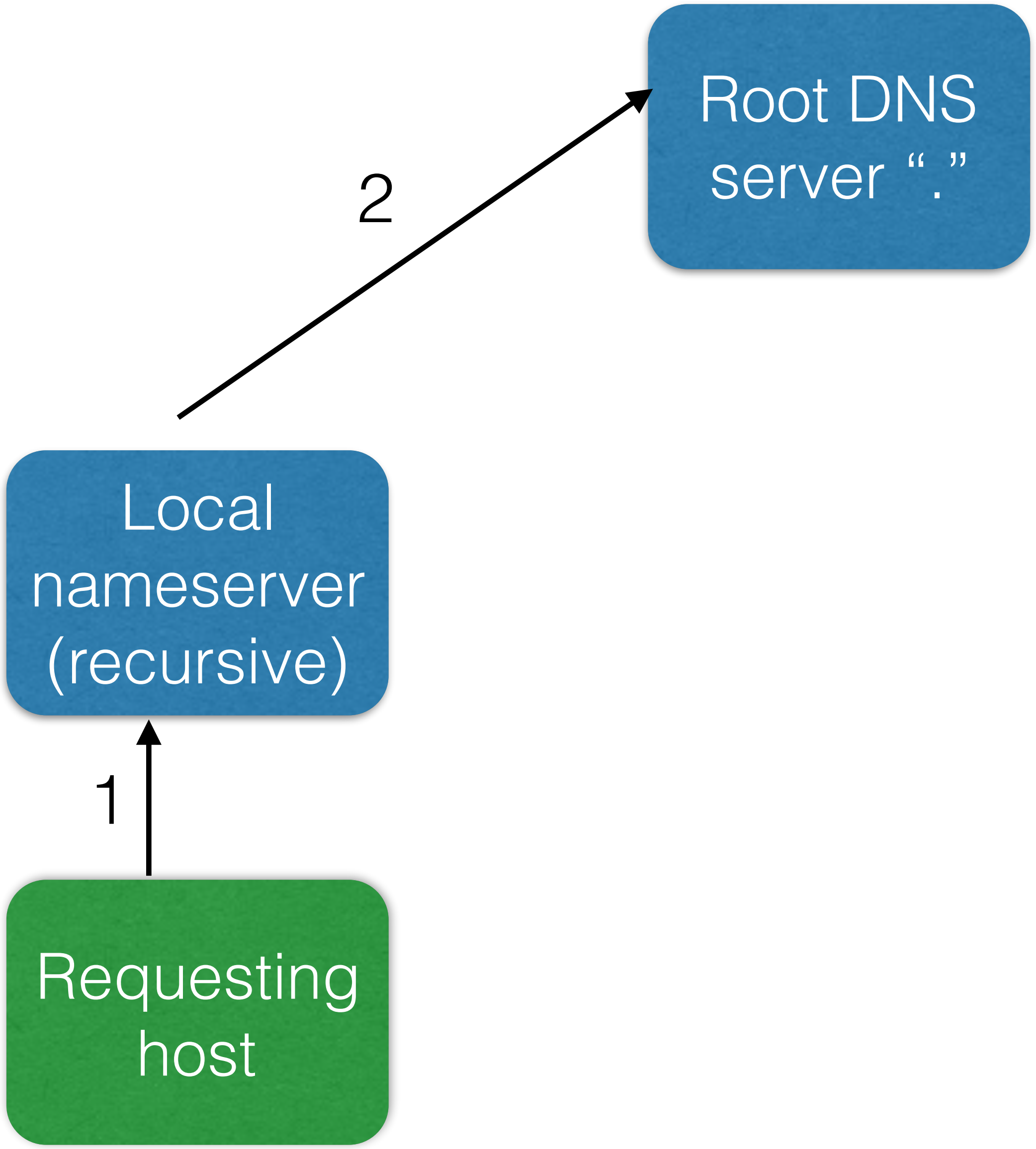

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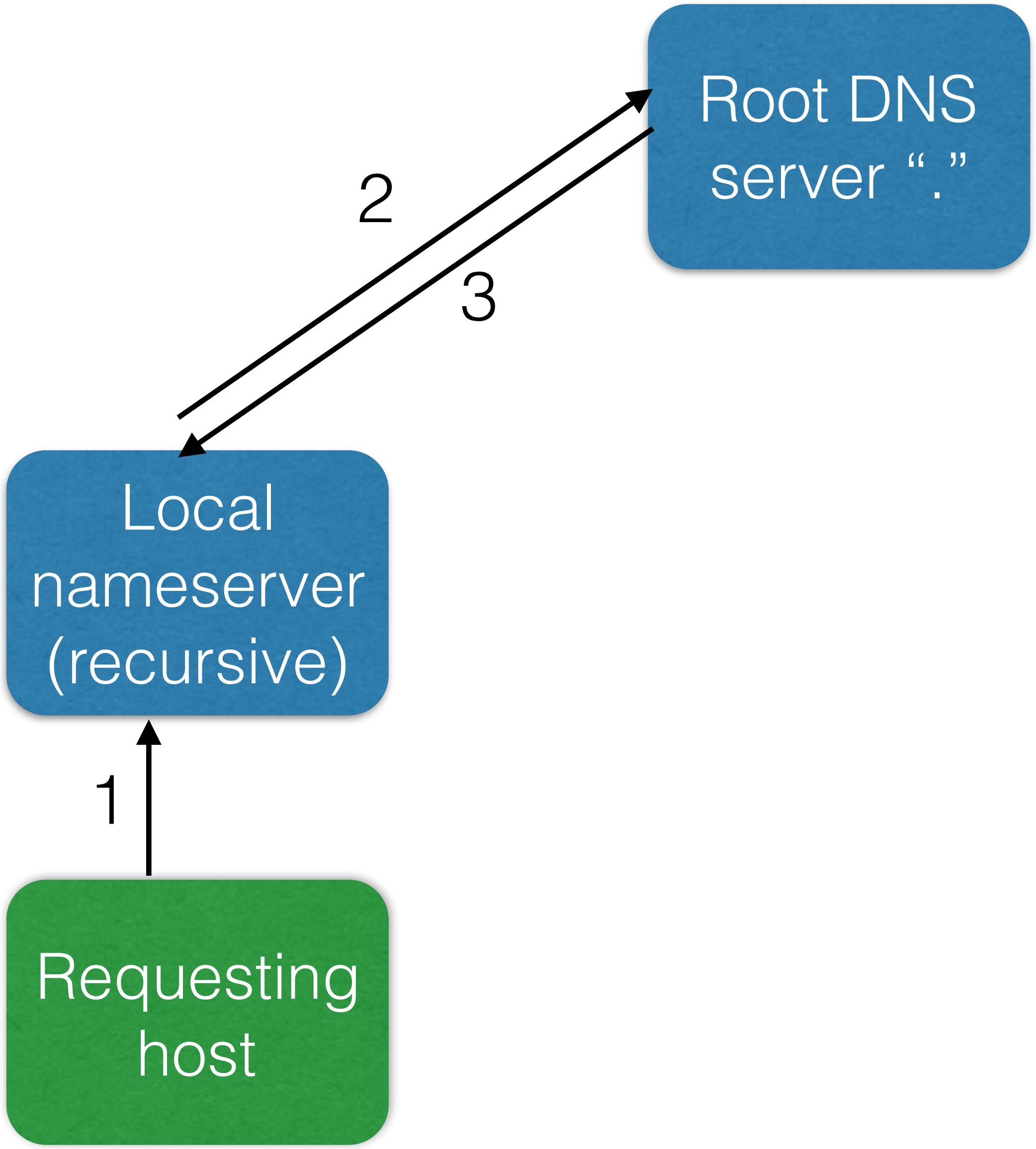

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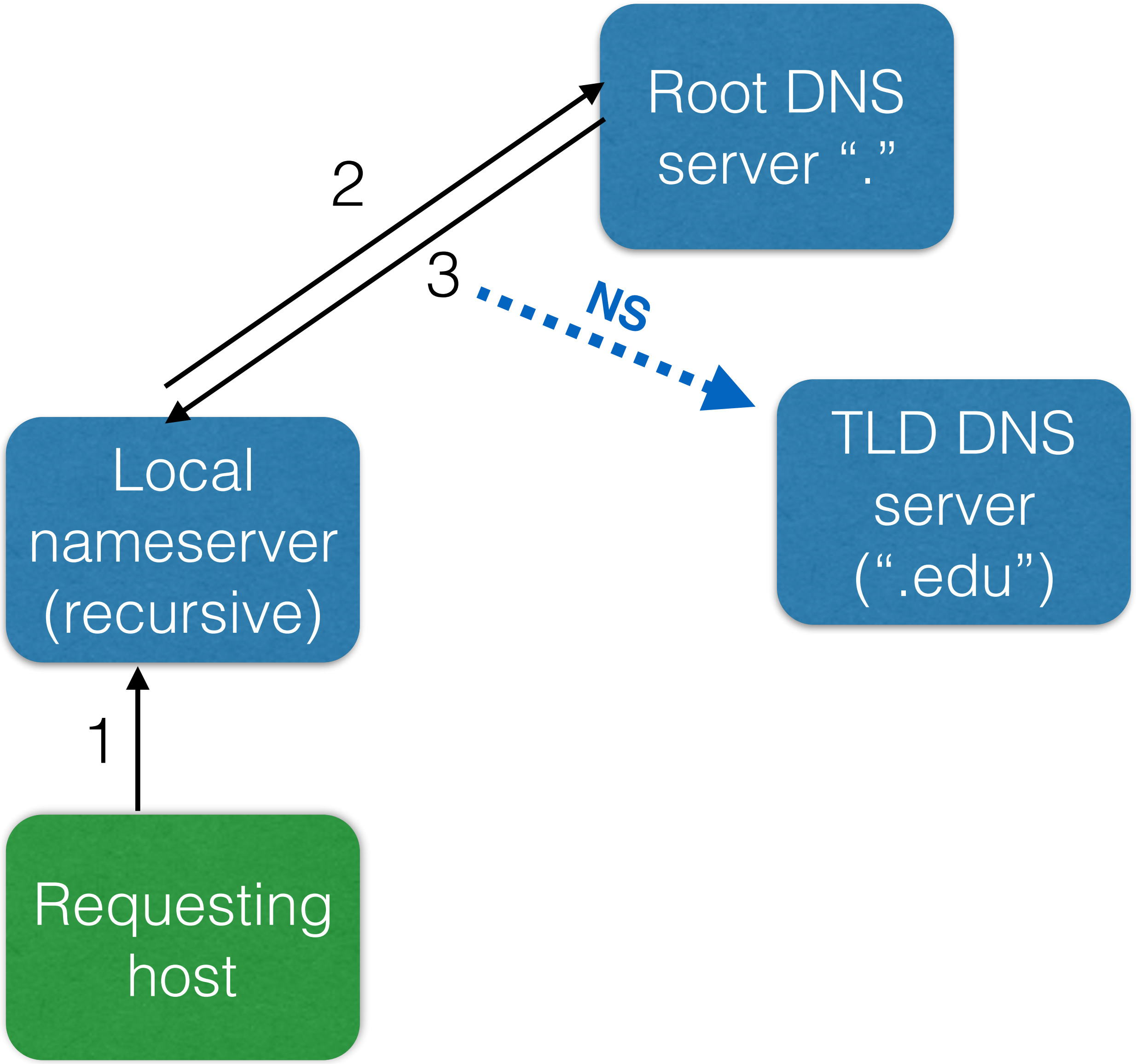

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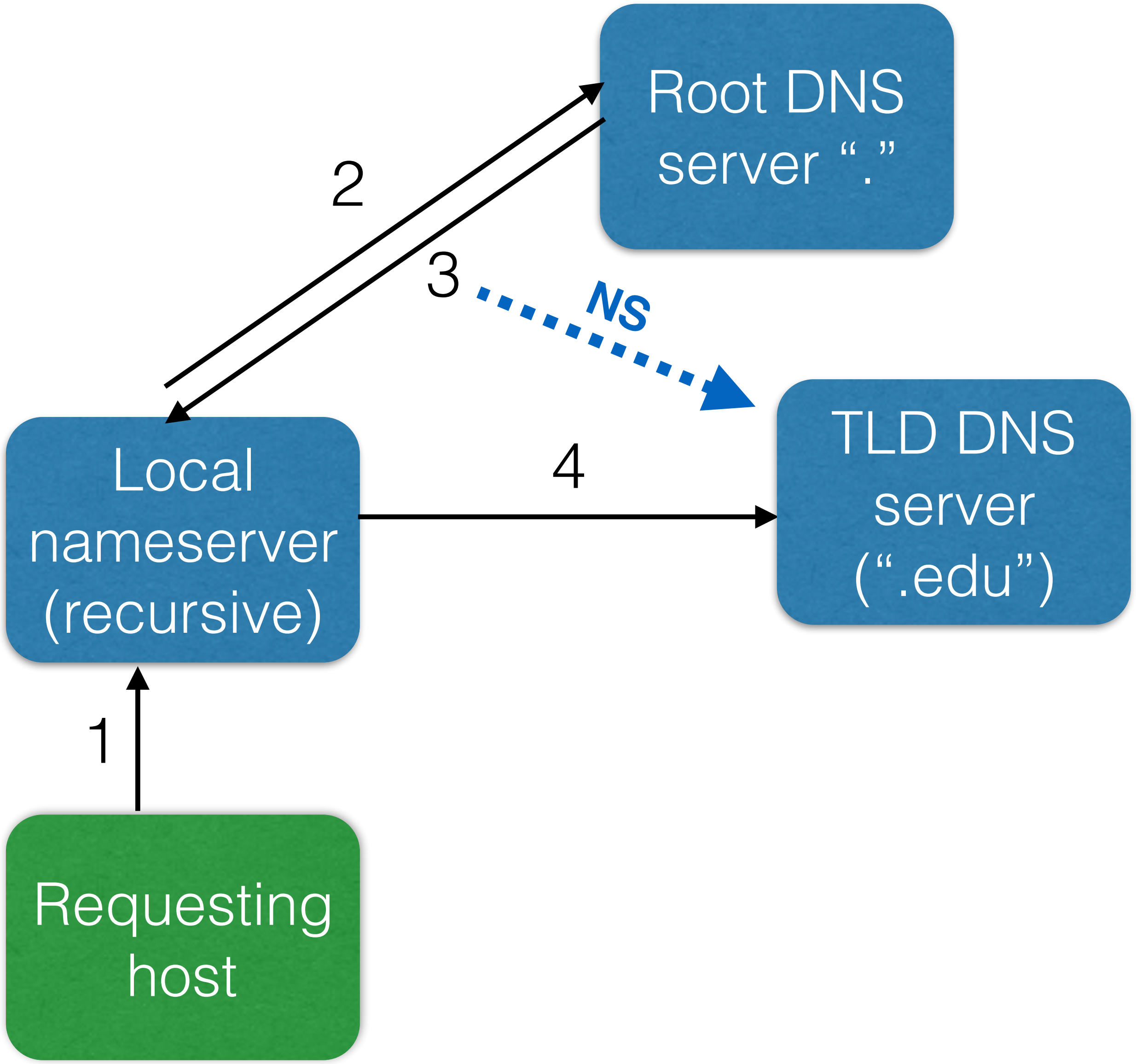

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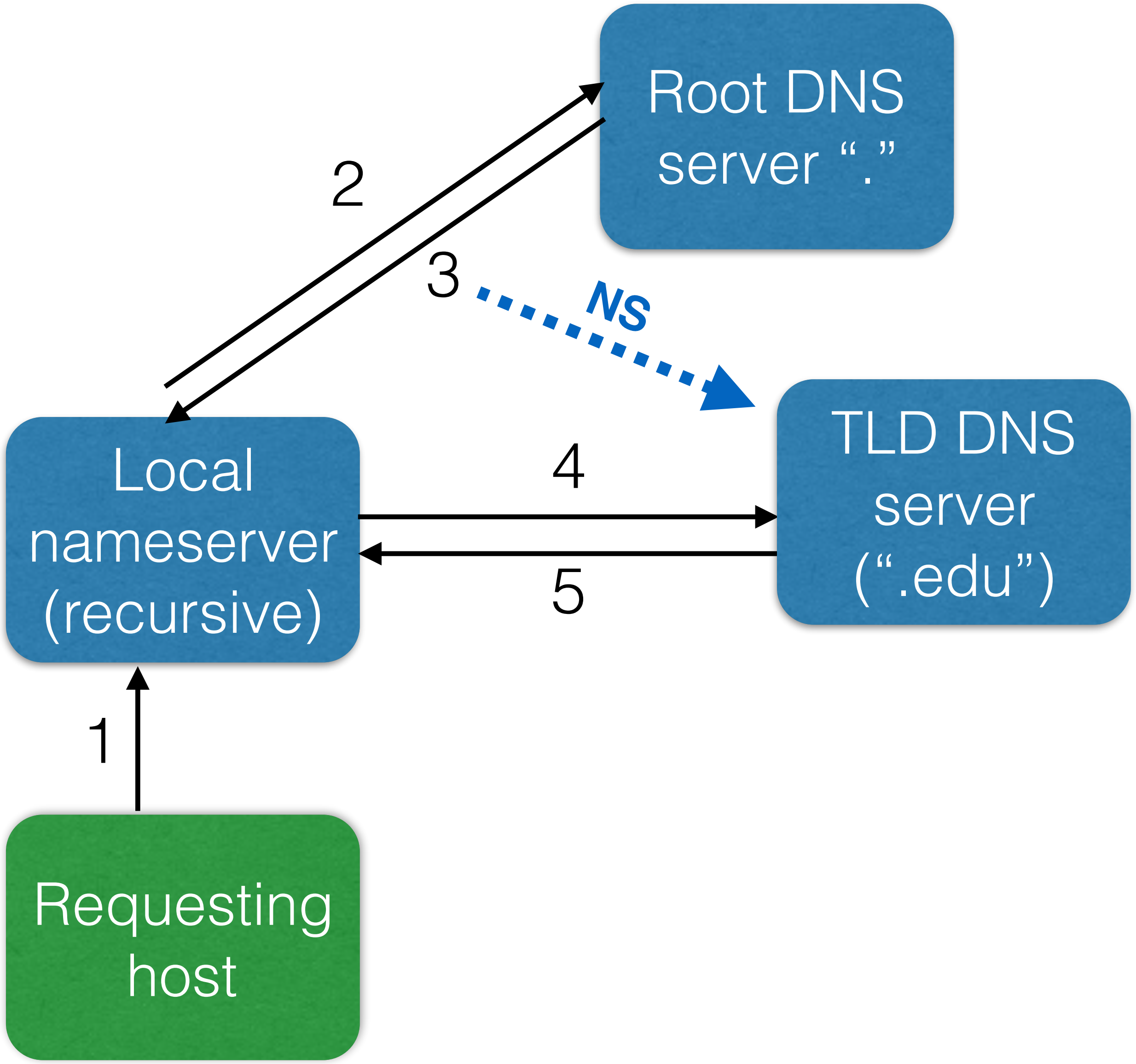

