

CMSC 330: Organization of Programming Languages

Lets, Tuples, Records

Let Expressions

- Enable binding variables in other expressions
 - These are different from the `let` *definitions* we've been using at the top-level
- They are expressions, so they have a value
- Syntax
 - `let x = e1 in e2`
 - *x* is a *bound variable*
 - *e1* is the *binding expression*
 - *e2* is the *body expression*

Let Expressions

- Syntax

- `let x = e1 in e2`

- Evaluation

- Evaluate **e1** to **v1**
 - Substitute **v1** for **x** in **e2** yielding new expression **e2'**
 - Evaluate **e2'** to **v2**
 - Result of evaluation is **v2**

Example

```
let x = 3+4 in 3*x
➤ let x = 7 in 3*x
➤ 3*7
➤ 21
```

Let Expression Example

`let x = 3+27 in x*3`

- `3+27 : int`
- `x*3 : int` (assuming `x:int`)
- **so** `let x = 3+27 in x*3 : int`

Let Definitions vs. Let Expressions

- At the top-level, we write
 - `let x = e;; (* no in e2 part *)`
 - This is called a let *definition*, not a let *expression*
 - Because it doesn't, itself, evaluate to anything
- Omitting `in` means “from now on”:
`# let pi = 3.14;;`
`(* pi is now bound in the rest of the top-level scope *)`

Top-level expressions

- We can write any expression at top-level, too

– **e**;;

– This says to evaluate **e** and then ignore the result

- Equivalent to `let _ = e;;`
- Useful when **e** has a side effect, such as reading/writing a file, printing to the screen, etc.

```
let x = 37;;  
let y = x + 5;;  
print_int y;;  
print_string "\n";;
```

- When run, outputs 42 to the screen

Let Expressions: Scope

- In `let x = e1 in e2`, variable **x** is *not* visible outside of **e2**

```
let pi = 3.14 in pi *. 3.0 *. 3.0;;  
print_float pi;;
```

error: **pi** not bound

bind **pi** (only) in body of **let**
(which is `pi *. 3.0 *. 3.0`)

Binding in other languages

- Compare to similar usage in Java/C

```
let pi = 3.14 in
  pi *. 3.0 *. 3.0;;
pi;; (* pi unbound! *)
```

```
{
  float pi = 3.14;

  pi * 3.0 * 3.0;
}
pi; /* pi unbound! */
```


Examples – Scope of Let bindings

- $x;;$
 - (* Unbound value x *)
- $\text{let } x = 1 \text{ in } x + 1;;$
 - (* 2 *)
- $\text{let } x = x \text{ in } x + 1;;$
 - (* Unbound value x *)

Examples – Scope of Let bindings

- `let x = 1 in (x + 1 + x) ;;`
 - `(* 3 *)`
- `(let x = 1 in x + 1) ;; x;;`
 - `(* Unbound value x *)`
- `let x = 4 in (let x = x + 1 in x) ;;`
 - `(* 5 *)`

Shadowing Names

- **Shadowing** is rebinding a name in an inner scope to have a different meaning
 - May or may not be allowed by the language

C

```
int i;  
  
void f(float i) {  
    {  
        char *i = NULL;  
        ...  
    }  
}
```

Java

```
void h(int i) {  
    {  
        float i; // not allowed  
        ...  
    }  
}
```

OCaml

```
let x = 3;;  
let g x = x + 3;;
```

Shadowing, by the Semantics

- What if **e2** is also a **let** for **x** ?
 - Substitution will **stop** at the **e2** of a shadowing **x**

Example

```
let x = 3+4 in let x = 3*x in x+1
```


```
➤ let x = 7 in let x = 3*x in x+1
```

```
➤ let x = 3*7 in x+1
```

```
➤ let x = 21 in x+1
```

```
➤ 21+1
```

```
➤ 22
```



Not substituted,
since it is shadowed
by the inner let

Let Expressions in Functions

- You can use **let** inside of functions for local vars

```
let area r =  
  let pi = 3.14 in  
  pi *. r *. r
```

