CMSC 132: OBJECT-ORIENTED PROGRAMMING II

Heaps & Priority Queues

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Complete Binary Trees

- An binary tree (height h) where
  - Perfect tree to level h-1
  - Leaves at level h are as far left as possible
Complete Binary Trees

Basic complete tree shape

Not Allowed
Heaps

- Two key properties
  - Complete binary tree (shape property)
  - Value at node (value property)
    - **Minheap** $\rightarrow$ smaller than or equal to values in subtrees ($X \leq Y$, $X \leq Z$)
    - **Maxheap** $\rightarrow$ larger than or equal to values in subtrees ($X \geq Y$, $X \geq Z$)
  - We will use minheap in our discussion
Heap (min) & Non-heap Examples

Heaps

Non-heaps
Heap Properties

• Heaps are balanced trees
  • Height = log₂(n) = O(log(n))
• Can find smallest/largest element easily
  • Always at top of heap!
  • Heap can track either min or max, but not both
Heap

- Key operations
  - Insert (X)
  - getSmallest()

- Key applications
  - Heapsort
  - Priority queue
Heap Operations – Insert( X )

• Algorithm
  • Add X to end of tree
  • While (X < parent)
    Swap X with parent  // X bubbles up tree

• Complexity
  • # of swaps proportional to height of tree
  • O( log(n) )
Heap Insert Example

- Insert (20)

1) Insert to end of tree
2) Compare to parent, swap if parent key larger
3) Insert complete
Heap Insert Example

• Insert (8)

1) Insert to end of tree

2) Compare to parent, swap if parent key larger

3) Insert complete
Heap Operation – getSmallest()

- **Algorithm**
  - Get smallest node at root
  - Replace root with X at end of tree
  - While ( \( X > \text{child} \) )
    - Swap X with smallest child // X drops down tree
  - Return smallest node

- **Complexity**
  - # swaps proportional to height of tree
  - \( O(\log(n)) \)
Heap GetSmallest Example

- `getSmallest()`

1) Replace root with end of tree
2) Compare node to children, if larger swap with smallest child
3) Repeat swap if needed
Heap GetSmallest Example

• getSmallest ()

1) Replace root with end of tree
2) Compare node to children, if larger swap with smallest child
3) Repeat swap if needed
Heap Implementation

- Can implement heap as array
  - Store nodes in array elements
  - Assign location (index) for elements using formula

- Observations
  - Compact representation
  - Edges are implicit (no storage required)
  - Works well for complete trees (no wasted space)

(a) Heap represented as a tree
(b) Heap represented as an array
Heap Implementation

- \lfloor \cdot \rfloor \rightarrow \text{floor} (\text{e.g., } 1.7 \rightarrow 1, 2 \rightarrow 2)

- Calculating node locations
  - Array index \( i \) starts at 0
  - Parent\((i) = \lfloor (i - 1) / 2 \rfloor\)
  - LeftChild\((i) = 2 \times i + 1\)
  - RightChild\((i) = 2 \times i + 2\)

(a) Heap represented as a tree

(b) Heap represented as an array
Heap Implementation

- Example
  - Parent(1) = \( \left\lfloor \frac{1 - 1}{2} \right\rfloor = \left\lfloor \frac{0}{2} \right\rfloor = 0 \)
  - Parent(2) = \( \left\lfloor \frac{2 - 1}{2} \right\rfloor = \left\lfloor \frac{1}{2} \right\rfloor = 0 \)
  - Parent(3) = \( \left\lfloor \frac{3 - 1}{2} \right\rfloor = \left\lfloor \frac{2}{2} \right\rfloor = 1 \)
  - Parent(4) = \( \left\lfloor \frac{4 - 1}{2} \right\rfloor = \left\lfloor \frac{3}{2} \right\rfloor = 1 \)
  - Parent(5) = \( \left\lfloor \frac{5 - 1}{2} \right\rfloor = \left\lfloor \frac{4}{2} \right\rfloor = 2 \)
Heap Implementation

• Example
  • LeftChild(0) = 2 × 0 +1 = 1
  • LeftChild(1) = 2 × 1 +1 = 3
  • LeftChild(2) = 2 × 2 +1 = 5
Heap Implementation

- Example
  - RightChild(0) = 2 × 0 + 2 = 2
  - RightChild(1) = 2 × 1 + 2 = 4
Heap Application – Heapsort

- Use heaps to sort values
  - Heap keeps track of smallest/largest element in heap
- Algorithm
  1. Create heap
  2. Insert values in heap
  3. Remove values from heap (in ascending/descending order)
- Complexity
  - $O(n \log(n))$
Heapsort Example

• Input
  • 11, 5, 13, 6, 1
• View heap during insert, removal
  • As tree
  • As array
Heapsort – Insert Values