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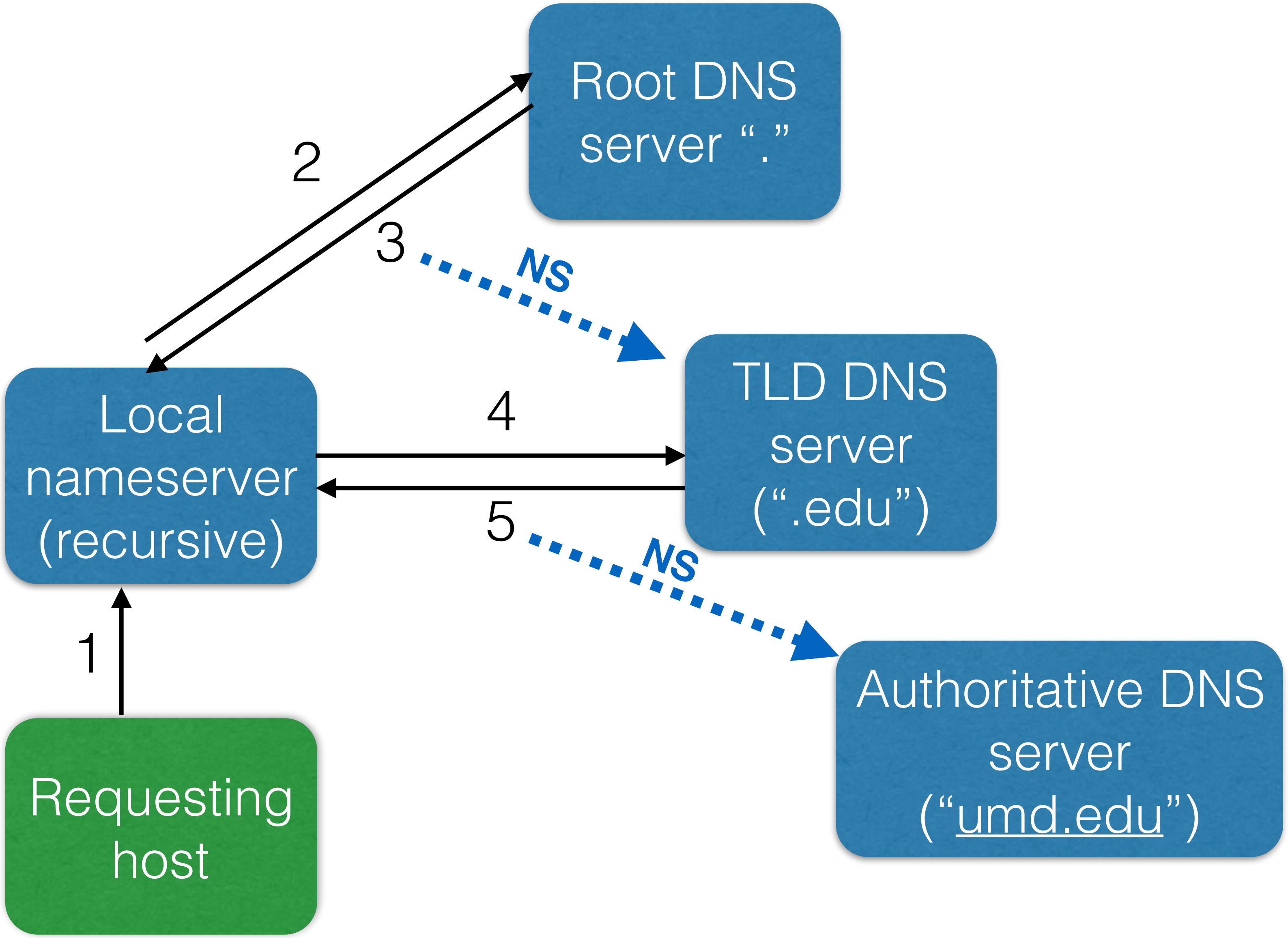

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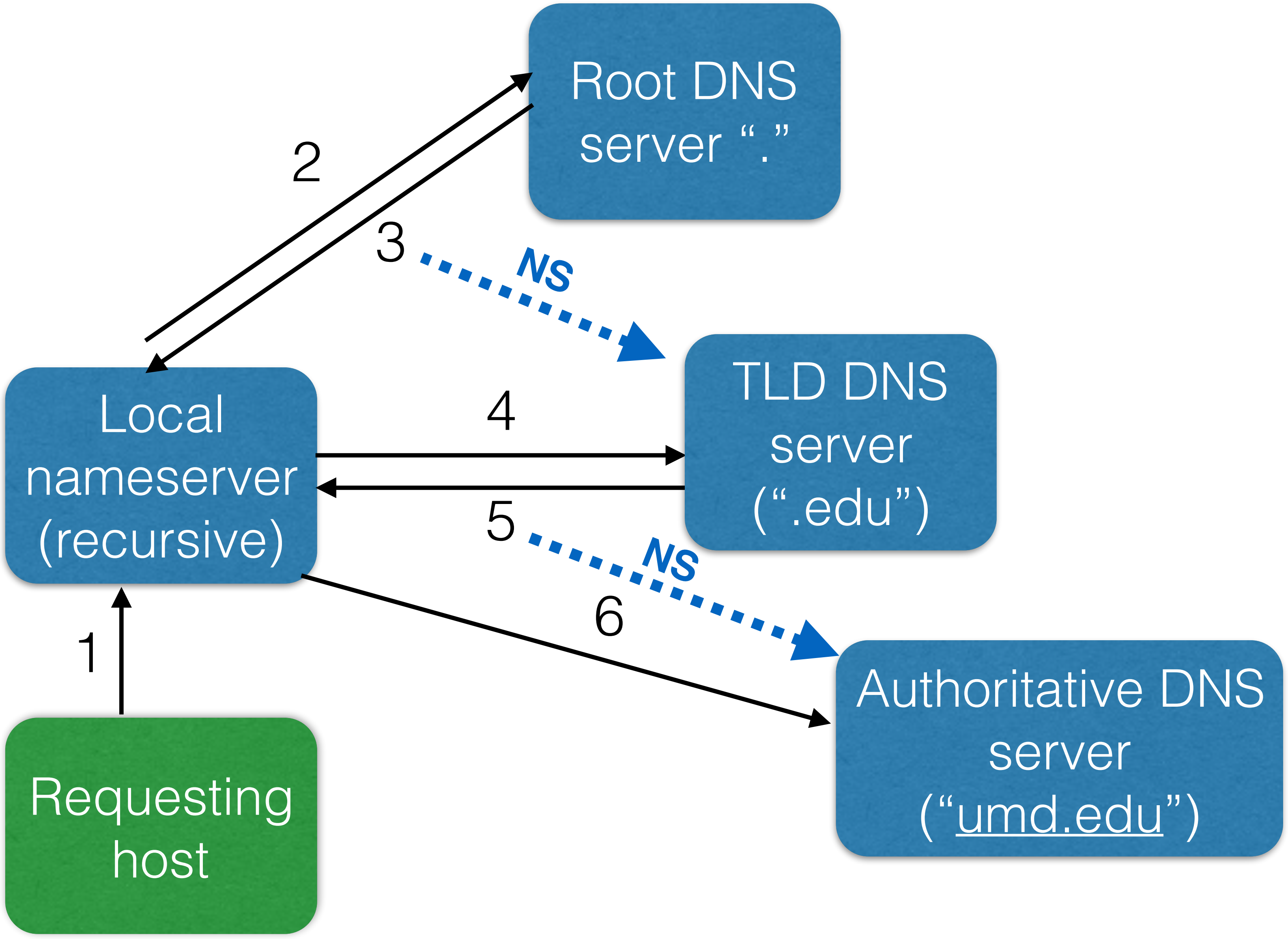

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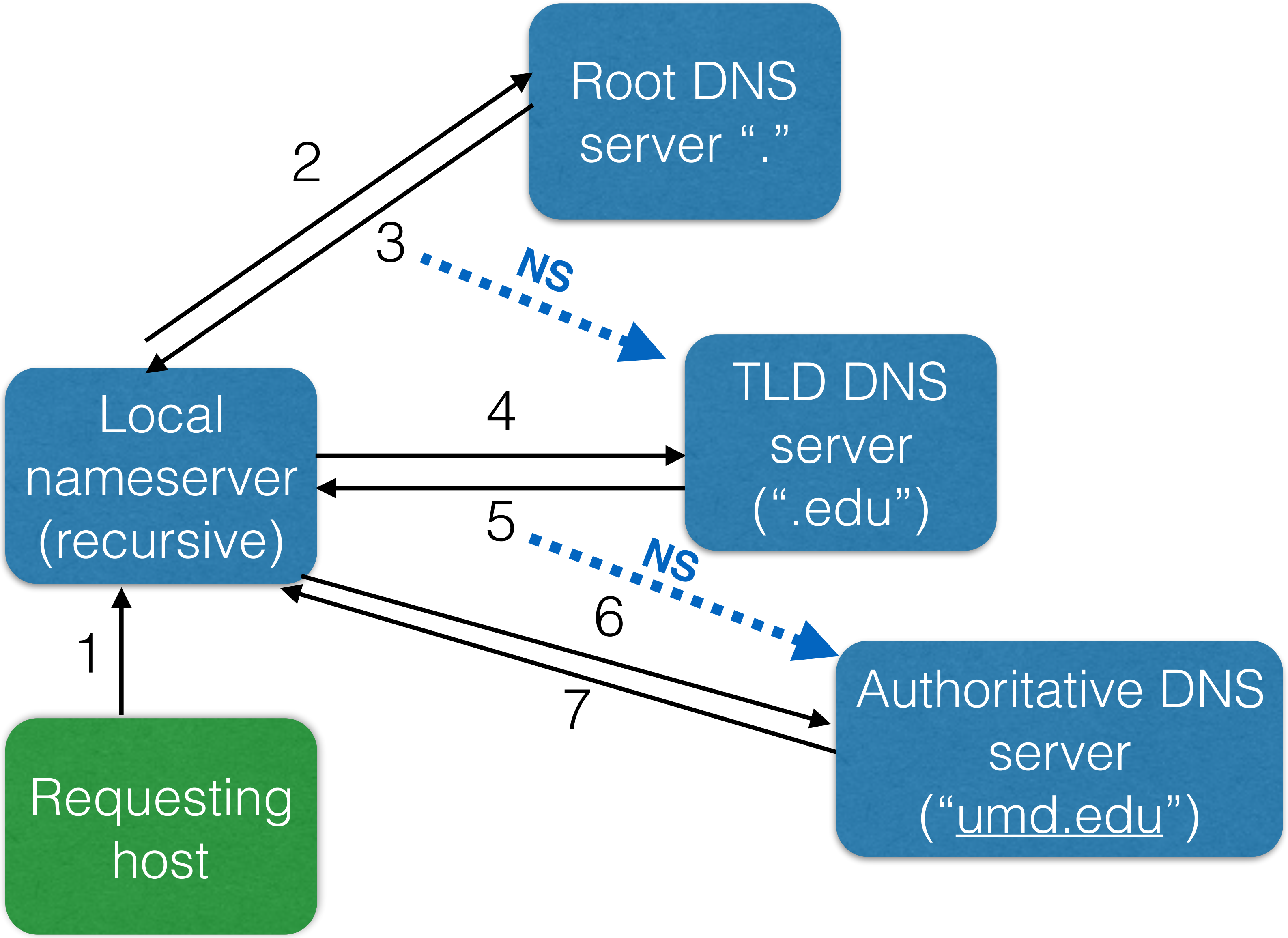

Root DNS server "."

TLD DNS server (".edu")

Authoritative DNS server ("umd.edu")

Local nameserver (recursive)

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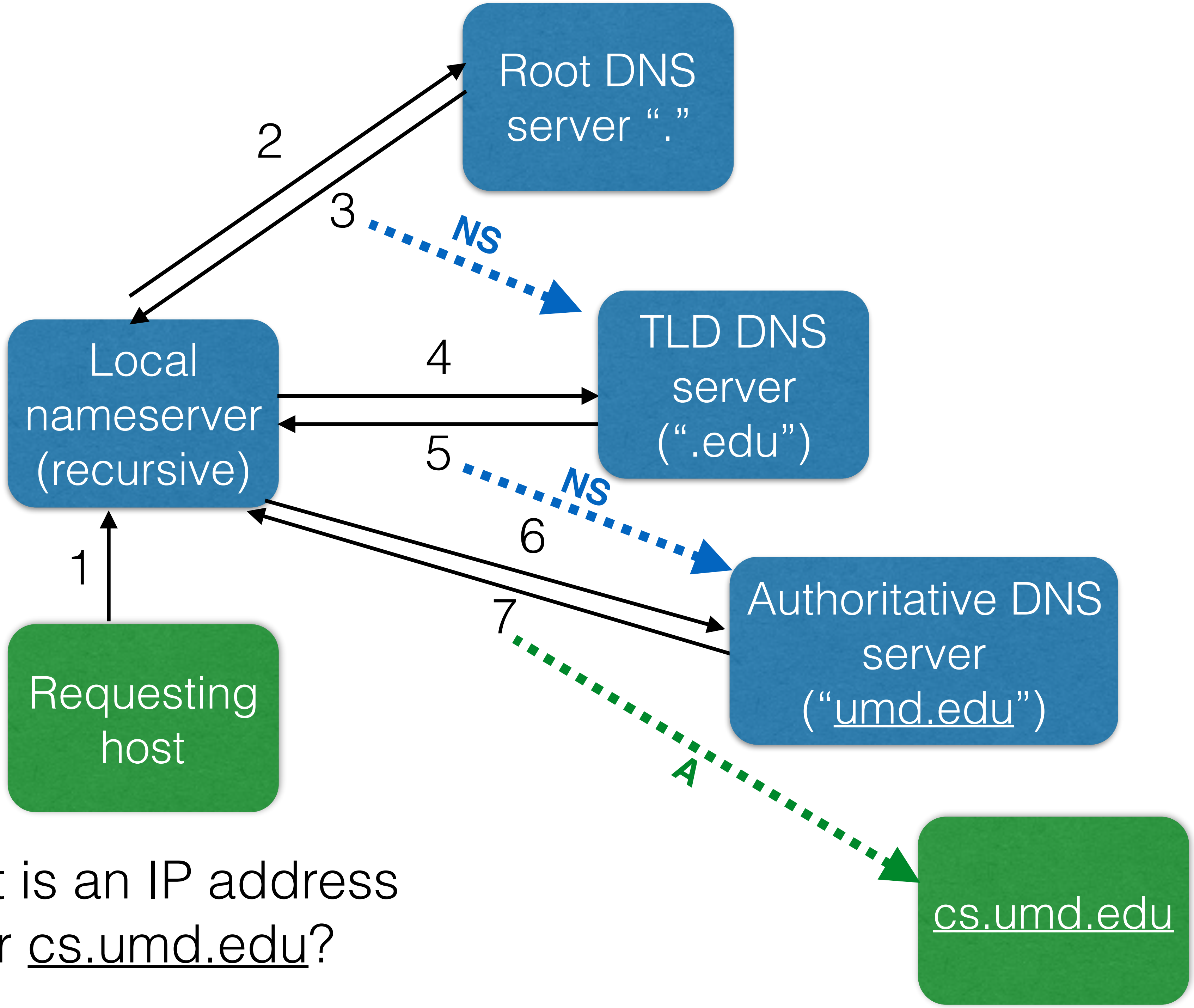

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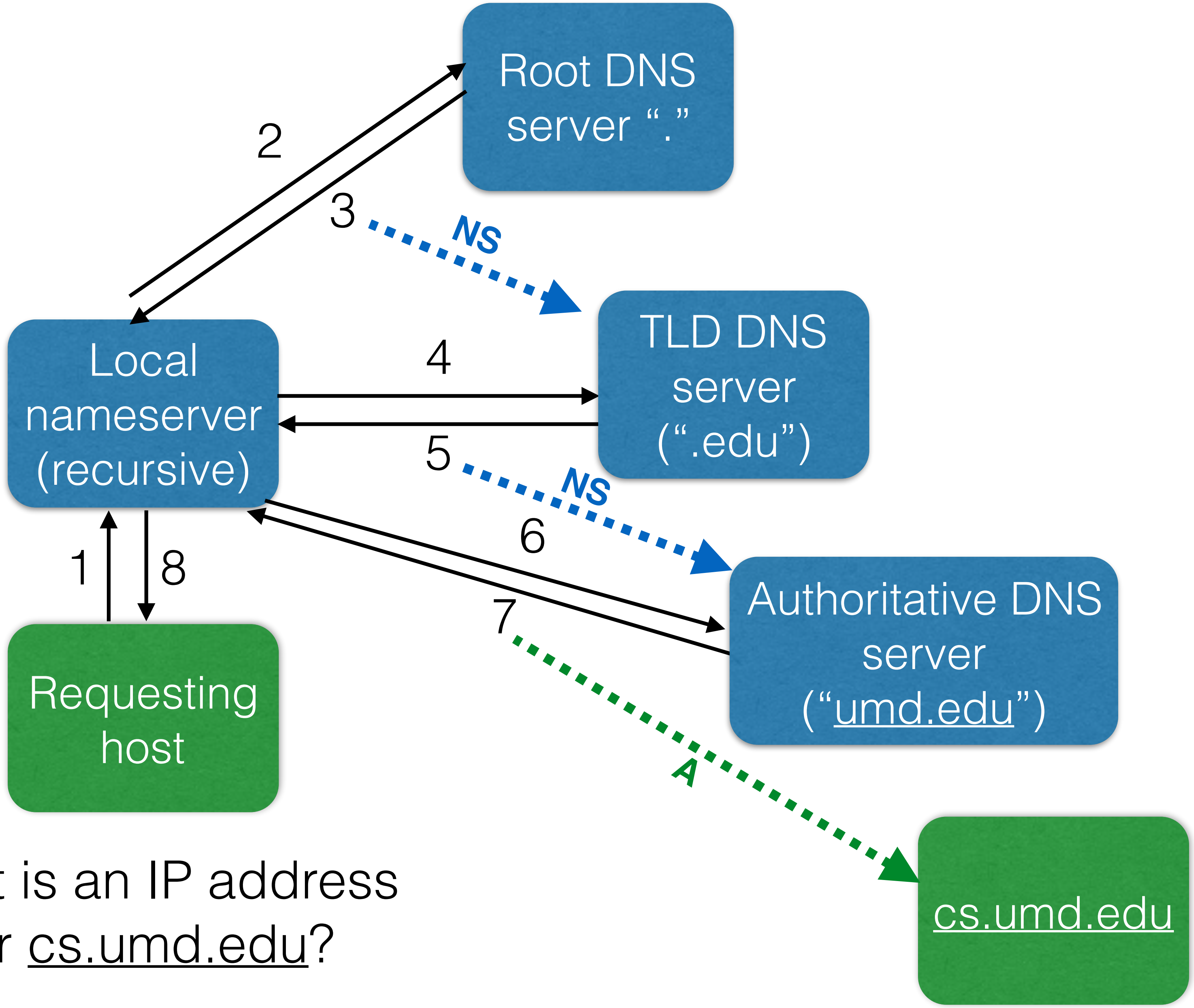

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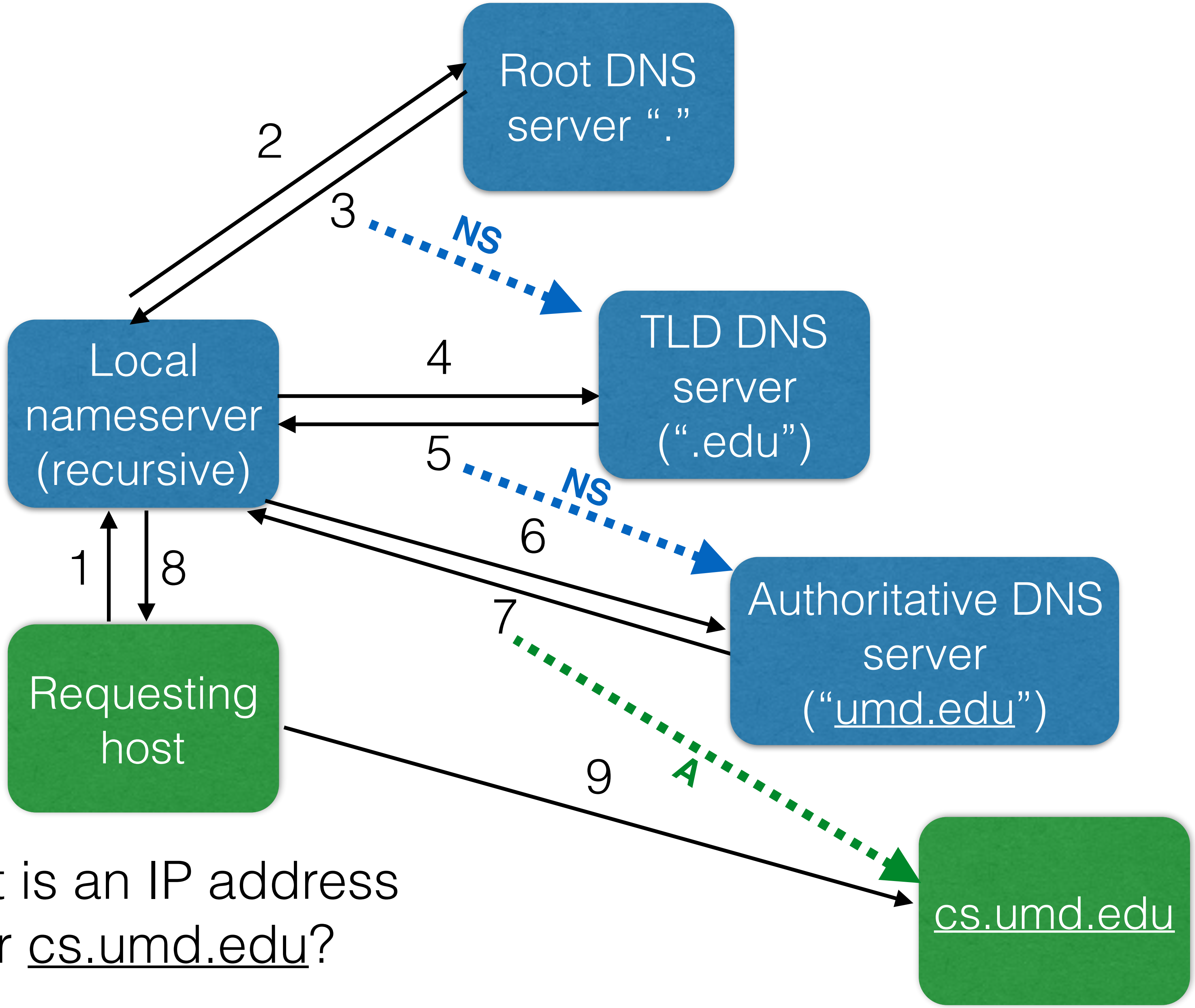

