



Lecture 12: Fat-tree and Dragonfly Networks

Abhinav Bhatele, Department of Computer Science



UNIVERSITY OF
MARYLAND

Summary of last lecture

- Key requirements of HPC networks
 - extremely low latency, high bandwidth, scalable
 - low network diameter, high bisection bandwidth
- Torus networks (less common now)
 - Network diameter grows as $O(\sqrt[3]{N})$ where N is the number of nodes
- Different types of routing algorithms:
 - Shortest path vs. non-minimal
 - Static vs. dynamic

Fat-tree network


- Most popular network topology
 - Low network diameter, high bandwidth



Fat-tree network

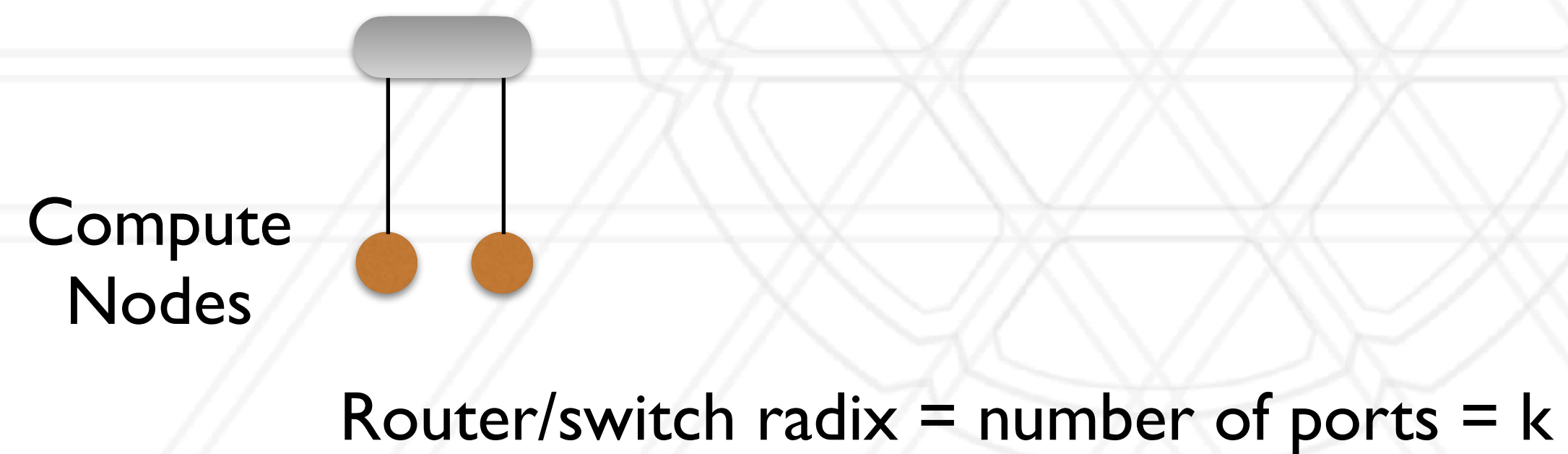
- Most popular network topology
 - Low network diameter, high bandwidth

