

250H-Discussion: Algorithm Design Techniques

2/9/26





Brute Force

- Attempting all possible solutions, seeing which one is the “best”
- Guaranteed to give you the best possible solution
- Not always efficient
 - Exponential runtimes
 - P VS NP
- We get better solutions using design techniques
 - Narrowing down solution space (LPs)
 - Memoization
 - Heuristics



Greedy Algorithms

- An algorithm is greedy if it builds up a solution in small steps, each a “locally” optimum solution satisfying some criteria
- Although greedy algorithms are not optimal for all problems, they can produce optimal solutions for some
 - Typically very intuitive and easy to prove correctness
 - We won’t be going over the proof techniques



Example 1) Making Change

- Suppose you are a cashier and you need to make change for someone who just paid you
- They are annoying and ask for the minimum number of coins in return
- What approach do you take?



Taking the biggest coin

- Naturally, taking the coin with the most value lowers the remaining amount the most



Taking the biggest coin

- Naturally, taking the coin with the most value lowers the remaining amount the most
- This is the algorithm! take the largest possible coin available and subtract its value from the total, repeating until total is 0



Taking the biggest coin

Suppose the change you need to make is 87 cents

- Take as many quarters as you can
 - 1 – 62 remaining
 - 2 – 37 remaining
 - 3 – 12 remaining
 - Can't take anymore
- Take as many dimes as you can
 - 1 – 2 remaining
 - Can't take anymore
- Take as many nickels
 - Can't take any!
- Take as many pennies
 - 2 pennies



Coin values

Vote: The algorithm we gave before

1. GIVES OPTIMAL SOLUTIONS FOR ALL COIN VALUES
2. DOESN'T GIVE OPTIMAL SOLUTIONS FOR ALL COIN VALUES



Coin values

- Not all coin values give you the optimal solution with this algorithm



Coin values

- Not all coin values give you the optimal solution with this algorithm
- {1, 3, 4,}



Coin values

- Not all coin values give you the optimal solution with this algorithm
- {1, 3, 4,}
- {1, 15, 25}



Coin values

- Not all coin values give you the optimal solution with this algorithm
- $\{1, 3, 4\} - 6$
 - Greedy = $\{4, 1, 1\}$
 - Optimal = $\{3, 3\}$
- $\{1, 15, 25\} - 35$
 - Greedy = $\{25, 1 \text{ (x10)}\}$
 - Optimal = $\{15, 15, 1, 1, 1, 1, 1\}$



Brute Force?

What is a brute force algorithm for this?



Example 2) Interval Scheduling

- We have n requests, labeled $1, \dots, n$, to use a room between times a - b . (i.e $[a, b]$).
- Each request has a starting time, $S(i)$, and a finish time, $F(i)$.
- We want to schedule as many of these as possible, without creating any overlap

Task 5

Task 3

Task 2

Task 4

Task 1



Answer:

We schedule Tasks 2, 3, and 4



Strategy/Heuristic

- How can we make a selection at each step based on the problem's properties
- What are some ideas?
 - Select the one with the earliest start time?



Strategy/Heuristic

- How can we make a selection at each step based on the problem's properties
- What are some ideas?
 - Select the one with the earliest start time?
 - Select the one shortest in length?



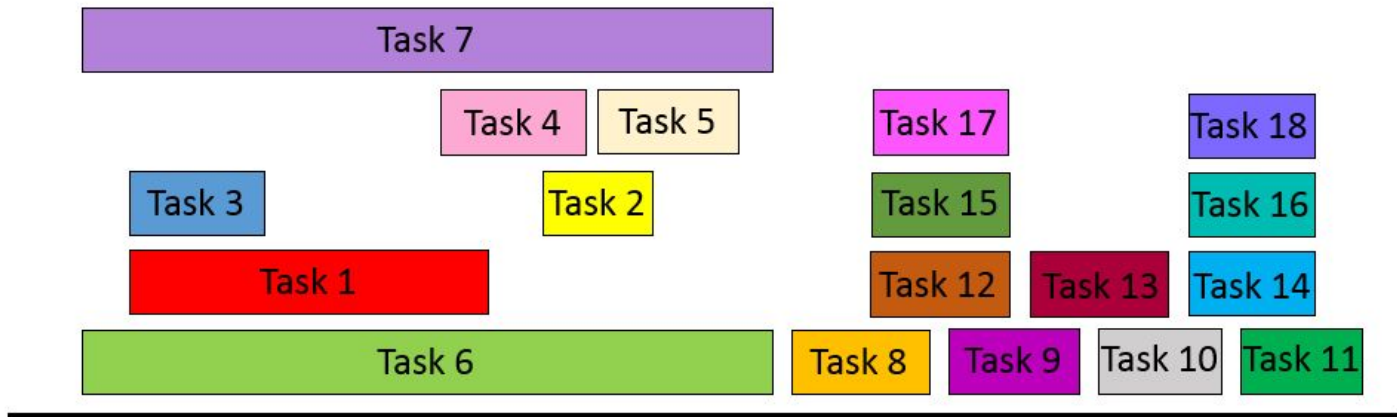
Strategy/Heuristic

- How can we make a selection at each step based on the problem's properties
- What are some ideas?
 - Select the one with the earliest start time?
 - Nope!
 - Select the one shortest in length?
 - Nope!
 - Select the one that finishes first?
 - Works!