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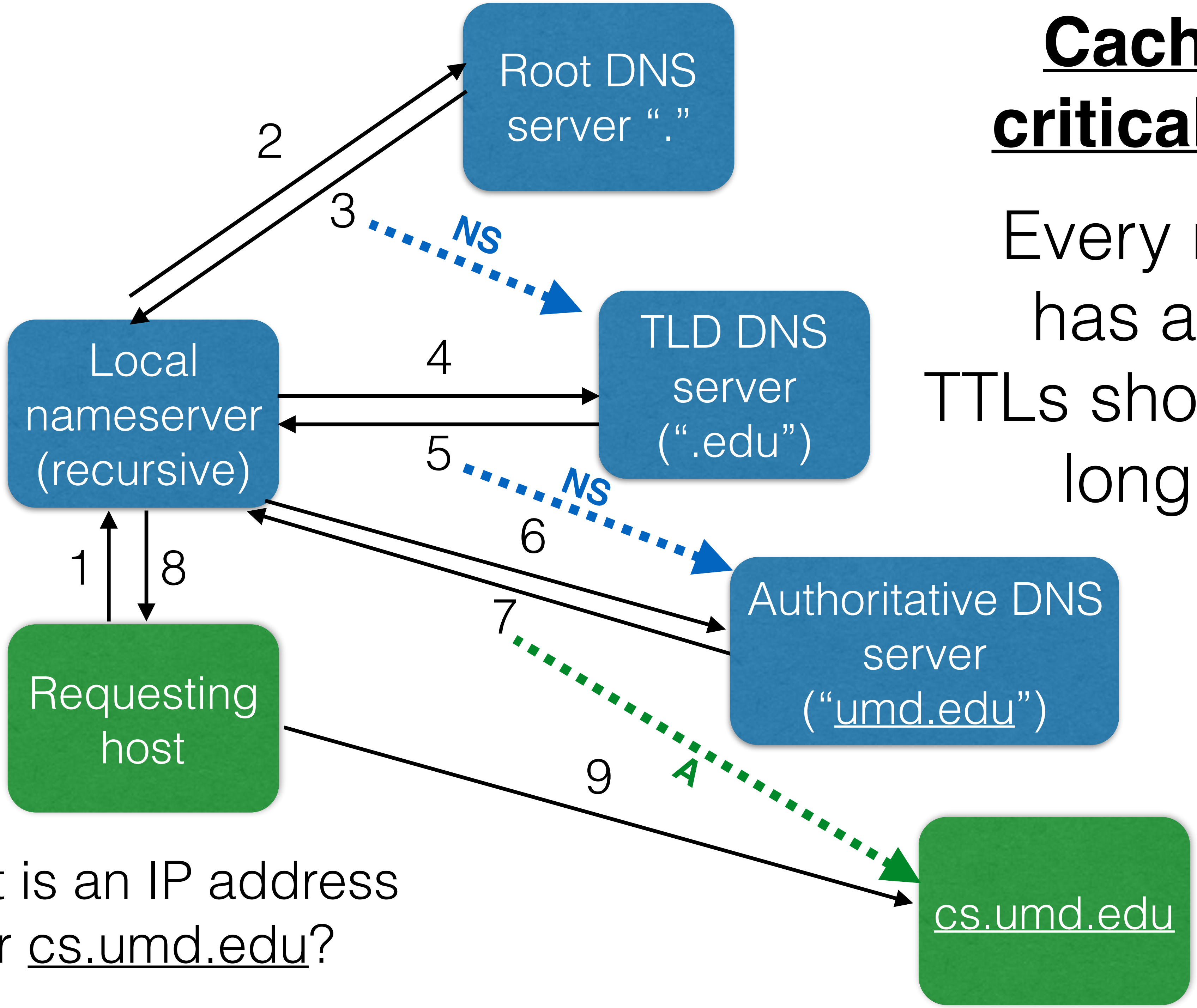


What is an IP address for cs.umd.edu?

Caching responses is critical to DNS's success

Every response (3,5,7,8) has a time-to-live (TTL). TTLs should be reasonably long (days), but some are minutes.

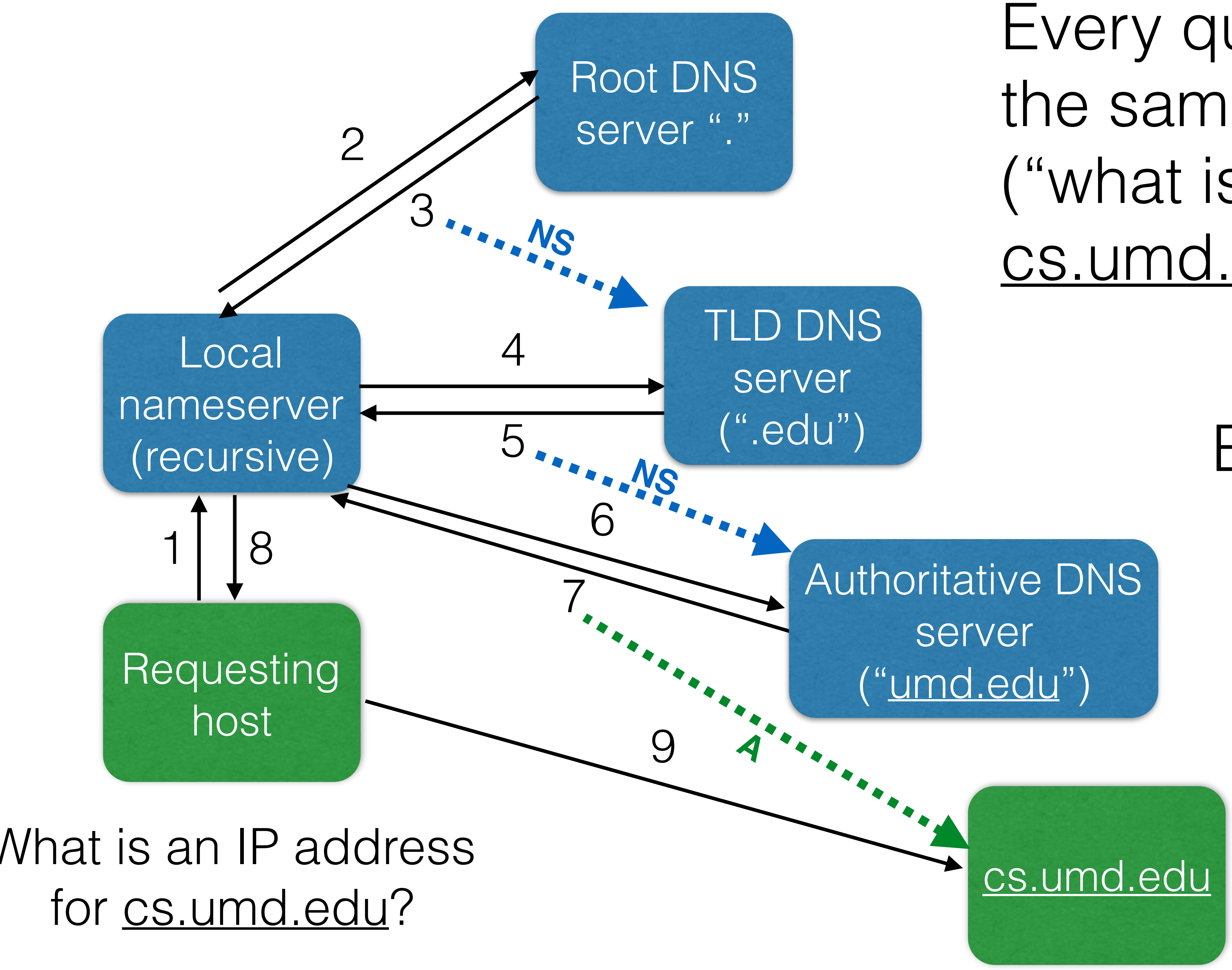
How do they know these IP addresses?

- Local DNS server: host learned this via DHCP
- A parent knows its children: part of the registration process
- Root nameserver: *hardcoded* into the local DNS server (and every DNS server)
  - 13 root servers (logically): A-root, B-root, …, M-root
  - These IP addresses change *very* infrequently
  - **UMD runs D-root.**
    - IP address changed beginning of 2013!!
    - For the most part, the change-over went alright, but Lots of weird things happened — ask me some time.

Caching

- Central to DNS’s success
- Also central to attacks
- “Cache poisoning”: filling a victim’s cache with false information

Queries



Vhat is an IP address for cs.umd.edu?

Every query (2,4,6) has the same request in it ("what is the IP address for cs.umd.edu?")

But **different**:

- dst IP (port = 53)
- query ID

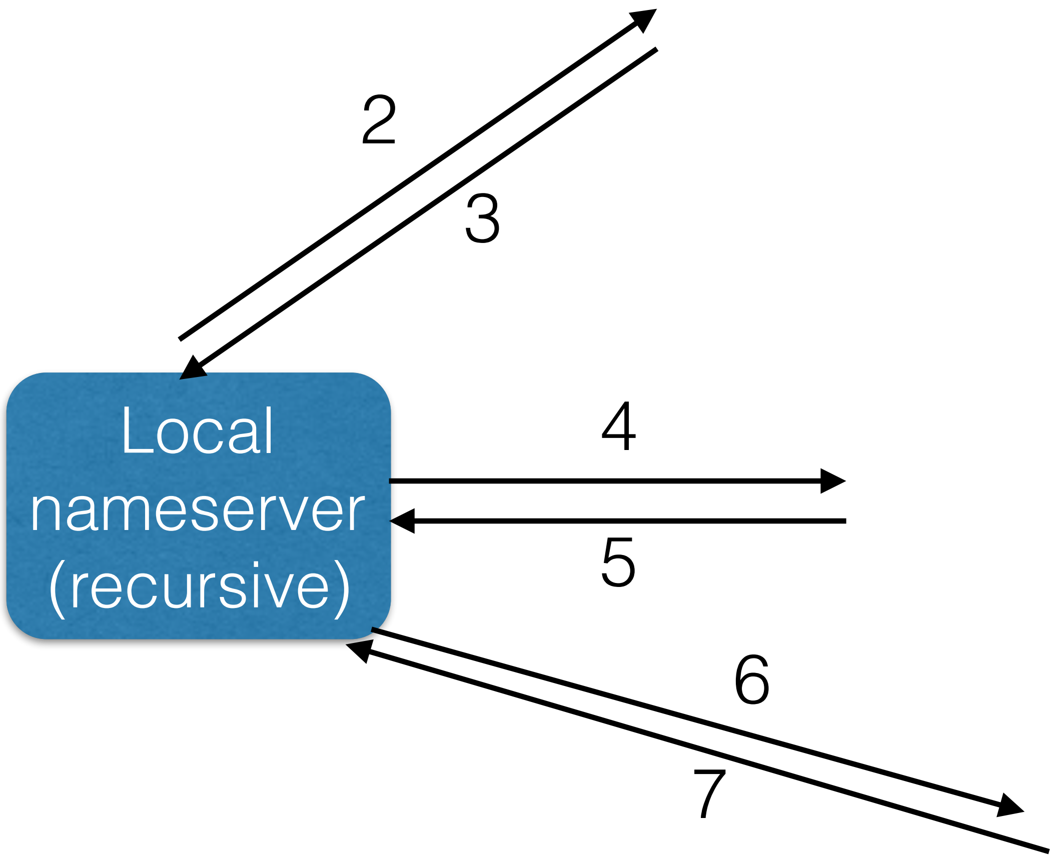
What’s in a response?

- Many things, but for the attacks we’re concerned with…

- A record: gives “the authoritative response for the IP address of this hostname”

- NS record: describes “this is the name of the nameserver who should know more about how to answer this query than I do”
  - Often also contains “glue” records (IP addresses of those name servers to avoid chicken and egg problems)
  - Resolver will generally cache all of this information

Query IDs



- The local resolver has a lot of incoming/outgoing queries at any point in time.

- To determine which response maps to which queries, it uses a *query ID*

- Query ID: 16-bit field in the DNS header
    - Requester sets it to whatever it wants
    - Responder must provide the same value in its response

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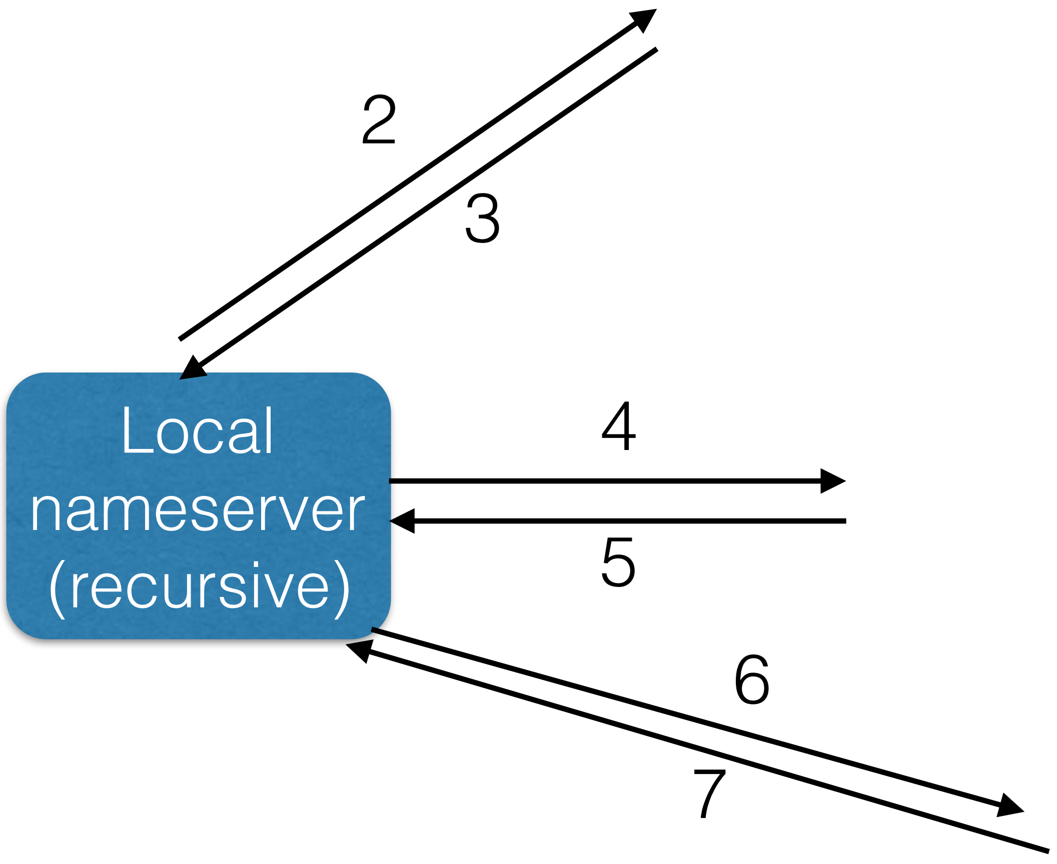

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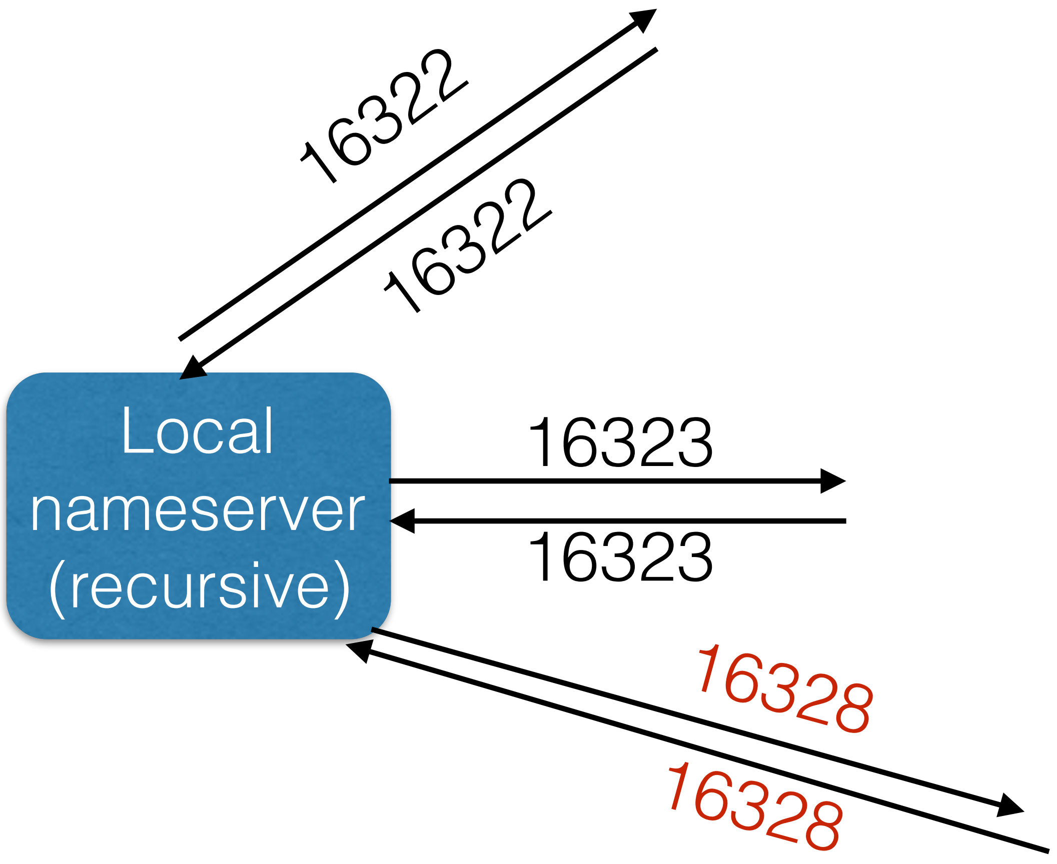
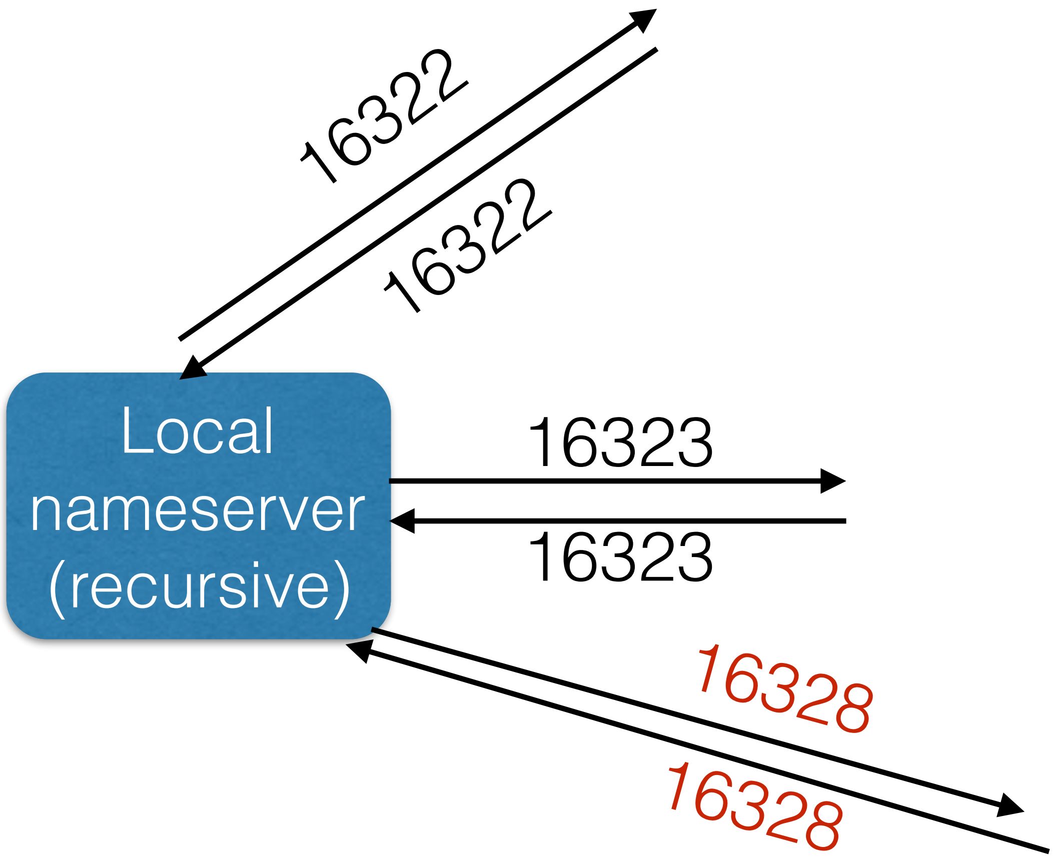
How would you implement query IDs at a resolver?