- And you can use many **lets** in sequence

```
let area d =  
  let pi = 3.14 in  
  let r = d /. 2.0 in  
  pi *. r *. r
```

Shadowing (of Locals) Discouraged

- You can use shadowing to simulate mutation (variable update)

```
let rec f x n =  
  if x = 0 then 1  
  else  
    let x = x - 1 in (* shadowed *)  
    n * (f x n)
```

- But avoiding shadowing can be clearer, so we **recommend not using it**
 - With no shadowing, if you see a variable x , you know it hasn't been "changed," no matter where it appears
 - if you want to "update" n , use a new name $n1$, n' , etc.

Nested Let Expressions

- Uses of **let** can be nested in OCaml
 - Nested bound variables (**pi** and **r**) invisible outside
- Similar scoping possibilities C and Java


```
let res =  
  (let area =  
    (let pi = 3.14 in  
      let r = 3.0 in  
        pi *. r *. r) in  
    area /. 2.0) ;;
```

```
float res;  
{ float area;  
  { float pi = 3.14  
    float r = 3.0;  
    area = pi * r *  
r;  
  }  
  res = area / 2.0;  
}
```

Nested Let Style: Generally Avoid

- Oftentimes a nested binding can be rewritten in a more linear style
 - Easier to understand
- Can go too far: namespace pollution
 - Avoiding adding unnecessary variable bindings to top-level

```
let res =  
  (let area =  
    (let pi = 3.14 in  
      let r = 3.0 in  
        pi *. r *. r) in  
    area /. 2.0);;
```



```
let res =  
  let pi = 3.14 in  
  let r = 3.0 in  
  let area = pi *. r *. r  
in  
  area /. 2.0;;
```

```
let pi = 3.14;;  
let r = 3.0;;  
let area = pi *. r *. r;;  
let res = area /. 2.0;;
```


Quiz 1

Which of these is **not** an expression that evaluates to 3?

A. `let x=3`

B. `let x=2 in x+1`

C. `let x=3 in x`

D. `3`

Quiz 1

Which of these is **not** an expression that evaluates to 3?

A. `let x=3` ---> not an expression

B. `let x=2 in x+1`

C. `let x=3 in x`

D. `3`

Quiz 2: What does this evaluate to?

```
let x = 2 in  
x = 3
```

- A. 3
- B. 2
- C. true
- D. false

Quiz 2: What does this evaluate to?

```
let x = 2 in  
x = 3
```

- A. 3
- B. 2
- C. true
- D. false

Quiz 3: What does this evaluate to?

```
let x = 3 in  
let y = x+2 in  
let x = 8 in  
x+y
```

- A. 13
- B. 8
- C. 11
- D. 18

Quiz 3: What does this evaluate to?

```
let x = 3 in  
let y = x+2 in  
let x = 8 in  
x+y
```

- A. 13
- B. 8
- C. 11
- D. 18

let Specializes match

More general form of let allows patterns:

- `let p = e1 in e2`
 - where *p* is a pattern. If *e1* fails to match that pattern then an exception is thrown

This pattern form of `let` is equivalent to

- `match e1 with p -> e2`

Examples

- `let [x] = [1] in 1::x` (* evaluates to [1;1] *)
- `let h::_ = [1;2;3] in h` (* evaluates to 1 *)
- `let () = print_int 5 in 3` (* evaluates to 3 *)

Tuples

- **Constructed** using `(e1, ..., en)`
- **Deconstructed** using pattern matching
 - Patterns involve parens and commas, e.g., `(p1, p2, ...)`
- Tuples are similar to C structs
 - But without field labels
 - Allocated on the heap
- Tuples can be heterogenous
 - Unlike lists, which must be homogenous
 - `(1, ["string1"; "string2"])` is a valid tuple

