

AMSC 607 / CMSC 764 Homework 6, Fall 2010
Due October 26, before class begins.

8. (20) Write a Matlab program that uses a feasible direction method to solve the linear programming problem

$$\begin{aligned} \min_{\mathbf{x}} \quad & \mathbf{c}^T \mathbf{x} \\ \mathbf{A} \mathbf{x} = & \mathbf{b}, \\ \mathbf{x} \geq & \mathbf{0}, \end{aligned}$$

where $\mathbf{x}, \mathbf{c} \in \mathcal{R}^n$ and $\mathbf{b} \in \mathcal{R}^m$ with $m < n$. Assume a constraint qualification. Also assume that the initial point is a vertex (i.e., exactly n active constraints), and step from vertex to vertex, as in the simplex method.

Write a Matlab function `xopt = lpfeasdir(A,b,c,x)`. The parameters to your feasible direction algorithm are \mathbf{A} , \mathbf{b} , \mathbf{c} , and an initial feasible point \mathbf{x} .

- Use `qrupdate`, `qrinsert`, and/or `qrdelete` (instead of the \mathbf{B} and \mathbf{N} method in the notes) to update a factorization of the matrix \mathbf{A}_W corresponding to the currently active constraints.
- At each iteration, one row of \mathbf{A}_W is replaced by another.
- The next point is $\mathbf{x} + \alpha \mathbf{p}$, where \mathbf{p} is determined from solving the system involving a column of the identity matrix, and α defines the longest step that is possible without violating any of the constraints. The constraint that we hit becomes the added one.
- Stop when there is no feasible downhill direction.

Test your program on this linear programming problem (Griva, Nash, and Sofer, p.221):

```
A = [2 4 2 1 0 0 0 0
      3 5 4 0 1 0 0 0
      1 0 0 0 0 1 0 0
      0 1 0 0 0 0 1 0
      0 0 1 0 0 0 0 1];
b = [50; 80; 20; 20; 20];
c = [-5; 10; 15; 3; 0; 2; 0; 0];
% x0 is the initial point.
x0 = [zeros(3,1);b];
```

You can check your answer using Matlab's simplex algorithm for linear programming:

```
options = optimset('largescale','off');
[x,fval,exitflag,output,lamba] = ...
    linprog(c,[],[],A,b,zeros(8,1),inf*ones(8,1),x0,options)
```

Grading: 20 points for the efficient implementation of the algorithm as a bug-free MATLAB function, with good documentation for the calling sequence and the algorithm. “Efficient” means not using an order of magnitude more computation than necessary.

Notes

- Let A and B be matrices, and let c be a vector. Make sure you understand why the statements $A*(B*c)$ and $A \setminus (B*c)$ take much less time than $A*B*c$ and $A \setminus B * c$, and then use this knowledge in your programming.
 - Make sure that each iteration of your algorithm uses only $O(mn + n^2)$ multiplications. This is possible if you compute a QR factorization once, at the beginning of the algorithm, and then, for each iteration, use the updating functions rather than refactoring, or computing an inverse, or solving a linear system involving a general matrix.
 - The grader will test your program on a larger problem. See the sample problem generator on the homepage.
 - Your program does not need to handle error conditions such as violation of the constraint qualification, infeasibility, etc. This is just an exercise to understand the algorithm. Use a trusted routine such as `linprog.m` if you ever really need to solve such a problem.
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