### Introduction to Parallel Computing (CMSC416 / CMSC616)

# **Designing Parallel Programs**



### Abhinav Bhatele, Department of Computer Science



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### • Decide the serial algorithm first



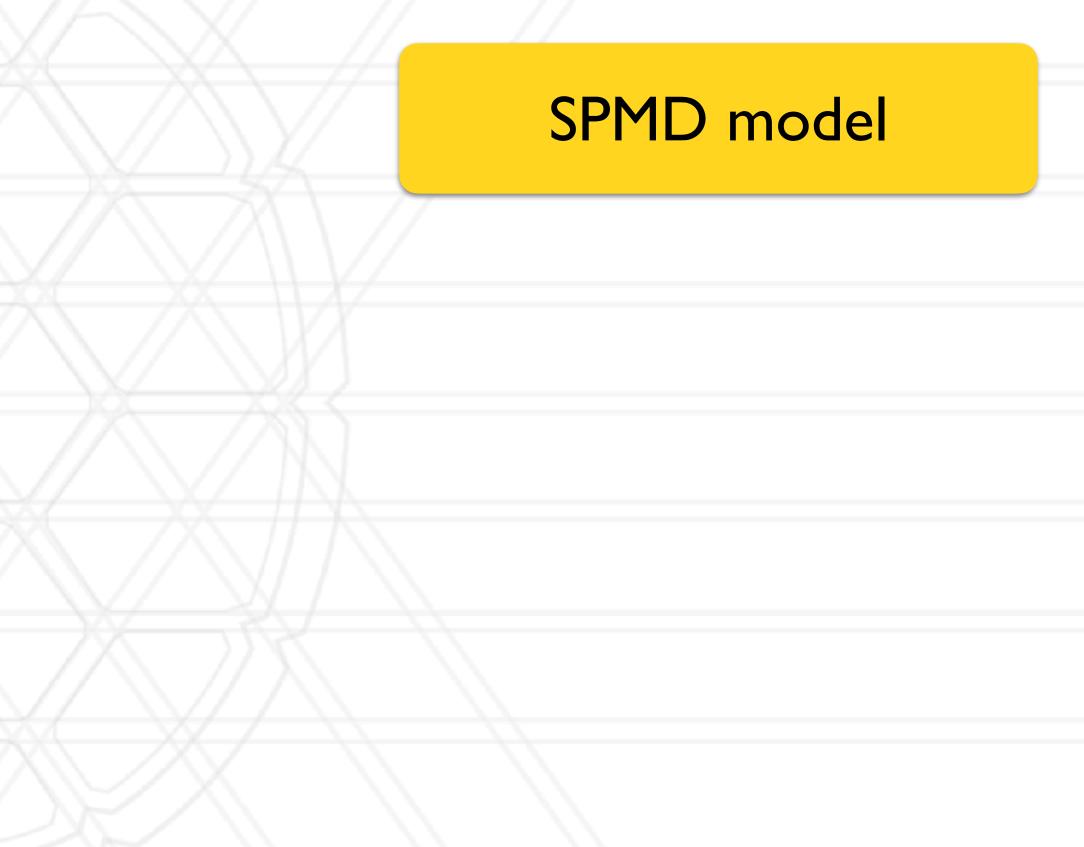
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- Decide the serial algorithm first
- Data: how to distribute data among threads/processes?
  - Data locality: assignment of data to specific processes to minimize data movement



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### SPMD model

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- Data: how to distribute data among threads/processes?
  - Data locality: assignment of data to specific processes to minimize data movement
- Computation: how to divide work among threads/processes?





### SPMD model

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- Decide the serial algorithm first
- Data: how to distribute data among threads/processes?
  - Data locality: assignment of data to specific processes to minimize data movement
- Computation: how to divide work among threads/processes?
- Figure out how often communication will be needed



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### SPMD model

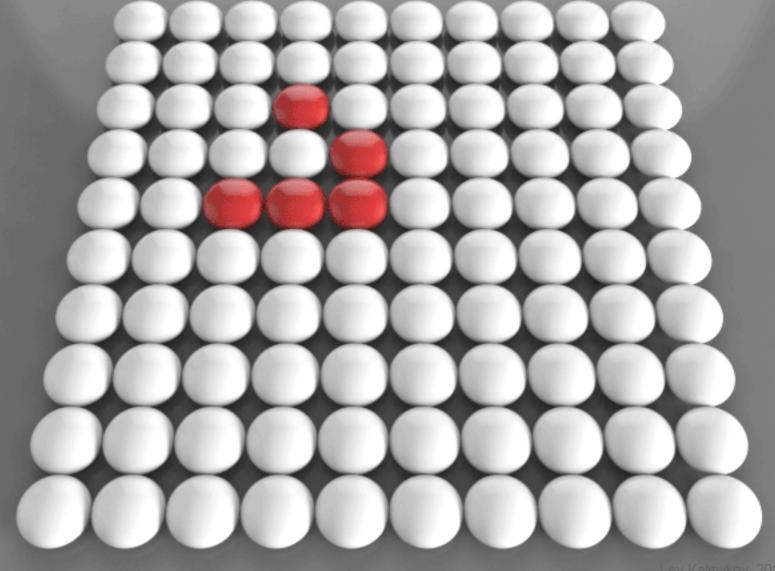
# Conway's Game of Life

- Two-dimensional grid of (square) cells
- Each cell can be in one of two states: live or dead
- Every cell only interacts with its eight nearest neighbors
- In every generation (or iteration or time step), there are some rules that decide if a cell will continue to live or die or be born (dead -> live)

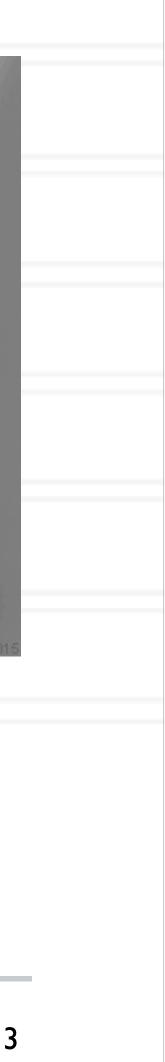
https://en.wikipedia.org/wiki/Conway's\_Game\_of\_Life







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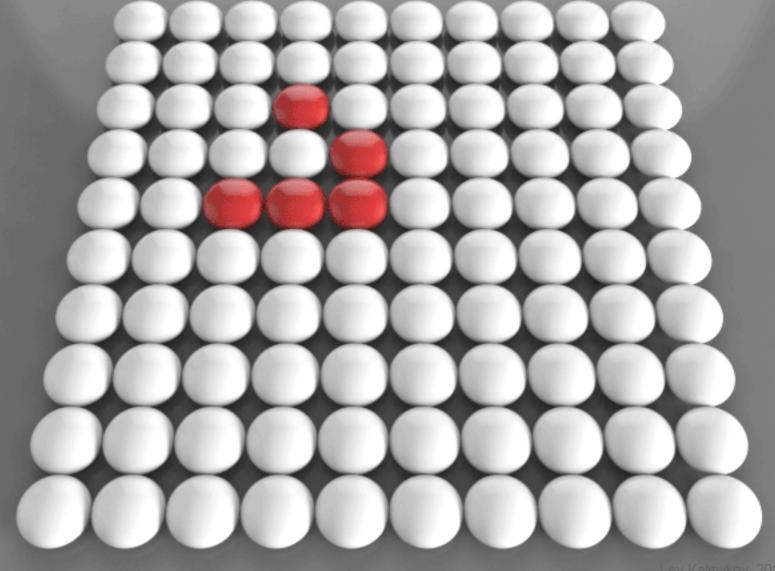
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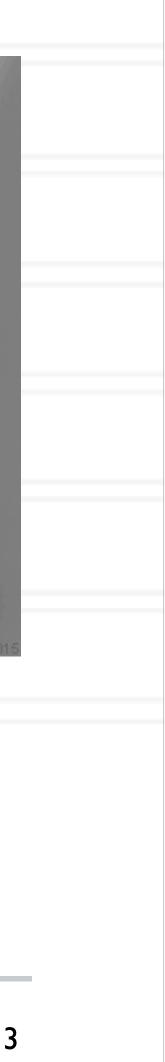
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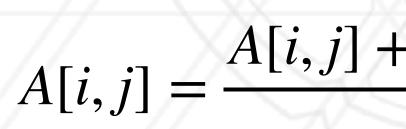


