#### Introduction to Parallel Computing (CMSC416)

# **Designing Parallel Programs**



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## Announcements

- Zaratan accounts have been created for everyone
- ELMS should be visible to everyone now so far only Announcements
- When emailing me, please cc the TAs also
  - Emails are on the class website: https://www.cs.umd.edu/class/spring2023/cmsc416/index.shtml
  - Prefix [CMSC416] to your email subject
- Office hours posted on class website Zoom info for online hours in ELMS
- Assignment 0 will (likely) be posted on Feb. 9 and will be due a week later
  - Not graded, 0 points



• Some of you will have to activate your TERPConnect account, and should have received an email on how to do that



## Getting started with zaratan

### • 380 nodes with AMD "Milan" processors (128 cores/node)

### • 20 nodes with four NVIDIA AI00 GPUs (in addition to I28 cores/node)



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ssh username@login.zaratan.umd.edu



# Writing parallel programs

- Decide the serial algorithm first
- Data: how to distribute data among threads/processes?
  - Data locality: assignment of data to specific processes to minimize data movement
- Computation: how to divide work among threads/processes?
- Figure out how often communication will be needed





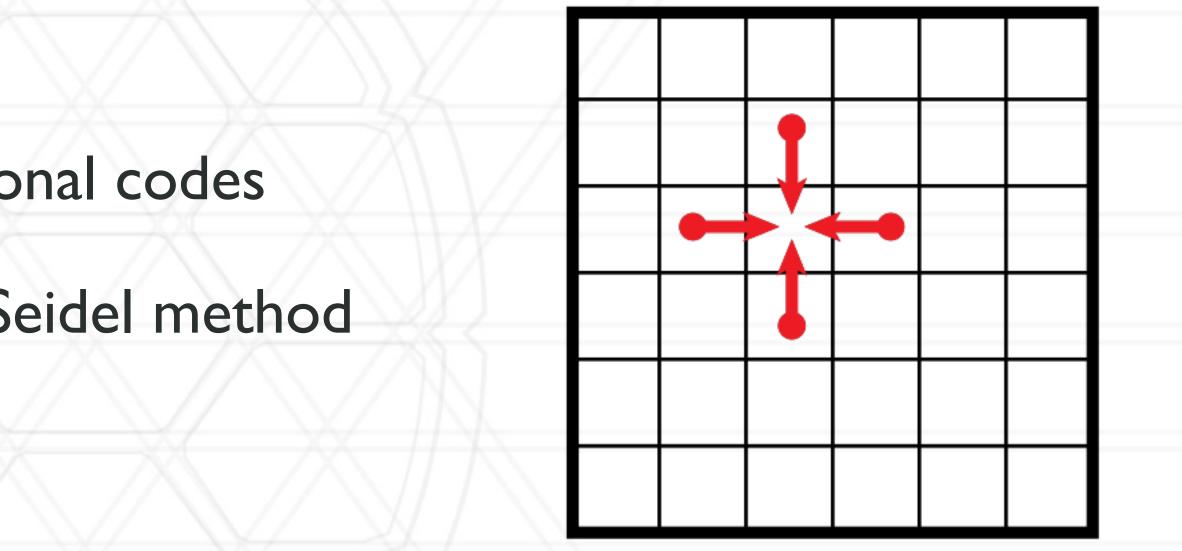


## **Two-dimensional stencil computation**

- Commonly found kernel in computational codes
- Heat diffusion, Jacobi method, Gauss-Seidel method

A[i, j] -A[i, j]





$$+ A[i - 1, j] + A[i + 1, j] + A[i, j - 1] + A[i, j + 1]$$
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## Serial code

for(int t=0; t<num\_steps; t++) {</pre>

for(i ...)
for(j ...)
 A\_new[i, j] = (A[I, j] + A[i-1, j]

// copy contents of A\_new into A
...