Compute
Nodes



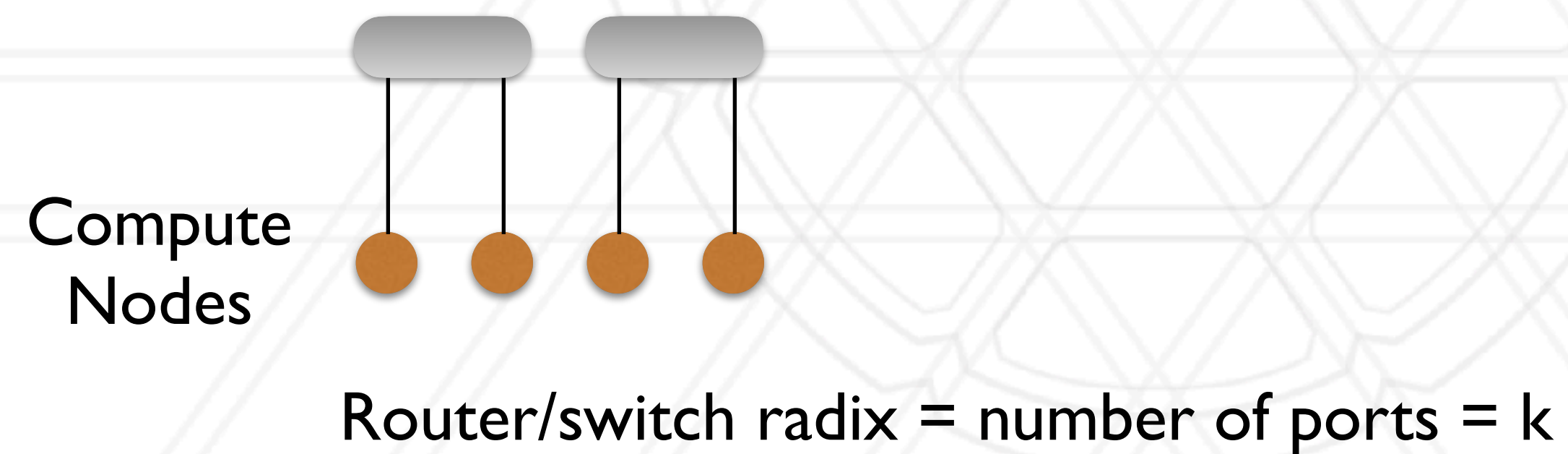
Fat-tree network

- Most popular network topology
 - Low network diameter, high bandwidth



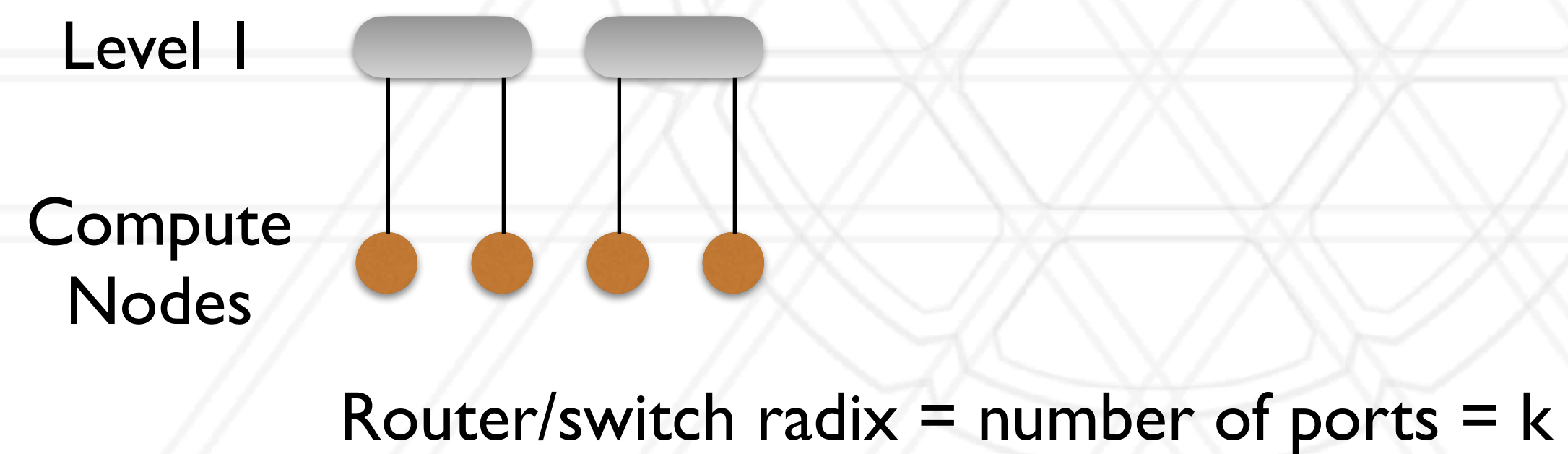
Fat-tree network

- Most popular network topology
 - Low network diameter, high bandwidth



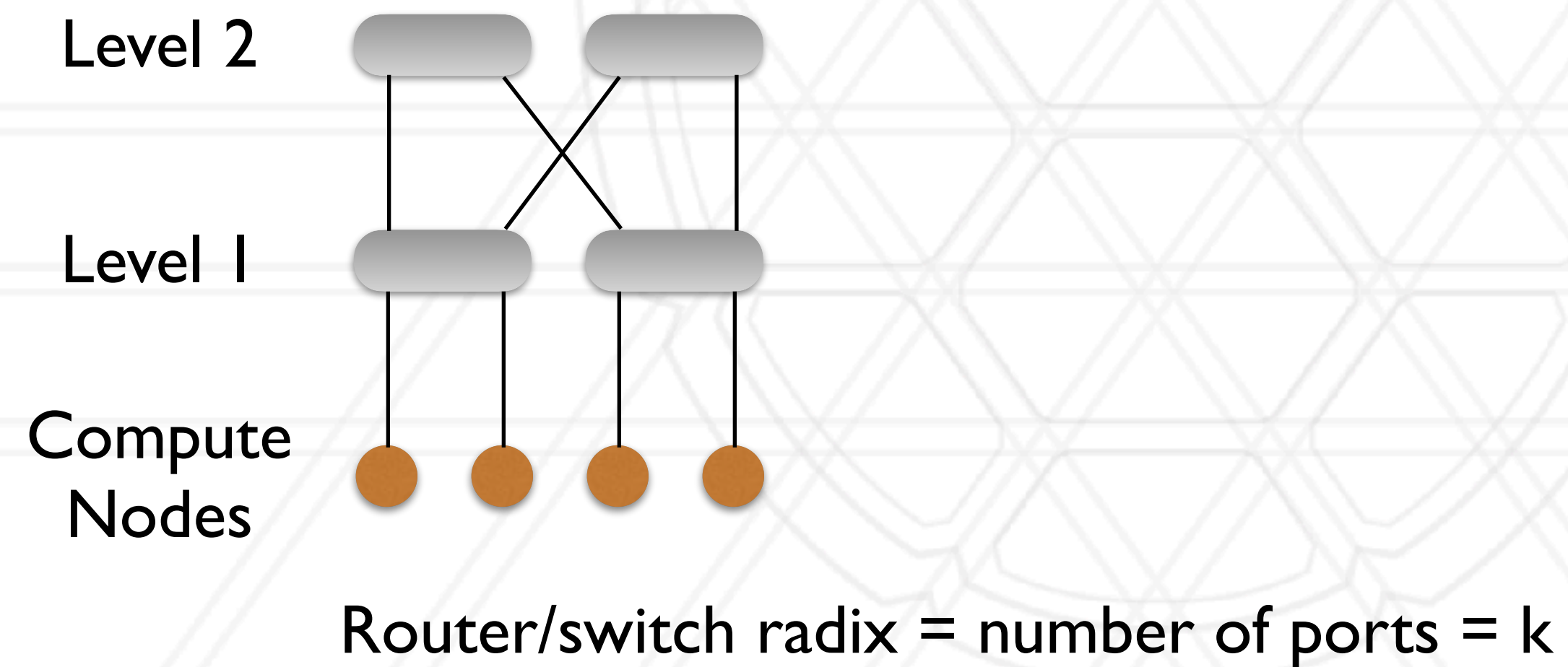
Fat-tree network

- Most popular network topology
 - Low network diameter, high bandwidth



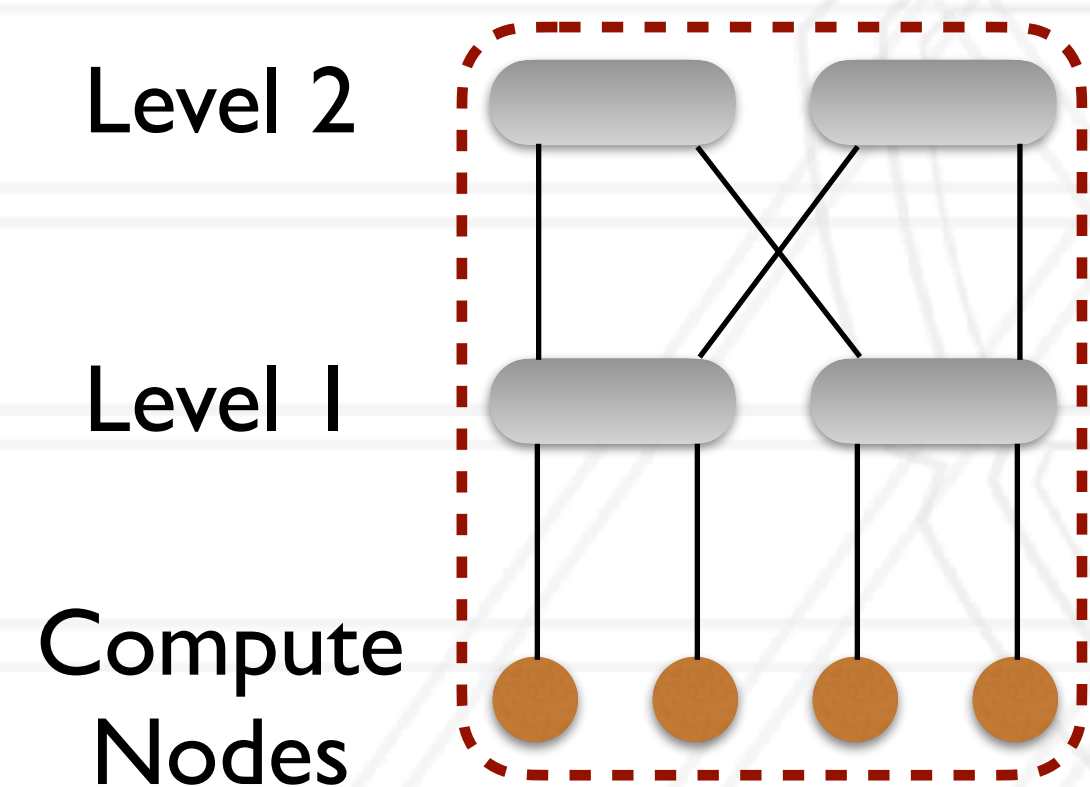
Fat-tree network

- Most popular network topology
 - Low network diameter, high bandwidth



Fat-tree network

- Most popular network topology
 - Low network diameter, high bandwidth

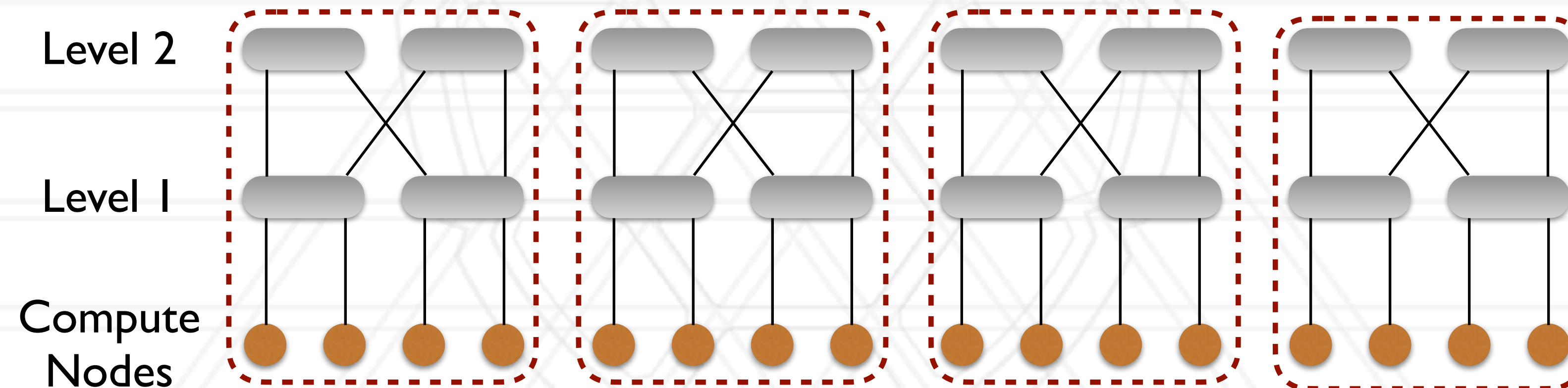


Router/switch radix = number of ports = k

