The Emptier-Filler Game

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We describe several games between



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There will be a bin with numbers in it.

If the bin is ever empty then E wins.

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We describe several games between E: The Emptier F: The Filler.

There will be a bin with numbers in it.

- If the bin is ever empty then E wins.
- ▶ If game goes forever and bin is always nonempty then F wins.

1) F puts a **finite** multiset of \mathbb{N} into the bin. (e.g., bin has $\{1, 1, 1, 2, 3, 4, 9, 9, 18, 18\}$.

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 F puts a finite multiset of N into the bin. (e.g., bin has {1, 1, 1, 2, 3, 4, 9, 9, 18, 18}.
E takes out ONE number *n* (e.g., 18).

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- (e.g., bin has $\{1, 1, 1, 2, 3, 4, 9, 9, 18, 18\}$.
- 2) E takes out ONE number n (e.g., 18).
- 3) F puts in as many numbers as he wants that are < n (e.g., replace 18 with 99,999,999 17's and 5000 16's.)

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Which player has the winning strategy? What is that strategy. Breakout Rooms!



E wins!

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Strategy for E Keep removing the largest number in the box.



E wins! **Strategy for E** Keep removing the largest number in the box. **Why does this work?** Lets prove it by induction! But on what?

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Answer on next slide.

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Then we associate to the position the ordered pair (r, n).

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What happens if E removes a ball of rank r and F puts in LOTS of balls of lower rank?

- If $n \ge 1$ then the ordered pair is now (r, n-1).
- If n = 0 then their are no balls of rank r. Let the highest rank be r' < r. Assume there are n' balls of rank r'. Then the ordered pair is now (r', n').

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Consider the following funky ordering on ordered pairs.

$$(0,0) < (0,1) < (0,2) < \cdots < (1,0) < (1,1) < (1,2) < \cdots$$

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This is the ordering to use since this quantity always decreases.

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Proof By induction on the funky ordering.

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As noted in the last slide if E removes the top ranked ball and F puts in as many balls of lower rank, then the resulting position is associated to (n', r') < (n, r).

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Formal Proof

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End of Proof

But Can we do induction on this funky ordering?

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 \mathbb{N} It only used that if you start at some *n* and march downward you will **in a finite number of steps** get to 0. In fact, just *n* steps.

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Funky Ordering If you start at (n, r) and march downward will you get to (0, 0) in a finite number of steps? Discuss.

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Def An ordering is **well ordered** if when you start at any element *x* and march downward you will get to a MIN element in a finite number of steps.

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Upshot You can do induction on any well ordered ordering.

Does the Strategy Matter?

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Does the Strategy Matter?

What if E plays differently? One can show that no matter what E does, she wins!



What if E plays differently? One can show that no matter what E does, she wins!

How to prove that? By an induction on a an even funkier ordering. We won't be doing that.

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X is any of \mathbb{Z} , $\mathbb{Q}^{\geq 0}$, $\mathbb{N} + \mathbb{N}$, $\mathbb{N} + \mathbb{N} + \cdots$, $\mathbb{N} + \mathbb{Z}$, $\mathbb{N} + \mathbb{N}^*$. 1) F puts a **finite** multiset of X into the bin.

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- 1) F puts a **finite** multiset of X into the bin.
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For each of
$$X = \mathbb{Z}$$
, $X = \mathbb{Q}$, $X = \mathbb{N} + N$, $X = \mathbb{N} + \mathbb{N} + \cdots$, $X = \mathbb{N} + \mathbb{Z}$, $X = \mathbb{N} + \mathbb{N}^*$ who wins?

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Breakout Rooms!

 $\mathbf{X} = \mathbb{N}$

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$X = \mathbb{N}$ E wins–Always remove the largest element

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 $\begin{array}{l} \boldsymbol{X} = \mathbb{N} \text{ E wins-Always remove the largest element} \\ \boldsymbol{X} = \mathbb{Z} \text{ F wins-If E removes } n, \text{ F puts in } n-1. \\ \boldsymbol{X} = \mathbb{Q} \text{ F wins-If E removes } n, \text{ F puts in } \frac{n}{2}. \\ \boldsymbol{X} = \mathbb{N} + \mathbb{N} \end{array}$

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Key When you remove the 0 in second copy of \mathbb{N} you have to replace it with some element of the first \mathbb{N} . So eventually all elements are in first \mathbb{N} .

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How to make this rigorous? Ind on the number of copies of \mathbb{N} . $X = \mathbb{N} + \mathbb{Z}$ F wins. Bin initially has 0 in \mathbb{Z} , then always replace *n* by n - 1 in \mathbb{Z} $X = \mathbb{N} + \mathbb{N}^*$ F wins. Bin initially has 0 in \mathbb{N}^* , then always replace *n* by n - 1 in \mathbb{N}^* **Question** Let X be a set and \leq be an ordering on it. Let the (X, \leq) -game be the game as above where we put elements of X in the bin.

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Question Let X be a set and \leq be an ordering on it. Let the (X, \leq) -game be the game as above where we put elements of X in the bin.

In the following sentence fill in the ??? E can win the (X, \preceq) -game if and only if (X, \preceq) has property ???.

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Breakout Rooms!

Def (X, \preceq) is **well ordered** if there are NO infinite decreasing sequences.

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Def (X, \leq) is well ordered if there are NO infinite decreasing sequences.

E can win the (X, \preceq) -game if and only if (X, \preceq) is well ordered.