Sort by Finish Times

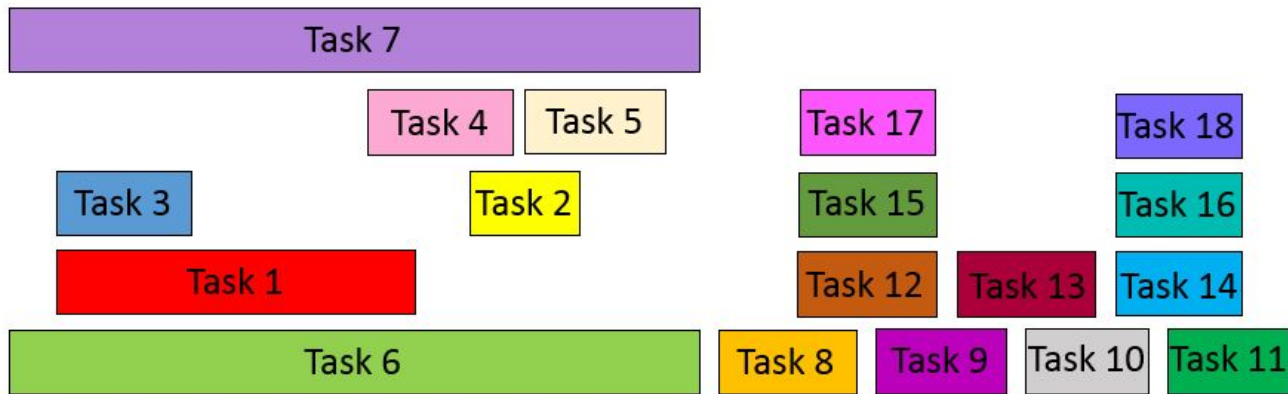
- Intuition:
 - This ensures that we have the most possible **REMAINING TIME** to schedule the rest of the problems





Sort by Finish Times

- Intuition:
 - This ensures that we have the most possible **REMAINING TIME** to schedule the rest of the problems
- Answer: Task 3, 4, 5, 8, 9, and 11





Expanding this Problem - Weighted I.S

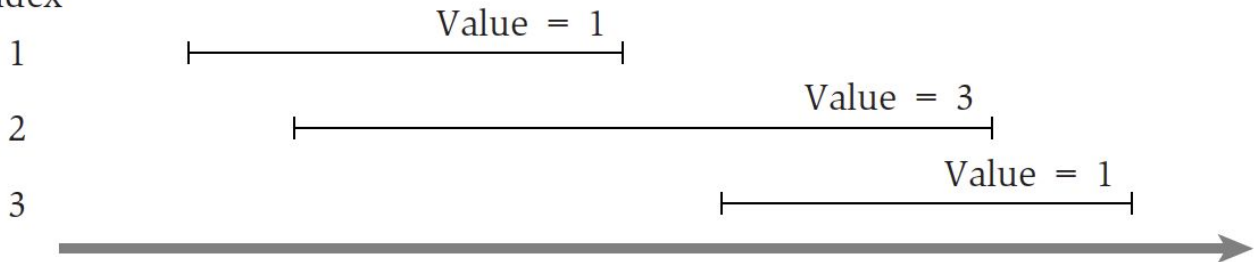
- What if we assign weights to each of these requests
 - This represents some requests being more important than others
 - Each request now has a weight, $w(i)$.
 - We want to maximize weight!
- Would our algorithm work now?



Expanding this Problem - Weighted I.S

- What if we assign weights to each of these requests
 - This represents some requests being more important than others
 - Each request now has a weight, $w(i)$.
 - We want to maximize weight!
- Would our algorithm work now?
 - No

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Proving Optimality of Greedy Algorithms

- Suppose we have access to an optimal solution, O .
 - Our greedy solution: A
-
1. Assume A is not optimal
 2. Compare it to O
 3. Derive a contradiction by the properties of our algorithm

Or

1. Compare A and O at each step
2. Show that at each step A does **AT LEAST** as well as O