

1. (10) Suppose you want to solve a linear system of equations $Ax = b$, $n = 100,000$, and have your choice of 2 preconditioners for GMRES:

- M_1 creates a $\hat{G}_1$ with 5 small clusters of eigenvalues. 2 seconds and 10^6 storage locations are required to form $M_1^{-1}z$ for an arbitrary vector z .
- M_2 creates a $\hat{G}_2$ with 10 small clusters of eigenvalues. 1 second and 5×10^5 storage locations is required to form $M_2^{-1}z$ for an arbitrary vector z .

It takes .25 seconds and 10^6 storage locations to form Az . Which preconditioner would you advise using in order to compute an approximate solution to the problem? Why?

Answer: First, note that $\hat{G} = I - M^{-1}N = M^{-1}(M - N) = M^{-1}A$, so each multiplication by $\hat{G}$ requires one multiplication by A and one solve of a linear system involving M .

Consider M_1 . There are 5 clusters of eigenvalues, so we expect a good answer in 5 iterations. Each iteration costs 2.25 sec (since the GMRES overhead will be microseconds at most when $n = 10^5$), so the time will be about 11.25 sec. The storage required will be about 10^6 for multiplication by A , 10^6 for multiplication by M_1^{-1} , and 5×10^5 for the P matrix.

Consider M_2 . There are 10 clusters of eigenvalues, so we expect a good answer in 10 iterations. Each iteration costs 1.25 sec, so the time will be about 12.5 sec. The storage required will be about 10^6 for multiplication by A , 5×10^5 for multiplication by M_1^{-1} , and 10×10^5 for the P matrix.

So M_1 should be faster, and M_2 should require less storage. M_1 is overall probably the better choice.

2. (10) The following is the Lanczos algorithm, closely related to CG. As usual, A is a matrix with nz nonzero elements and b is a vector of length n . How much storage does the algorithm use? How many multiplications and divisions does it perform? (Express your answers in terms of the parameters n , k , and nz .)

	r=b;
n mult	beta = norm(b);
	vsav=[];
	beta0=beta;
	v=zeros(n,1);
	for i=1:k,
	vold=v;
n div.	v=r/beta;
nz mult	av=A*v;
n mult	alpha=v'*av;
$2n$ mult	r=av-alpha*v-beta*vold;
n mult	beta=norm(r);
	vsav=[vsav v];
	alphasav(i)=alpha;
	betasav(i)=beta;
	end

Answer: $n + (5n + nz)k$ multiplications and divisions.

Storage:

- A takes $3nz$ locations (Matlab sparse matrix).
- b , r , v , $vold$, av each take n locations.
- $vsav$ takes nk locations.
- $alphasav$, $betasav$ each take k locations.

Total: $3nz + (k + 5)n + 2k + O(1)$ locations.