Query IDs used to increment



- Global query ID value
- Map outstanding query ID to local state of who to respond to (the client)
- Basically:
  new Packet(queryID++)

Query IDs used to increment



- Global query ID value
- Map outstanding query ID to local state of who to respond to (the client)
- Basically:
  new Packet(queryID++)

How would you attack this?

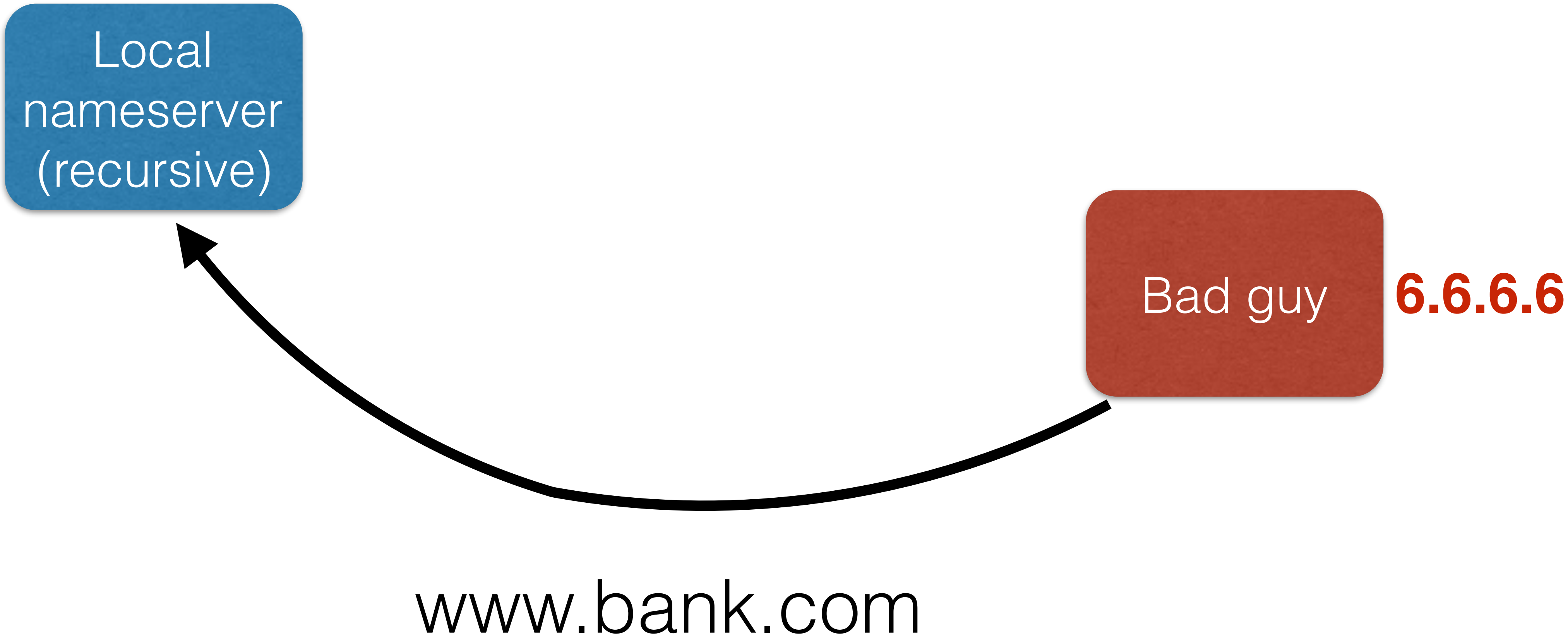
Cache poisoning

Local nameserver (recursive)

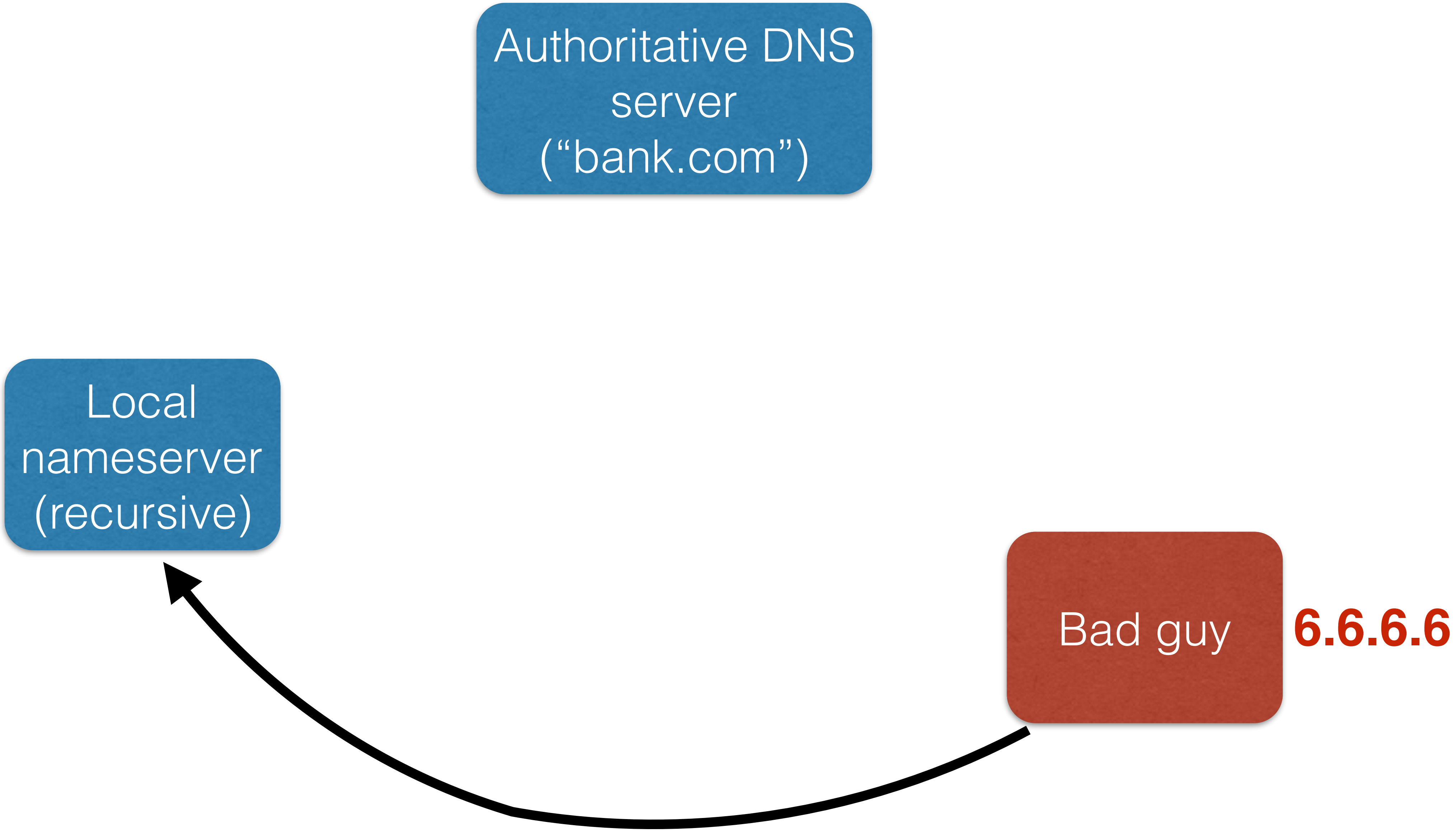
Bad guy

6.6.6.6

Cache poisoning

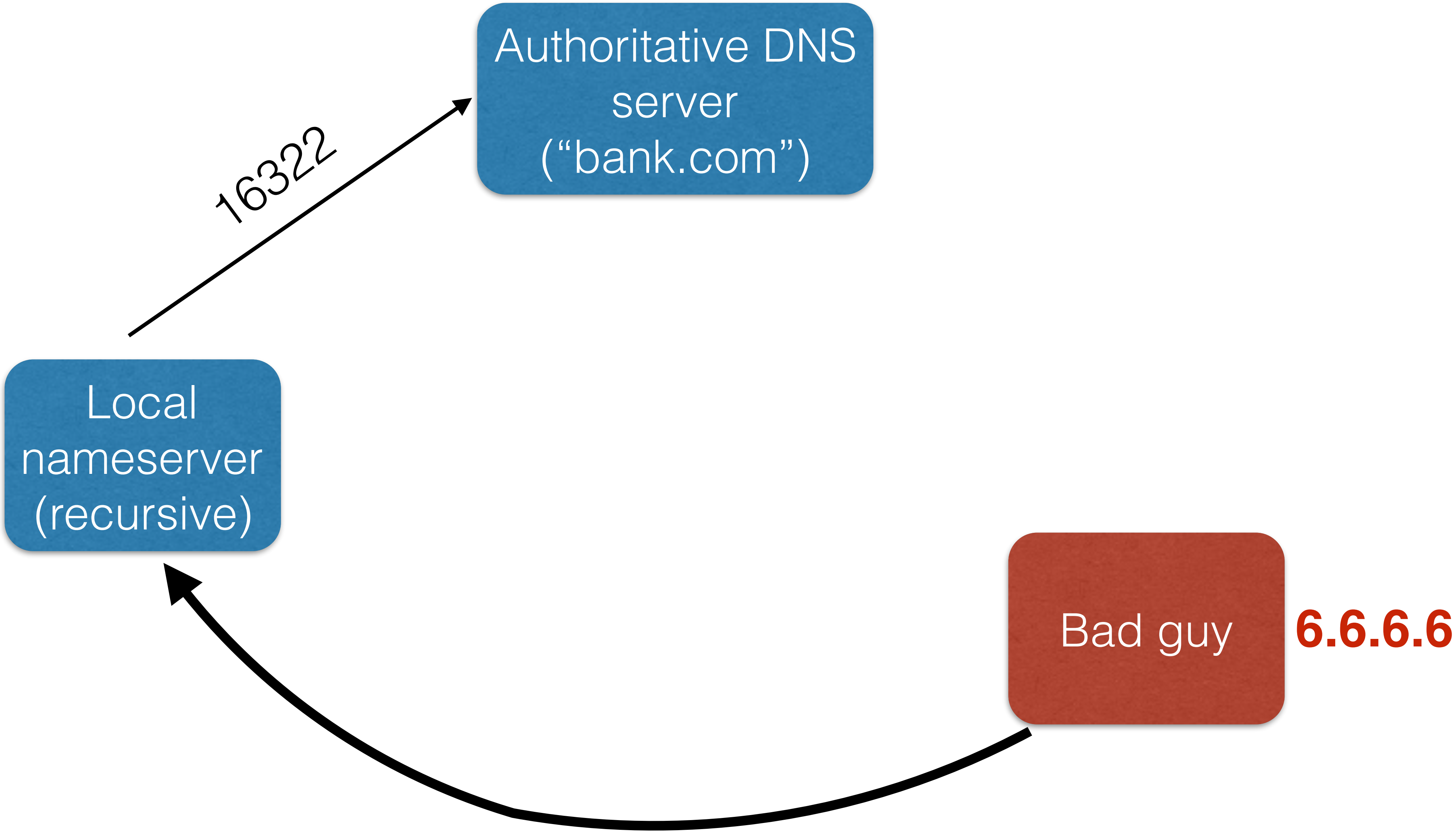


Cache poisoning


Authoritative DNS server ("bank.com")
Local nameserver (recursive)
Bad guy
6.6.6.6

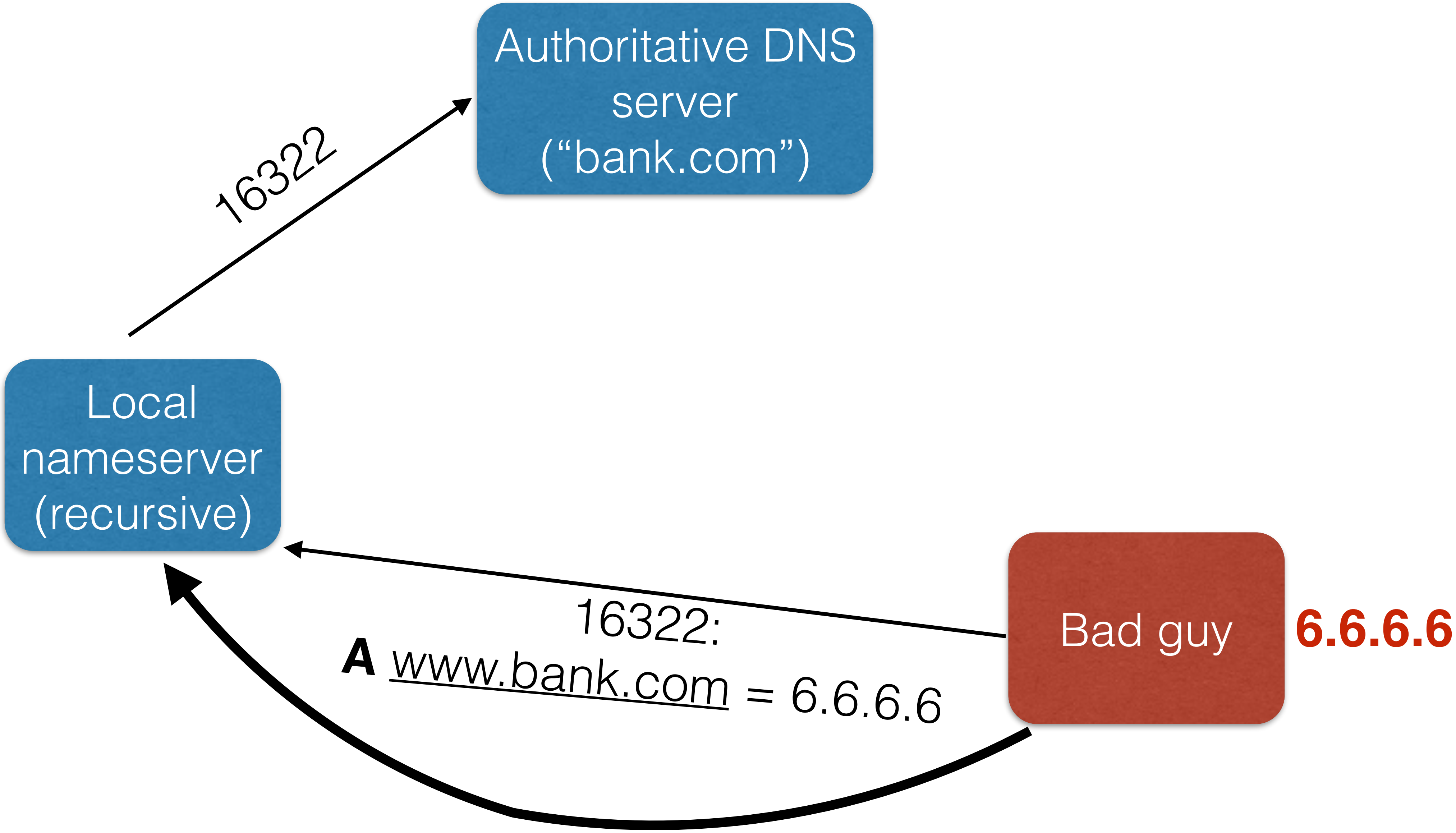
www.bank.com

Cache poisoning


Authoritative DNS server ("bank.com")
16322
Local nameserver (recursive)
Bad guy
6.6.6.6

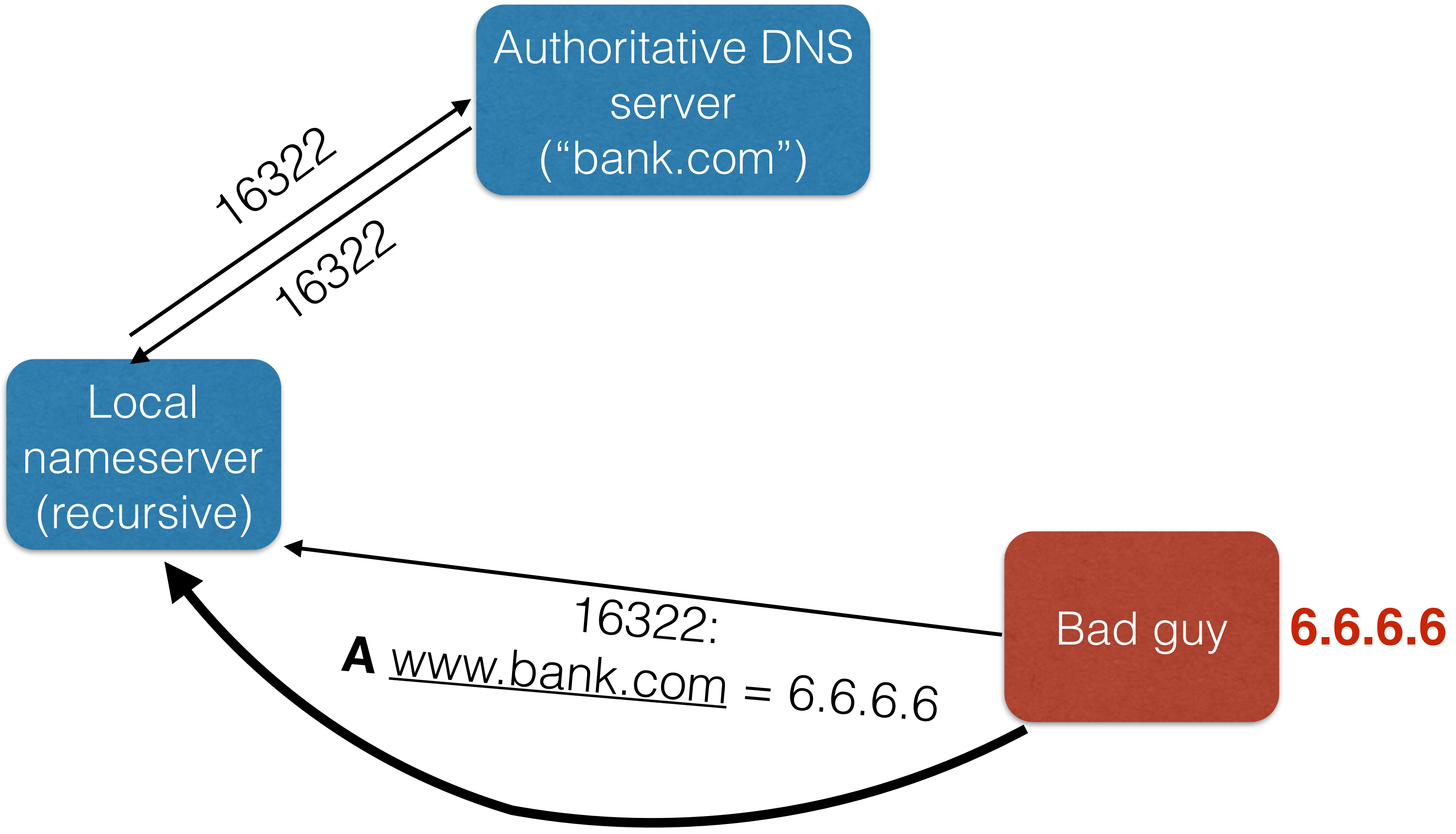
www.bank.com

Cache poisoning


Authoritative DNS server ("bank.com")
16322
Local nameserver (recursive)
16322:
A www.bank.com = 6.6.6.6
Bad guy
6.6.6.6