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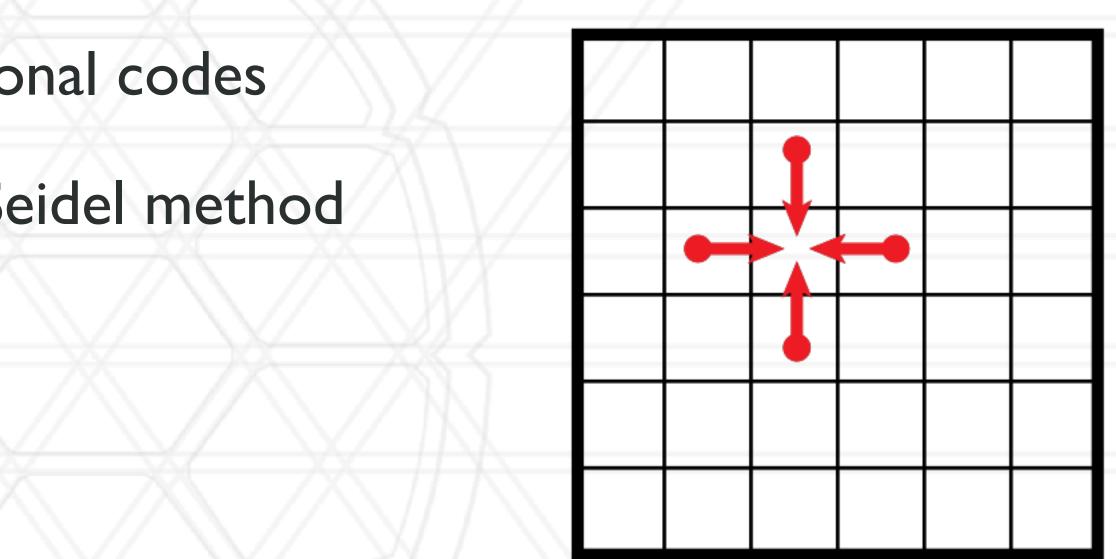
### **Two-dimensional stencil computation**

- Commonly found kernel in computational codes
- Heat diffusion, Jacobi method, Gauss-Seidel method





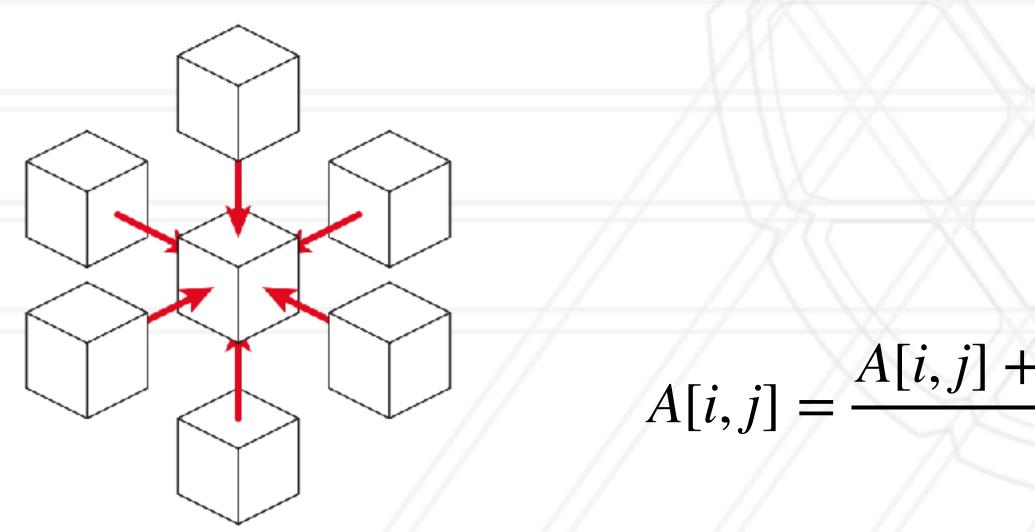
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### $A[i,j] = \frac{A[i,j] + A[i-1,j] + A[i+1,j] + A[i,j-1] + A[i,j+1]}{A[i,j] + A[i,j] + A$

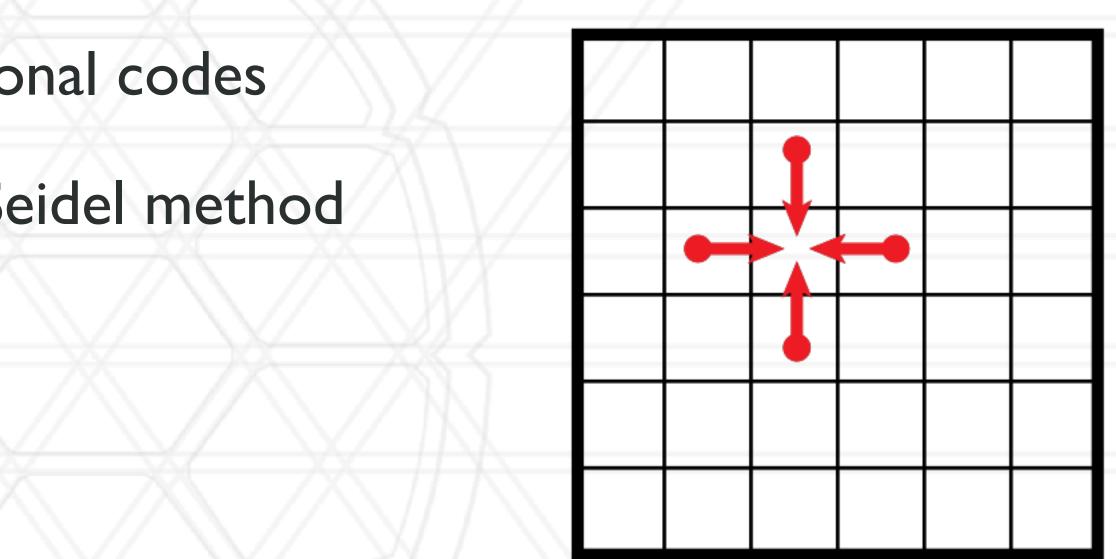
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### $A[i,j] = \frac{A[i,j] + A[i-1,j] + A[i+1,j] + A[i,j-1] + A[i,j+1]}{5}$

### Serial code

for(int t=0; t<num\_steps; t++) {
 ...</pre>

// copy contents of A\_new into A
...



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### + A[i+1, j] + A[i, j-1] + A[i, j+1]) \* 0.2



### Serial code

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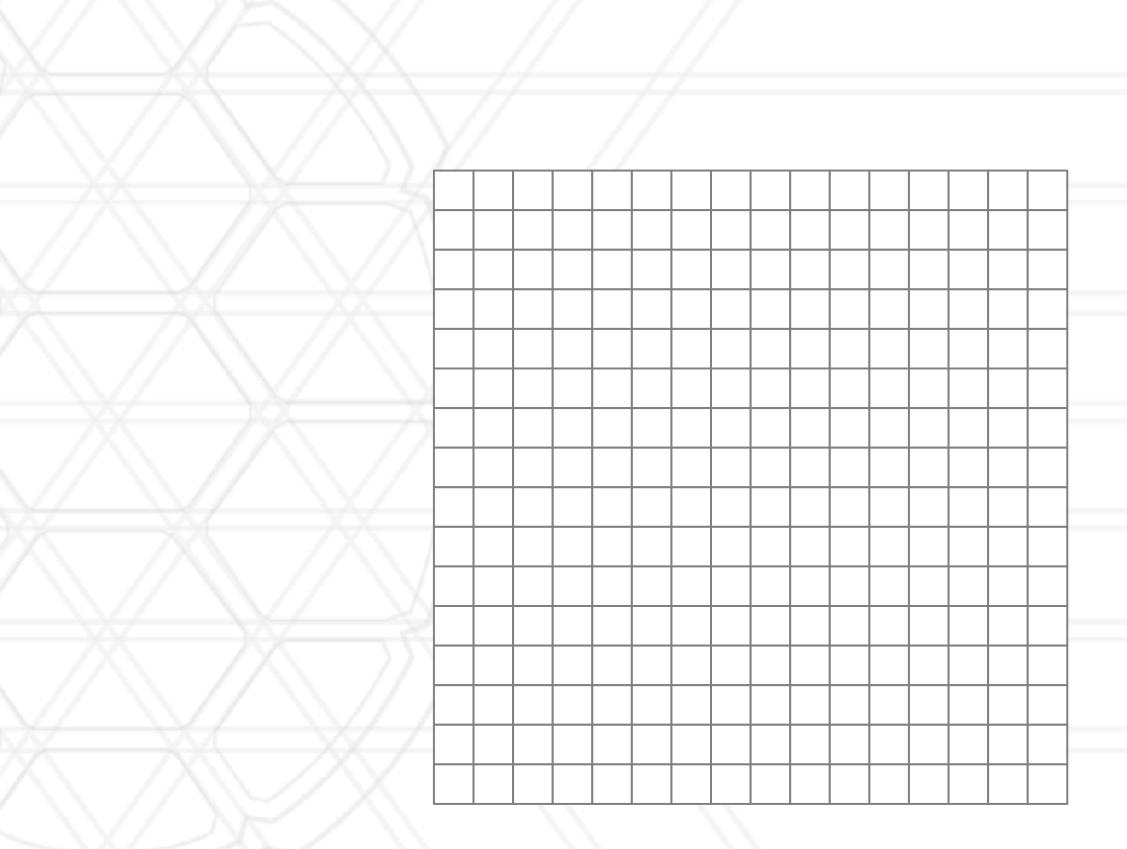
### + A[i+1, j] + A[i, j-1] + A[i, j+1]) \* 0.2

