START RECORDING

# Constructive Induction

CMSC 250

# Introductory Example

- We already know that

$$\sum_{i=1}^{n} i = \frac{n(n+1)}{2} = \frac{n^2}{2} + \frac{n}{2}$$

- **But how? Who told us this?**
- This is **not** how math works; we usually do **not** know the answer ahead of time!

# Making a Good Guess with Calculus

- Calculus tells us that (discrete) sums are approximations of (continuous) integrals.
- Then, we can observe that:

$$\sum_{i=1}^{n} i \approx \int_{1}^{n} x\,dx = \frac{1}{2}n^2 + c, \qquad c \in \mathbb{R}$$

- So we know that the sum ought to be *some* quadratic function of $n$.

# Making a Good Guess with CS

- Another way to guess the quadratic form would be with plotting!
- Suppose $f(n) = \sum_{i=1}^{n} i$. Then:

  - $f(0) = \sum_{i=1}^{0} i = 0$
  - $f(1) = \sum_{i=1}^{1} i = 1$
  - $f(2) = \sum_{i=1}^{2} i = 1 + 2 = 3$
  - $f(3) = \sum_{i=1}^{3} i = 1 + 2 + 3 = 6$
  - …
  - $f(30) = \sum_{i=1}^{30} i = 1 + 2 + \cdots + 30 = 465$

- We can then fit a curve and see the quadratic curve by ourselves!

# Making a Good Guess

- We saw that the sum is some quadratic polynomial. This is all we know!
- So $\sum_{i=1}^{n} i$ is some $poly(n)$ with degree 2, i.e

$$\sum_{i=1}^{n} i = An^2 + Bn + C, \qquad A, B, C \in \mathbb{R}$$

- ***How to determine A, B, and C?***

# General Logic

- Solve **as if** you had an inductive proof (so IB, IH, IS)
- For every step, we will establish **conditions** on A, B,C **such that** the relevant step is correct.
    - Contrast this with **directly proving** that every step is correct.

# Constant $C$

- IB: LHS is $\sum_{i=1}^{0} i = 0$. For RHS to be equal to LHS we need:

$$An^2 + Bn + C = 0 \Rightarrow C = 0$$

- So we already know that $C = 0$.

# Co-efficients $A, B$

- IH: Assume that the proposition holds for $n \geq 0$. Then:

$$\sum_{i=1}^{n} i = An^2 + Bn$$

- IS: We want to prove that

$$\left(\sum_{i=1}^{n} i = An^2 + Bn\right) \Rightarrow \left(\sum_{i=1}^{n+1} i = A(n+1)^2 + B(n+1)\right)$$

# Co-efficients $A, B$

- IH: Assume that the proposition holds for $n \geq 0$. Then:

$$\sum_{i=1}^{n} i = An^2 + Bn \rightarrow P(n)$$

- IS: We want to prove that

$$\underbrace{\left(\sum_{i=1}^{n} i = An^2 + Bn\right)}_{P(n)} \Rightarrow \underbrace{\left(\sum_{i=1}^{n+1} i = A(n+1)^2 + B(n+1)\right)}_{P(n+1)}$$

# Co-efficients $A, B$

$$\sum_{i=1}^{n+1} i = \sum_{i=1}^{n} i + (n+1) \stackrel{\text{IH}}{=} An^2 + Bn + (n+1)$$

- We have to equate this to $A(n+1)^2 + B(n+1)$, since this is what we're trying to prove:

$$An^2 + Bn + (n+1) = A(n+1)^2 + B(n+1) \Rightarrow$$
$$\cancel{An^2} + \cancel{Bn} + (n+1) = \cancel{An^2} + 2An + A + \cancel{Bn} + B \Rightarrow$$
$$n + 1 = 2An + (A + B)$$

# Co-efficients $A, B$

$$n + 1 = 2An + (A + B)$$

- This is an equality between polynomials of $k$, so equating the co-efficients yields:

$$1 = 2A$$
$$A + B = 1$$

# Co-efficients $A, B$

$$n + 1 = 2An + (A + B)$$

- This is an equality between polynomials in $n$, so equating the co-efficients yields:

$$1 = 2A$$
$$A + B = 1$$

- Note: The IS did not end up with **TRUE**, but with conditions on A,B **for it to be TRUE.**

# All Our Constraints

1. $C = 0$
2. $A + B = 1$
3. $2 \cdot A = 1$

- Algebra yields $A = B = {}^{1}/_{2}$, so:

$$\sum_{i=0}^{n} i = \frac{1}{2}n^2 + \frac{1}{2}n + 0 = \frac{n(n+1)}{2}$$

# What if Our Guess is Wrong (Over)?

1. Suppose we guess

$$\sum_{i=1}^{n} i = A \cdot n^3 + B \cdot n^2 + C \cdot n + D$$

2. **This still works**, we will just find $A = 0$ (try it at home!)