Pod = group of switches = $k/2$ switches

Fat-tree network

- Most popular network topology
 - Low network diameter, high bandwidth

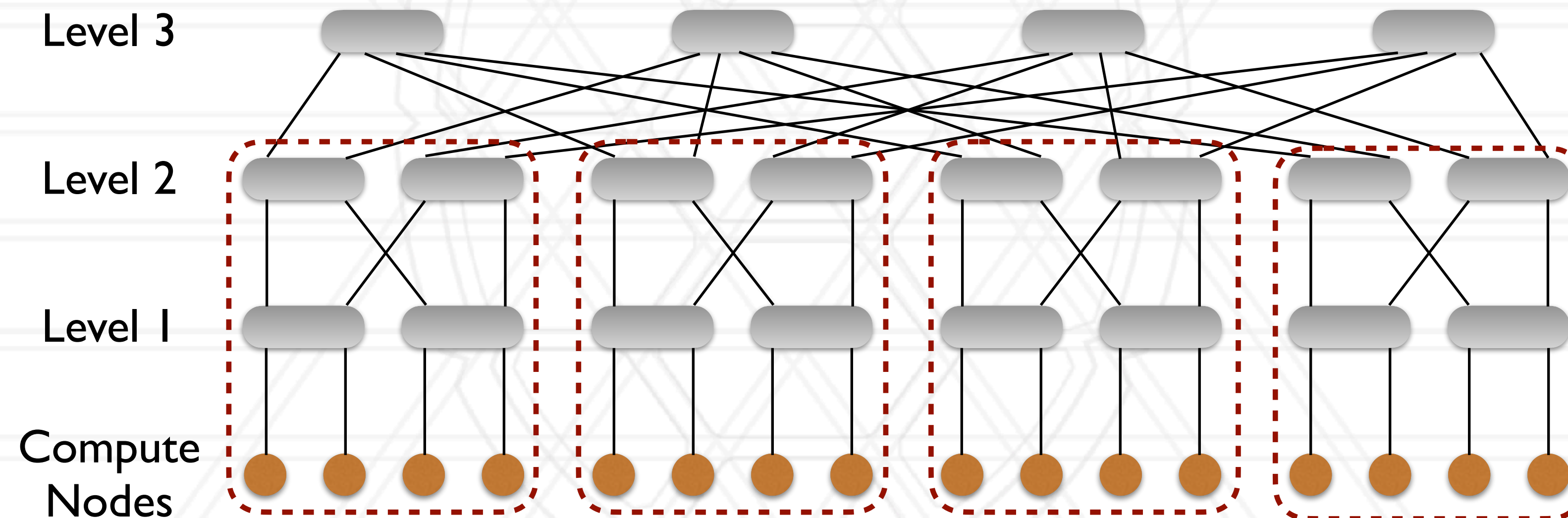


Router/switch radix = number of ports = k

Pod = group of switches = $k/2$ switches

Fat-tree network

- Most popular network topology
 - Low network diameter, high bandwidth

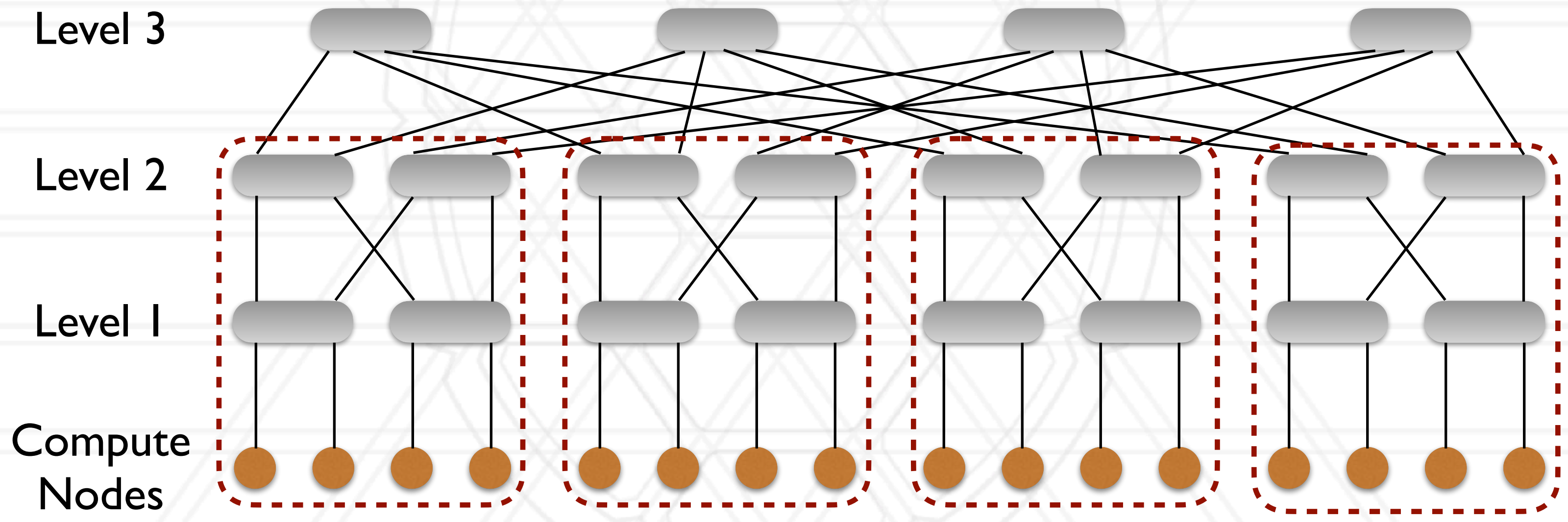


Router/switch radix = number of ports = k

Pod = group of switches = $k/2$ switches

Fat-tree network

- Most popular network topology
 - Low network diameter, high bandwidth

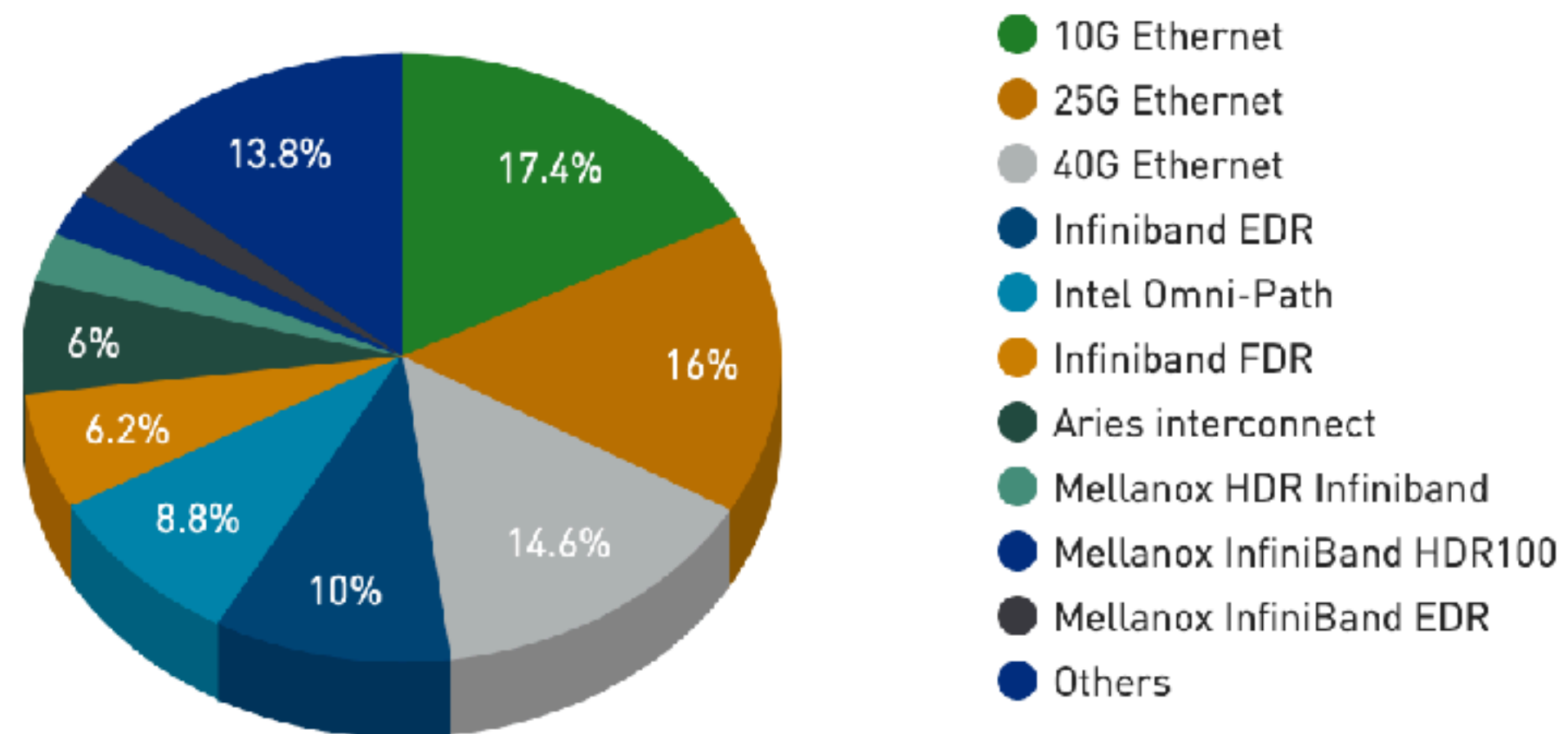


Router/switch radix = number of ports = k
Pod = group of switches = $k/2$ switches Max. number of pods = k