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#### A\_new[i, j] = (A[I, j] + A[i-1, j] + A[i+1, j] + A[i, j-1] + A[i, j+1]) \* 0.2



## **2D stencil computation in parallel**

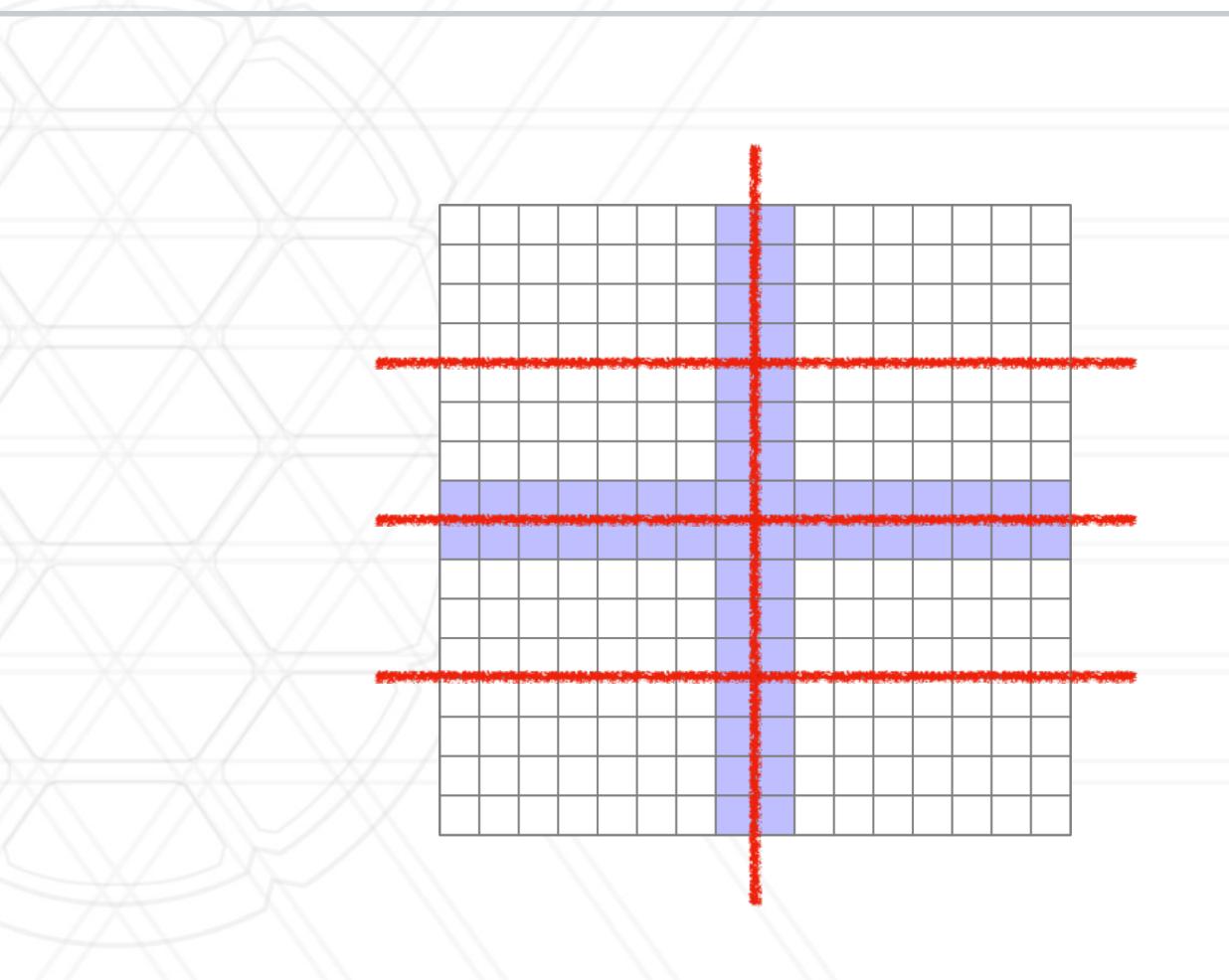
### ID decomposition

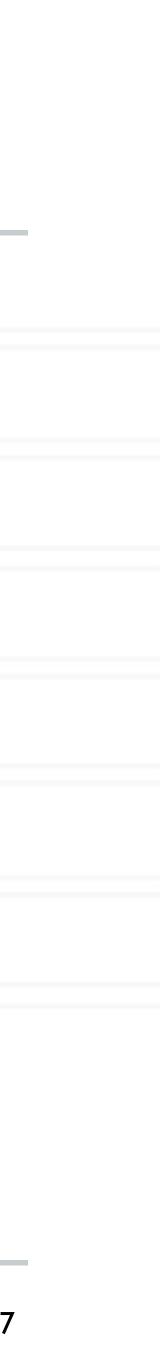
• Divide rows (or columns) among processes

