CMSC 754: Lecture EX1 Review for the Midterm Exam

Midterm Exam: Will be Thu, Mar 12 in class. The exam will be closed-book and closed-notes, but you are allowed one sheet (front and back) of notes.

We have discussed a number of algorithms and data structures that are useful for solving geometric problems. Here is a list of the main topics we have presented so far.

- **Convex hulls:** Convexity and convex hulls algorithms including Graham's scan $(O(n \log n))$, divideand-conquer $(O(n \log n) \text{ time})$, Jarvis march (O(hn) time), and Chan's algorithm $(O(n \log h) \text{ time})$. The latter two algorithms are examples of *output sensitive* algorithms, whose running time is a function of output size (h being the number of vertices on the final hull).
- Line Segment Intersection: We presented an $O((n+m)\log n)$ time plane sweep algorithm, where m is the number of intersection points. This introduced the important algorithmdesign concept called *plane-sweep*.
- **Intersection of Halfplanes:** We presented an $O(n \log n)$ time divide-and-conquer algorithm. The key is a merging step that intersects two convex polygons using plane-sweep. (This was not discussed in class, but it appeared in the lecture notes.)
- **Polygon Triangulation:** We presented a simple $O(n \log n)$ time for polygon triangulation. This was based on two pieces: an O(n)-time algorithm for triangulating an x-monotone polygon and an $O(n \log n)$ time algorithm to decompose a simple polygon into x-monotone pieces. Both algorithms are based on plane-sweep. (Note that there exists an O(n)-time algorithm for simple polygon triangulation, but it is quite complicated.)
- **Point-line duality:** We demonstrated a *dual transformation* that maps a point (a, b) in \mathbb{R}^2 to the (nonvertical) line y = ax b and vice versa. We discussed how numerous affine (point-line) properties are preserved by this transformation. We observed that problems involving lines/points can be converted to an equivalent problem involving points/lines through this transformation. As an example, we showed that computing lower and upper envelopes of a set lines can be reduced to the problem of computing the upper and lower convex hulls of the dual points.
- Linear Programming: We presented a randomized incremental algorithm for linear programming in spaces of constant dimension d. The algorithm runs in O(d!n) expected time. The randomized incremental method involves inserting halfplanes at random and then updating the optimal point with each new insertion. We introduced the technique of backwards analysis, which analyzing the expected time of each step under the assumption that the last insertion is random. In the homework assignment, we explored a number of applications of low-dimensional linear programming.
- **Trapezoidal Maps:** We introduced the useful planar decomposition of a collection of line segments called the *trapezoidal map*. We presented a randomized incremental $O(n \log n)$ expected-time algorithm for computing the trapezoidal map of n line segments.

- **Planar Point Location:** By building a *history DAG* for the trapezoidal-map construction process, it is possible to answer point location queries in $O(\log n)$ time and O(n) space. This data structure records the history of the rebucketing decisions made by the construction algorithm.
- **DCEL (Not covered):** The material on the doubly-connected edge list (DCEL) from the last lecture will *not* be covered on the exam.