## Why do we keep two copies of A?





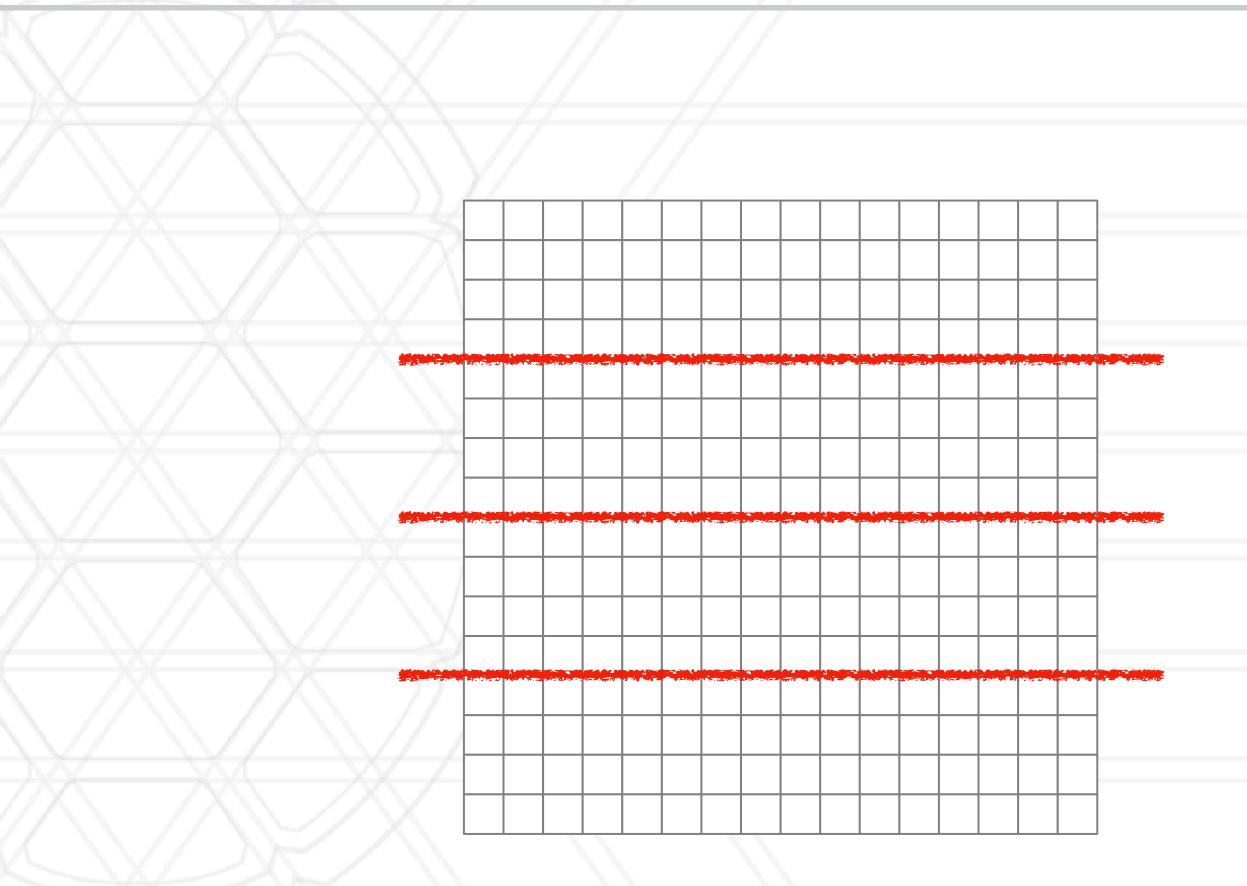
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### ID decomposition

• Divide rows (or columns) among processes

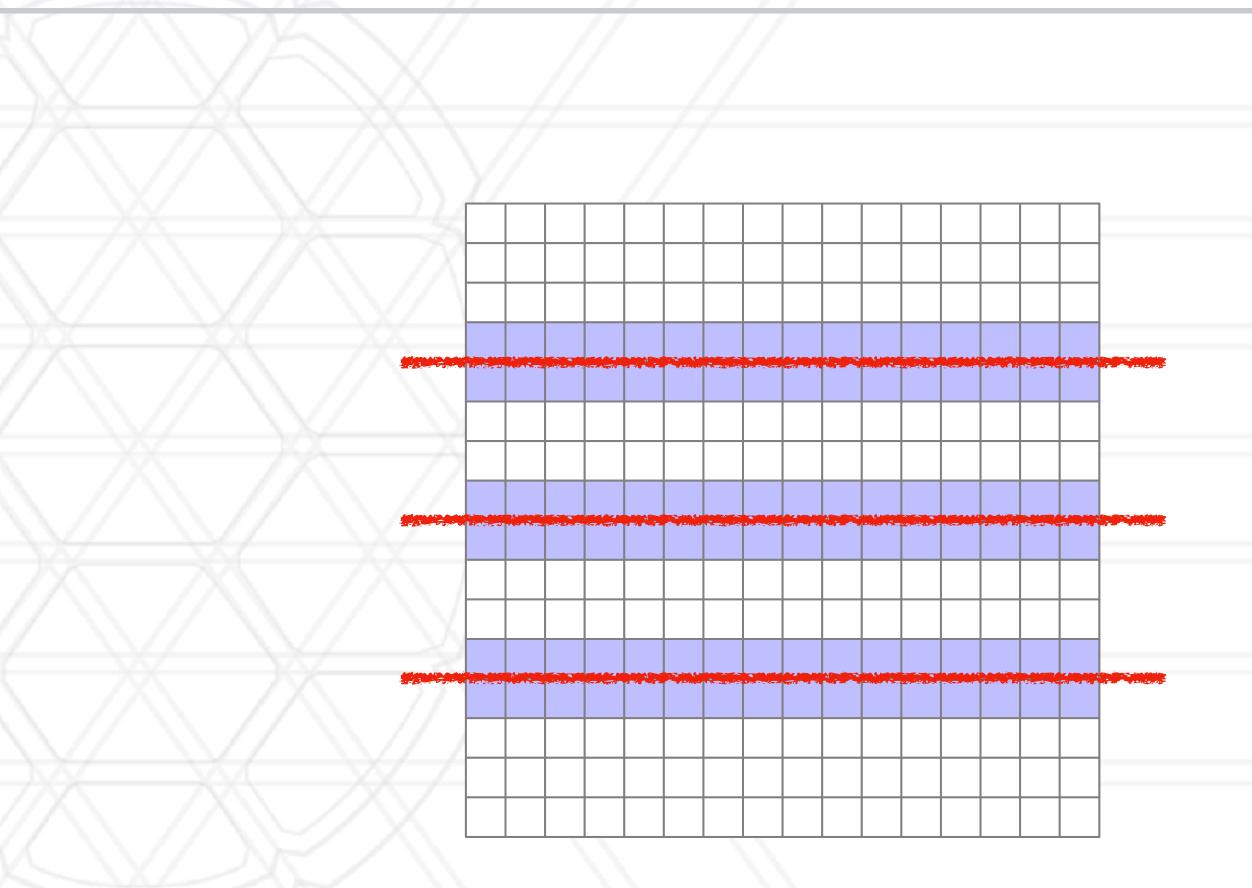




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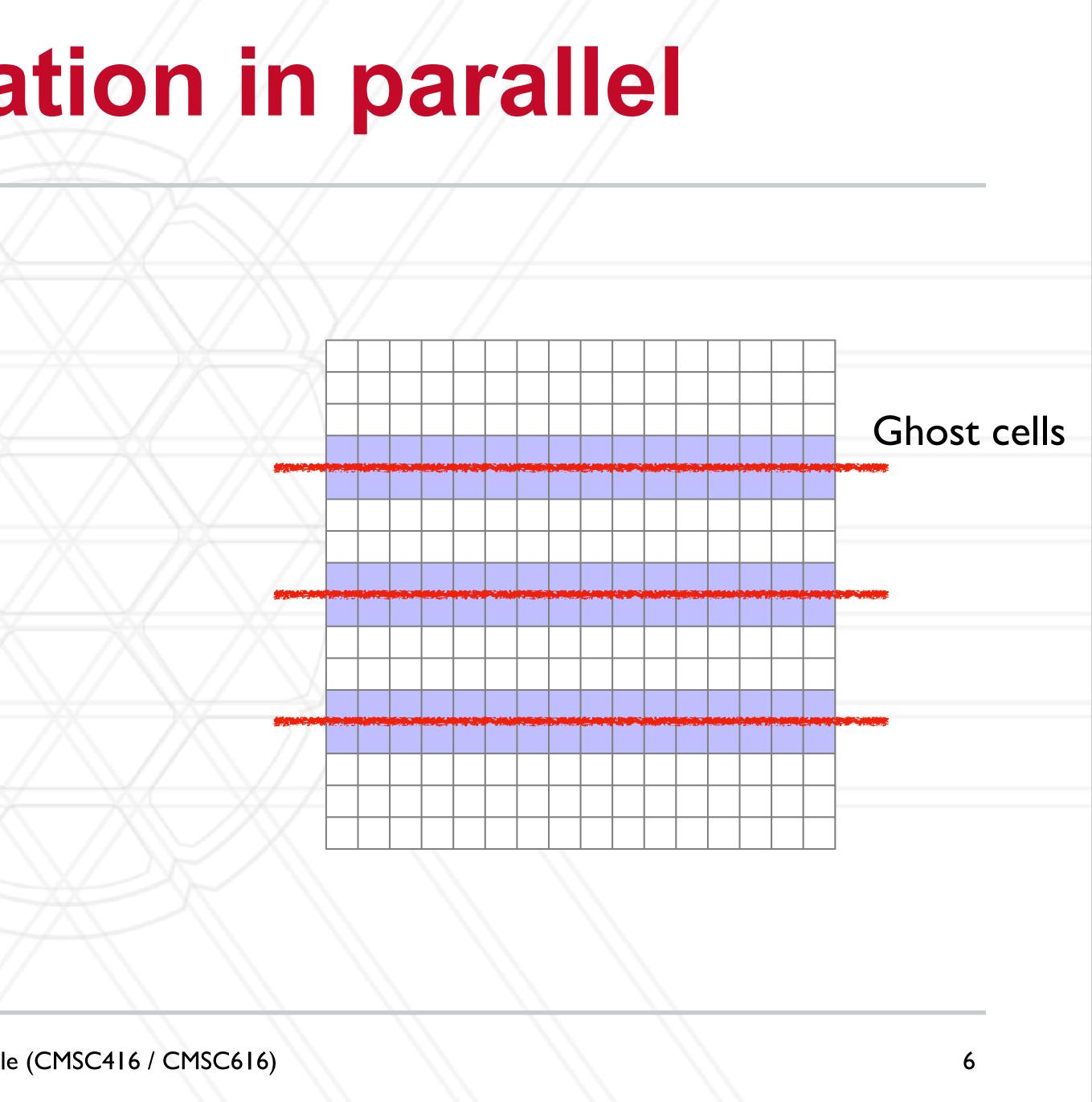