(a) Insert 11
   11

(b) Insert 5
   11
   5

(c) Rebuild heap
   11
   5
   11

(d) Insert 13
   5
   11
   13

(e) Insert 6
   5
   11
   13
   6

(f) Rebuild heap
   5
   11
   13
   6
   11

(g) Insert 1
   5
   6
   11
   13

(h) Rebuild heap
   1
   11
   6
   13
   5
   11
   6
Heapsort – Remove Values

(a) Print root = 1

1

11 6

5 13

(b) Rebuild heap

6

11

13 6

5 13

(c) Print root = 5

5

11

6 13

(d) Rebuild heap

11

6

13 11 13

6 13

(e) Print root = 6

6

11

13

(f) Rebuild heap

13

11

11 13

13

(g) Print root = 11

11

11

13

(h) Rebuild heap

13

(f) Print root = 13

13

Done
Heapsort – Insert in to Array 1

- **Input**
  - 11, 5, 13, 6, 1

```
Index = 0 1 2 3 4
Insert 11 11
```
Heapsort – Insert in to Array 2

- Input
  - 11, 5, 13, 6, 1

<table>
<thead>
<tr>
<th>Index</th>
<th>0</th>
<th>1</th>
<th>2</th>
<th>3</th>
<th>4</th>
</tr>
</thead>
<tbody>
<tr>
<td>Insert 5</td>
<td>11</td>
<td>5</td>
<td></td>
<td></td>
<td></td>
</tr>
</tbody>
</table>

Insert 5

<table>
<thead>
<tr>
<th>Swap</th>
<th>5</th>
<th>11</th>
<th></th>
<th></th>
<th></th>
</tr>
</thead>
</table>

Swap
Heapsort – Insert in to Array 3

• Input
  • 11, 5, 13, 6, 1

Index = 0 1 2 3 4

Insert 13 5 11 13
Heapsort – Insert in to Array 4

- Input
  - 11, 5, 13, 6, 1

<table>
<thead>
<tr>
<th>Index</th>
<th>0</th>
<th>1</th>
<th>2</th>
<th>3</th>
<th>4</th>
</tr>
</thead>
<tbody>
<tr>
<td>Insert 6</td>
<td>5</td>
<td>11</td>
<td>13</td>
<td>6</td>
<td></td>
</tr>
<tr>
<td>Swap</td>
<td>5</td>
<td>6</td>
<td>13</td>
<td>11</td>
<td></td>
</tr>
</tbody>
</table>

...
# Heapsort – Remove from Array 1

- **Input**
  - 11, 5, 13, 6, 1

<table>
<thead>
<tr>
<th>Index</th>
<th>0</th>
<th>1</th>
<th>2</th>
<th>3</th>
<th>4</th>
</tr>
</thead>
<tbody>
<tr>
<td>Remove root</td>
<td>1</td>
<td>5</td>
<td>13</td>
<td>11</td>
<td>6</td>
</tr>
<tr>
<td>Replace</td>
<td>6</td>
<td>5</td>
<td>13</td>
<td>11</td>
<td></td>
</tr>
<tr>
<td>Swap w/ child</td>
<td>5</td>
<td>6</td>
<td>13</td>
<td>11</td>
<td></td>
</tr>
</tbody>
</table>
Heapsort – Remove from Array 2

- Input
  - 11, 5, 13, 6, 1

Index =

<table>
<thead>
<tr>
<th>Index</th>
<th>0</th>
<th>1</th>
<th>2</th>
<th>3</th>
<th>4</th>
</tr>
</thead>
<tbody>
<tr>
<td>Remove root</td>
<td>5</td>
<td>6</td>
<td>13</td>
<td>11</td>
<td></td>
</tr>
<tr>
<td>Replace</td>
<td>11</td>
<td>6</td>
<td>13</td>
<td></td>
<td></td>
</tr>
<tr>
<td>Swap w/ child</td>
<td>6</td>
<td>11</td>
<td>13</td>
<td></td>
<td></td>
</tr>
</tbody>
</table>
Heap Application – Priority Queue

- Queue
  - Linear data structure
  - First-in First-out (FIFO)
  - Implement as array / linked list

Enqueue  Dequeue
Heap Application – Priority Queue

- Priority queue
  - Elements are assigned priority value
  - Higher priority elements are taken out first
  - Implement as heap
    - Enqueue ⇒ insert( )
    - Dequeue ⇒ getSmallest( )
**Priority Queue**

- **Properties**
  - Lower value = higher priority
  - Heap keeps highest priority items in front

- **Complexity**
  - Enqueue \(\Rightarrow\) `insert( )` = \(O(\log(n))\)
  - Dequeue \(\Rightarrow\) `getSmallest( )` = \(O(\log(n))\)
  - For any heap
Heap vs. Binary Search Tree

- Binary search tree
  - Keeps values in sorted order
  - Find any value
    - $O(\log(n))$ for balanced tree
    - $O(n)$ for degenerate tree (worst case)

- Heap
  - Keeps smaller values in front
  - Find minimum value
    - $O(\log(n))$ for any heap
About Heap Implementation

- Implementing a heap ([Video I](#)/ [Video II](#)) This videos illustrates the process a programmer (Prof. Bill Pugh in this case) goes through while implementing code. This video was filmed in Dr. Bill Pugh's lecture. Keep in mind that in this video some bugs might be present in the implementation as the testing phase has not been completed yet.