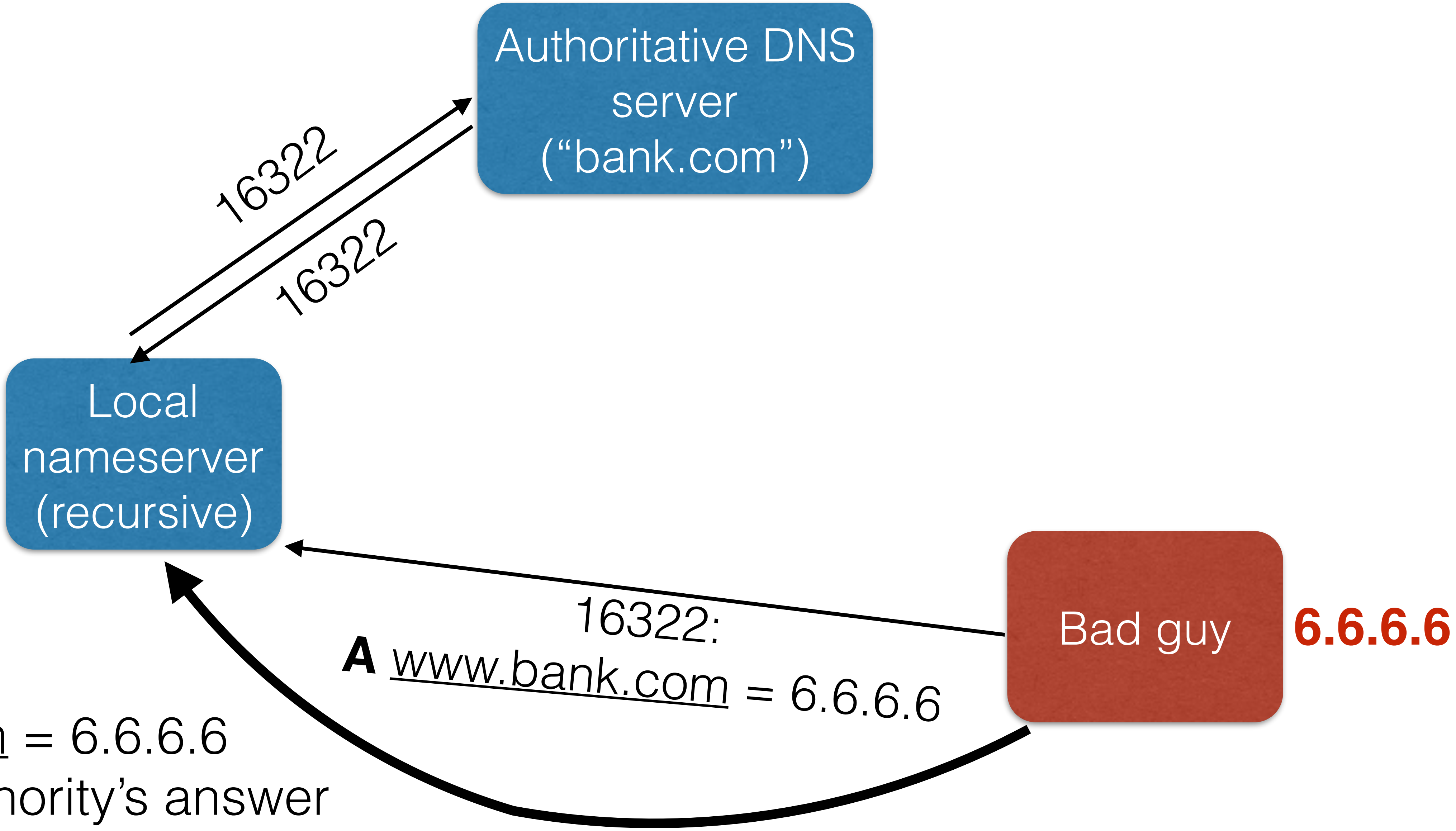
www.bank.com

Cache poisoning


Authoritative DNS server ("bank.com")
16322
16322
Local nameserver (recursive)
Bad guy
6.6.6.6
16322:
A www.bank.com = 6.6.6.6

www.bank.com

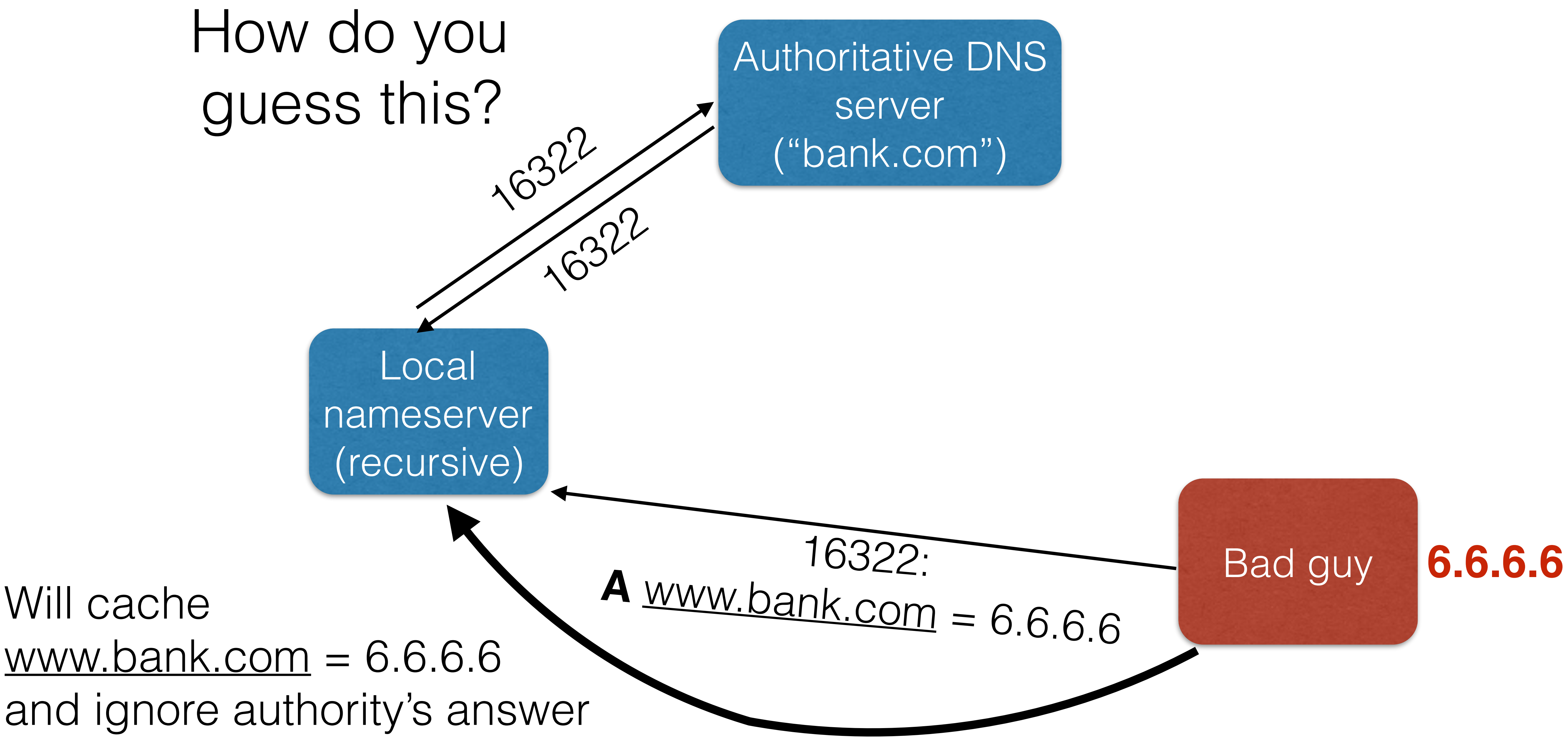
Cache poisoning


Authoritative DNS server ("bank.com")
16322
16322
Local nameserver (recursive)
16322:
A www.bank.com = 6.6.6.6
Bad guy
6.6.6.6

Will cache
www.bank.com = 6.6.6.6
and ignore authority's answer

www.bank.com

Cache poisoning



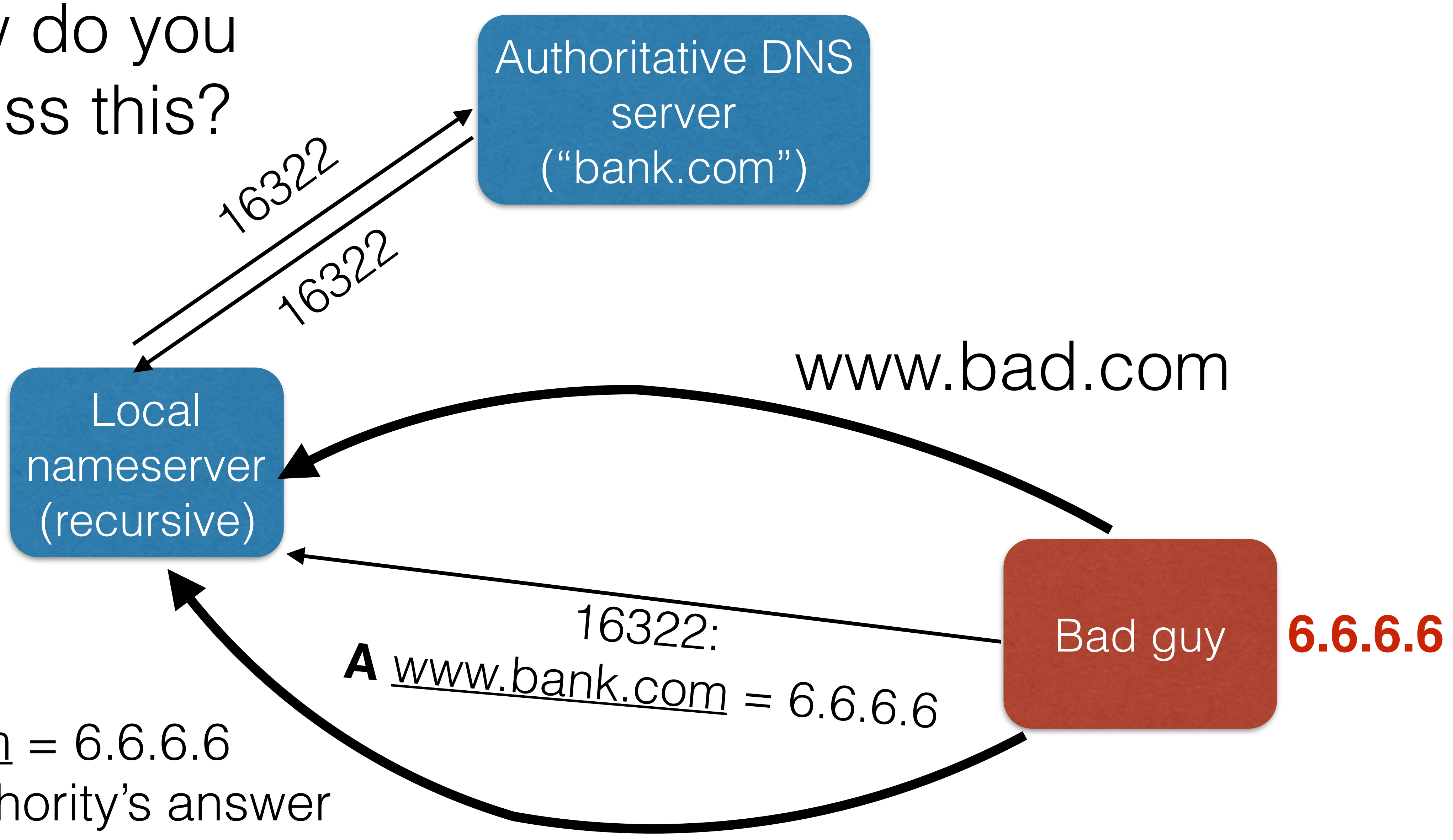
How do you guess this?

Authoritative DNS server ("bank.com")
16322
16322
Local nameserver (recursive)
Bad guy
6.6.6.6
16322:
A www.bank.com = 6.6.6.6

Will cache www.bank.com = 6.6.6.6 and ignore authority's answer

www.bank.com

Cache poisoning

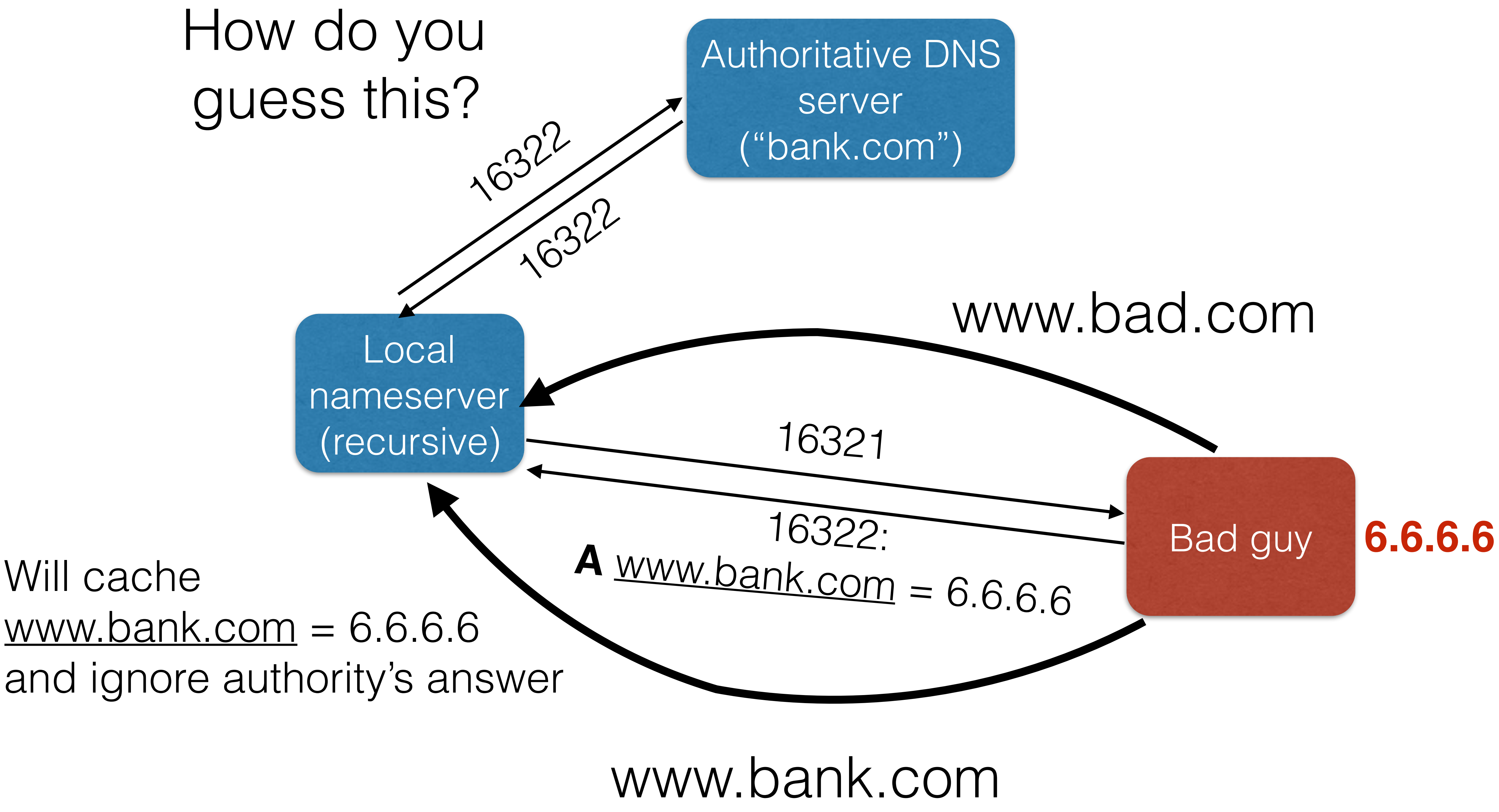

Authoritative DNS server ("bank.com")
16322
16322
Local nameserver (recursive)
www.bad.com
Bad guy
6.6.6.6
16322:
A www.bank.com = 6.6.6.6

How do you guess this?

Will cache
www.bank.com = 6.6.6.6
and ignore authority's answer

www.bank.com

Cache poisoning



How do you guess this?

Authoritative DNS server ("bank.com")

16322

16322

Local nameserver (recursive)

www.bad.com

16321

16322: **A** www.bank.com = 6.6.6.6

Bad guy

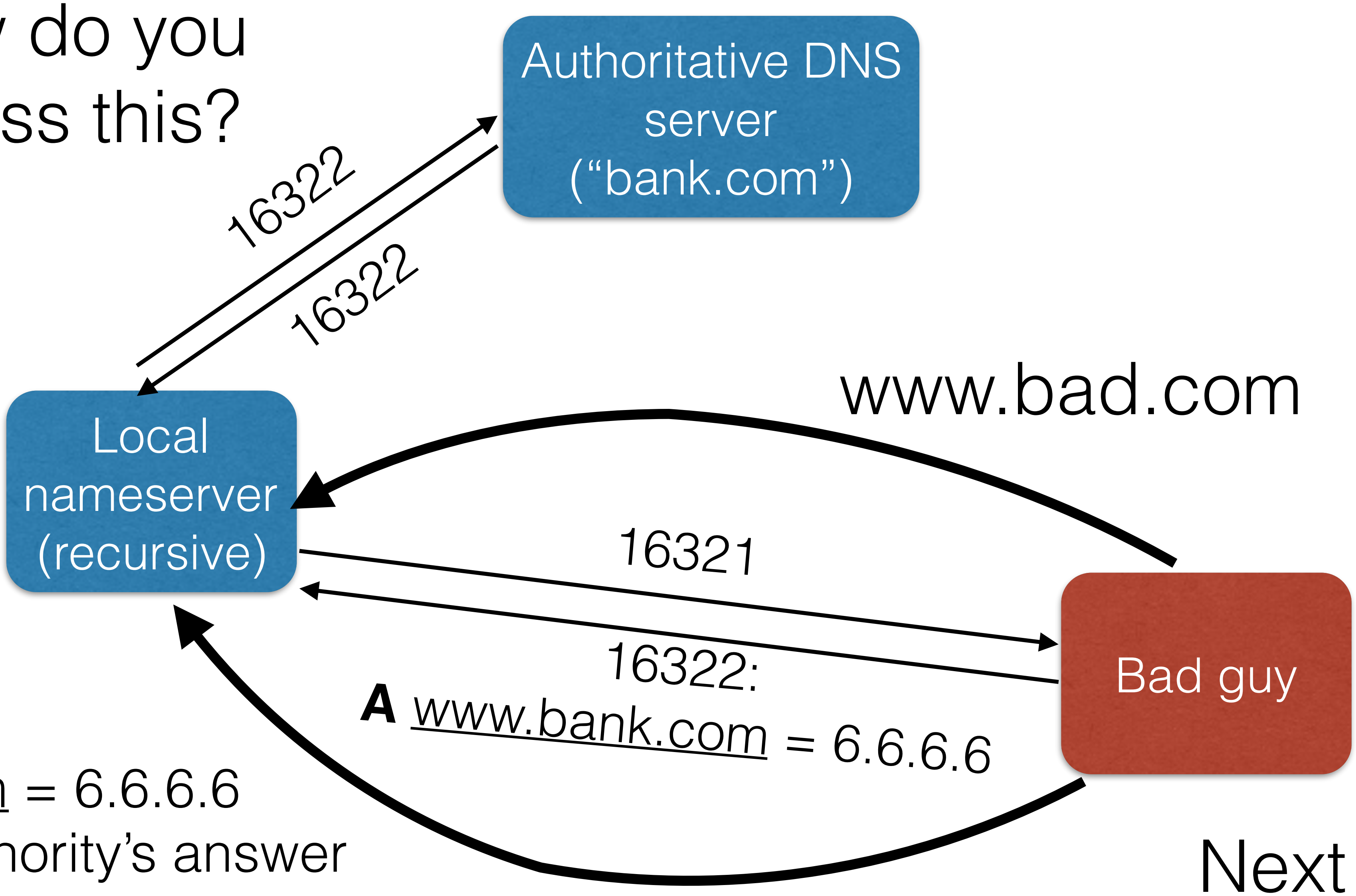
6.6.6.6

Will cache
www.bank.com = 6.6.6.6
and ignore authority's answer

www.bank.com

Cache poisoning

How do you guess this?



