Closed book. Closed notes. No electronic device.

1.

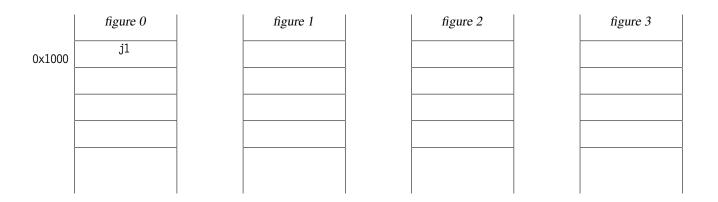
ASLR Buffer overflow Canary chroot() CodeRed CSRF Clickjacking Cookie Cookie elevation Document virus Effective uid Flash cookie Kerckhoff's principle Least privilege Metamorphic Mitigation Polymorphic Ransomware Real uid Rootkit Same origin policy seccomp Session elevation Setuid bit Session cookie seteuid() setuid() URL's scope Virus Web storage Worm X-Frame-Options XSS Stored XSS Reflected XSS

For each description below, give **one** term that *best* describes it. In many cases, but not all, the term will be in the table at left. In each case, give only **one** answer of **at most** 4 words (otherwise you get zero).

- 1. This determines whether the effective uid of a process changes when it starts executing a file.
- 2. This can change the real uid of a process.
- 3. This infects a running program.
- 4. This infects a stored program.
- 5. This detects stack smashing.
- 6. This kind of virus alters its code at each generation.
- 7. This kind of virus alters its memory image at each generation but its code does not change.
- 8. Virus contained in a macro of a Word file.
- 9. Virus contained in the autoplay file of a USB stick.
- 10. Information that a website gets concerning a client's platform.
- 11. State that is created at browsers by http headers.
- 12. Website-specific state at browsers that can be sent to websites only via scripts.
- 13. State that can be created by a browser and shared across other browsers running on the same OS.
- 14. This determines which cookies to include in a http request.
- 15. A website increases the privileges of an existing cookie.
- 16. A Linux system call that prevents the executing process from opening any more files.
- 17. This is a way to have framebusting without resorting to javascript.
- 18. A browser has pages a.com and b.com open. Page a.com sends a request with an attack script to b.com; the latter sends a response containing that script; the browser executes the script.
- 19. This is a way to limit the filesystem available to a process to a subtree of the original filesystem.

2

```
int f(int arg) {
                                                 Assume the Ubuntu environment of project 1.
    int j3;
    char buf[16];
                                                 This program is compiled without Stack Guard and executed
    FILE *badfile:
                                                 such that variable j1 is always at address 0x000001000.
s1: badfile = fopen("badfile", "r");
    fread(buf, sizeof(char), 40, badfile);
                                                 Below, "draw the stack layout" means indicate, in the referred
s2: return arg + 1;
                                                 figure, the contents of the stack from address 0x1000 to the top of
}
                                                 stack, and give the addresses (in hex) of the contents at the side.
                                                 Use "&g" to indicate the address of function g(). The stack grows
int g(int arg) { ... }
                                                 downward.
                                                 Below, "f() returns to g()" means that when control returns from
int main(int argc, char **argv) {
                                                 f(), it starts executing g() with any argument. And "f() returns
    int j1;
                                                 to g(4)" means that when control returns from f(), it starts exe-
    int j2[4];
                                                 cuting g() with an argument of 4.
    j1 = 123;
    return f(j1);
}
```



2a. In figure 0, draw the stack layout when control comes to s1 (i.e., "badfile = fopen(...)" is to be executed next).

2b Assume badfile is such that f() returns to g(). In *figure 1*, draw the stack layout when control comes to s2 (i.e., "return arg+1" is to be executed next). (Use "&g" to denote the address of function g().)

2c Assume badfile is such that f() returns to g(4). In *figure 2*, draw the stack layout when control comes to s2.

2d Does your answer for part 2a change if Address Randomization is turned on. Explain briefly.

2e Does your answer for part 2a change if the stack is made Non-Executable. Explain briefly.

2f Now assume the program is compiled with Stack Guard. In *figure 3*, draw the stack layout when control comes to s1. Can file badfile be initialized so that func() returns to g(). Explain briefly.

- **3** Here are two files owned by root:
 - /passwd.txt: text file that contains user passwords. Root has read-write-execute access. All other users have no access.
 - /chpwd: executable file that users can run to change their passwords. Root has read-write-execute access. All other users have write-execute access. The setuid bit is set (so it is a set-root-uid file).

Does this configuration allow an ordinary (i.e., non-root) user to delete passwd.txt? If no, explain briefly.

If yes, briefly give the steps of the attack.

- 4 Website a. com has an SQL table Users with a row for every user and columns Name, Pwd, Age and others.
 - A user changes its password with the GET request http://a.com/chpwd.php?a1 = <name>&a2 = <opwd>&a3 = <npwd>.

The server handles this with

chpwd.php:

\$db->sql_query("UPDATE Users SET Pwd='\$npwd' WHERE (Name='\$name' AND Pwd='\$opwd');");

• A user gets the names of all users with the GET request http://a.com/getusers.php

• Among the users are Ted and Bob (these are their Name entries). Ted does not know Bob's Pwd value.

3a. Give the path of a GET request that Ted can issue in order to change Bob's password to fqr123.

3b. Rewrite the chpwd code to use the *prepare* construct.

- **5** This problem concerns a browser c1, website s1, and attacker website s2.
 - c1 clicks http://s1/p1.html. In the response, s1 sets a cookie for domain s1.
 - Then c1 clicks http://s2/p2.html.

.

• Then p1.html regularly issues POST requests to s1. Each POST request contains the cookie value in its data. The server treats a request as valid iff the cookie value (in the request header) matches the value in the data.

For each of the following cases, answer whether p2.html can send a POST request to s1 that the latter treats as valid. Write "YES" if it can, and "NO" if it cannot. (Below, "unguessable" is equivalent to "randomly generated".)

s1-cookie name	s1-cookie value	your answer
guessable	guessable	
guessable	unguessable	
unguessable	guessable	
unguessable	unguessable	

6 When deposit() is called with user pointing to a valid User instance, it should update the user's balance or return 0 upon error. Does the above function achieve that. If not, fix the function.

```
struct User {
    unsigned int balance;
};
unsigned int deposit(struct User *user, unsigned int amount) {
    user->balance += amount;
    return user->balance;
}
```

7 The following function is in a program owned by root and executable by any user. ustr points to a string supplied by the user. secret points to a string that is stored within the program and should not be disclosed to the user. The function prints a message indicating the result of strcmp(ustr, secret). Can the function expose the secret string. If so, fix the function.

```
int f(char* ustr, char* secret) {
    printf(ustr);
    if (strcmp(ustr, secret) <= 0)
        printf(" <= secret");
    else
        printf(" > secret");
}
```