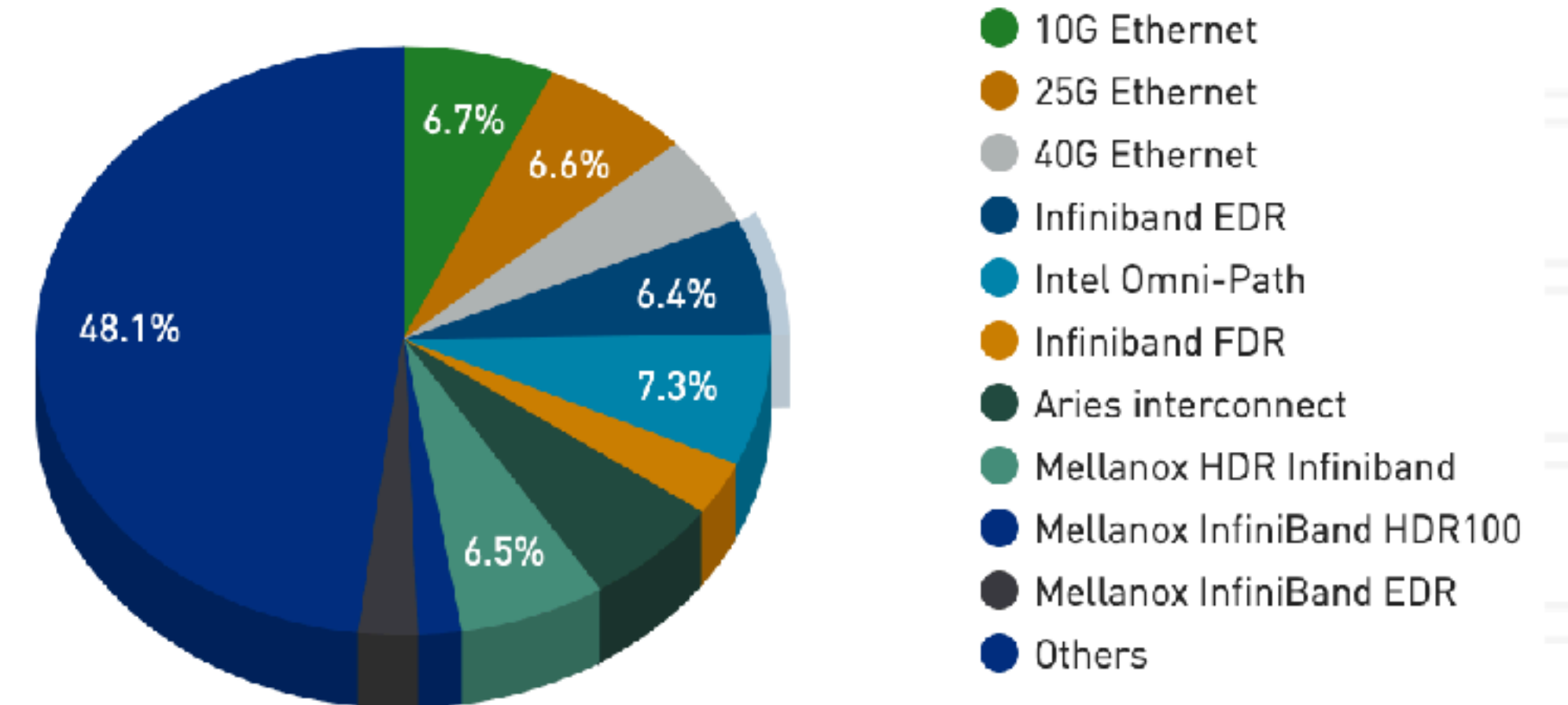
Fat-tree networks on the top500 list

- Infiniband EDR/FDR/HDR, Intel Omni-Path

Interconnect System Share



Interconnect Performance Share

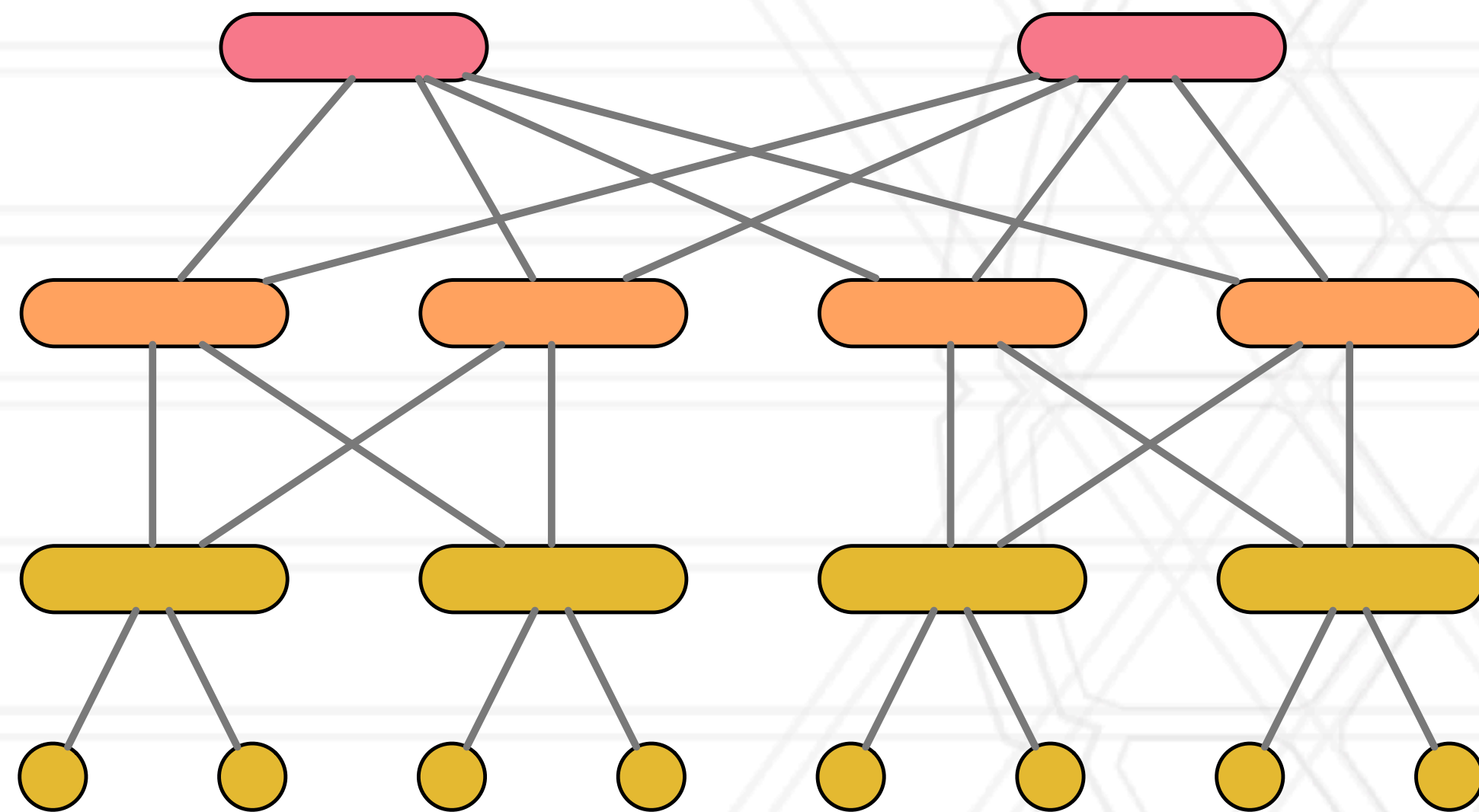


<https://www.top500.org/statistics/list>, November 2020

Routing on a fat-tree

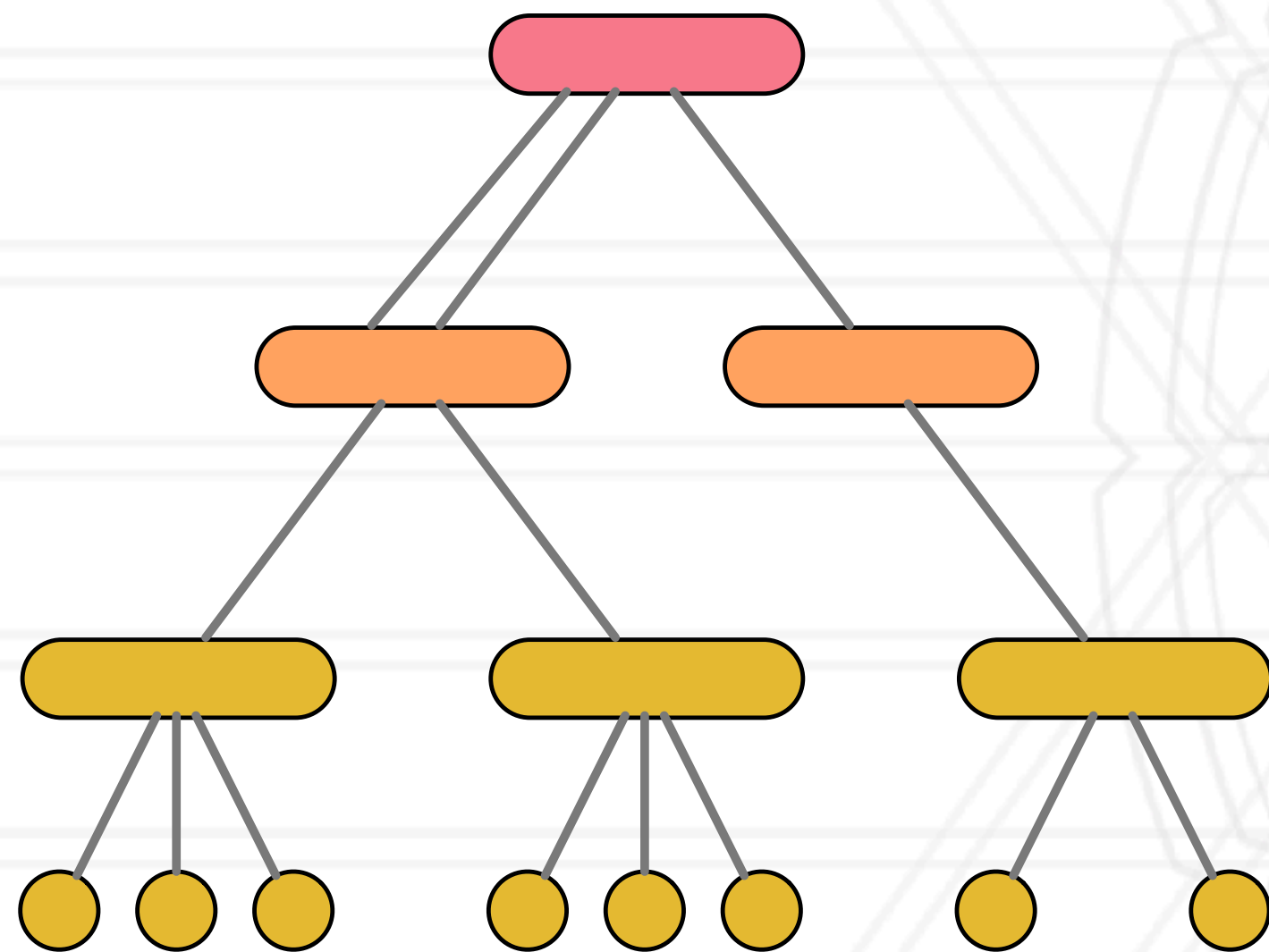
- Until recently, most fat-tree installations used static routing
 - Destination-mod-k (D-mod-k) routing
- Adaptive routing is now starting to be used