### 2D decomposition

 Divide both rows and columns (2d blocks) among processes







## **Prefix sum**

- Calculate partial sums of elements in a
- Also called a "scan" sometimes

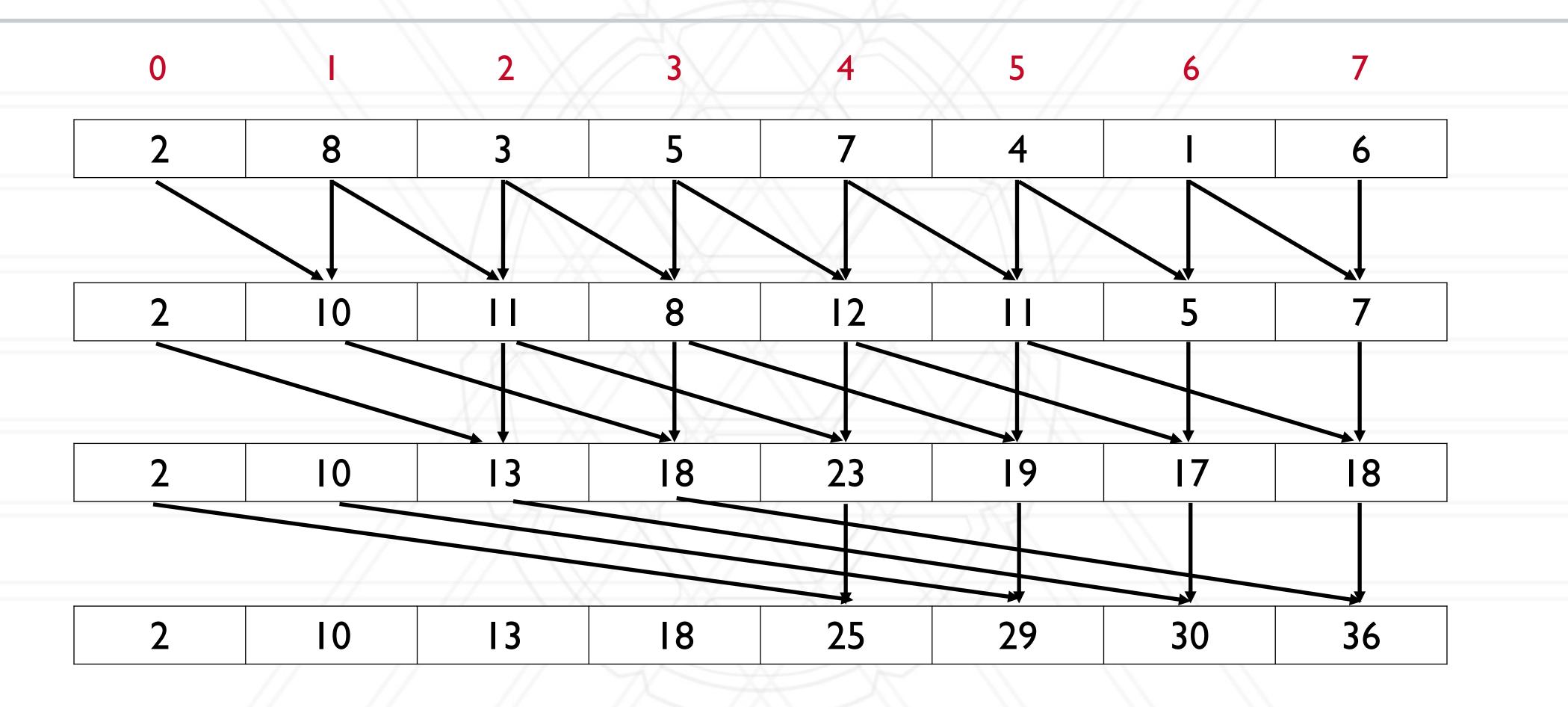
```
pSum[0] = A[0]
for(i=1; i<N; i++) {
    pSum[i] = pSum[i-1] + A[i]
```



in arra								
A		2	3	4	5	6	• • •	
pSum	I	3	6	10	15	21	• • •	



## Parallel prefix sum







## In practice

- You have N numbers and P processes, N >> P
- Assign a N/P block to each process