# What if Our Guess is Wrong (Under)?

1. Suppose we guess

$$\sum_{i=1}^{n} i = A \cdot n + B$$

2. **This does not work (infeasible equation),** no $A, B \in \mathbb{R}$ will satisfy the constraints (try it at home!)

# Another Example (with Bounds!)

- Let $a$ be a sequence defined as follows:

$$a_n = \begin{cases} 2, & n = 0 \\ 50, & n = 1 \\ 10a_{n-1} + 3a_{n-2}, & n \geq 2 \end{cases}$$

- Task: Find an upper bound for $a_n$.

# Another Example (with Bounds!)

- Let $a$ be a sequence defined as follows:

$$a_n = \begin{cases} 2, & n = 0 \\ 50, & n = 1 \\ 10a_{n-1} + 3a_{n-2}, & n \geq 2 \end{cases}$$

- Task: Find an upper bound for $a_n$.
- What kind of inductive structure am I expecting?

Weak

Strong

# Another Example (with Bounds!)

- Let $a$ be a sequence defined as follows:

$$a_n = \begin{cases} 2, & n = 0 \\ 50, & n = 1 \\ 10a_{n-1} + 3a_{n-2}, & n \geq 2 \end{cases}$$

- Task: Find an upper bound for $a_n$.
- What kind of inductive structure am I expecting?

Weak | Strong

An inductive base with > 1 elements and a recursive rule with references to two prior terms hints towards strong induction…

# Key Step

$$a_n = \begin{cases} 2, & n = 0 \\ 50, & n = 1 \\ 10a_{n-1} + 3a_{n-2}, & n \geq 2 \end{cases}$$

- Because of our experience with sequences like Fibonacci, Tribonacci that all have this form, we **suspect**:

$$a_n \leq C \cdot D^n, \quad C, D \in \mathbb{R}$$

# Constraints on $C$

- IB:
  - $a_0 \leq C \cdot D^0 \Leftrightarrow 2 \leq C$
  - $a_1 \leq C \cdot D^1 \Leftrightarrow 50 \leq C \cdot D$

# Inductive Hypothesis

- IB:
  - $a_0 \leq C \cdot D^0 \Leftrightarrow 2 \leq C$
  - $a_1 \leq C \cdot D^1 \Leftrightarrow 50 \leq C \cdot D$
- IH: Let $n \geq 1$. Assume that $(\forall i \in \{0, 1, 2, \ldots n\})[a_i \leq C \cdot D^i]$

# Inductive Step

- IB:
  - $a_0 \leq C \cdot D^0 \Leftrightarrow 2 \leq C$
  - $a_1 \leq C \cdot D^1 \Leftrightarrow 50 \leq C \cdot D$
- IH: Let $n \geq 1$. Assume that $\forall i \in \{0, 1, 2, \ldots n\},\ a_i \leq C \cdot D^i$ .
- IS:

$$(\forall i \in \{0, 1, 2, \ldots n\})[a_i \leq C \cdot D^i] \Rightarrow \left(a_{n+1} \leq C \cdot D^{n+1}\right)$$

# Inductive Step

- IS:

$$(\forall i \in \{0, 1, 2, \dots n\})[a_i \leq C \cdot D^i] \Rightarrow \left(a_{n+1} \leq C \cdot D^{n+1}\right)$$

- From the definition of $a$, we have $a_{n+1} = 10a_n + 3a_{n-1}$. Therefore,

$$a_{n+1} = 10a_n + 3a_{n-1} \leq 10 \cdot C \cdot D^n + 3 \cdot C \cdot D^{n-1} \text{ (By IH)}$$

- Want $10 \cdot C \cdot D^n + 3 \cdot C \cdot D^{n-1} \leq C \cdot D^{n+1}$

# Inductive Step

- Want

$$10 \cdot \cancel{C} \cdot D^n + 3 \cdot \cancel{C} \cdot D^{n-1} \leq \cancel{C} \cdot D^{n+1} \Leftrightarrow$$
$$10 \cdot D^n + 3 \cdot D^{n-1} \leq D^{n+1}$$