Will cache www.bank.com = 6.6.6.6 and ignore authority's answer

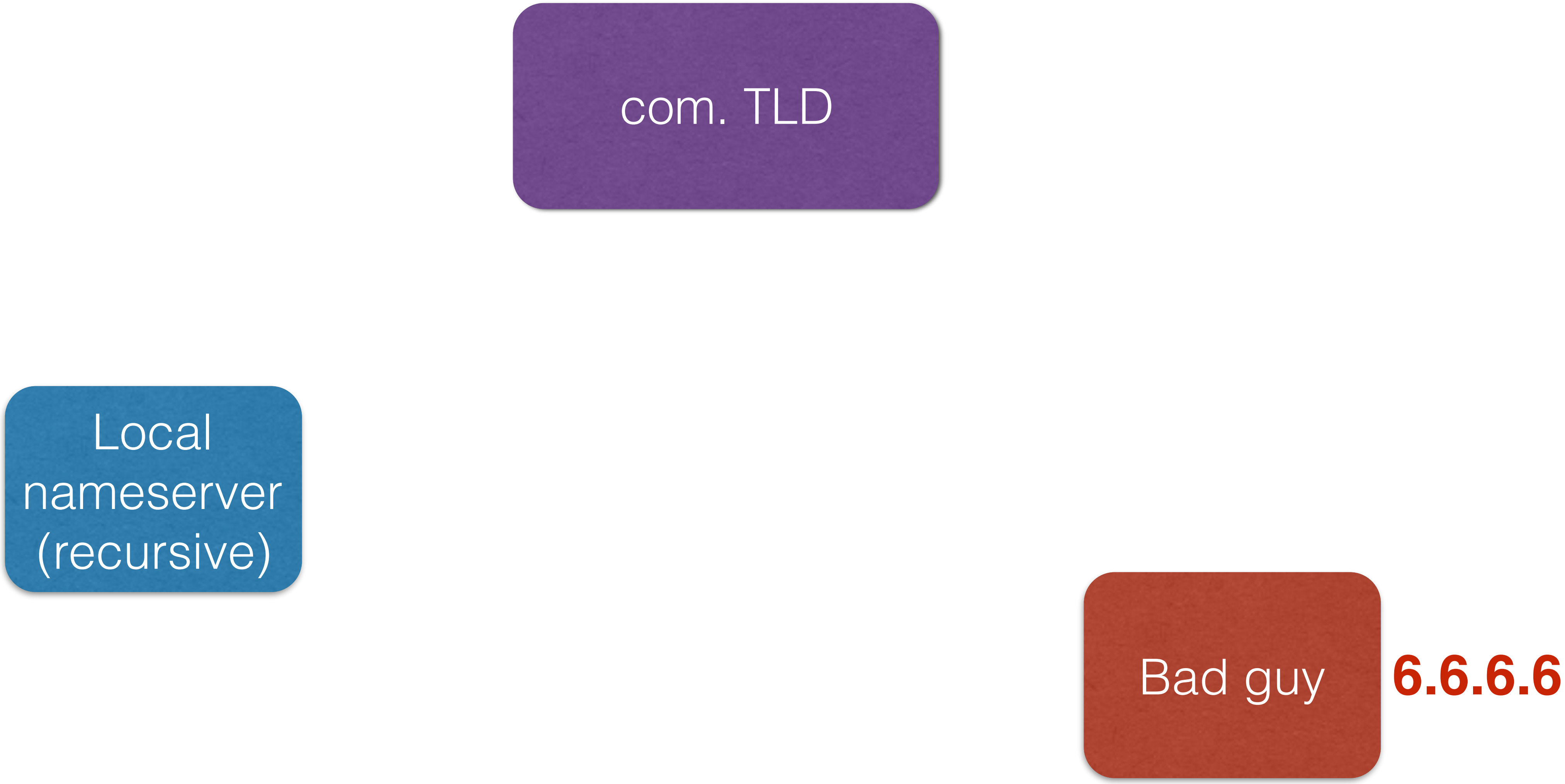
Next is likely 16322

Details of getting the attack to work

- Must guess query ID: ask for it, and go from there
  - Partial fix: randomize query IDs
  - Problem: small space
  - Attack: issue a Lot of query IDs
- Must guess source port number
  - Typically constant for a given server (often always 53)
- The answer must not already be in the cache
  - It will avoid issuing a query in the first place

Cache poisoning

Can we do more harm than a single record?



com. TLD

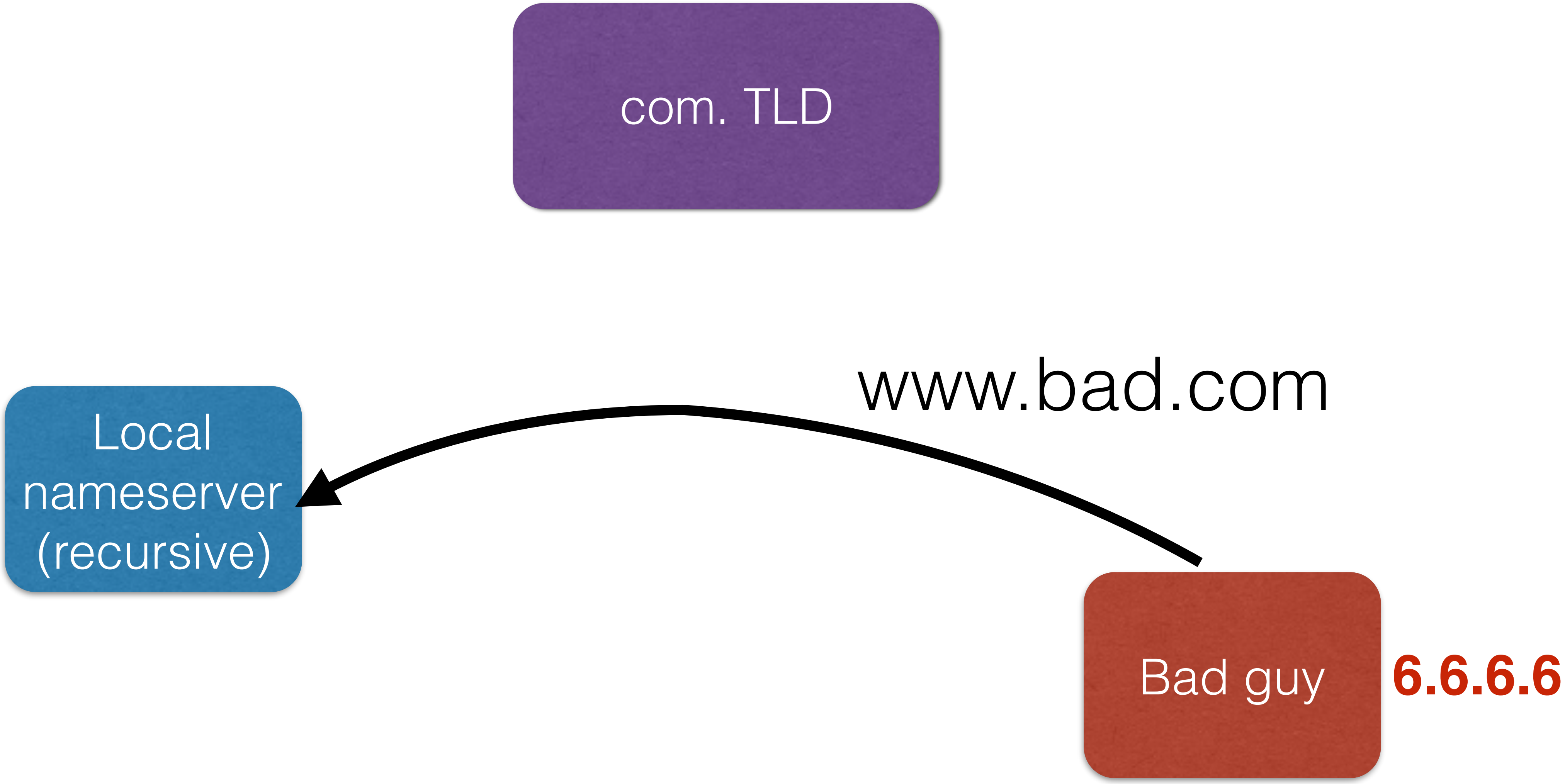
Local nameserver (recursive)

Bad guy

6.6.6.6

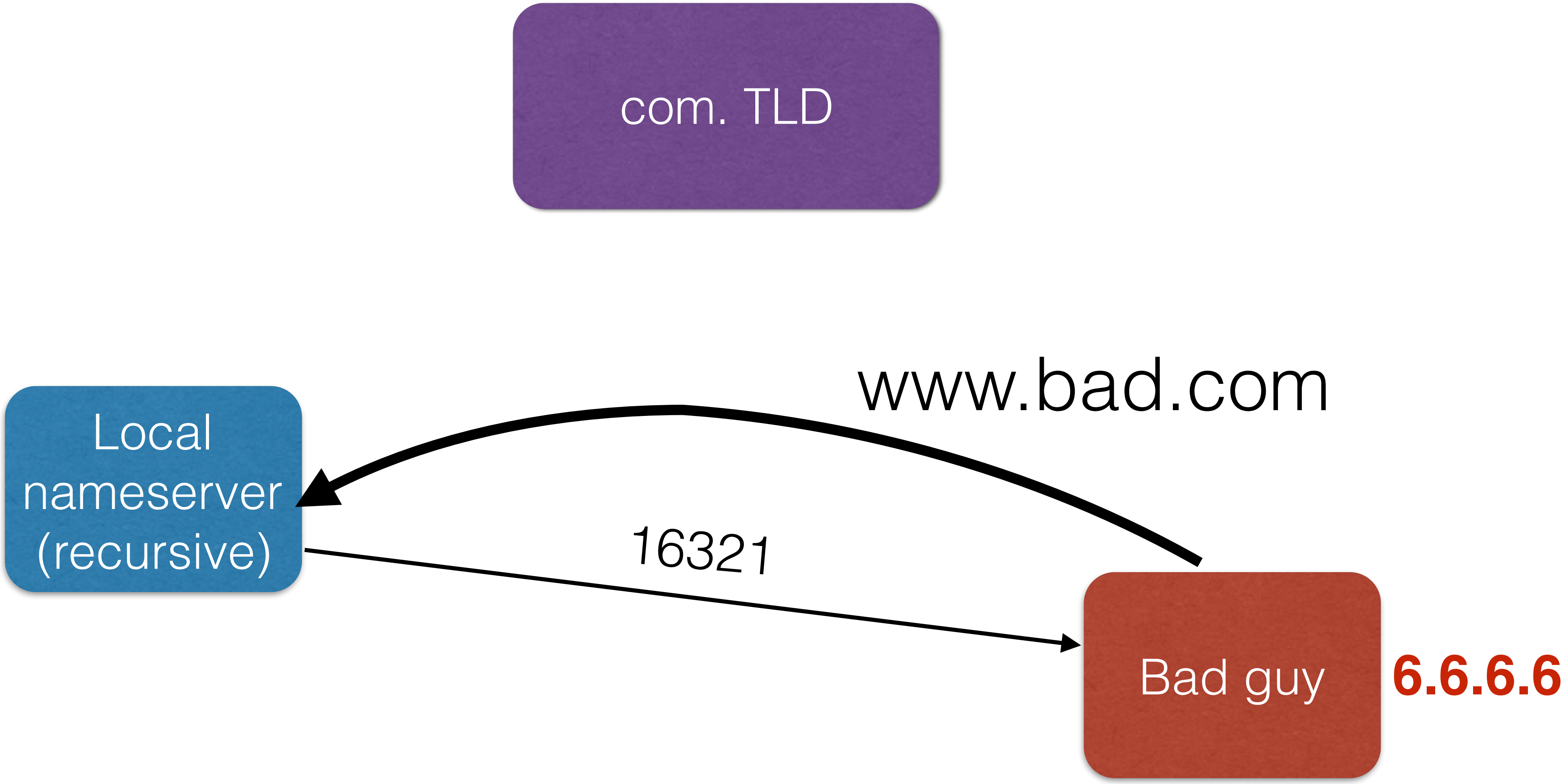
Cache poisoning

Can we do more harm than a single record?



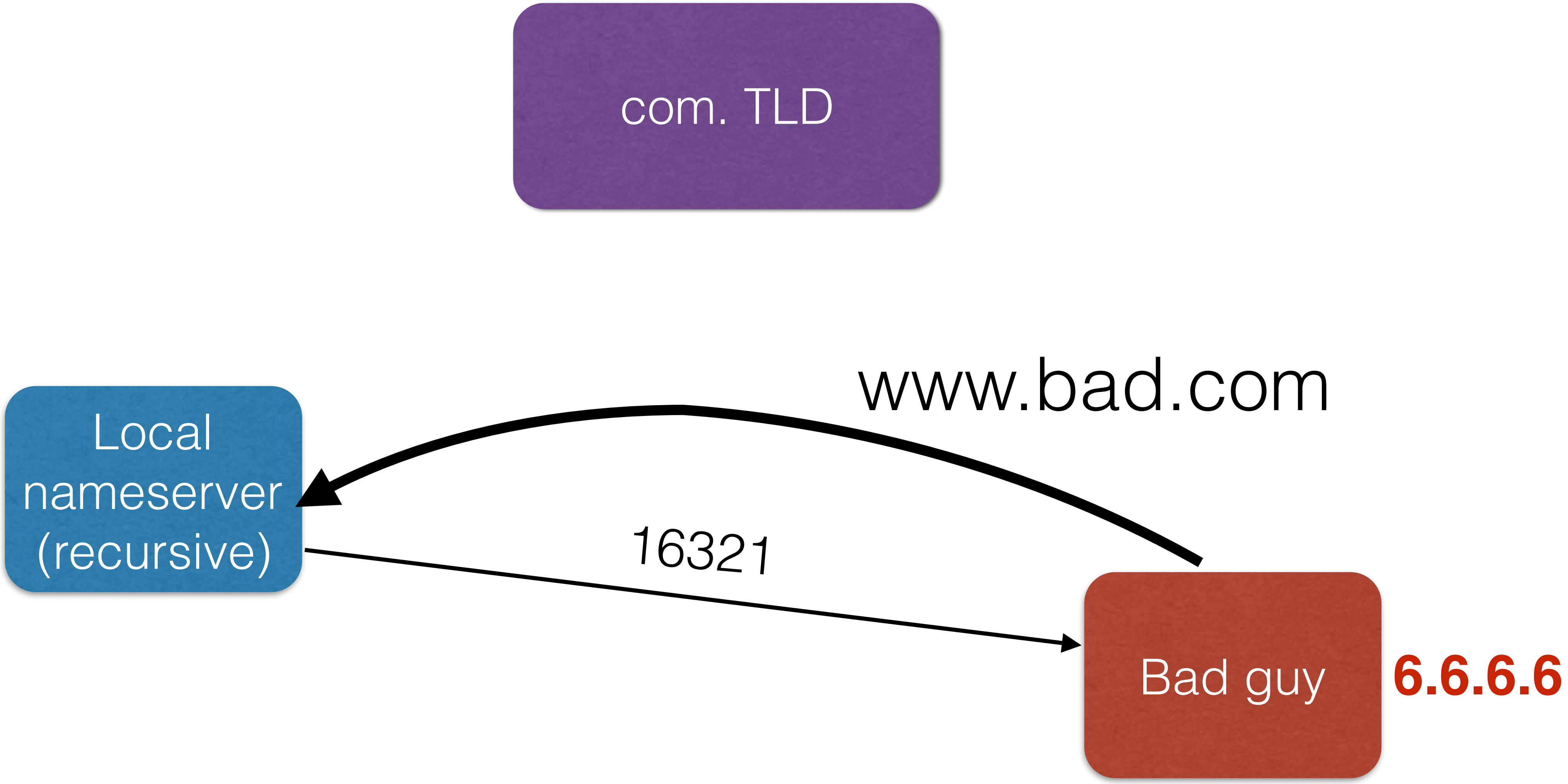
Cache poisoning

Can we do more harm than a single record?



Cache poisoning

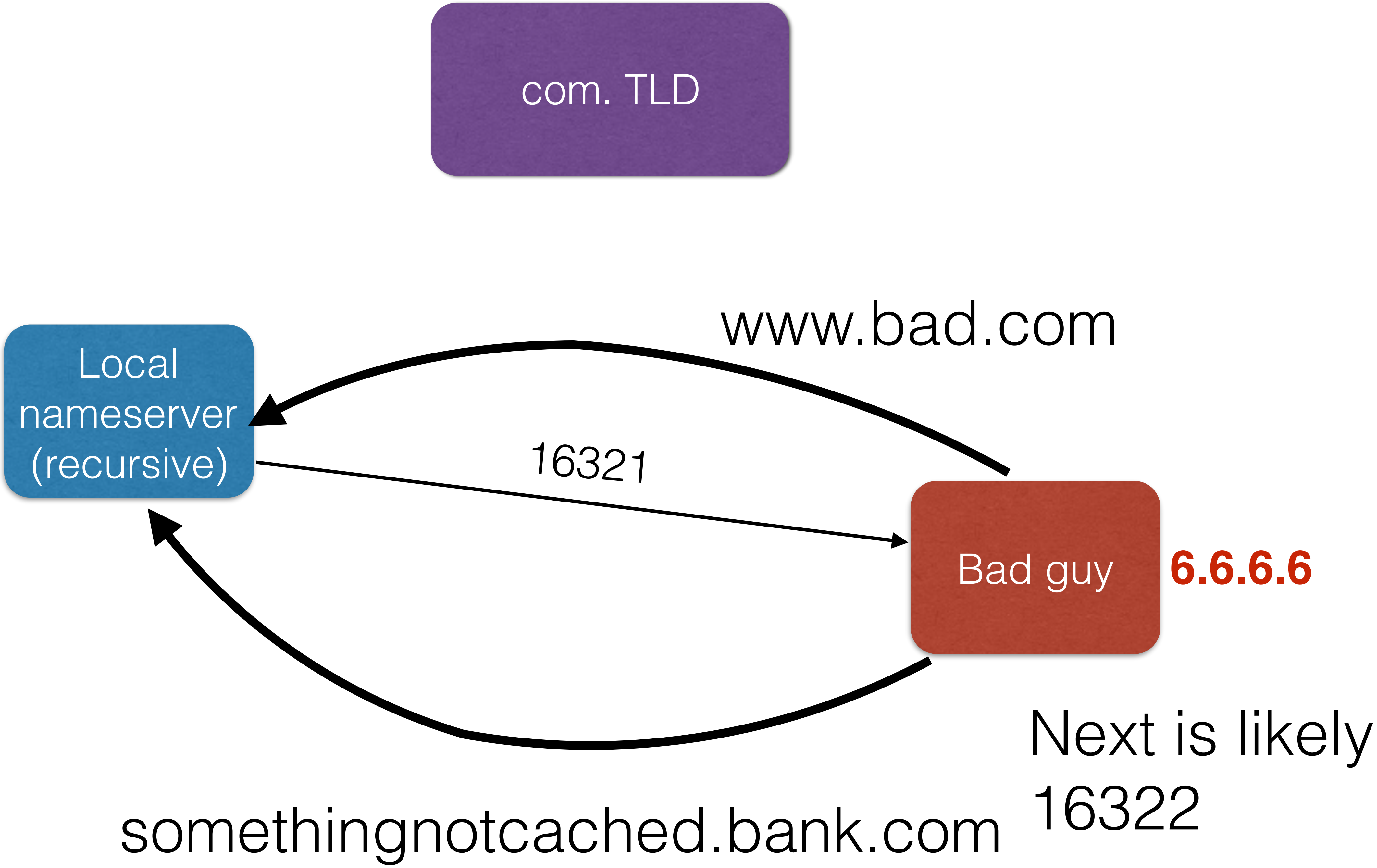
Can we do more harm than a single record?



