

1. (6) The following is the CG algorithm for solving the linear system $\mathbf{Ax} = \mathbf{b}$. As usual, \mathbf{A} is a symmetric positive definite matrix with nz nonzero elements and \mathbf{b} is a vector of length n . How much storage does the algorithm use? How many multiplications and divisions does it perform? (Express your answers in terms of the parameters n , K , and nz .)

Let $\mathbf{r} = \mathbf{b}$, $\mathbf{x} = \mathbf{0}$, and let $\gamma = \mathbf{r}^T \mathbf{r}$, $\mathbf{p} = \mathbf{r}$.
for $k = 1, \dots, K$

$$\mathbf{q} = \mathbf{A}\mathbf{p}$$

$$\alpha = \gamma / (\mathbf{p}^T \mathbf{q})$$

$$\mathbf{x} = \mathbf{x} + \alpha \mathbf{p}$$

$$\mathbf{r} = \mathbf{r} - \alpha \mathbf{q}$$

$$\hat{\gamma} = \mathbf{r}^T \mathbf{r}$$

$$\beta = \hat{\gamma} / \gamma, \gamma = \hat{\gamma}$$

$$\mathbf{p} = \mathbf{r} + \beta \mathbf{p}$$

end (for k)

Answer: We need storage for matrix \mathbf{A} and vectors \mathbf{x} , \mathbf{r} , \mathbf{b} , \mathbf{p} , \mathbf{q} , for a total of $nz + 5n$ floating-point numbers plus a few scalars, and (in MATLAB) $2nz$ integers.

In each iteration, we need one matrix-vector product, 2 inner products (for α and $\hat{\gamma}$), and 3 vector updates (\mathbf{x} , \mathbf{r} , \mathbf{p}) for a total of $K(nz + 5n) + n$ multiplications plus $2K$ divisions.

2a. (6) Suppose we use CG to solve the linear system $\widehat{\mathbf{G}}\mathbf{x}^* = \mathbf{c}$, and we know that all of the eigenvalues of $\widehat{\mathbf{G}}$ are in the interval $[1, 10^4]$. Give an upper bound on the number of iterations it will take to reduce the error

$$\|\mathbf{x}^{(k)} - \mathbf{x}^*\|_{\widehat{\mathbf{G}}}$$

by a factor of 10^5 .

Answer: Let $e_k = \|\mathbf{x}^{(k)} - \mathbf{x}^*\|_{\widehat{\mathbf{G}}}$. Let κ be the ratio of the largest eigenvalue of $\widehat{\mathbf{G}}$ to the smallest. The convergence theorem for CG says

$$e_k \leq 2 \left(\frac{\sqrt{\kappa} - 1}{\sqrt{\kappa} + 1} \right)^k e_0.$$

We know that $\kappa \leq 10^4$. Therefore,

$$e_k \leq 2 \left(\frac{100 - 1}{100 + 1} \right)^k e_0.$$

We need

$$2 \left(\frac{100 - 1}{100 + 1} \right)^k \leq 10^5,$$

or

$$\left(\frac{100 - 1}{100 + 1} \right)^k \leq \frac{10^5}{2}.$$

Taking logs of both sides (base e or base 10 would be fine for a calculator) gives

$$k(\log 99 - \log 101) \leq \log(10^5/2).$$

Since k must be an integer, we set

$$k \geq -\lceil \frac{\log(10^5/2)}{(\log 99 - \log 101)} \rceil,$$

and, with a calculator we could find that we need at most $k = 541$ iterations. (This only makes sense if $n \geq 541$; otherwise, without rounding error, the iteration will stop after n iterations with the exact solution.)

2b. (8) Apply one step of SOR ($\omega = 1.5$) to the linear system

$$\begin{bmatrix} 2 & -1 \\ -1 & 2 \end{bmatrix} \begin{bmatrix} x_1^* \\ x_2^* \end{bmatrix} = \begin{bmatrix} 4 \\ 1 \end{bmatrix}$$

with a starting guess of $x_1^{(0)} = x_2^{(0)} = 1$. What is $\mathbf{x}^{(1)}$? Will the iteration converge to the true solution? Justify your answer.

Answer: For Gauss-Seidel,

$$\begin{aligned} x_1^{(1)}_{GS} &= (4 + 1)/2 = 2.5 \\ x_2^{(1)}_{GS} &= (1 + 2.5)/2 = 1.75 \end{aligned}$$

Therefore, the SOR iterate is

$$\mathbf{x}^{(1)} = (1 - \omega)\mathbf{x}^{(0)} + \omega\mathbf{x}_{GS}^{(1)} = -0.5 \begin{bmatrix} 1 \\ 1 \end{bmatrix} + 1.5 \begin{bmatrix} 2.50 \\ 1.75 \end{bmatrix} = \begin{bmatrix} 3.250 \\ 2.125 \end{bmatrix}.$$

The eigenvalues of \mathbf{A} are roots of

$$\det(\mathbf{A} - \lambda\mathbf{I}) = 0.$$

Therefore,

$$(2 - \lambda)^2 - 1 = 0$$

or

$$2 - \lambda = \pm 1,$$

so

$$\lambda = 1, 3.$$

Therefore matrix \mathbf{A} is symmetric and positive definite. (Gershgorin's theorem would also tell us this.)

The "Convergence of SIMs" handout (<http://www.cs.umd.edu/users/oleary/c661/simnotes.pdf>) says that SOR is convergent for this ω when \mathbf{A} is symmetric and positive definite.