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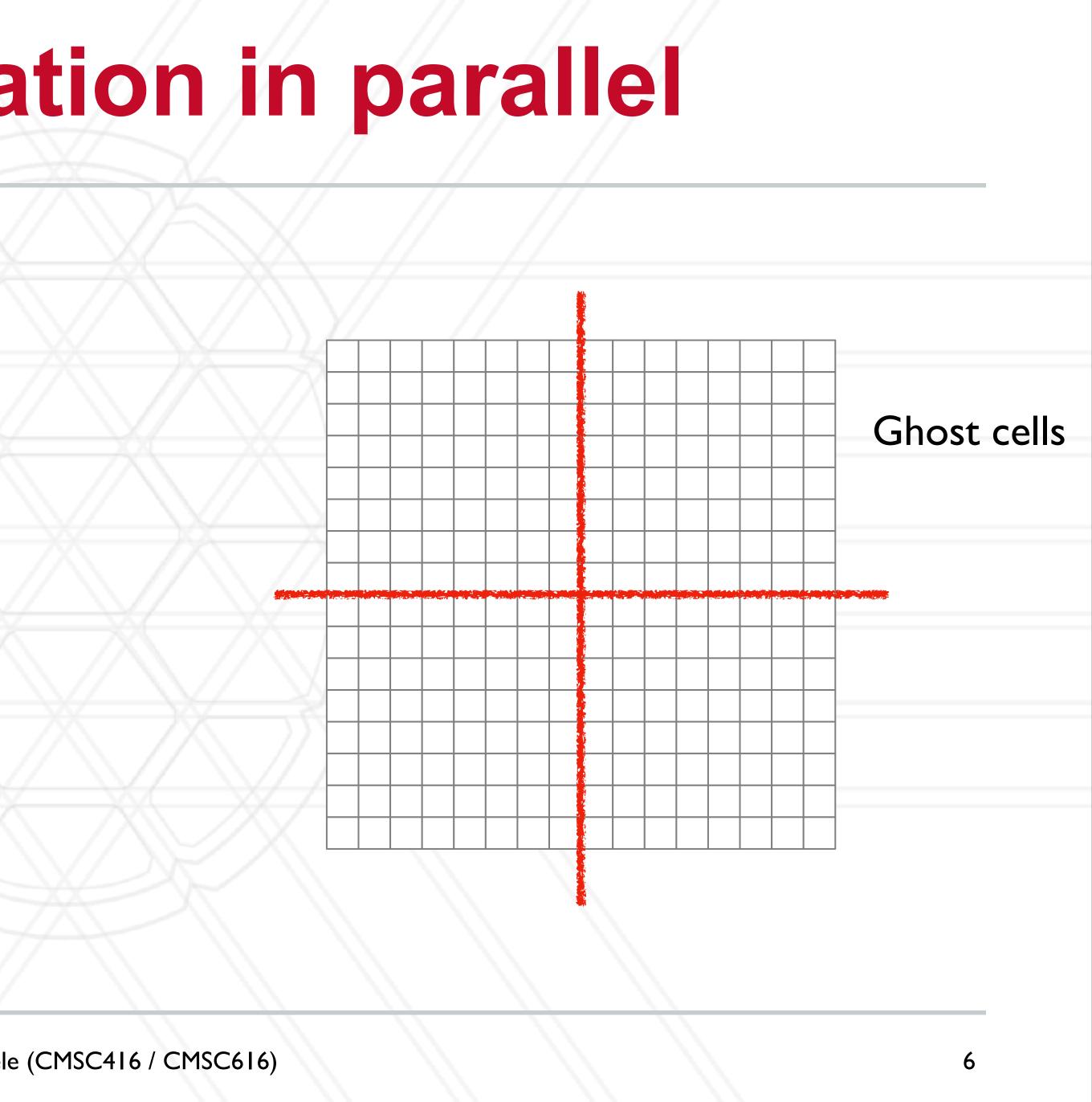
### ID decomposition

• Divide rows (or columns) among processes

### • 2D decomposition

• Divide both rows and columns (2d blocks) among processes





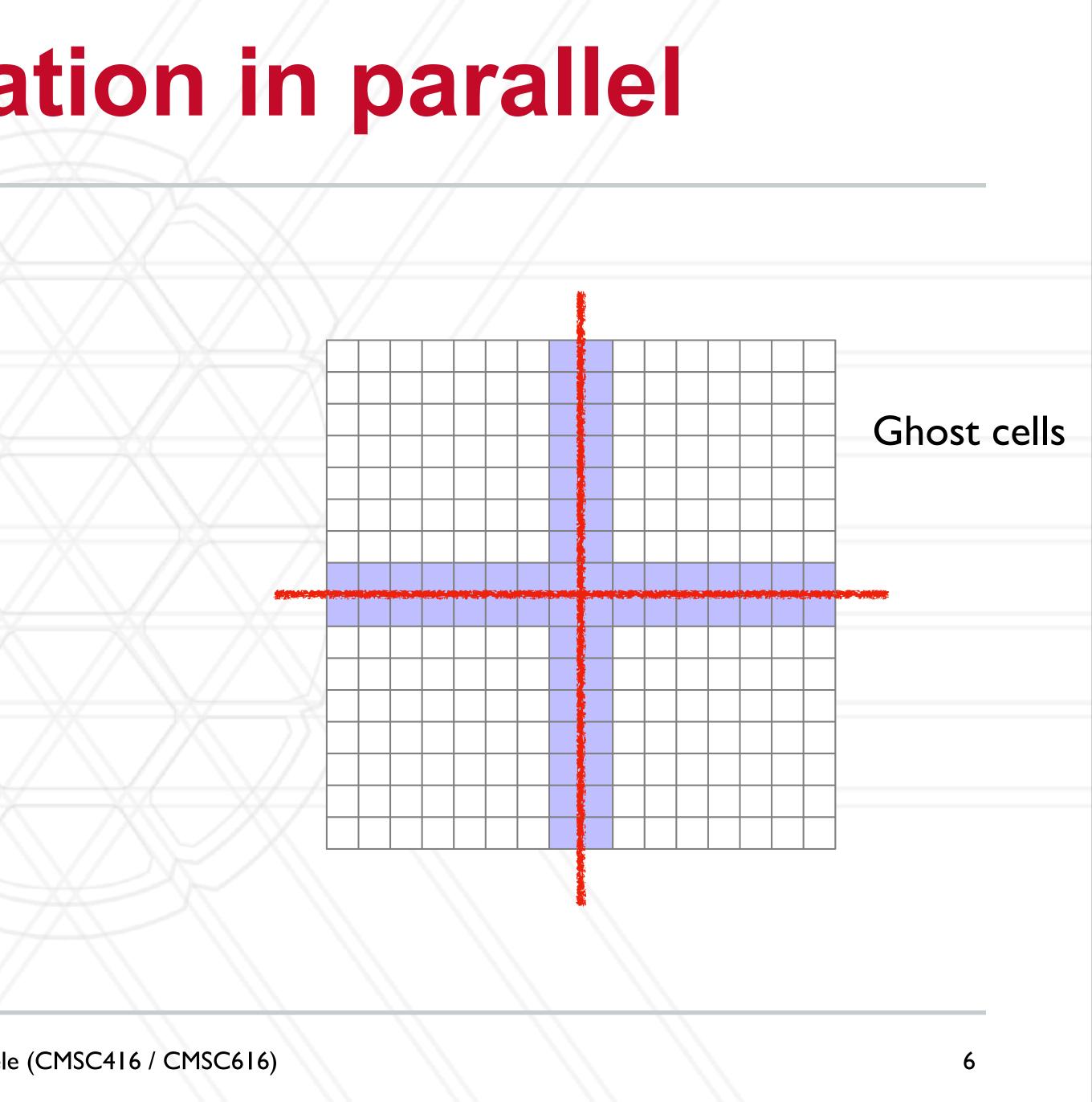
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### **Prefix sum**