Next is likely 16322

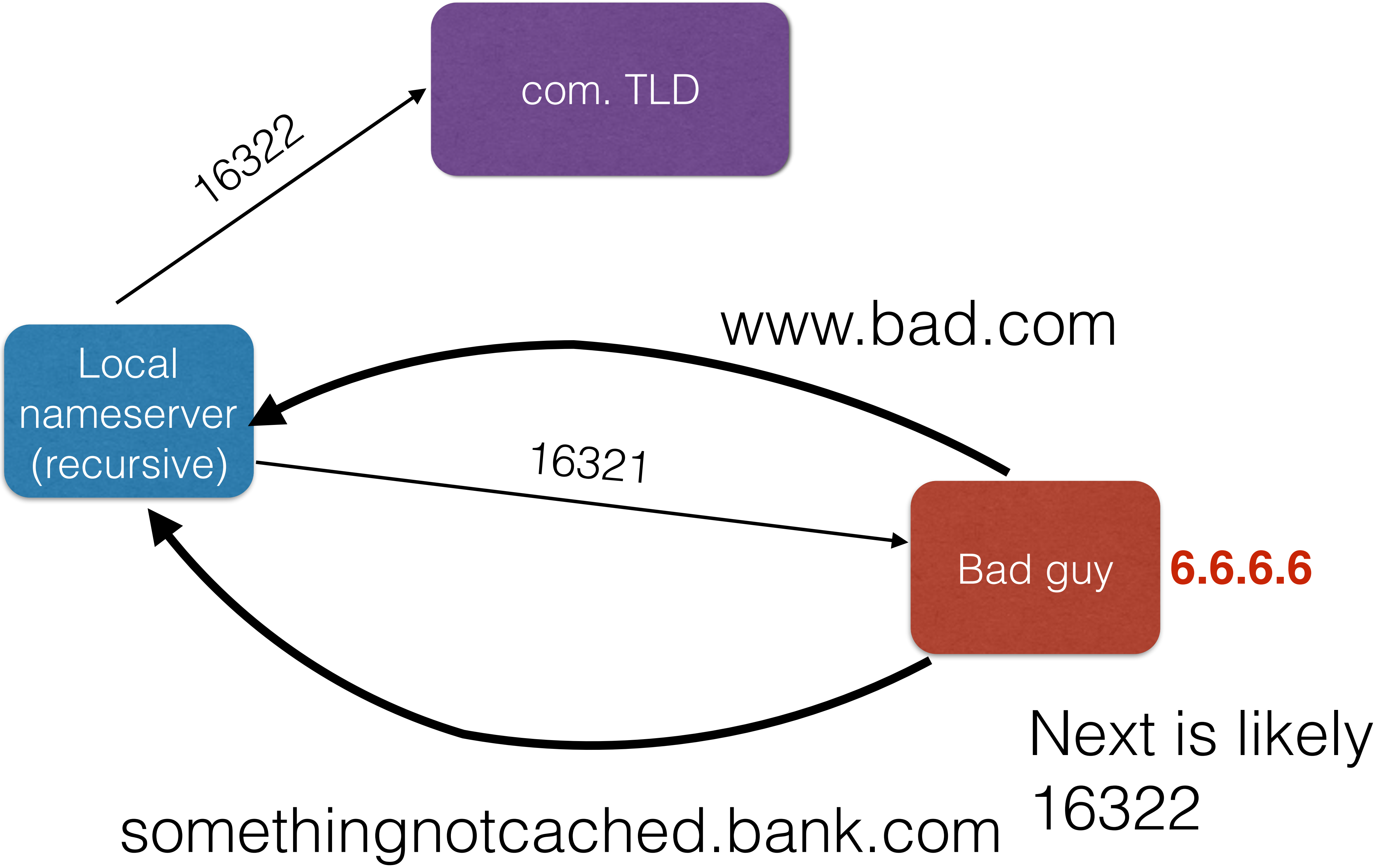
Cache poisoning

Can we do more harm than a single record?



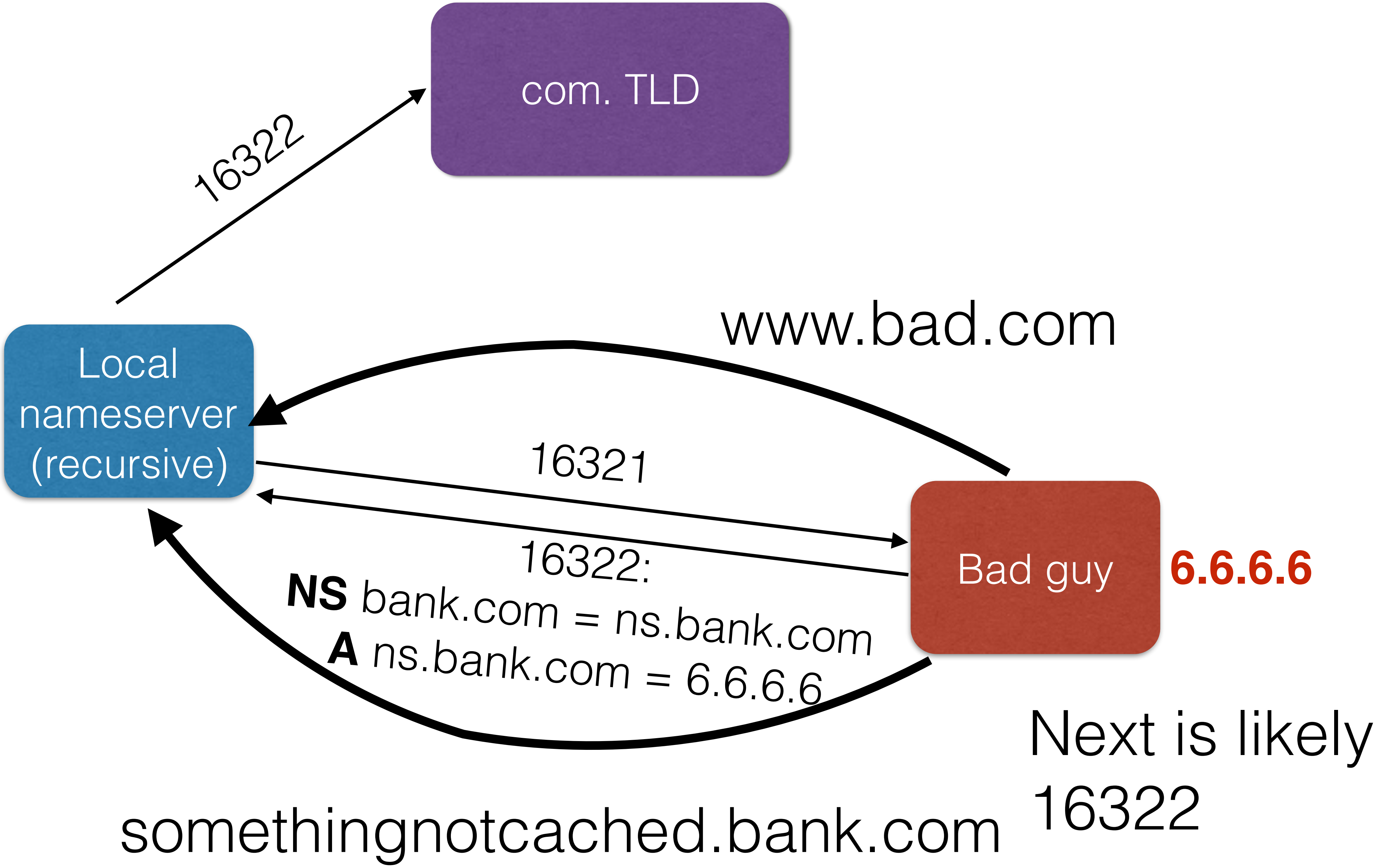
Cache poisoning

Can we do more harm than a single record?



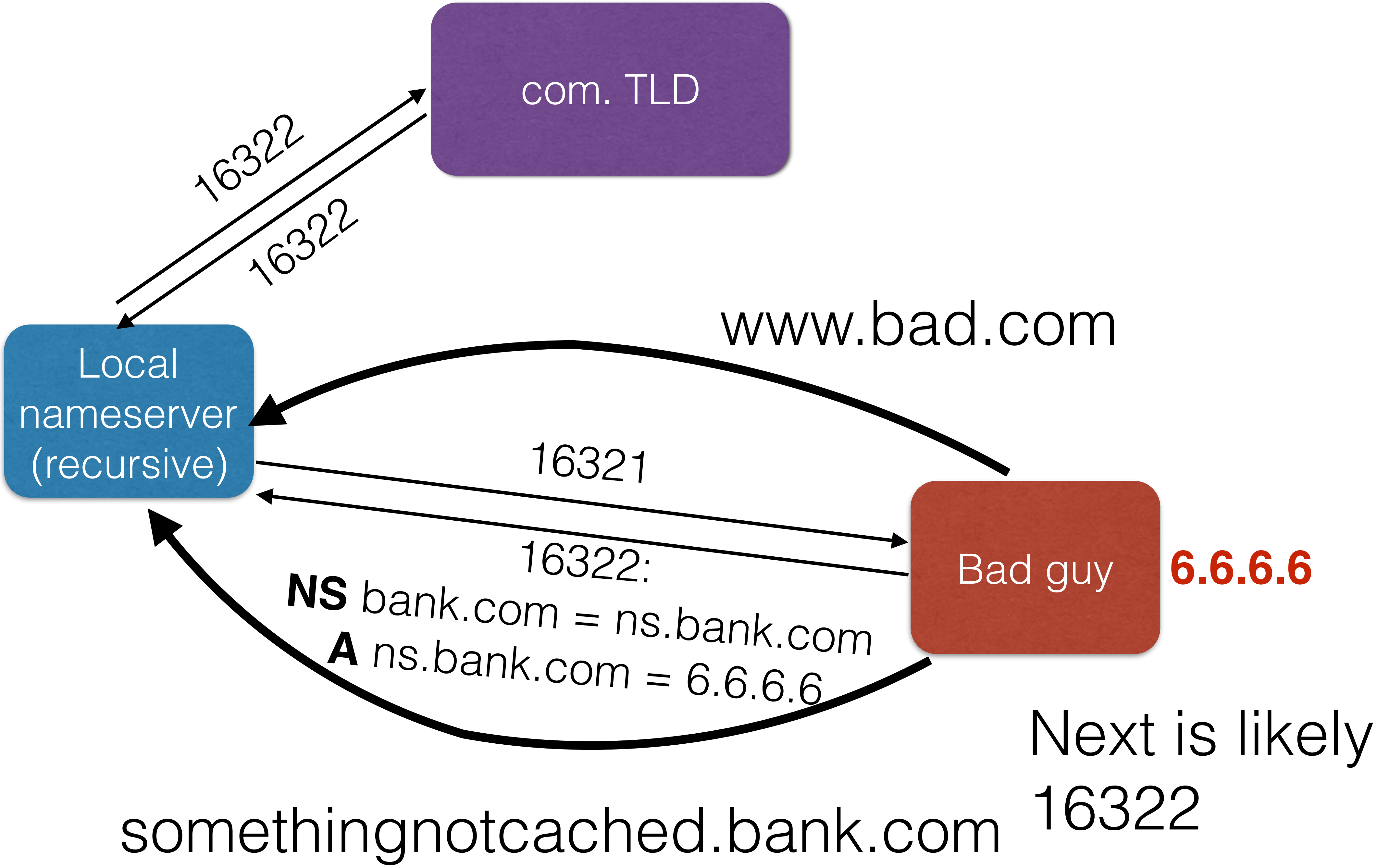
Cache poisoning

Can we do more harm than a single record?



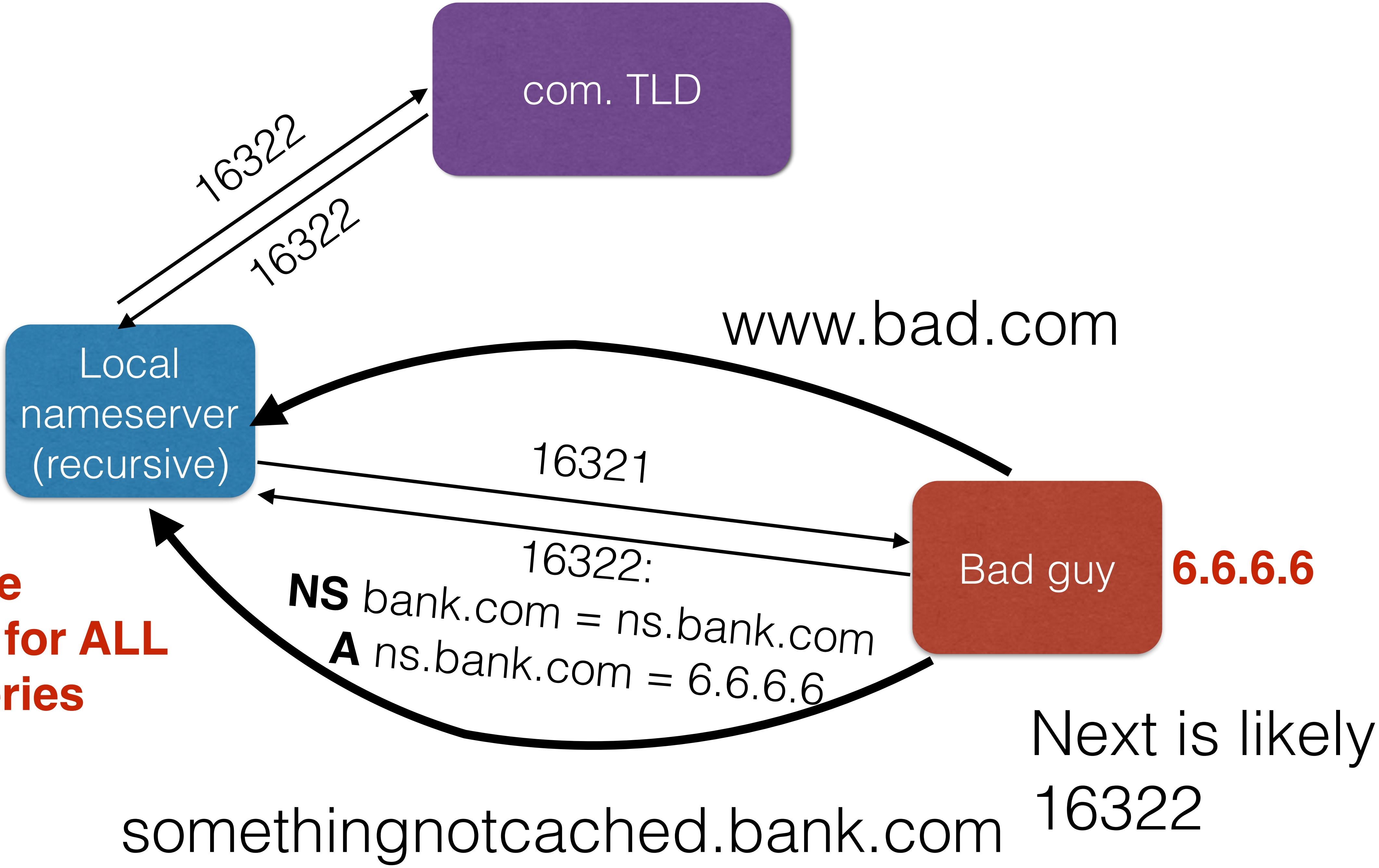
Cache poisoning

Can we do more harm than a single record?


com. TLD
16322
16322
Local nameserver (recursive)
www.bad.com
16321
16322:
NS bank.com = ns.bank.com
A ns.bank.com = 6.6.6.6
Bad guy
6.6.6.6
Next is likely 16322
somethingnotcached.bank.com

Cache poisoning

Can we do more harm than a single record?


com. TLD
16322
16322
Local nameserver (recursive)
www.bad.com
16321
16322:
NS bank.com = ns.bank.com
A ns.bank.com = 6.6.6.6
Bad guy
6.6.6.6
somethingnotcached.bank.com

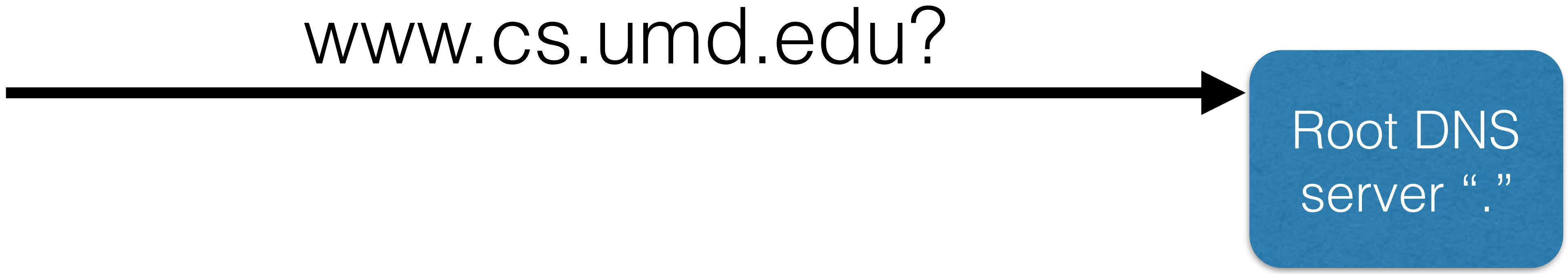
Next is likely 16322

Will cache "the person to ask for ALL bank.com queries is 6.6.6.6"

Solutions?

- Randomizing query ID?
  - Not sufficient alone: only 16 bits of entropy
- Randomize source port, as well
  - There’s no reason for it stay constant
  - Gets us another 16 bits of entropy
- DNSSEC?

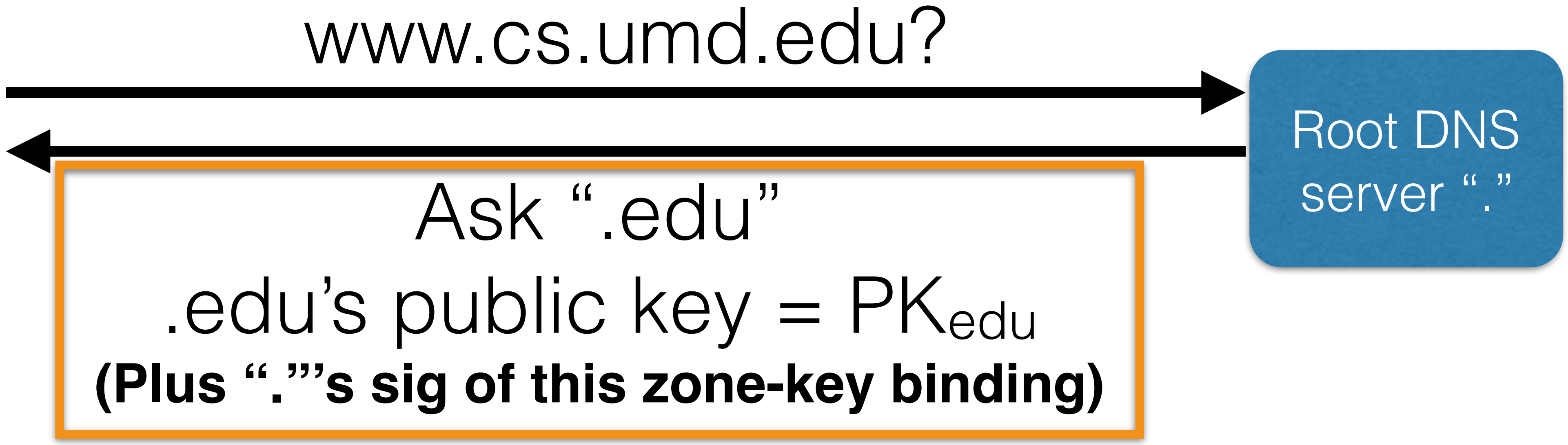
DNSSEC



www.cs.umd.edu?

Root DNS server “.”

DNSSEC



www.cs.umd.edu?

