Announcements

- Reading Chapter 12
- Project #4a is Due Wed at 6:00 PM

Filesystems

Raw Disks can be viewed as:

- a linear array of fixed sized units of allocation, called blocks
 - assume that blocks are error free (for now)
 - typical block size is 512 to 4096 bytes
- can update a block in place, but must write the entire block
- can access any block in any desired order
 - blocks must be read as a unit
 - for performance reasons may care about "near" vs. "far" blocks (but that is covered in a future lecture)

A Filesystem:

- provides a hierarchical namespace via directories
- permits files of variable size to be stored
- provides disk protection by restricting access to files based on permissions

Allocation Methods

- How do we select a free disk block to use?
- Contiguous allocation
 - allocate a contiguous chunk of space to a file
 - directory entry indicates the starting block and the length of the file
 - easy to implement, but
 - how to satisfy a given sized request from a list of free holes?
 - two options
 - first fit (find the first gap that fits)
 - best fit (find the smallest gaps that is large enough)
 - What happens if one wants to append to file?
 - from time to time, one will need to repack files

Linked Allocation

- Each file is a linked list of disk blocks, blocks can be located anywhere
 - Directory contains a pointer to the first and last block of a file
 - Each block contains a pointer to the next block
 - This is essentially a linked-list data structure

Problems:

- Best for sequential access data structures
 - requires sequential access whether you want to or not!
- Reliability one bad sector and all portions of your file downstream are lost

Useful fix:

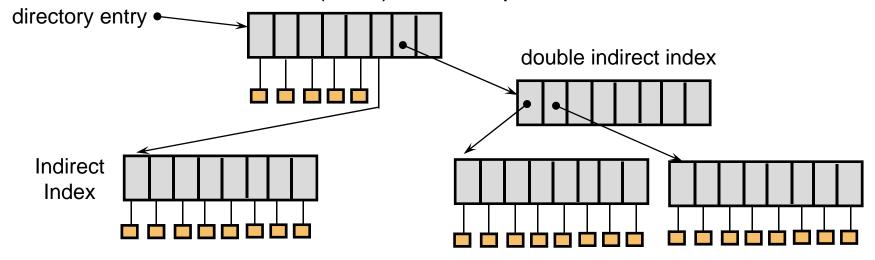
- Maintain a separate data structure just to keep track of linked lists
- Data-structure includes pointers to actual blocks

Indexed Allocation

- Bring all pointers together in an *index block*
 - Each file has its own index block ith entry of index block points to ith block making up the file
- How large to make an index block?
 - To avoid a fixed maximum file size, index block must be extensible
- Linked scheme:
 - maintain a linked list of indexed blocks
- Multilevel index:
 - Index block can point to other index blocks (which point to index blocks), which point to files
- Hybrid multi-level index
 - first n blocks are from a fixed index
 - next m blocks from an indirect index
 - next o blocks from a double indirect index

Hybrid Multi-level Index (UNIX)

- Observations
 - most files are small
 - most of the space on the disk is consumed by large files
- Want a flexible way to support different sized
 - assume 4096 byte block
 - first 12 blocks (48 KB) are from a fixed index
 - next 1024 blocks (4 MB) from an indirect index
 - next 1024² blocks (4 GB) from a double indirect index
 - final 1024³ blocks (4 TB) from a triple indirect index

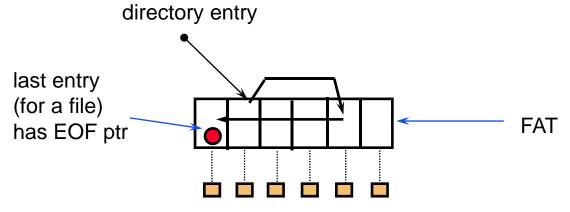


CMSC 412 - F11 (lect 17)

6

Modified Linked Allocation (FAT)

- Section of disk contains a table
 - called the file allocate table (FAT)
 - used in MS-DOS
- Directory entry contains the block number of the first block in the file
- Table entry contains the number of the next block in the file
- Last block has a end-of-file value as a table entry



ith block corresponds to the ith FAT entry

Performance Issues

FAT

- simple, easy to implement
- faster to traverse than linked allocation
- random access requires following links
- files can't have holes in them

Hybrid indirect

- fast access to any part of the file
- files can have holes in them
- more complex

Free Space Management

- How do we find a disk block to allocate?
- Bit Vectors
 - array of bits (one per block) that indicates if a block is free
 - compact so can keep in memory
 - 100 GB disk, 4K blocks -> 6MB per disk (0.003%)
 - easy to find long runs of free blocks
- Linked lists
 - each disk block contains the pointer to the next free block
 - pointer to first free block is keep in a special location on disk
- Run length encoding (called counting in book)
 - pointer to first free block is keep in a special location on disk
 - each free block also includes a count of the number of consecutive blocks that are free