- Calculate sums of prefixes (running totals) of elements (numbers) in an array
- Also called a "scan" sometimes

```
pSum[0] = A[0]
for(i=1; i<N; i++) {</pre>
    pSum[i] = pSum[i-1] + A[i]
```







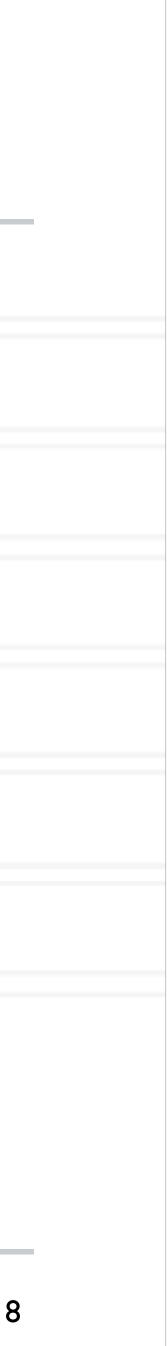
	I	2	3	4	5	6	•••
ım	Ι	3	6	10	15	21	•••



2	8	3	5



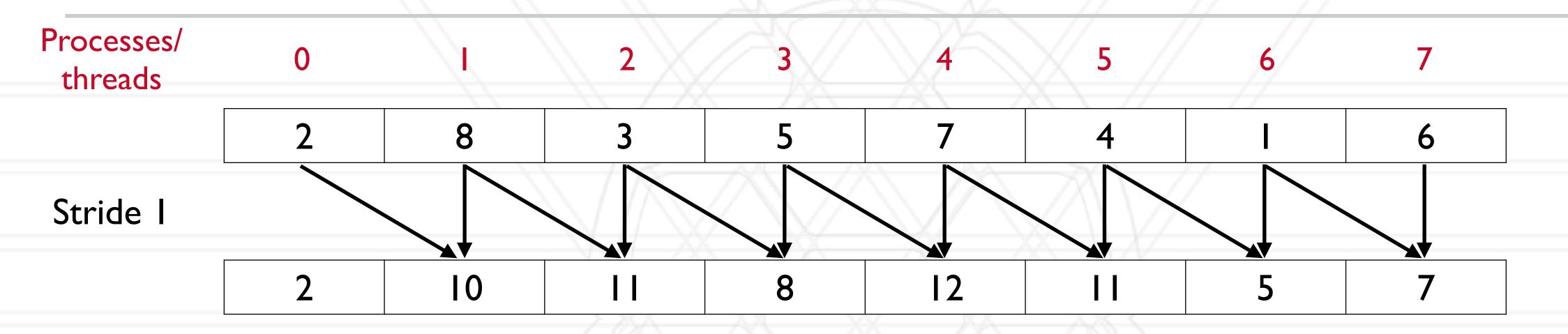
7	▲	6	
	4	6	



rocesses/ threads	0		2	3	4	5	6	7
	2	8	3	5	7	4		6

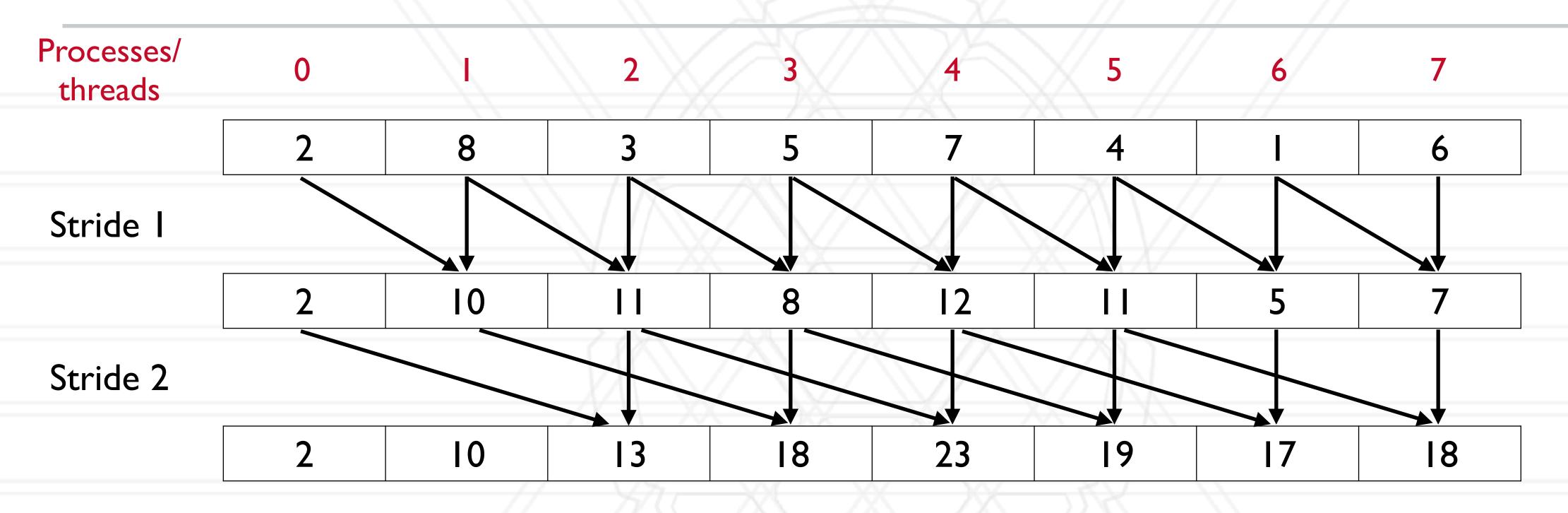




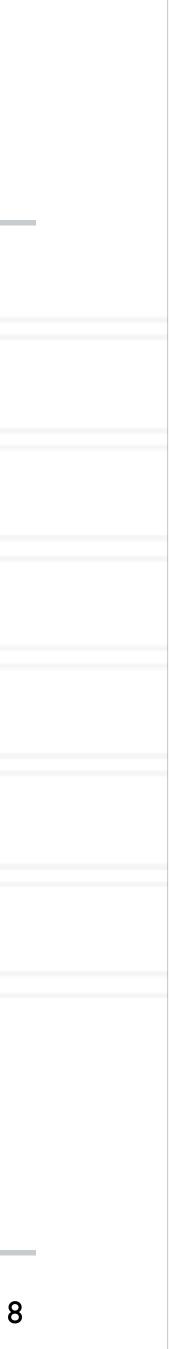


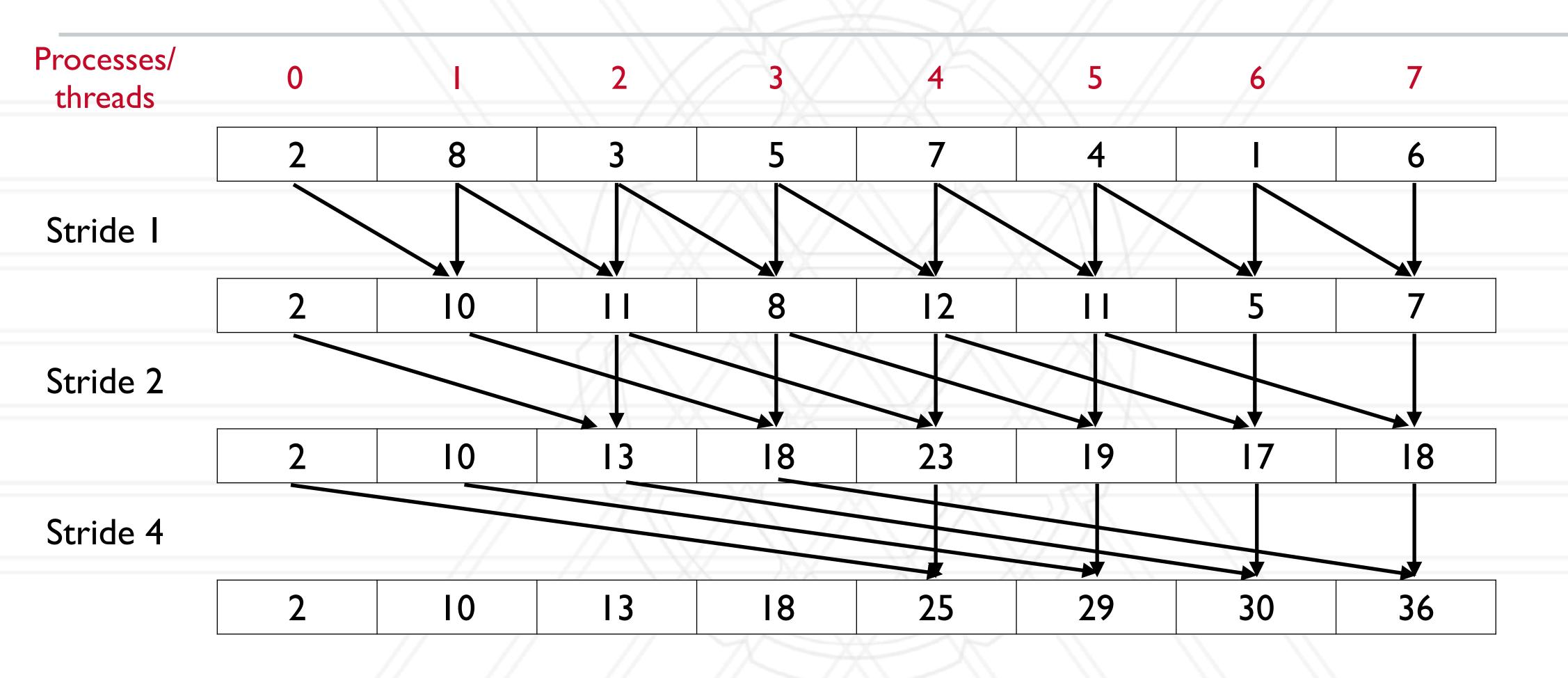




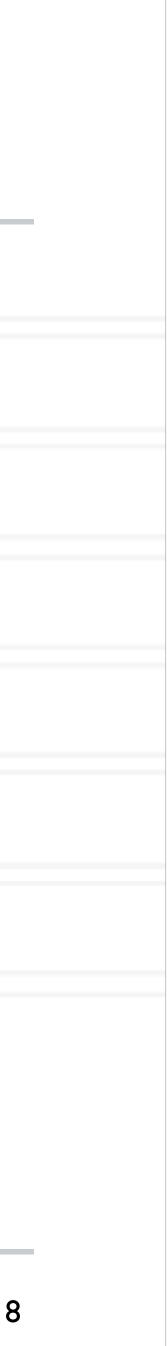




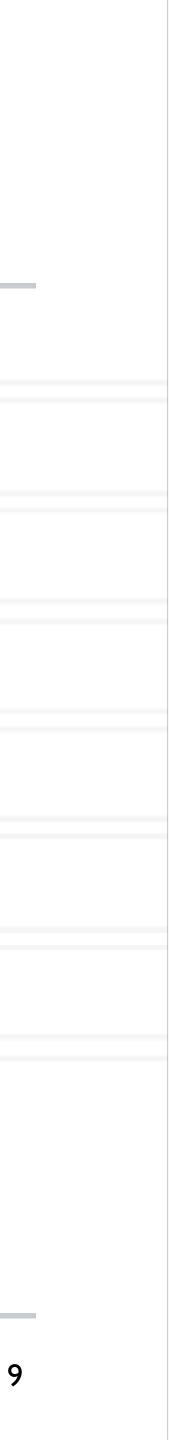






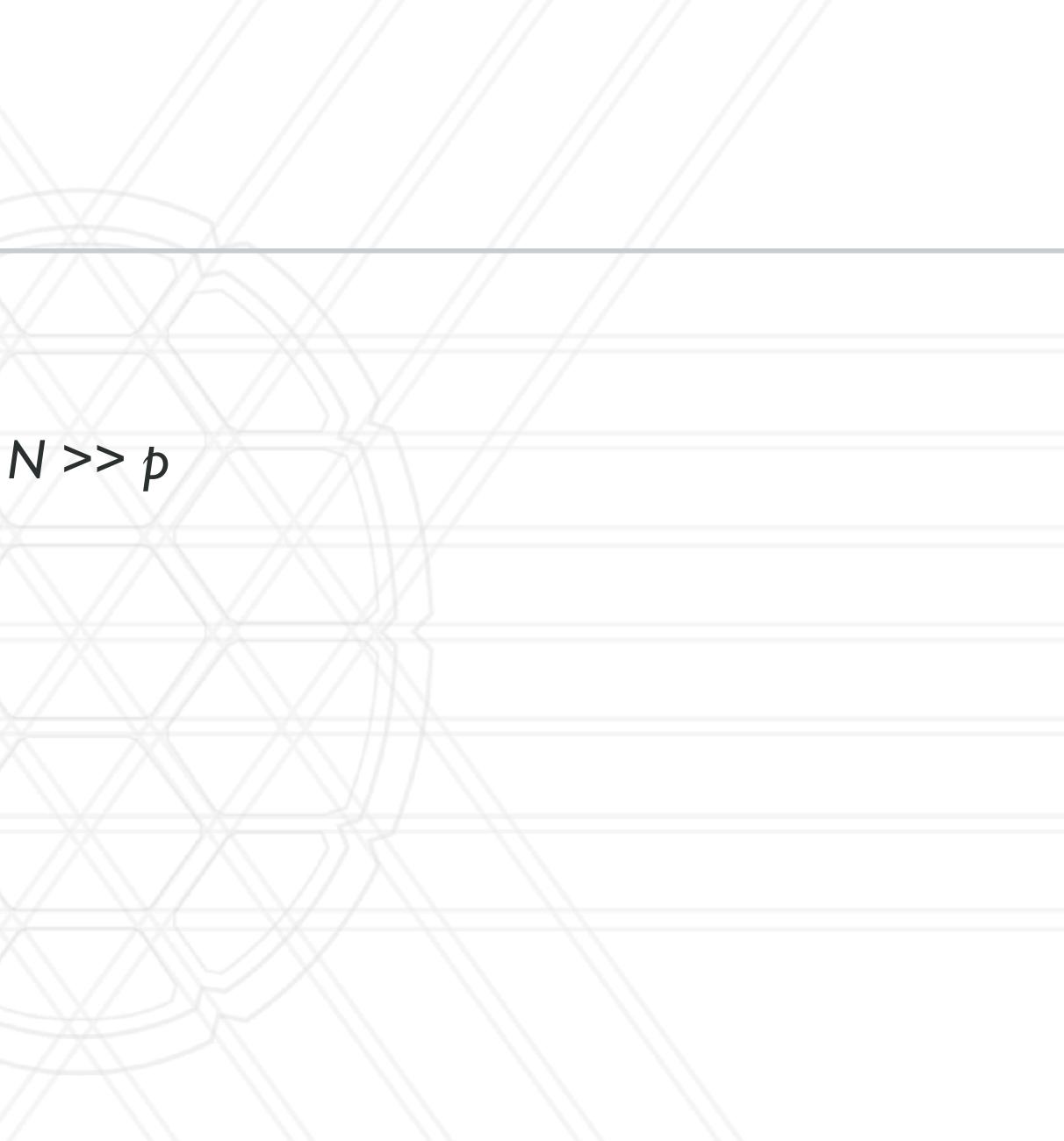






### You have N numbers and p processes, N >> p

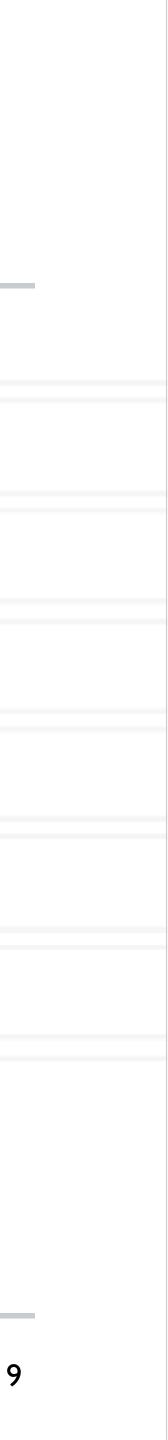






- You have N numbers and p processes, N >> p
- Assign a N/p block to each process
  - Do the serial prefix sum calculation for the blocks owned on each process locally





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  - Do the serial prefix sum calculation for the blocks owned on each process locally
- local block)
  - Last element from sending process is added to all elements in receiving process' sub-block



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• Then do parallel algorithm with partial prefix sums (using the last element from each