Do calculation for the blocks on each process locally

Then do parallel algorithm with partial prefix sums





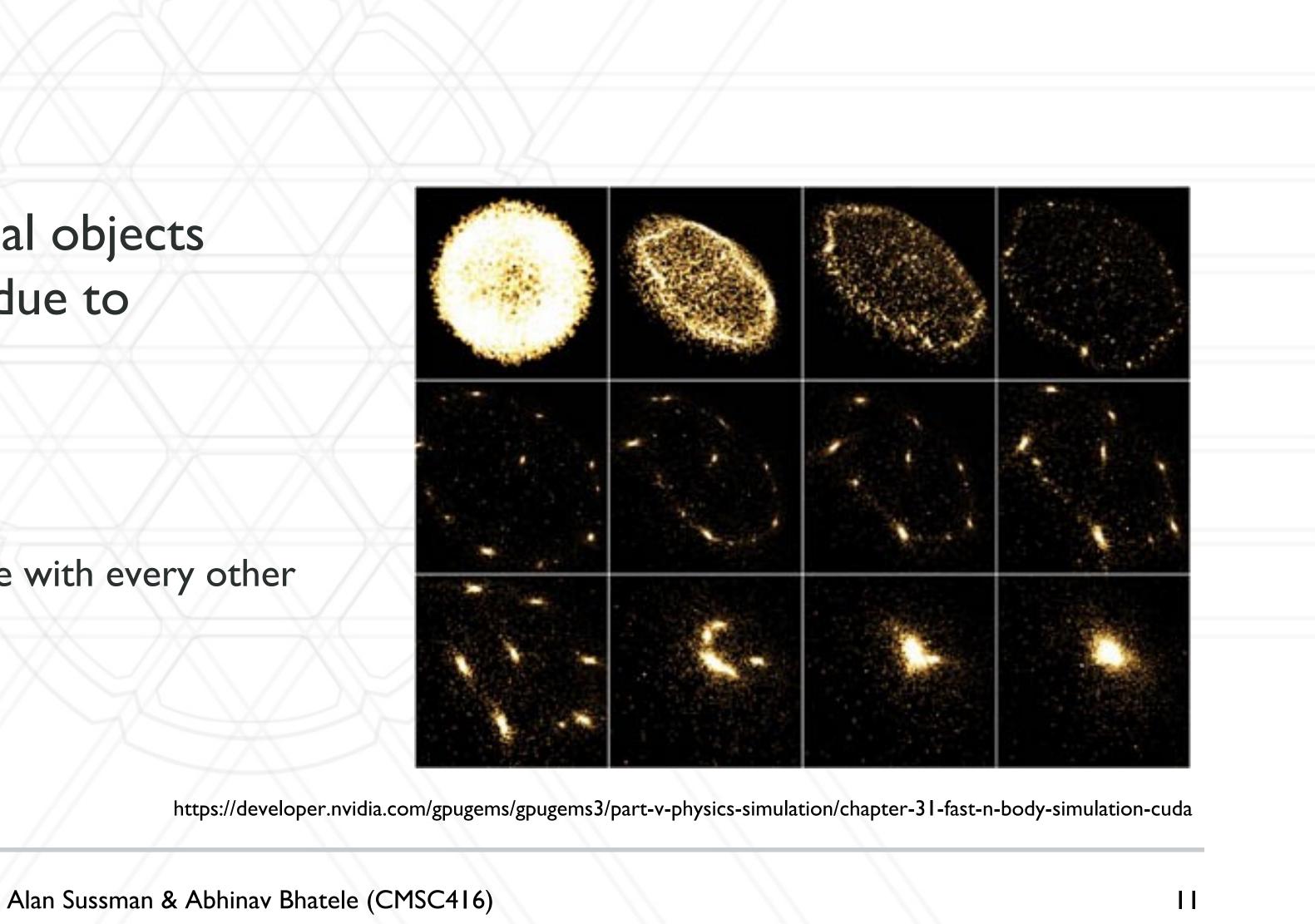


# The n-body problem

- Simulate the motion of celestial objects interacting with one another due to gravitational forces
- Naive algorithm:  $O(n^2)$ 
  - Every body calculates forces pair-wise with every other body (particle)