Variations on a full bandwidth fat-tree



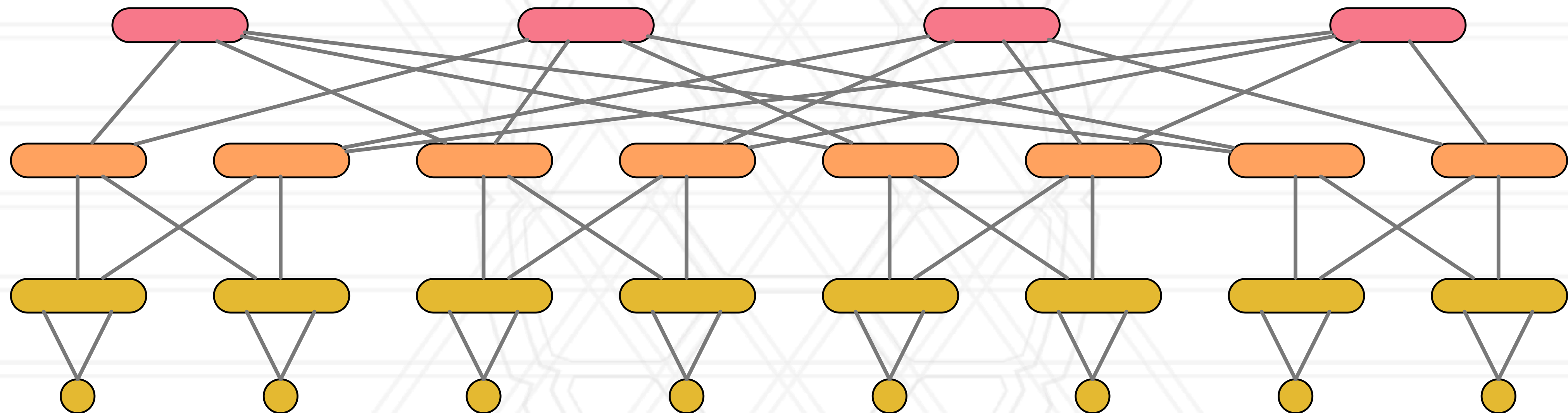
Single-rail single-plane fat-tree

Variations on a full bandwidth fat-tree



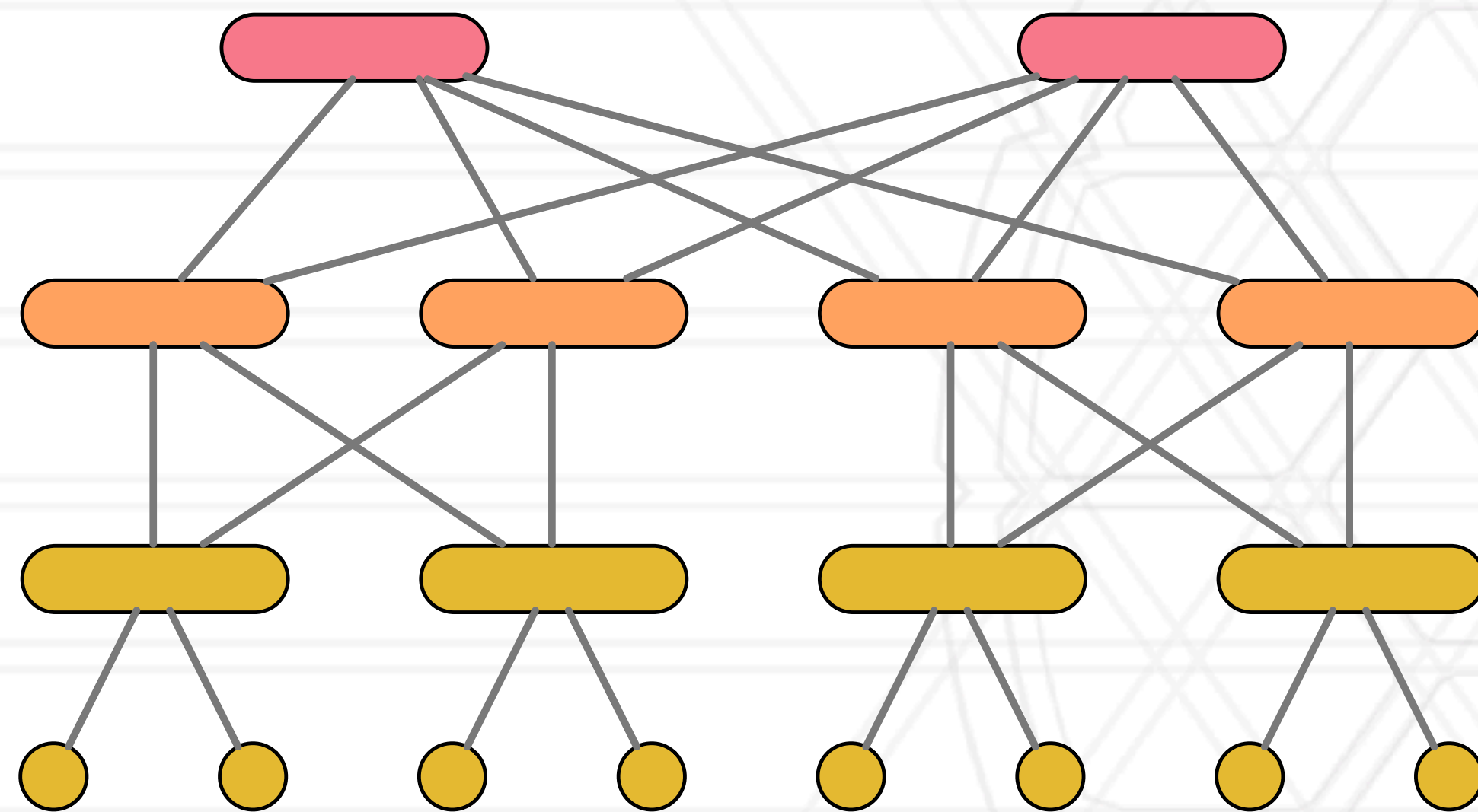
Single-rail single-plane fat-tree (tapered)