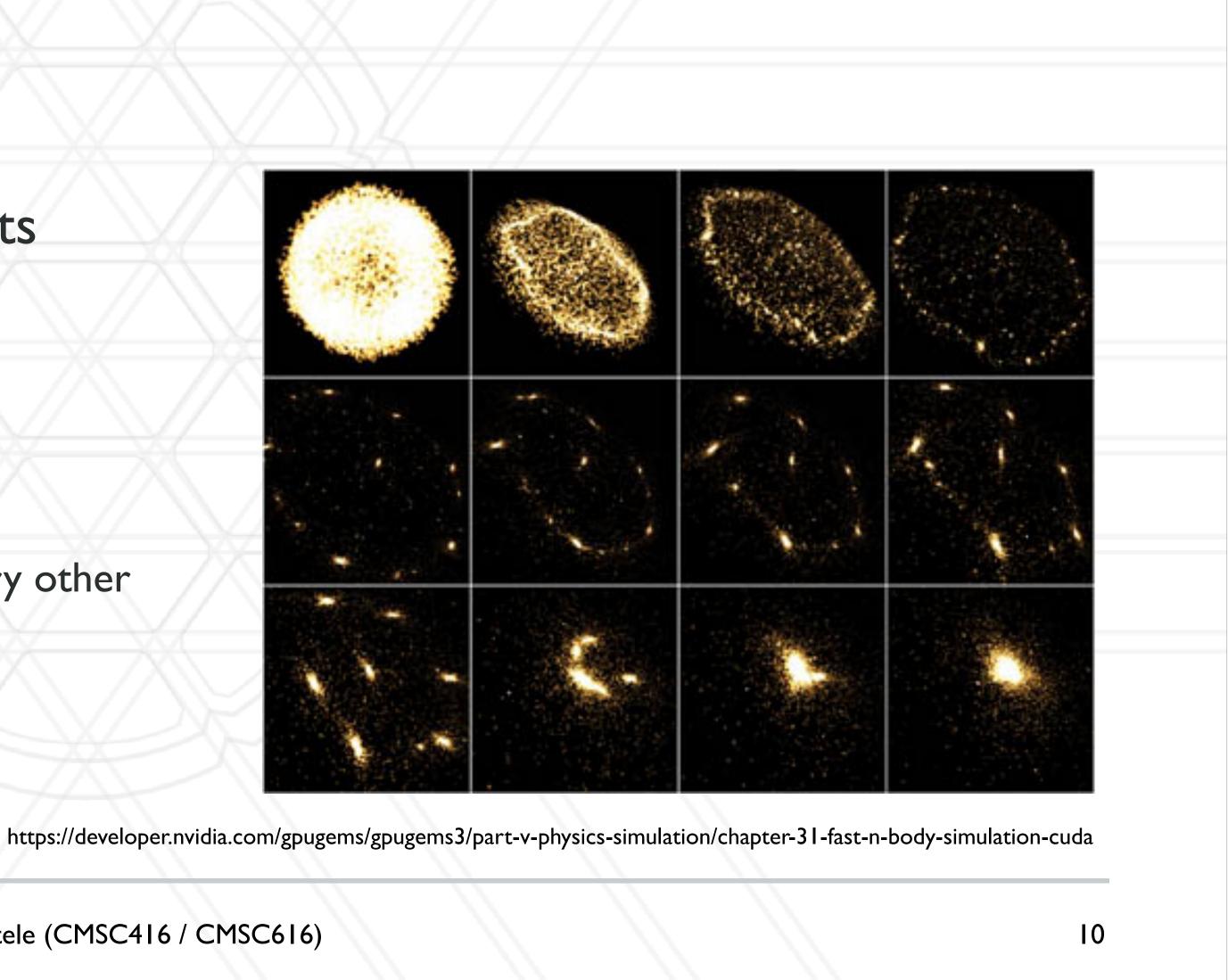


# The *n*-body problem

- Simulate the motion of celestial objects interacting with one another due to gravitational forces
- Naive algorithm:  $O(n^2)$ 
  - Every body calculates forces pair-wise with every other body (particle)





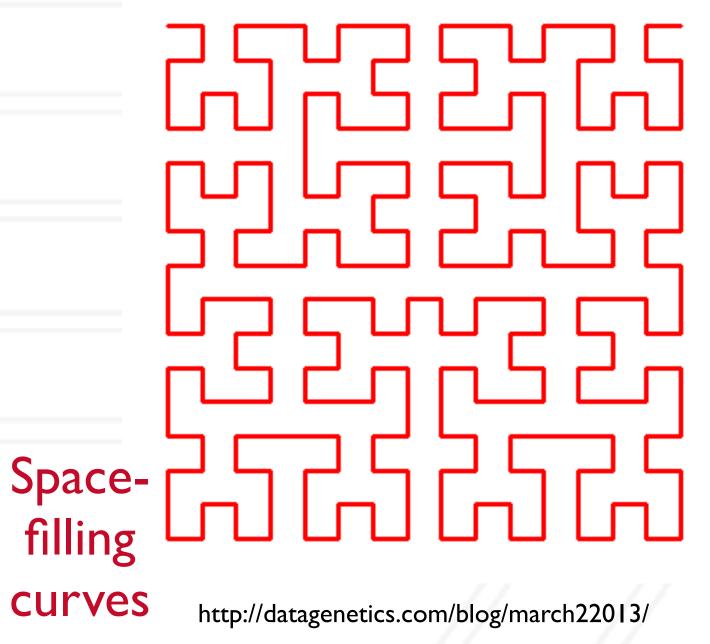


- Naive approach: Assign n/p particles to each process
- Other approaches?



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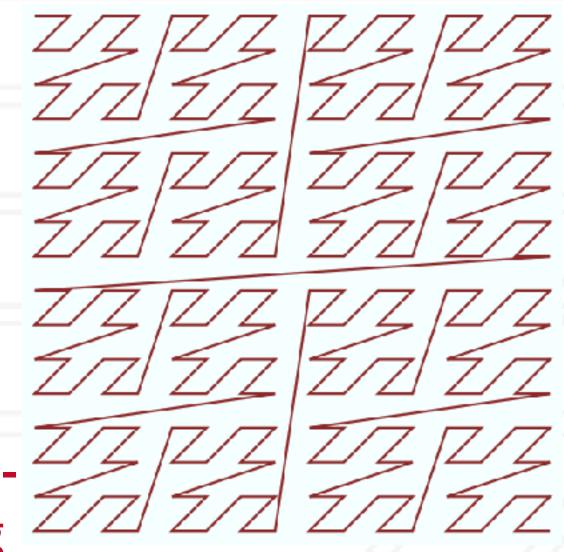


https://en.wikipedia.org/wiki/Z-order\_curve



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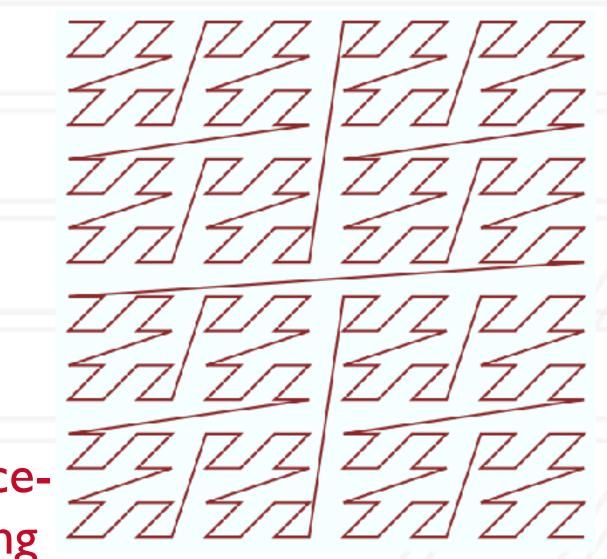


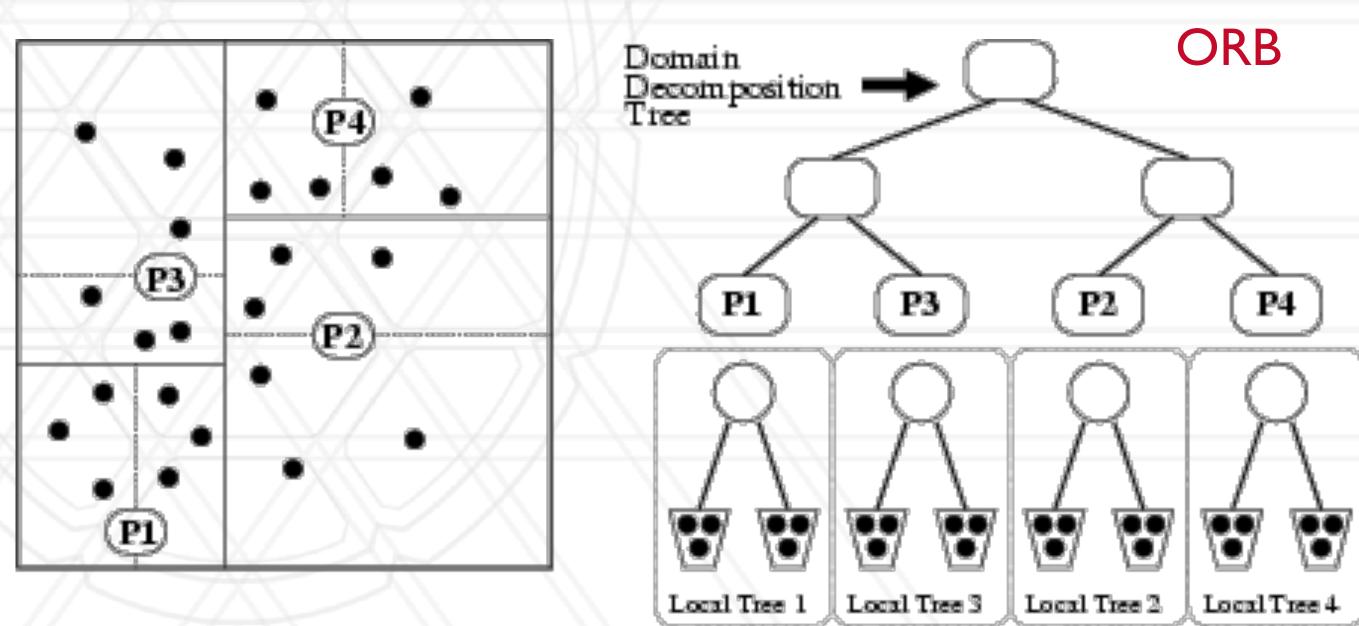
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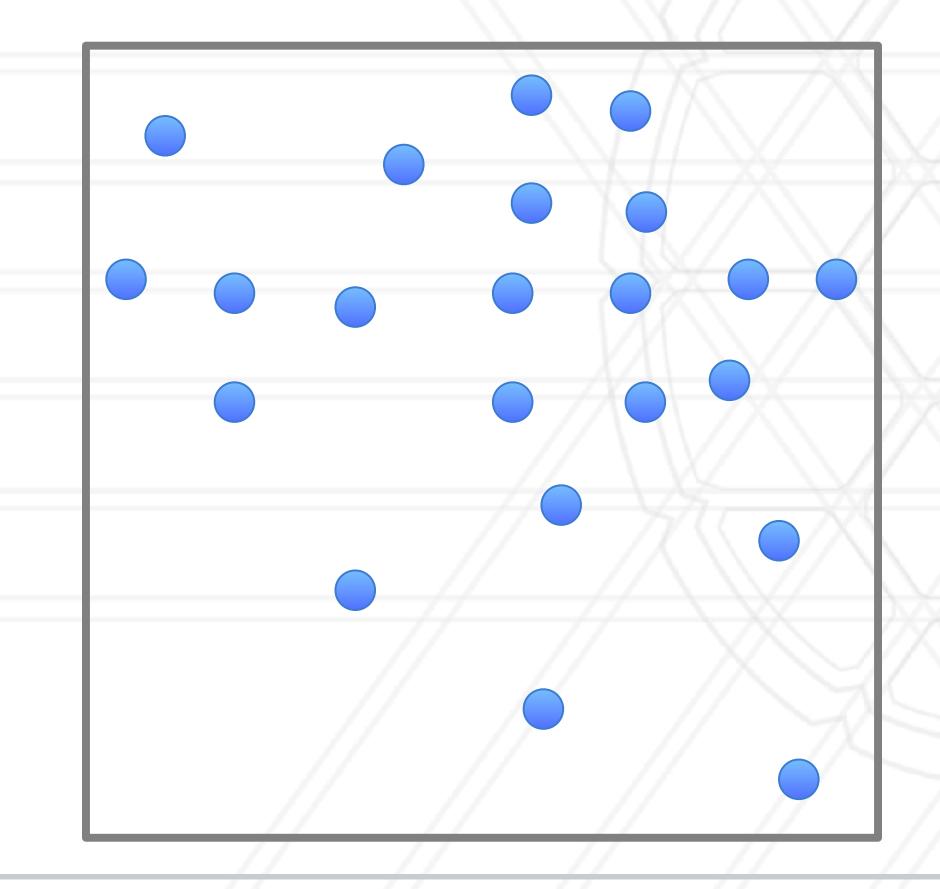