- Dividing both sides by $D^{n-1}$ yields:

$$10D + 3 \leq D^2$$

# All Constraints

1. $2 \leq C$
2. $50 \leq C \cdot D$
3. $10D + 3 \leq D^2$

- We deal with constraint 3 first.
  - Smallest $D \in \mathbb{R}^{>0}$ that satisfies it:

# All Constraints

1. $2 \leq C$
2. $50 \leq C \cdot D$
3. $10D + 3 \leq D^2$

- We deal with constraint 3 first.
  - Smallest $D \in \mathbb{R}^{>0}$ that satisfies it: NO, WE ARE BUSY PEOPLE AND WE DON'T WANT TO SPEND TIME SOLVING $D^2 - 10D - 3 \geq 0$
  - Smallest $D \in \mathbb{N}$ that satisfies it: $D = \cdots ???$ (FIND ONE REAL QUICK, PLZ)

# All Constraints

1. $2 \leq C$
2. $50 \leq C \cdot D$
3. $10D + 3 \leq D^2$

- We deal with constraint 3 first.
  - Smallest $D \in \mathbb{R}^{>0}$ that satisfies it: NO, WE ARE BUSY PEOPLE AND WE DON'T WANT TO SPEND TIME SOLVING $D^2 - 10D - 3 \geq 0$
  - Smallest $D \in \mathbb{N}$ that satisfies it: $D = \cdots ???$ (FIND ONE REAL QUICK, PLZ)

$D = 11$ works!

# All Constraints

1. $2 \leq C$
2. $50 \leq C \cdot D$
3. $10D + 3 \leq D^2$

- Constraint (3) satisfied when $D \geq 11$ (just discussed)
- Since we want to find **tight** bounds for $a_n$, to minimize $C$, we select

$D = 11$ and from constraint (2) we have: $50 \leq C \cdot 11 \Leftrightarrow C \geq 4.55 \Rightarrow$
$C_{min} = 4.55$

# All Constraints

1. $2 \leq C$
2. $50 \leq C \cdot D$
3. $10D + 3 \leq D^2$

- Constraint (3) satisfied when $D \geq 11$ (just discussed)
- Since we want to find **tight** bounds for $a_n$, to minimize $C$, we select

$D = 11$ and from constraint (2) we have: $50 \leq C \cdot 11 \Leftrightarrow C \geq 4.55 \Rightarrow C_{min} = 4.55$

- Conclusion:

$$a_n \leq 4.55 \cdot 11^n$$

# Work on This

- A slight modification on the previous sequence:

$$a_n = \begin{cases} 10, & n = 0 \\ 50, & n = 1 \\ 10a_{n-1} + 3a_{n-2}, & n \geq 2 \end{cases}$$

- Assuming that we still suspect $a_n \leq C \cdot D^n$, **you** solve for the new $C, D$ **right now**!

# Work on This

- A slight modification on the previous sequence:

$$a_n = \begin{cases} 10, & n = 0 \\ 50, & n = 1 \\ 10a_{n-1} + 3a_{n-2}, & n \geq 2 \end{cases}$$

- Assuming that we still suspect $a_n \leq C \cdot D^n$, solve for the new $C, D$!
- Your solution ought to be $C = 10, D = 11$. What do you observe?

# Coin Problem

- In Celestia, there are only $7c$ and $10c$ coins.
- We want to find the *least monetary amount* payable **exclusively** with such coins!
- In quantifiers (all quantifications assumed over $\mathbb{N}$)

$$(\forall n \geq A)(\exists n_1, n_2)[n = 7n_1 + 10n_2]$$

- **Goal: Find constraints on A via constructive induction!**
- IB: ???

# Coin Problem

- In Celestia, there are only $7c$ and $10c$ coins.
- We want to find the *least monetary amount* payable **exclusively** with such coins!
- In quantifiers (all quantifications assumed over $\mathbb{N}$)

$$(\forall n \geq A)(\exists n_1, n_2)[n = 7n_1 + 10n_2]$$

- **Goal: Find constraints on A via constructive induction!**
- IB: Defer for later!!!

# Coin Problem

- In Celestia, there are only $7c$ and $10c$ coins.
- We want to find the *least monetary amount* payable **exclusively** with such coins!
- In quantifiers (all quantifications assumed over $\mathbb{N}$)

$$(\forall n \geq A)(\exists n_1, n_2)[n = 7n_1 + 10n_2]$$

- **Goal: Find constraints on A via constructive induction!**
- IB: Defer for later!!!
- IH: Assume that for $n \geq A$, $(\exists n_1, n_2)[n = 7 \cdot n_1 + 10n_2]$

# Coin Problem (IS)

- From the IH we have $(\exists n_1, n_2)[n = 7 \cdot n_1 + 10n_2]$
- How can we add/remove coins to get another cent?

