CMSC 330: Organization of Programming Languages

Tail Recursion and Continuation Passing Style (CPS)

Reverse

Pushes a stack frame on each recursive call

```
rev [1;2;3]

→ (rev [2;3]) @ [1]

→ ((rev [3]) @ [2]) @ [1]

→ (((rev []) @ [3]) @ [2]) @ [1]

→ (([] @ [3]) @ [2]) @ [1]

→ ([3] @ [2]) @ [1]

→ [3;2] @ [1]

→ [3;2;1]
```

A Clever Version of Reverse

No need to push a frame for each call!

```
rev [1;2;3] →
rev_helper [1;2;3] [] →
rev_helper [2;3] [1] →
rev_helper [3] [2;1] →
rev_helper [] [3;2;1] →
[3;2;1]
```

Tail Recursion

- Whenever a function ends with a recursive call, it is called tail recursive
 - Its "tail" is recursive
- Tail recursive functions can be implemented without requiring a stack frame for each call
 - No intermediate variables need to be saved, so the compiler overwrites them
- Typical pattern is to use an accumulator to build up the result, and return it in the base case

Compare rev and rev_helper

```
let rec rev 1 =
   match 1 with
   [] -> []
   | (x::xs) -> (rev xs) @ [x]
```

Waits for recursive call's result to compute final result

```
let rec rev_helper l a =
  match l with
  [] -> a
  | (x::xs) -> rev_helper xs (x::a)
```

final result is the result of the recursive call

True/false: map is tail-recursive.

A. True

B. False

True/false: map is tail-recursive.

True/false: fold is tail-recursive

```
let rec fold f a = function
[] -> a
| (h::t) -> fold f (f a h) t
```

A. True

B. False

True/false: fold is tail-recursive

```
let rec fold f a = function
  [] -> a
  [ (h::t) -> fold f (f a h) t
```

A. True

B. False

True/false: fold_right is tail-recursive

```
let rec fold_right f l a =
  match l with
  [] -> a
  | (h::t) -> f h (fold_right f t a)
```

True/false: fold_right is tail-recursive

```
let rec fold_right f l a =
  match l with
  [] -> a
  | (h::t) -> f h (fold_right f t a)
```

Tail Recursion is Important

- Pushing a call frame for each recursive call when operating on a list is dangerous
 - One stack frame for each list element
 - Big list = stack overflow!
- So: favor tail recursion when inputs could be large (i.e., recursion could be deep). E.g.,
 - Prefer List.fold_left to List.fold_right
 - · Library documentation should indicate tail recursion, or not
 - Convert recursive functions to be tail recursive

Tail Recursion Pattern (1 argument)

```
let func x =
 let rec helper arg acc =
  if (base case) then acc
  else
    let arg' = (argument to recursive call)
    let acc' = (updated accumulator)
    helper arg' acc' in (* end of helper fun *)
 helper x (initial val of accumulator)
"
```

Tail Recursion Pattern with fact

```
let fact x =
 let rec helper arg acc =
  if arg = 0 then acc
  else
    let arg' = arg - 1 in
    let acc' = acc * arg in
    helper arg' acc' in (* end of helper fun *)
 helper x 1
"
```

Tail Recursion Pattern with rev

```
let rev x =
                                   Can generalize to
 let rec rev helper arg acc =
                                   more than one
  match arg with [] -> acc
                                   argument, and
                                   multiple cases for
  | h::t ->
                                   each recursive call
    let arg' = t in
    let acc' = h::acc in
    rev helper arg' acc' in (* end of helper fun *)
 rev helper x []
"
```

True/false: this is a tail-recursive map

```
let map f l =
  let rec helper l a =
    match l with
    [] -> a
    | h::t -> helper t ((f h)::a)
  in helper l []
```

True/false: this is a tail-recursive map

```
let map f l =
  let rec helper l a =
    match l with
    [] -> a
    | h::t -> helper t ((f h)::a)
  in helper l []
```

A. True

B. False (elements are reversed)

A Tail Recursive map

```
let map f l =
  let rec helper l a =
    match l with
    [] -> a
    | h::t -> helper t ((f h)::a)
  in rev (helper l [])
```

Could instead change (f h)::a to be a@(f h)

Q: Why is the above implementation a better choice?

A: O(n) running time, not $O(n^2)$ (where n is length of list)

How far does this generalize?

 A function that is tail-recursive returns at most once (to its caller) when completely finished

 Is it possible to convert an arbitrary program into an equivalent one, except where no call ever returns?

Yes. This is called continuation-passing style