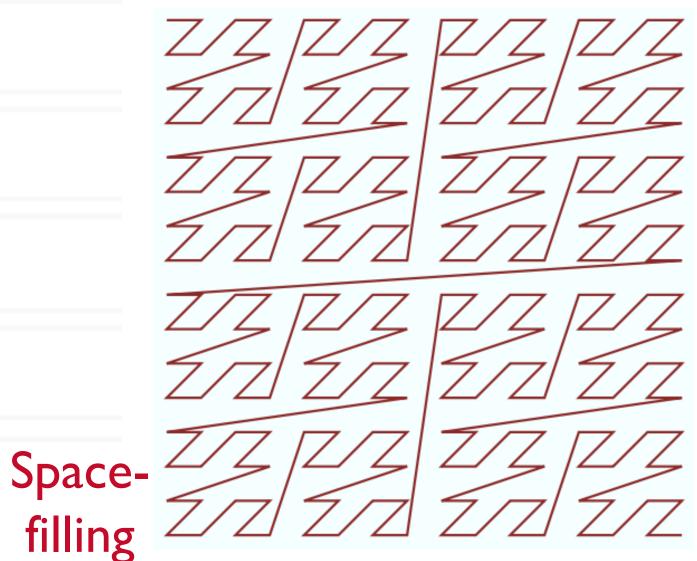


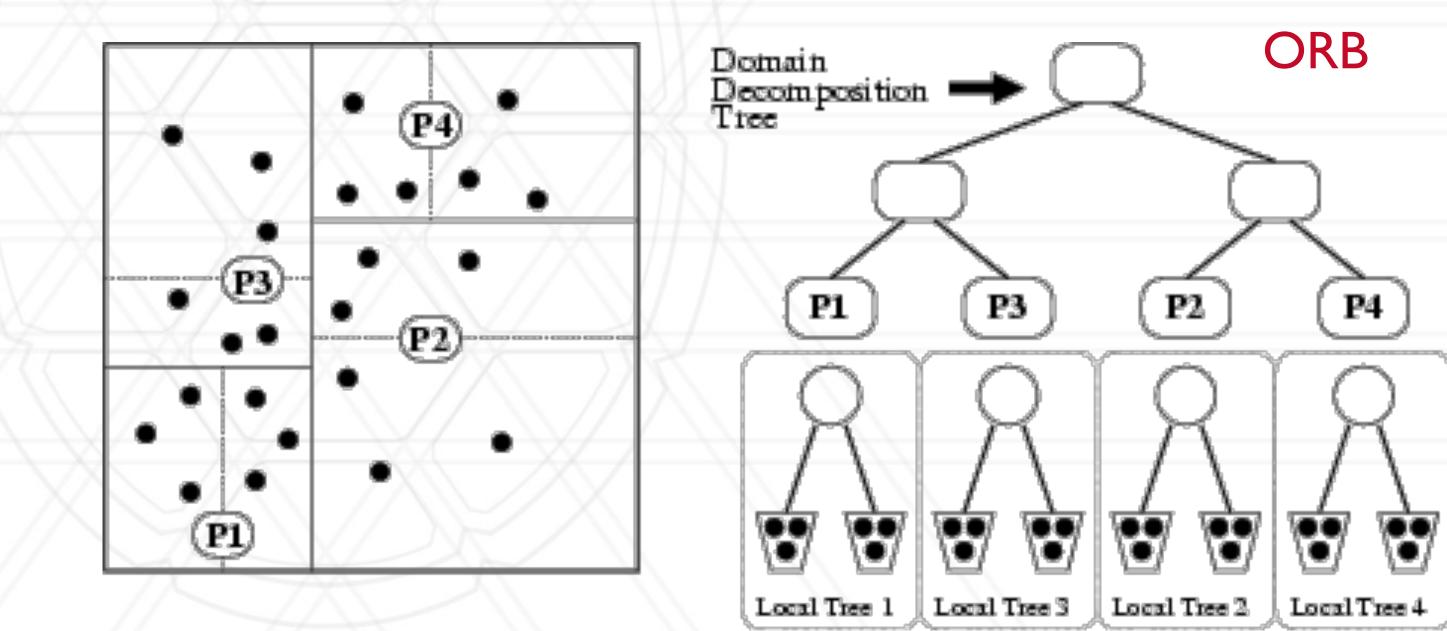




## Data distribution in *n*-body problems

- Naive approach: Assign n/p particles to each process
- Other approaches?





