

# Errata

Last updated April 12, 2022

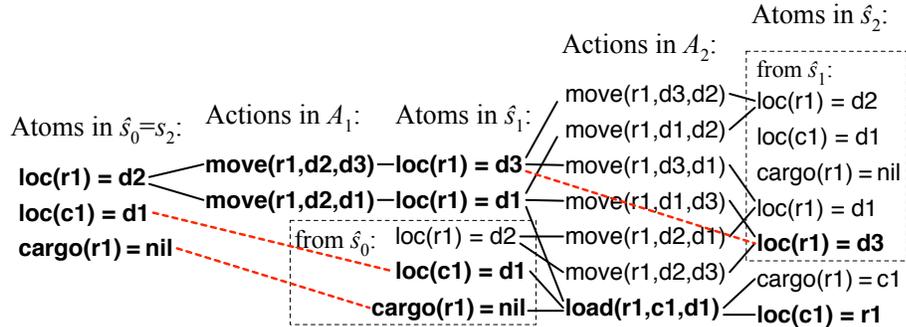
Errata for Ghallab, Nau, and Traverso, *Automated Planning and Acting*, Cambridge University Press, 2016.

This is a work in progress. Some of the corrections are tentative and may be revised, and additional corrections will probably be added. The list refers to section numbers instead of page numbers because the printed book and the authors' manuscript have different pagination.

– Dana Nau, University of Maryland

**Section 2.2.7.** In the first bullet, change *Children* to *Frontier*.

**Section 2.3.2.** Figure 2.8 should be as shown below. The dashed lines indicate situations where an assertion in an r-state  $\hat{s}_i$  is used to r-satisfy the goal  $\hat{g}_{i+1}$  of a later iteration.



**Section 2.3.3.** In Step 1 of RPG-landmark, replace the phrase “and the only landmark is  $\phi$  itself, so return  $\phi$ ” with this:

and there are no intermediate states, so return  $\emptyset$ .

**Section 2.3.3.** At the beginning of Step 5, replace “For every landmark  $\phi'$  found in the previous step” with this:

For every landmark  $\phi'$  found in the previous step that is not subsumed by other landmarks found in the previous step,

**Example 2.28.** The last assignment statement should be

$$\pi \leftarrow \text{move}(r1, d3, d1), \text{load}(r1, c1, d1), \text{move}(r1, d1, d3)\rangle.$$

**Section 2.6.1.** In the first bullet of the bulleted list, replace the phrase “Run-Lookahead is a simple version of the receding-horizon approach ...” with this:

If Run-Lookahead is used with a Lookahead procedure that searches to a fixed depth, then this is an example of the receding-horizon approach ...

**Exercise 3.19.** Part (a) should be

What sequence of commands will Refine-lookahead, Refine-lazy-lookahead, and Refine-concurrent-lookahead execute?

**Definition 4.4.** The first sentence of the definition should be

A ground instance of  $(\mathcal{T}', \mathcal{C}')$  of  $(\mathcal{T}, \mathcal{C})$  is *consistent* if  $\mathcal{T}'$  satisfies  $\mathcal{C}'$  and does not specify two different values for a state variable at the same time.

**Example 4.5.** The second paragraph should be

The assertions  $[t_1, t_2]\text{loc}(r1) = \text{loc1}$  and  $[t_2, t_3]\text{loc}(r1) : (\text{loc1}, \text{loc2})$  are nonconflicting: they have no inconsistent instances.

**Example 4.11.** For consistency with Examples 4.12 and 4.17,  $\text{put}(k', r, c, p')$  and  $\text{take}(k', r, c, p')$  should be  $\text{load}(k', r, c, p')$  and  $\text{unload}(k', r, c, p')$ , respectively.

**Example 4.12.** In `m-move1`, change  $\text{navigate}(w', w)$  to  $\text{navigate}(r, w', w)$ .

**Section 4.2.1.** In the paragraph before the bullets at the end of the section, “ $(\mathcal{T}, \mathcal{C})$  is consistent or secure if each of its timelines is” should be:

$(\mathcal{T}, \mathcal{C})$  is consistent if each of its timelines is consistent, and  
 $(\mathcal{T}, \mathcal{C})$  is secure if each of its timelines is secure and they have no variables in common.

**Exercise 4.8.** “Exercise 4.4” should be “Exercise 4.3”.

**Section 5.2.3.** The definition of a reachability graph should be this:

$$\text{Graph}(s, \pi) = (\widehat{\gamma}(s, \pi), \{(s', s'') \mid s' \in \widehat{\gamma}(s, \pi) \text{ and } s'' \in \gamma(s', \pi(s'))\})$$

or perhaps more clearly, this:

$$\begin{aligned} \text{Graph}(s, \pi) &= (V, E), \text{ where} \\ V &= \widehat{\gamma}(s, \pi), \\ E &= \{(s', s'') \mid s' \in \widehat{\gamma}(s, \pi) \text{ and } s'' \in \gamma(s', \pi(s'))\} \end{aligned}$$

**Section 5.2.3.** The last line before Example 5.5 should be

We let  $\widehat{\Gamma}(s)$  be the set of all states that are *reachable* from  $s$ , i.e.,  
 $\Gamma(s) = \bigcup_{\pi} \widehat{\gamma}(s, \pi)$ .

**Definition 5.10.** Put another right paren after  $\widehat{\gamma}(s_0, \pi)$ .

**Exercise 5.3.** Change “a policy  $\pi$ ” to “a solution policy  $\pi$ ”.

**Exercise 5.7(b).** Remove “by drawing the And/Or search tree.”

**Section 6.2.1.** In Definition 6.3, a *solution* should be defined as a policy  $\pi$  for  $\Sigma$  such that  $\widehat{\gamma}(s_0, \pi) \cap S_g \neq \emptyset$ . This differs from the definition used in Chapter 5, which required that  $leaves(s_0, \pi) \cap S_g \neq \emptyset$ .

**Section 6.2.3.** The paragraph after Equation 6.3 should be

A closed policy  $\pi'$  *dominates* a close policy  $\pi$  if and only if  $V^{\pi'}(s) \leq V^{\pi}(s)$  at every state  $s$  where both  $\pi$  and  $\pi'$  are defined. A closed policy  $\pi^*$  is *optimal* if it dominates all other closed policies. At every state  $s$  where  $\pi^*$  is defined, it has a minimal expected cost:  $V^*(s) = \min_{\pi} V^{\pi}(s)$ . Under our assumption of probabilistic planning in a domain without dead ends,  $\pi^*$  is guaranteed to exist.

**Section 6.3.2.** In Algorithm 6.8, AO\*, insert the following lines after the first line:

```
global V, π, Envelope
π ← ∅
V(s0) ← V0(s0)
```

**Section 6.4.2.** In the fourth paragrah, RRF should be RFF.

**Section 6.4.2.** Algorithm 6.16, RFF, should be as follows:

```
RFF(Σ, s0, Sg, θ)
  π ← Det-Plan(Σd, s0, Sg)
  if π = failure then return failure
  while ∃s ∈  $\widehat{\gamma}(s_0, \pi) \setminus (\text{Dom}(\pi) \cup S_g)$  such that  $\text{Pr}(s|s_0, \pi) \geq \theta$ , do
    π' ← Det-Plan(Σd, s, Sg ∪ Targets(π, s))
    if π = failure then return failure
    π ← π ∪ {(s, a) ∈ π' | s ∉ Dom(π)}
```

**Exercises 6.12 and 6.14.** FF-Replan should be FS-Replan.