# Coin Problem (IS)

- From the IH we have $(\exists n_1, n_2)[n = 7 \cdot n_1 + 10n_2]$
- How can we add/remove coins to get another cent?
  1. $n_2 \geq 2$ : Remove two $10c$ coins, add three $7c$ coins

$$n + 1 = 7n_1 + 10n_2 + 1 = 7n_1 + 10n_2 + (21 - 20)$$
$$= 7(n_1 + 3) + 10(n_2 - 2)$$

# Coin Problem (IS)

- From the IH we have $(\exists n_1, n_2)[n = 7 \cdot n_1 + 10n_2]$
- How can we add/remove coins to get another cent?

1. $n_2 \geq 2$ : Remove two $10c$ coins, add three $7c$ coins

$$n + 1 = 7n_1 + 10n_2 + 1 = 7n_1 + 10n_2 + (21 - 20)$$
$$= 7(n_1 + 3) + 10(n_2 - 2)$$

2. $n_1 \geq 7$: Remove seven $7c$ coins, add five $10c$ coins

$$n + 1 = 7n_1 + 10n_2 + 1 = 7n_1 + 10n_2 + (50 - 49)$$
$$= 7(n_1 - 7) + 10(n_2 + 5)$$

# Coin Problem (IS)

3. $(n_1 \leq 6) \land (n_2 \leq 1)$: Max value is $6 \times 7 + 1 \times 10 = 52$, so $n \leq 52$.

# RECAP

- We've shown that if $n \geq 53$, then

$$((\exists n_1, n_2)[n = 7 \cdot n_1 + 10n_2]) \Rightarrow ((\exists n_1, n_2\ )[n + 1 = 7 \cdot n_1 + 10n_2\ ])$$

- For which $n$ do we know that $((\exists a, b \in \mathbb{N})[n = 7a + 10b]$?

$\forall n \geq 52$

$\forall n \geq 53$

Something Else

# RECAP

- We've shown that if $n \geq 53$, then

$$((\exists n_1, n_2)[n = 7 \cdot n_1 + 10n_2]) \Rightarrow ((\exists n_1, n_2)[n + 1 = 7 \cdot n_1 + 10n_2])$$

- For which $n$ do we know that $((\exists a, b \in \mathbb{N})[n = 7a + 10b]$?

$\forall n \geq 52$

$\forall n \geq 53$

Something Else

*Only the implication holds! We don't have any **hard truth** (base) about whether it EVER holds.*

# Coin Problem (IS)

*3.* $(n_1 \leq 6) \wedge (n_2 \leq 1)$: Max value is $6 \times 7 + 1 \times 10 = 52$, so $n \leq 52$.

- Condition: $A \geq 53$.
- **Now** I need a base case.
- $(\exists ? n_1, n_2 \in \mathbb{N})[53 = 7 \cdot n_1 + 10 n_2]$

Yes (which?)

No

# Coin Problem (IS)

*3.* $(n_1 \leq 6) \wedge (n_2 \leq 1)$: Max value is $6 \times 7 + 1 \times 10 = 52$, so $k \leq 52$.

- Condition: $A \geq 53$.
- **Now** I need a base case.
- $(\exists ? n_1, n_2 \in \mathbb{N})[53 = 7 \cdot n_1 + 10n_2]$

Yes (which?)

No

*Prove it at home (use cases)*

# Coin Problem (IS)

*3.* $(n_1 \leq 6) \wedge (n_2 \leq 1)$: Max value is $6 \times 7 + 1 \times 10 = 52$, so $k \leq 52$.

- Condition: $A \geq 53$.
- **Now** I need a base case.
- $(\exists ? n_1, n_2 \in \mathbb{N})[53 = 7 \cdot n_1 + 10n_2]$
- $(\exists ? n_1, n_2 \in \mathbb{N})[54 = 7 \cdot n_1 + 10n_2]$

Yes (which?)

No

# Coin Problem (IS)

*3.* $(n_1 \leq 6) \wedge (n_2 \leq 1)$: Max value is $6 \times 7 + 1 \times 10 = 52$, so $k \leq 52$.