Variations on a full bandwidth fat-tree



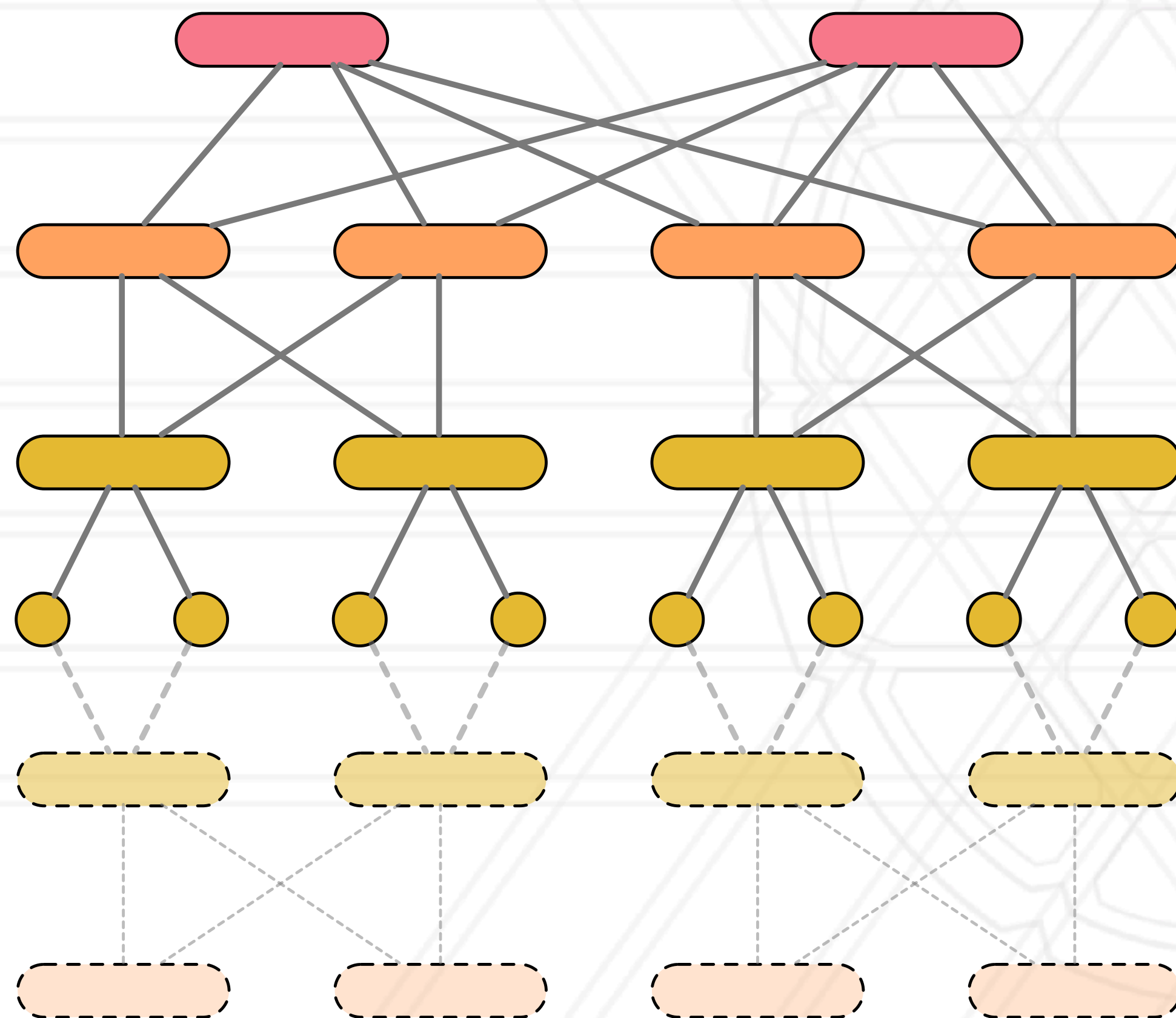
Dual-rail single-plane fat-tree

Variations on a full bandwidth fat-tree



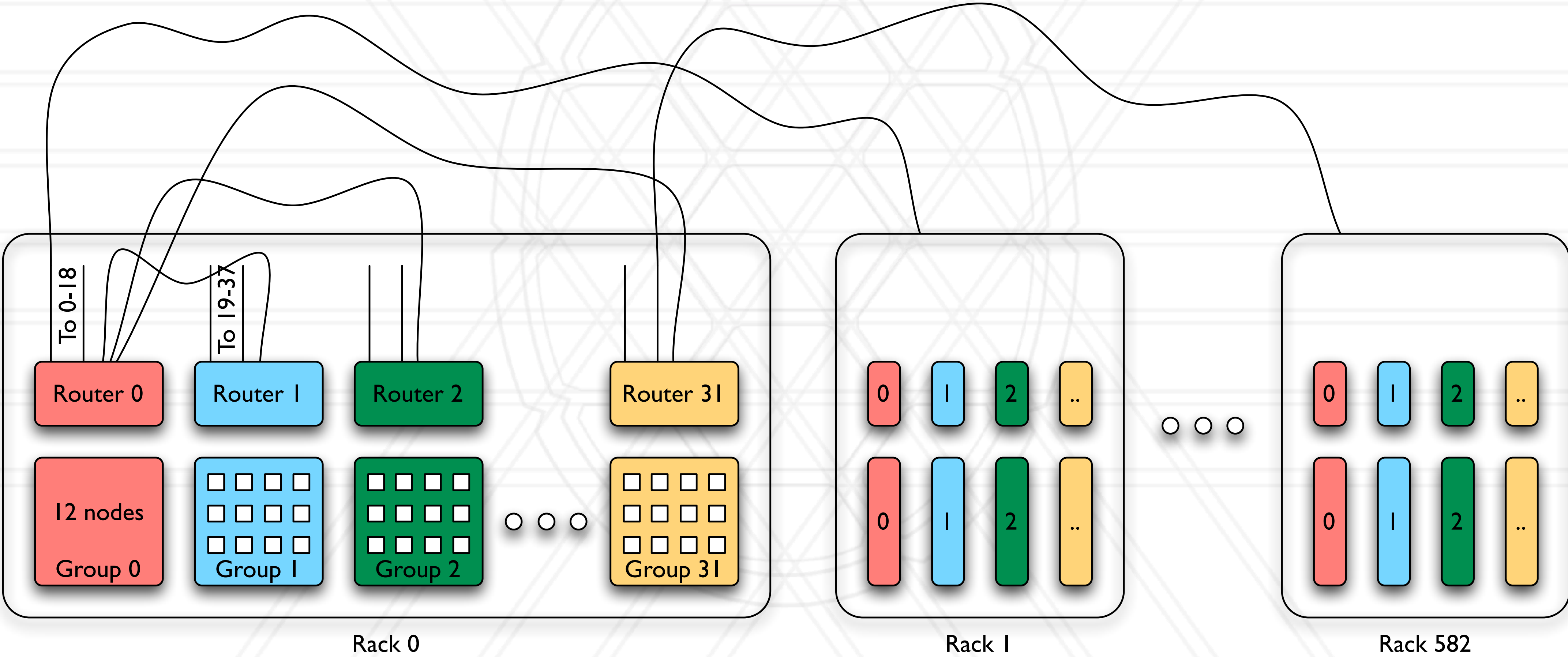
Single-rail single-plane fat-tree

Variations on a full bandwidth fat-tree



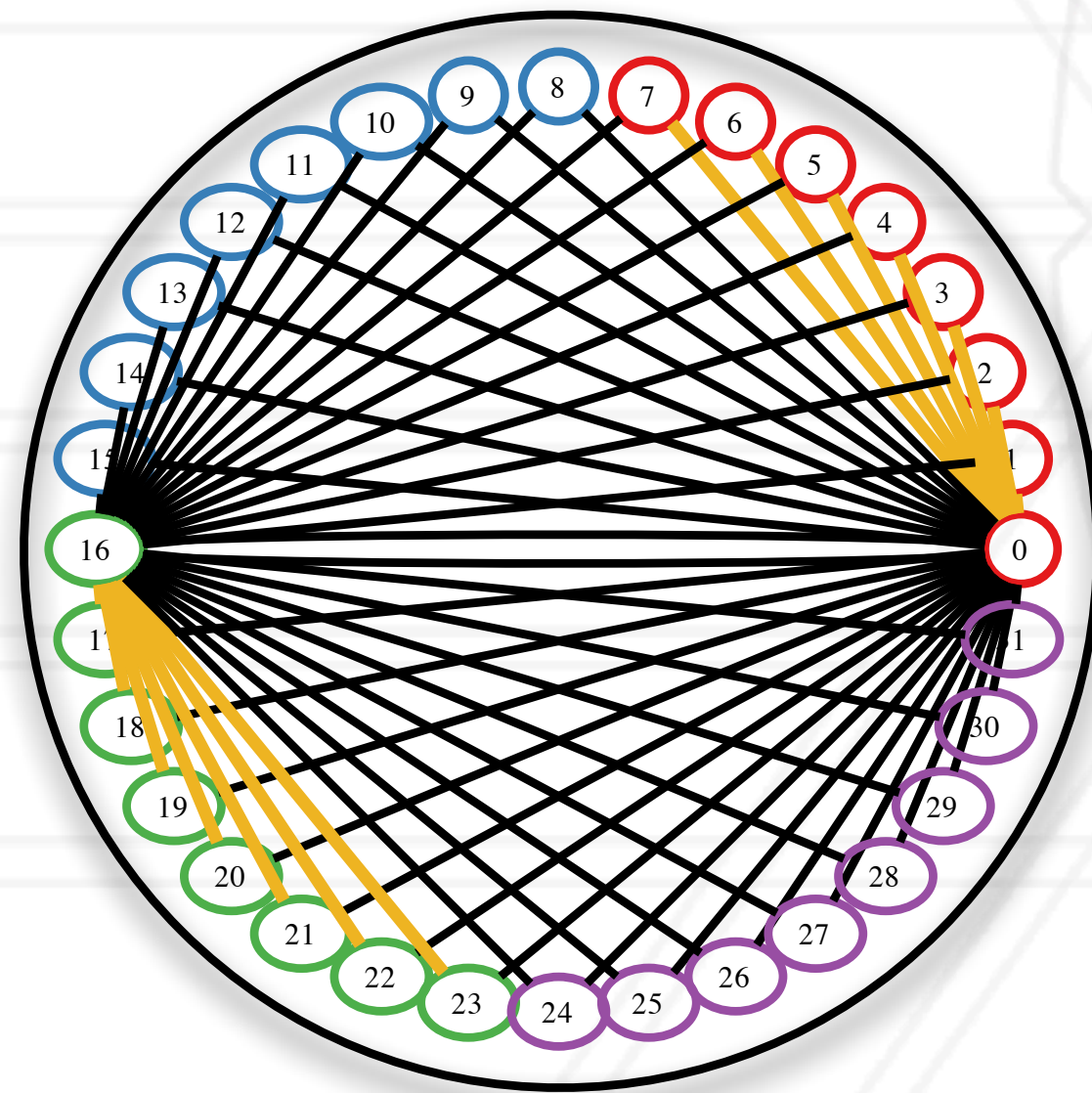
Dual-rail dual-plane fat-tree

Dragonfly network