http://charm.cs.uiuc.edu/workshops/charmWorkshop2011/slides/CharmWorkshop2011 apps ChaNGa.pdf

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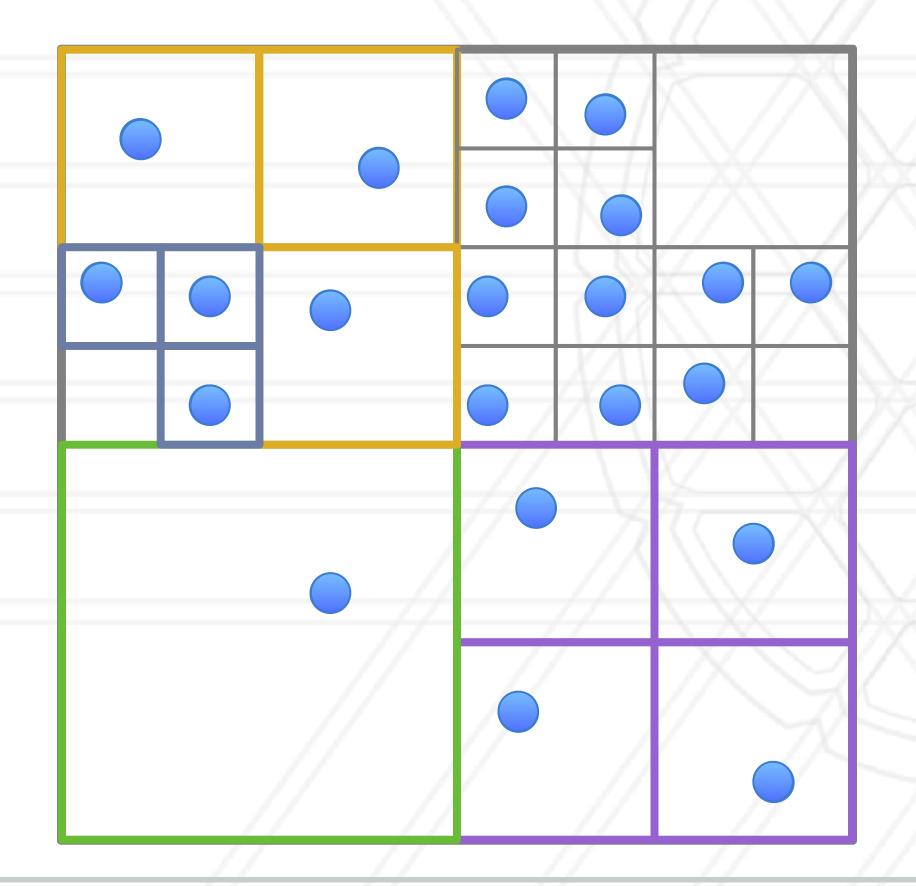
### • Let us consider a two-dimensional space with bodies/particles in it







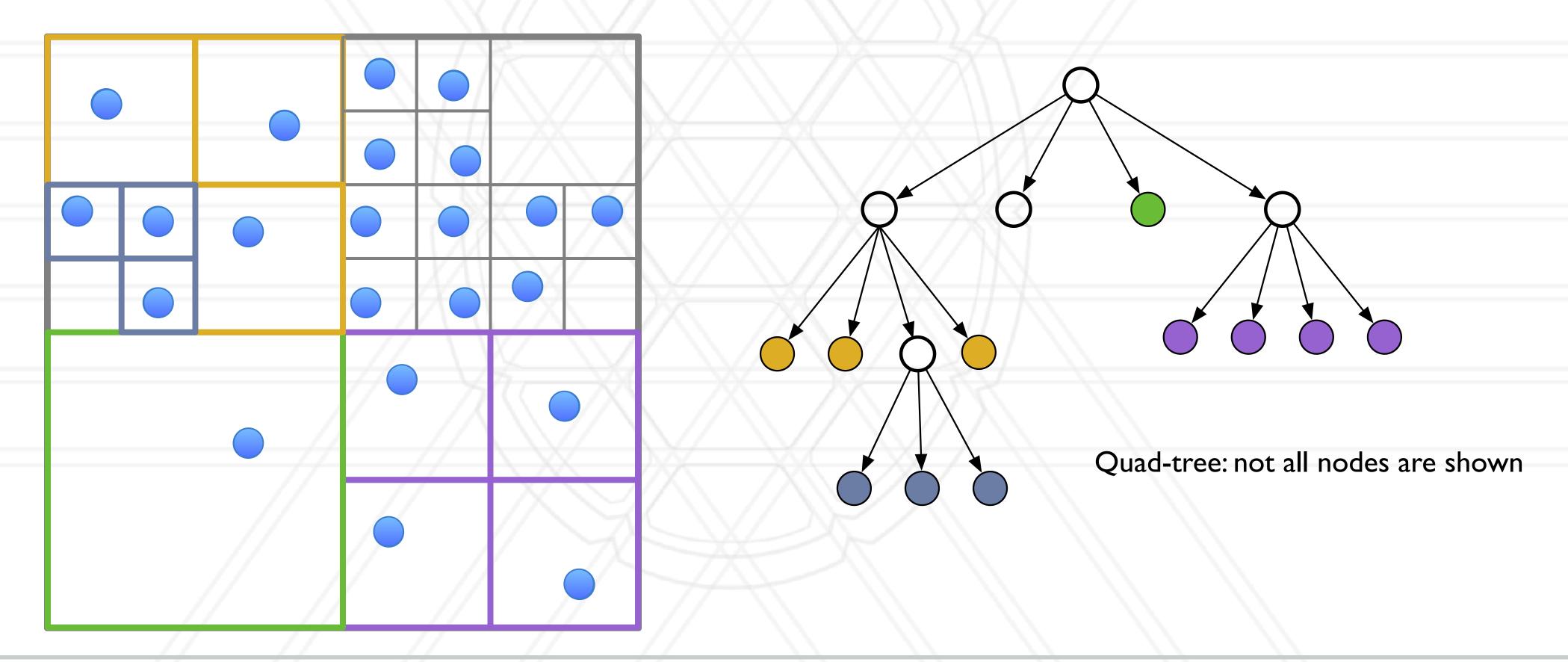
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## Load balance and grain size

- threads/ processes
  - Bring ratio of maximum to average load as close to 1.0 as possible
  - Secondary consideration: also load balance amount of communication
- Grain size: ratio of computation-to-communication
  - Coarse-grained (more computation) vs. fine-grained (more communication)



### • Load balance: try to balance the amount of work (computation) assigned to different



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