curves

http://datagenetics.com/blog/march22013/ https://en.wikipedia.org/wiki/Z-order\_curve



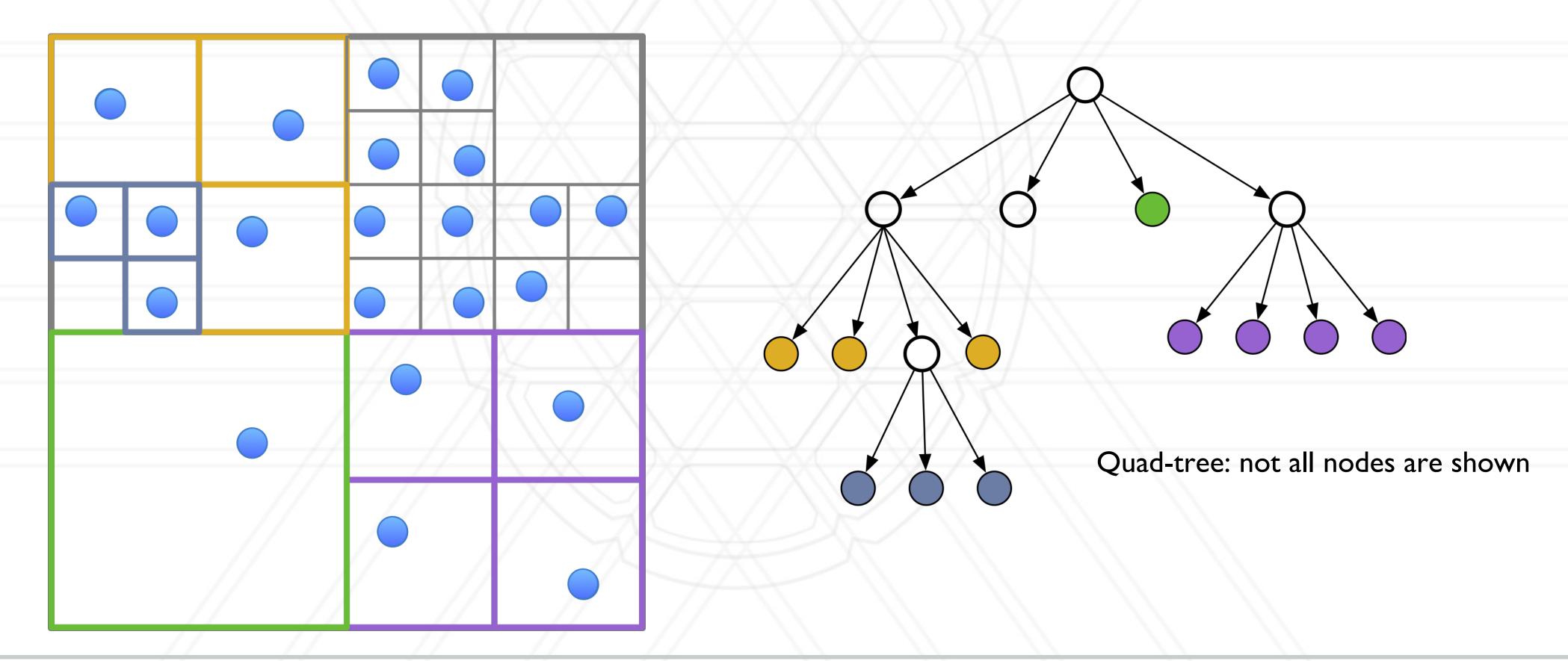
http://charm.cs.uiuc.edu/workshops/charmWorkshop2011/slides/CharmWorkshop2011 apps ChaNGa.pdf

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## Data distribution in *n*-body problems

#### Let us consider a two-dimensional space with bodies/particles in it







## Load balance and grain size

- threads/ processes
  - Bring ratio of maximum to average load as close to 1 as possible
  - Secondary consideration: also load balance amount of communication
- Grain size: ratio of computation-to-communication
  - Coarse-grained (more computation) vs. fine-grained (more communication)



#### • Load balance: try to balance the amount of work (computation) assigned to different







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