IBM PERCS network

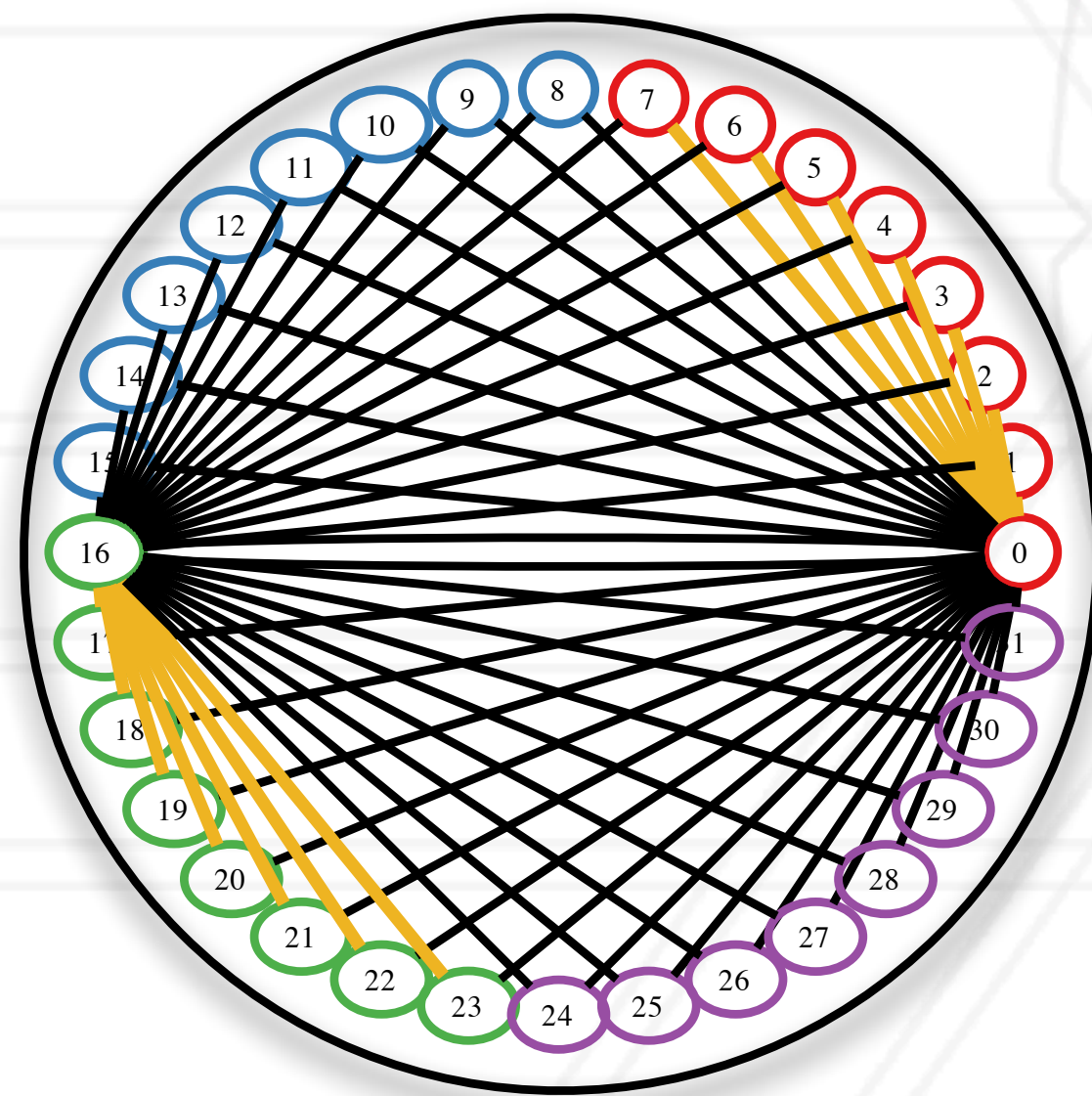
- All-to-all connections within each group



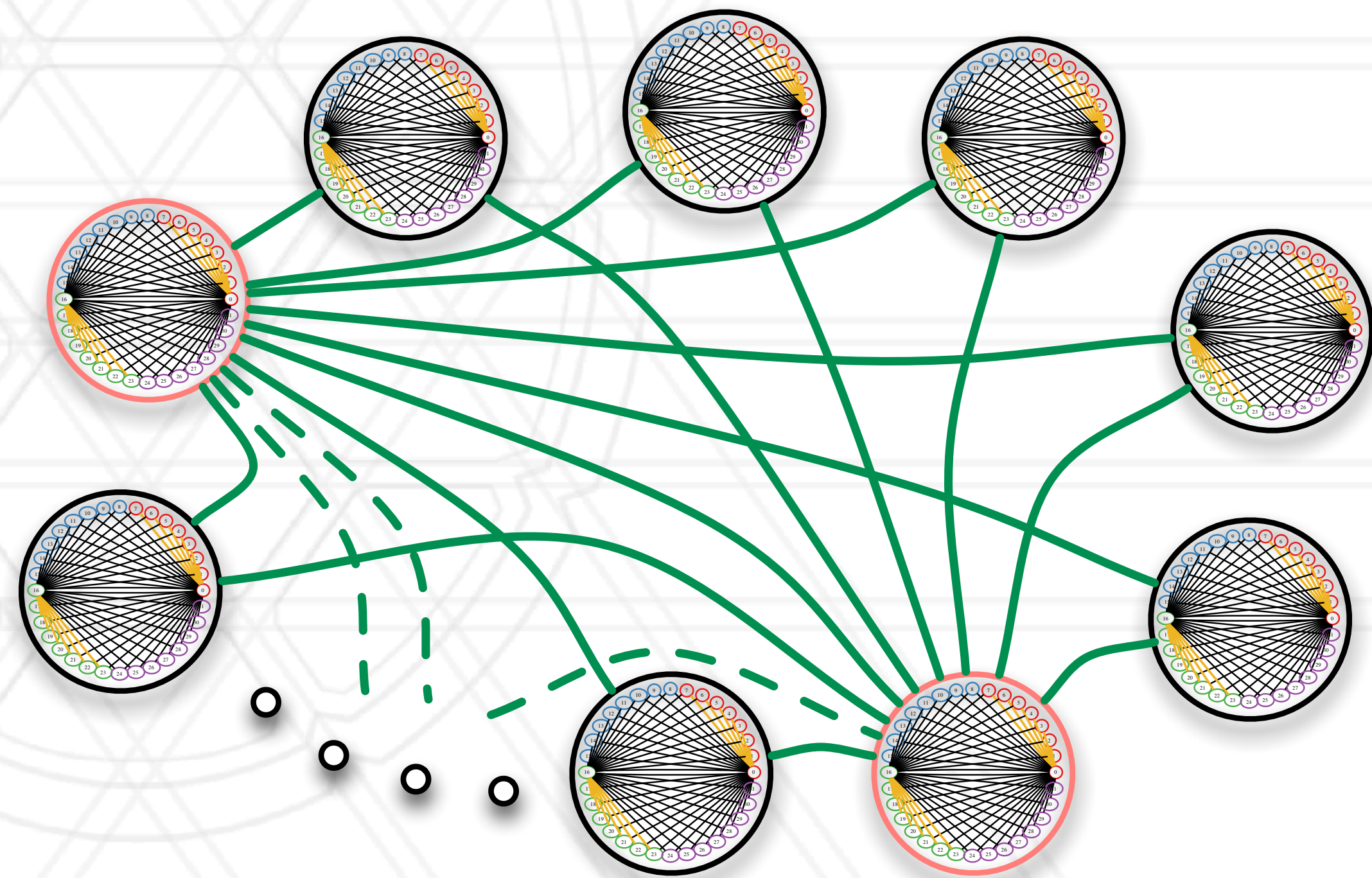
One supernode in the PERCS topology

IBM PERCS network

- All-to-all connections within each group



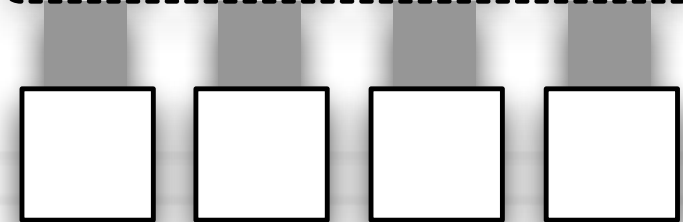
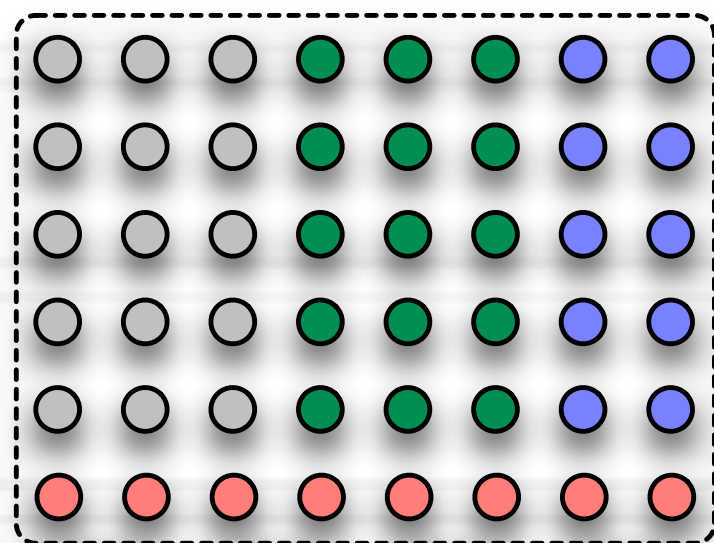
One supernode in the PERCS topology