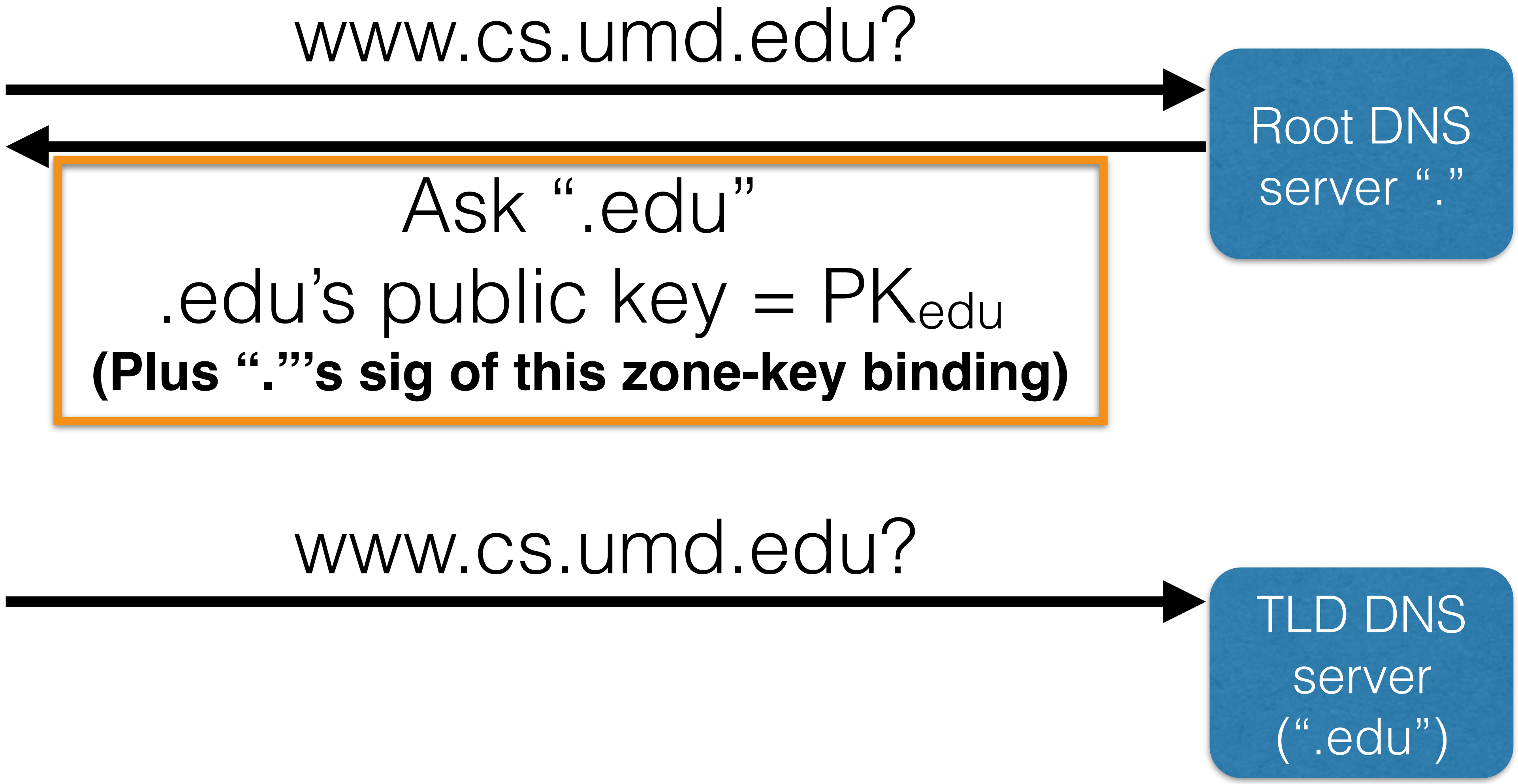
Root DNS server "."

Ask ".edu"

.edu's public key = PK_{edu}

(Plus ".""s sig of this zone-key binding)

DNSSEC



www.cs.umd.edu?

Root DNS server "."

Ask ".edu"

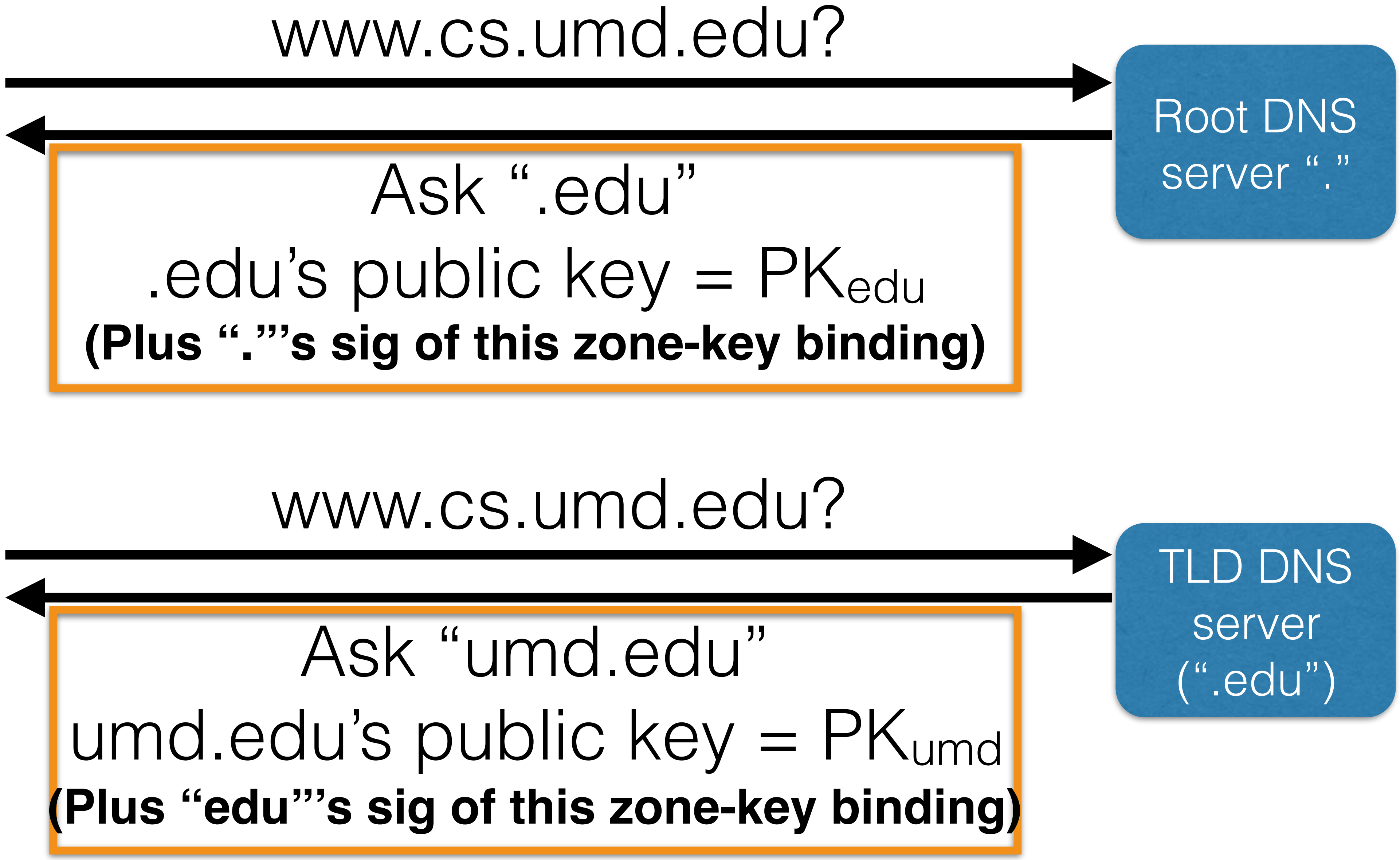
.edu's public key = PK_{edu}

(Plus "."'s sig of this zone-key binding)

www.cs.umd.edu?

TLD DNS server (".edu")

DNSSEC



www.cs.umd.edu?

Root DNS server “.”

Ask “.edu”

.edu’s public key = PK_{edu}

(Plus “.””s sig of this zone-key binding)

www.cs.umd.edu?

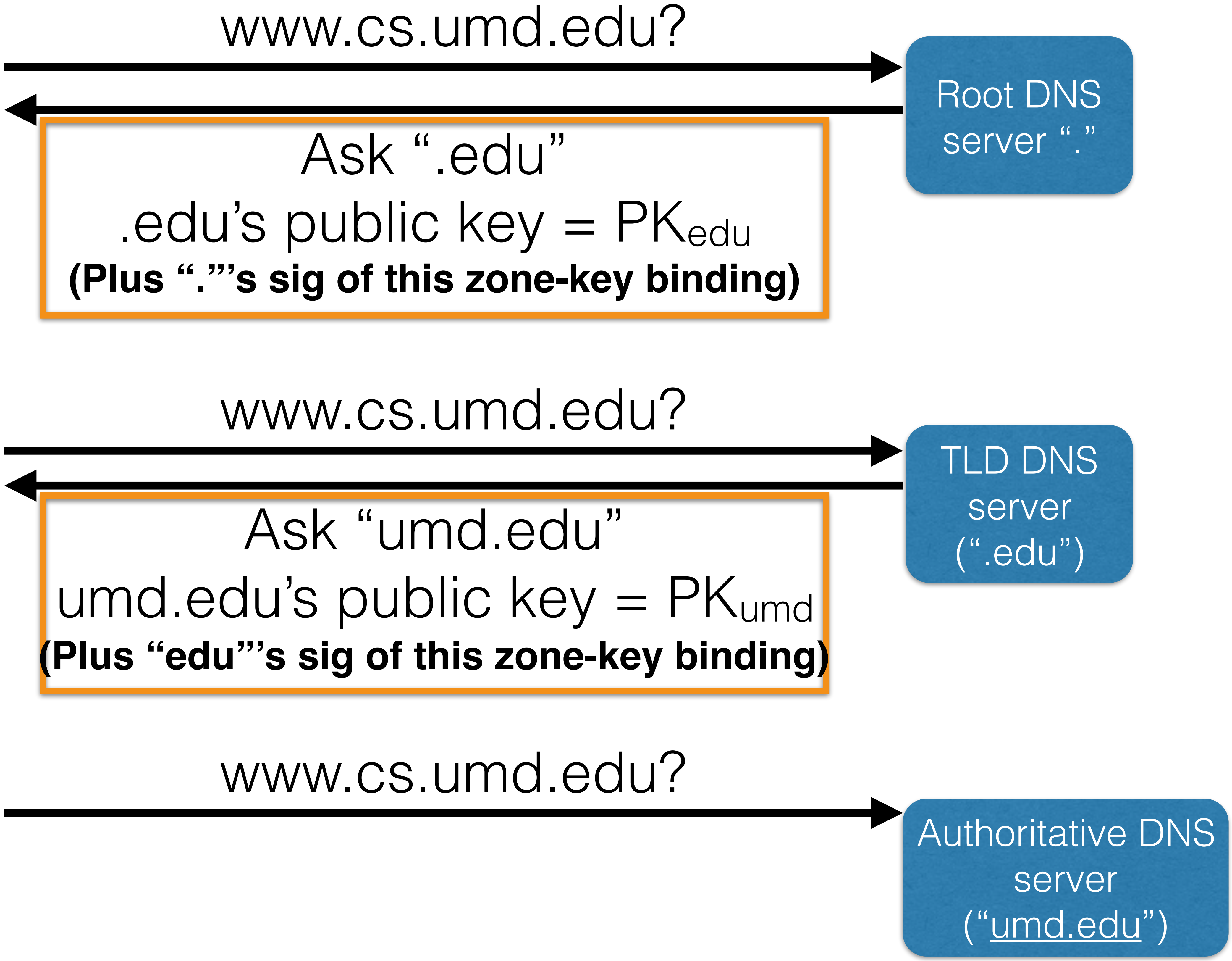
TLD DNS server (“.edu”)

Ask “umd.edu”

umd.edu’s public key = PK_{umd}

(Plus “edu””s sig of this zone-key binding)

DNSSEC



www.cs.umd.edu?

Root DNS server "."

Ask ".edu"

.edu's public key = PK_{edu}

(Plus ".""s sig of this zone-key binding)

www.cs.umd.edu?

TLD DNS server (".edu")

Ask "umd.edu"

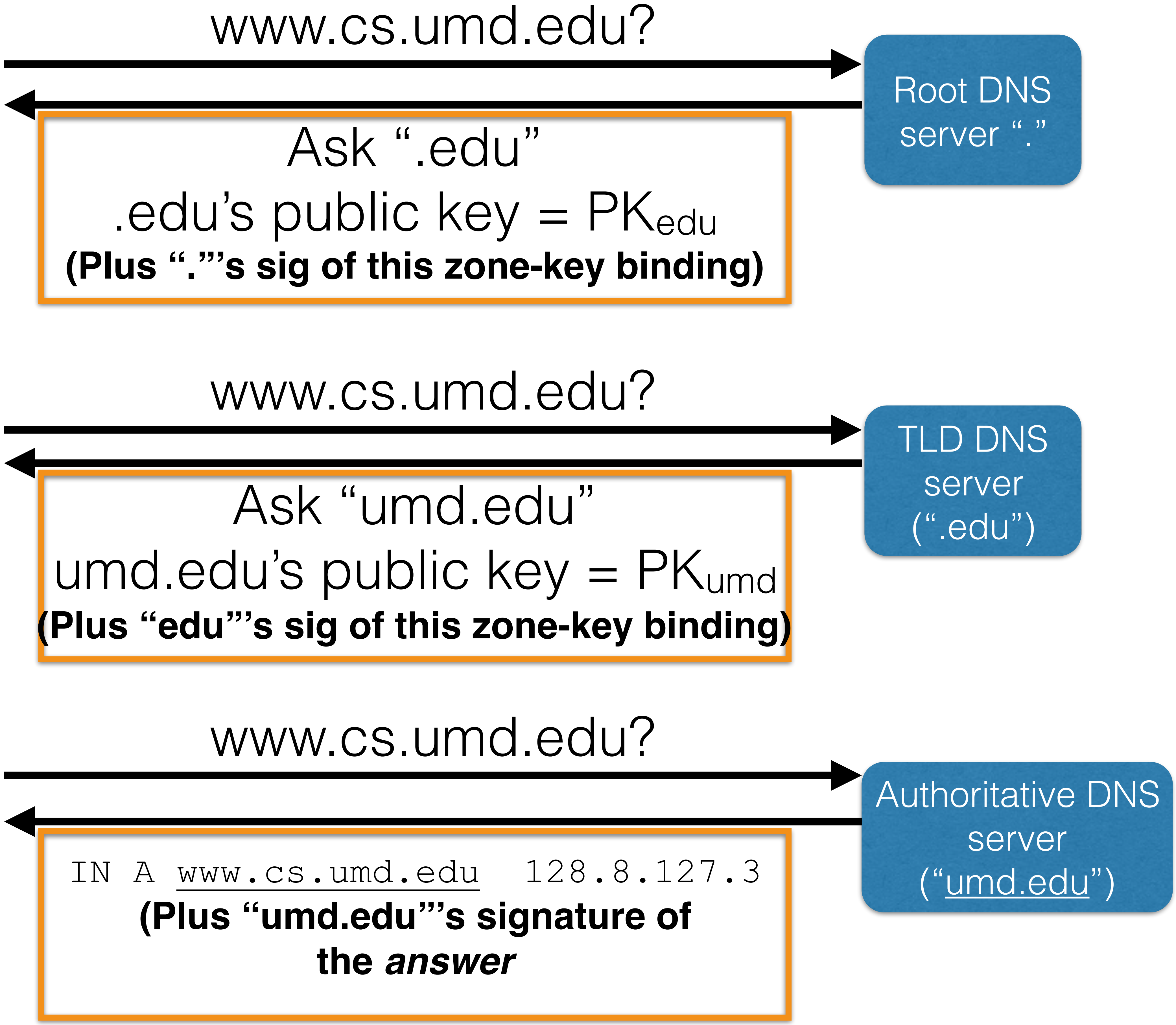
umd.edu's public key = PK_{umd}

(Plus "edu""s sig of this zone-key binding)

www.cs.umd.edu?

Authoritative DNS server ("umd.edu")

DNSSEC



www.cs.umd.edu?

Root DNS server "."

Ask ".edu"

.edu's public key = PK_{edu}

(Plus ".""s sig of this zone-key binding)

www.cs.umd.edu?

TLD DNS server (".edu")

Ask "umd.edu"

umd.edu's public key = PK_{umd}

(Plus "edu""s sig of this zone-key binding)

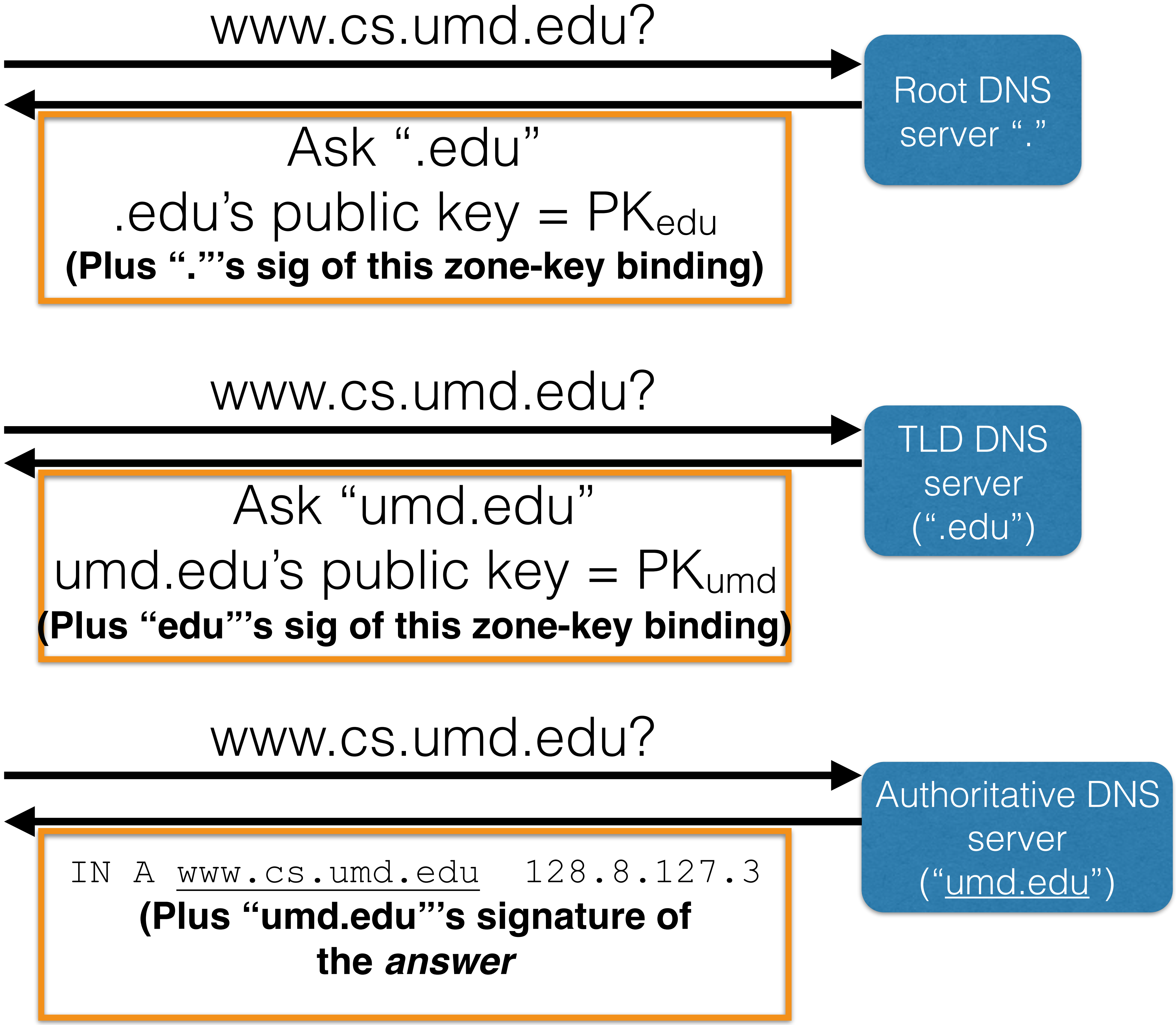
www.cs.umd.edu?

Authoritative DNS server ("umd.edu")

```
IN A www.cs.umd.edu  128.8.127.3
```

(Plus "umd.edu""s signature of the *answer*

DNSSEC



www.cs.umd.edu?

Root DNS server “.”

Ask “.edu”

.edu’s public key = PK_{edu}

(Plus “.”’s sig of this zone-key binding)

www.cs.umd.edu?

TLD DNS server (“.edu”)

Ask “umd.edu”

umd.edu’s public key = PK_{umd}

(Plus “edu”’s sig of this zone-key binding)

www.cs.umd.edu?

Authoritative DNS server (“umd.edu”)

```
IN A www.cs.umd.edu   128.8.127.3
```

(Plus “umd.edu”’s signature of the *answer*

Only the authoritative answer is signed

Properties of DNSSEC

- If everyone has deployed it, and if you know the root’s keys, then prevents spoofed responses
  - Very similar to PKIs in this sense
- But unlike PKIs, we still want authenticity despite the fact that not everyone has deployed DNSSEC
  - What if someone replies back without DNSSEC?
  - Ignore = secure but you can’t connect to a lot of hosts
  - Accept = can connect but insecure
- Back to our notion of incremental deployment
  - DNSSEC is not all that useful incrementally