- Condition: $A \geq 53$.
- **Now** I need a base case.
- $(\exists ? n_1, n_2 \in \mathbb{N})[53 = 7 \cdot n_1 + 10n_2]$
- $(\exists ? n_1, n_2 \in \mathbb{N})[54 = 7 \cdot n_1 + 10n_2]$

Yes (which?)

No

$n_1 = 2,$
$n_2 = 4$

# RECAP

- We've shown that if $n \geq 53$, then

$$((\exists n_1, n_2)[n = 7 \cdot n_1 + 10n_2]) \Rightarrow ((\exists n_1, n_2)[n + 1 = 7 \cdot n_1 + 10n_2 ])$$

- We've also shown that $(\exists r_1, r_2 \in \mathbb{N})[54 = 7r_1 + 10r_2 ]$

$$(r_1 = 2, r_2 = 4)$$

# RECAP

- We've shown that if $n \geq 53$, then

$$((\exists n_1, n_2)[n = 7 \cdot n_1 + 10n_2]) \Rightarrow ((\exists n_1, n_2)[n + 1 = 7 \cdot n_1 + 10n_2])$$

- We've also shown that $(\exists r_1, r_2 \in \mathbb{N})[54 = 7r_1 + 10r_2]$
  $(r_1 = 2, r_2 = 4)$
- What do we know NOW about the theorem?

| True for $n \geq 52$ | True for $n \geq 53$ | True for $n \geq 54$ | Nothing |
|---|---|---|---|

# RECAP

- We've shown that if $n \geq 53$, then

$$((\exists n_1, n_2)[n = 7 \cdot n_1 + 10n_2]) \Rightarrow ((\exists n_1, n_2)[n + 1 = 7 \cdot n_1 + 10n_2 ])$$

- We've also shown that $(\exists r_1, r_2 \in \mathbb{N})[54 = 7r_1 + 10r_2 ]$
$(r_1 = 2, r_2 = 4)$
- What do we know **NOW** about the theorem?

| True for $n \geq 52$ | True for $n \geq 53$ | True for $n \geq 54$ | Nothing |
|---|---|---|---|

# What is $A$?

- Recall the theorem (all quantifiers over $\mathbb{N}$ ):

$$(\forall n \geq A)(\exists n_1, n_2)[n = 7n_1 + 10n_2]$$

- Our goal was to find $A$.
- $A = 54$ works, and is **optimal, since** $A = 53$ does not work.

# Question

- Is the theorem true for any $n \leq 53$?

Yes
(which?)

No
(Why?)

# Question

- Is the theorem true for any $n \leq 53$?

Yes (which?)

No (Why?)

$$0, 7, 10, 14, 17, 20, 21, 24, 27, 28, 30, 31, 34, 35, 37, 38, 40,$$
$$41, 42, 44, 45, 47, 48, 49, 50, 51, 52$$

- Note that there are **gaps** between these integers!

# And Here's Another

- Let $a$ be a sequence defined as follows:

$$a_n = \begin{cases} 0, & n = 0 \\ 2, & n = 1 \\ a_{\left\lfloor \frac{n}{2} \right\rfloor} + a_{\left\lfloor \frac{n}{4} \right\rfloor} + 5n, & n \geq 2 \end{cases}$$

- Then, find $C \in \mathbb{R}$ such that

$$(\forall n \in \mathbb{N})[a_n \leq C \cdot n]$$

# And Here's Another

- Let $a$ be a sequence defined as follows:

$$a_n = \begin{cases} 0, & n = 0 \\ 2, & n = 1 \\ a_{\lfloor \frac{n}{2} \rfloor} + a_{\lfloor \frac{n}{4} \rfloor} + 5n, & n \geq 2 \end{cases}$$

- Then, find $C \in \mathbb{R}$ such that

$$(\forall n \in \mathbb{N})[a_n \leq C \cdot n]$$

*Recursions like this have* ***linear*** *upper bounds*

# And Here's Another

- Let $a$ be a sequence defined as follows:

$$a_n = \begin{cases} 0, & n = 0 \\ 2, & n = 1 \\ a_{\lfloor \frac{n}{2} \rfloor} + a_{\lfloor \frac{n}{4} \rfloor} + 5n, & n \geq 2 \end{cases}$$

- Then, find $C \in \mathbb{R}$ such that

$$(\forall n \in \mathbb{N})[a_n \leq C \cdot n]$$

*Recursions like this have* ***linear*** *upper bounds*

- We proceed via **strong induction** on $n$.