Cray Aries network

- Row and column all-to-all connections within each group

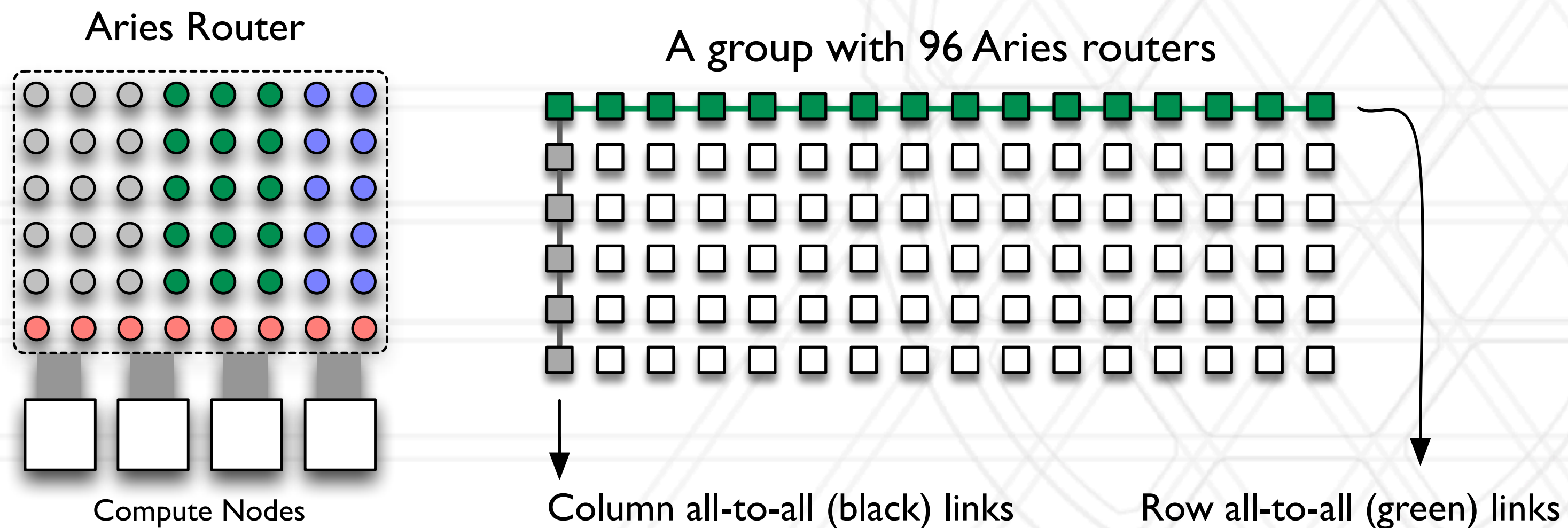
Aries Router



Compute Nodes

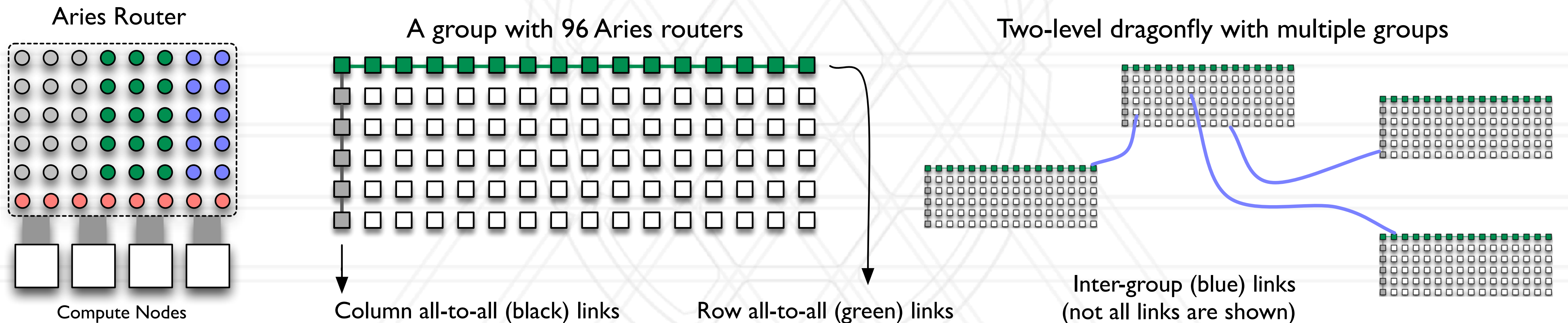
Cray Aries network

- Row and column all-to-all connections within each group



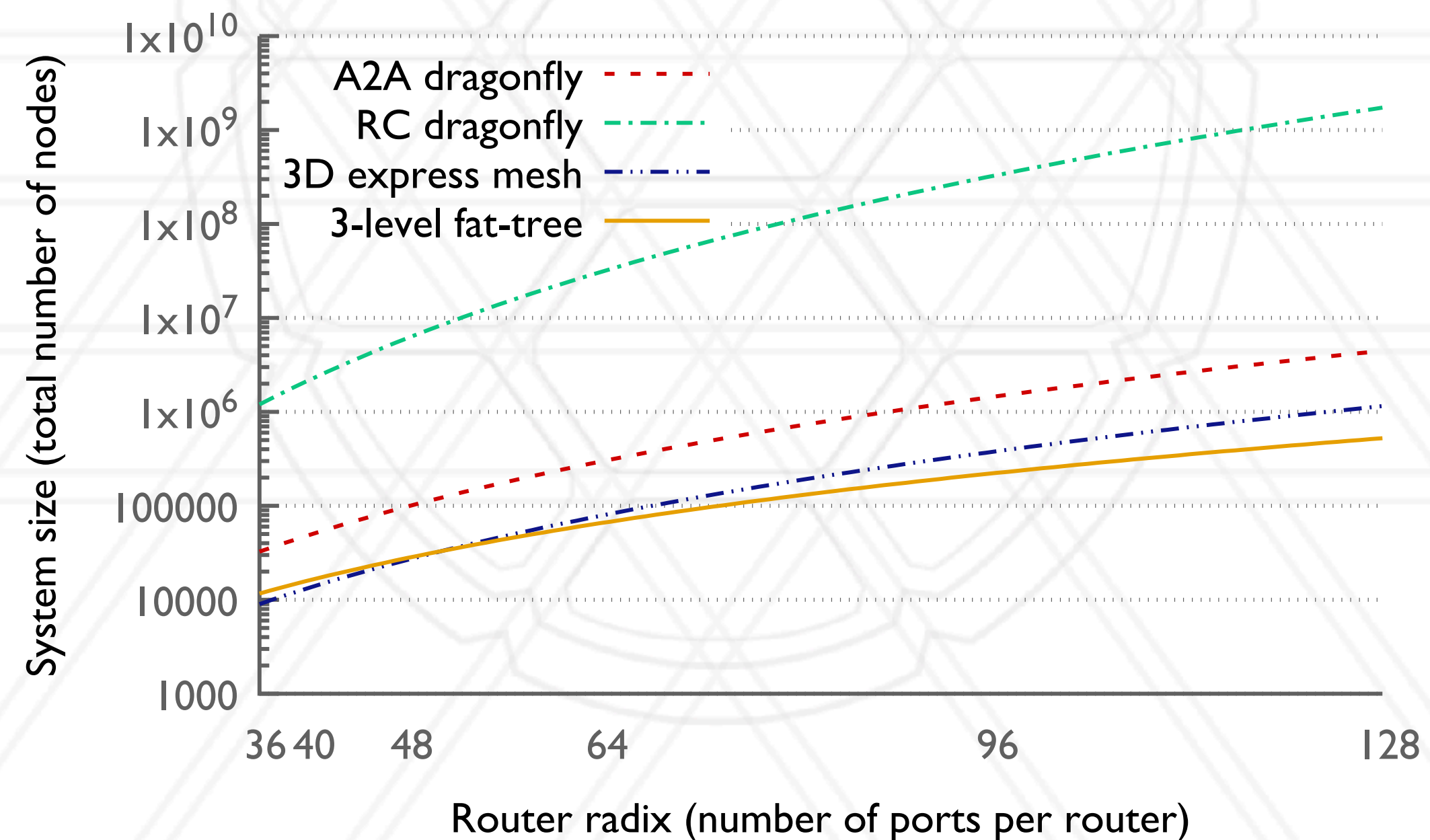
Cray Aries network

- Row and column all-to-all connections within each group



Network comparisons

Network topology	#nodes/router	#links/router	Maximum system size (#nodes)
All-to-all (A2A) dragonfly	$k/4$	$k/2$ (L), $k/4$ (G)	$(k/2 + 1)^2 \times (k/4 + 1) \times k/4$
Row-column (RC) dragonfly	$k/6$	$2k/3$ (L), $k/6$ (G)	$(k/3 + 1)^4 \times (k/6 + 1) \times k/6$
Express mesh (3D, gap=1)	$k/4$	$3k/4$	$(k/4 + 1)^3 \times k/4$
Fat-tree (three-level)	$k/2$	$k/2$	$k/2 \times k/2 \times k$



Questions?



UNIVERSITY OF
MARYLAND

Abhinav Bhatele

5218 Brendan Iribe Center (IRB) / College Park, MD 20742

phone: 301.405.4507 / e-mail: bhatele@cs.umd.edu