Tuple Types

- Tuple types use ***** to separate components
 - Type joins types of its components
- Examples
 - `(1, 2) :`
 - `(1, "string", 3.5) :`
 - `(1, ["a"; "b"], 'c') :`
 - `[(1,2)] :`
 - `[(1, 2); (3, 4)] :`
 - `[(1,2); (1,2,3)] :`

Tuple Types

- Tuple types use ***** to separate components
 - Type joins types of its components
- Examples
 - `(1, 2) : int * int`
 - `(1, "string", 3.5) : int * string * float`
 - `(1, ["a"; "b"], 'c') : int * string list * char`
 - `[(1,2)] : (int * int) list`
 - `[(1, 2); (3, 4)] : (int * int) list`
 - `[(1,2); (1,2,3)] : error`

Because the first list element has type `int * int`, but the second has type `int * int * int` – list elements must all be of the same type

Pattern Matching Tuples

```
# let plusThree t =  
  match t with  
    (x, y, z) -> x + y + z;;  
plusThree : int*int*int -> int = <fun>  
  
# let plusThree' (x, y, z) = x + y + z;;  
plusThree' : int*int*int -> int = <fun>  
  
# let addOne (x, y, z) = (x+1, y+1, z+1);;  
addOne : int*int*int -> int*int*int = <fun>  
  
# plusThree (addOne (3, 4, 5));;  
- : int = 15
```

Remember, **semicolon** for lists, **comma** for tuples

- `[1, 2] = [(1, 2)]` *which is a list of size one*
- `(1; 2)` *Warning: This expression should have type unit*

Tuples Are A Fixed Size

- This OCaml definition
 - `# let foo x = match x with
 (a, b) -> a + b
 | (a, b, c) -> a + b + c;;`
- Tuples of different size have different types
 - `(a, b)` has type: `'a * 'b * 'c`
 - `(a, b, c)` has type: `'a * 'b`

Records

- Records: identify elements by **name**
 - Elements of a tuple are identified by **position**
- Define a **record type** before defining record values

```
type date = { month: string; day: int; year: int }
```

- Define a **record value**

```
# let today = { day=16; year=2017; month="f"^"eb" };;  
  
today : date = { day=16; year=2017; month="feb" };;
```

Destructing Records

```
type date = { month: string; day: int; year: int }  
let today = { day=16; year=2017; month="feb" };;
```

- Access by field name or pattern matching

```
print_string today.month;; (* prints feb *)  
(* patterns *)  
let { month=_; day=d } = today in  
let { year } = today in  
let _ = print_int d in      (* prints 16 *)  
print_int year;;           (* prints 2017 *)
```

- Notes:
 - In record patterns, you can skip or reorder fields
 - You can use the field name as the bound variable

Quiz 4: What does this evaluate to?

```
let get (a,b) = a+b in  
get 1 2
```

- A. 3
- B. 2
- C. 1
- D. type error

Quiz 4: What does this evaluate to?

```
let get (a,b) = a+b in  
get 1 2
```

A. 3

B. 2

C. 1

D. type error – **get** takes one argument (a pair)

Quiz 5: What does this evaluate to?

```
let get x y =  
  match x with  
    (a,b) -> a+y  
in  
get (1,2) 1
```

- A. 3
- B. type error
- C. 2
- D. 1

Quiz 5: What does this evaluate to?

```
let get x y =  
  match x with  
    (a,b) -> a+y  
in  
get (1,2) 1
```

- A. 3
- B. type error
- C. 2
- D. 1

Quiz 6: What is the type of `shift`?

```
type point = {x:int; y:int}

let shift p =
  match p with
  { x=px; y=py } -> [px;py]
```

- A. `point -> int list`
- B. `int list -> int list`
- C. `point -> point`
- D. `point -> bool list`

Quiz 6: What is the type of `shift`?

```
type point = {x:int; y:int}

let shift p =
  match p with
  { x=px; y=py } -> [px;py]
```

- A. `point -> int list`
- B. `int list -> int list`
- C. `point -> point`
- D. `point -> bool list`