# And Here's Another

- Let $a$ be a sequence defined as follows:

$$a_n = \begin{cases} 0, & n = 0 \\ 2, & n = 1 \\ a_{\lfloor \frac{n}{2} \rfloor} + a_{\lfloor \frac{n}{4} \rfloor} + 5n, & n \geq 2 \end{cases}$$

- Then, find $C \in \mathbb{R}$ such that

$$(\forall n \in \mathbb{N})[a_n \leq C \cdot n]$$

- We proceed via **strong induction** on $n$.
- In fact, to make some of the math easier, we will assume the hypothesis until $P(n-1)$ and prove the step for $P(n)$ instead of $P(n+1)$

# Finding C

- IB:
  - For $n = 0, a_0 \leq C \cdot 0 \Leftrightarrow 0 \leq 0$. No constraints on $C$ yet!
  - For $n = 1, a_1 \leq C \cdot n \Leftrightarrow 2 \leq C$. Done. We have our first lower bound for $C$.

# Finding C

- IB:
  - For $n = 0, a_0 \leq C \cdot 0 \Leftrightarrow 0 \leq 0$. No constraints on $C$ yet!
  - For $n = 1, a_1 \leq C \cdot n \Leftrightarrow 2 \leq C$. Done. We have our first lower bound for *C*.
- IH: Let $n \geq \mathbf{2}.$ Then, assume $(\forall i \in \{0,1,2,\dots,n-1\})[P(i)]$, where $P(i)$ means $a_i \leq C \cdot i$

# Finding C

- IB:
  - For $n = 0, a_0 \leq C \cdot 0 \Leftrightarrow 0 \leq 0$. No constraints on $C$ yet!
  - For $n = 1, a_1 \leq C \cdot n \Leftrightarrow 2 \leq C$. Done. We have our first lower bound for *C*.
- IH: Let $n \geq 2$. Then, assume $(\forall i \in \{0,1,2,\dots,n-1\})[P(i)]$, where $P(i)$ means $a_i \leq C \cdot i$
- IS: We attempt to prove $(P(0) \wedge P(1) \wedge P(2) \wedge \cdots \wedge P(n-1)) \Rightarrow P(n)$:

$$\bigwedge_{i=0}^{i=n-1}(a_i \leq C \cdot i) \Rightarrow a_n \leq C \cdot n$$

# Finding C

- IS: We attempt to prove $(P(1) \wedge P(2) \wedge \cdots \wedge P(n-1)) \Rightarrow P(n)$:

$$\bigwedge_{i=0}^{i=n-1} (a_i \leq C \cdot i) \Rightarrow a_n \leq C \cdot n$$

- From the IH, and taking into consideration that $0 \leq \left\lfloor \frac{n}{4} \right\rfloor, \left\lfloor \frac{n}{2} \right\rfloor \leq n$, we have (next slide):

# Finding C

- From the IH, and taking into consideration that $0 \leq \left\lfloor \frac{n}{4} \right\rfloor, \left\lfloor \frac{n}{2} \right\rfloor \leq n$, we have:

$$\begin{cases} a_{\lfloor n/4 \rfloor} \leq C \cdot \lfloor n/4 \rfloor \leq C \cdot \frac{n}{4} \\ a_{\lfloor n/2 \rfloor} \leq C \cdot \lfloor n/2 \rfloor \leq C \cdot \frac{n}{2} \end{cases}$$

- $a_n = a_{\lfloor n/2 \rfloor} + a_{\lfloor n/4 \rfloor} + 5n \leq C \cdot \frac{n}{2} + C \cdot \frac{n}{4} + 5n = \frac{n*(3C+20)}{4}$

# Finding C

- We have:

$$a_n \leq \frac{n*(3C+20)}{4}$$

- We want:

$$a_n \leq C \cdot n$$

- Hence, we want a C such that:

$$\frac{n*(3C+20)}{4} \leq C \cdot n$$

# Finding C

$$\frac{\cancel{n}(3C + 20)}{4} \leq C \cdot \cancel{n} \overset{n \geq 1}{\Leftrightarrow}$$

$$\frac{(3C + 20)}{4} \leq C \Leftrightarrow$$

$$3C + 20 \leq 4C \Leftrightarrow$$

$$C \geq 20$$

$$\Rightarrow C_{min} = 20$$

# Constraints

- From the IB: $C \geq 2$
- From the IS: $C \geq 20$
- Since we want to minimize $C$, we set $C = 20$.

# STOP RECORDING