

# A Sane Proof that $COL_k \leq COL_3$

By William Gasarch

## 1 Introduction

Let  $A \leq B$  mean that there is a polynomial-time computable function  $f$  such that  $x \in A$  iff  $f(x) \in B$ .

**Def 1.1** Let  $k \geq 2$ .  $COL_k$  is the set of all graphs that are  $k$ -colorable

The following are well known.

- For all  $k \geq 2$ ,  $COL_k \leq SAT$  (this is by the Cook-Levin Theorem).
- For all  $k \geq 2$ , For all  $k \geq 3$ ,  $SAT \leq COL_k$ , hence  $COL_k$  is  $NP$ -complete.
- If  $a < b$  then  $COL_a \leq COL_b$  by an easy reduction (Take  $G$  and add  $K_{b-a}$  and an edge from every elements of  $K_{b-a}$  to the original graph.)

The proof that  $COL_3 \leq COL_4$  is very easy: just add a vertex to  $G$  and connect it to all the elements of  $G$ . Is  $COL_4 \leq COL_3$ ? Yes via

$$COL_4 \leq SAT \leq COL_3.$$

This is true but unsatisfying. One of my students said

*It's counterintuitive and makes me sad.*

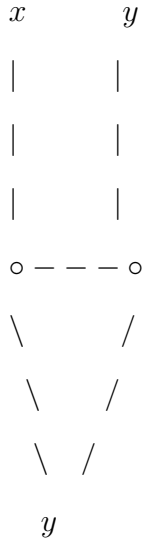
So we asked informally: Is there a SANE reduction  $COL_4 \leq COL_3$ . There is and we present it here. In fact we show  $COL_k \leq COL_3$ .

A sane proof is already known. Let  $HCOL_k$  is the set of all hypergraphs that are  $k$ -colorable Lovasz [?] showed

$$COL_k \leq HCOL_2 \leq COL_3.$$

Our proof does not use  $HCOL_2$  or any  $HCOL_k$ .

**Def 1.2**  $GAD(x, y, z)$  is the following graph. (The vertices that don't have labels are never referred to so we don't need to label them.)



We leave the proof of the following easy lemma to the reader.

**Lemma 1.3** *If  $GAD(x, y, z)$  is three colored and  $x, y$  get the same color, then  $z$  also gets that color.*

**Def 1.4**  $GAD(x_1, \dots, x_k, z)$  consists of  $GAD(x_1, x_2, y_1)$ ,  $GAD(y_1, x_3, y_2)$ ,  $GAD(y_2, x_4, y_3)$ ,  $\dots$ ,  $GAD(y_{k-3}, x_{k-1}, y_{k-2})$ , and  $GAD(y_{k-2}, x_k, z)$ . Note that, (1) not including  $x_1, \dots, x_k, z$ ,  $GAD(y_{k-2}, x_k, z)$  has  $3(k-2) + 1 = 3k - 5 \leq 3k$  vertices, and (2)  $5(k-1) = 5k - 5 \leq 5k$  edges.

We leave the proof of the following easy lemma to the reader.

**Lemma 1.5** *Let  $k \geq 2$ . If  $GAD(x_1, x_2, \dots, x_k, z)$  is three colored and  $x_1, \dots, x_k$  get the same color, then  $z$  also gets that color.*

**Theorem 1.6** *Let  $k \geq 2$ .  $COL_k \leq COL_3$  by a simple reduction. Let  $f$  be the reduction. If  $G$  has  $n$  vertices and  $e$  edges then  $f(G) = G'$  has  $\leq 2k^2n + 2ke$  vertices and  $\leq 3k^2n + 2ke$  edges.*

**Proof:** Let  $G$  have vertices  $v_1, \dots, v_n$  and edge set  $E$ . We construct  $G'$ :

1. Vertices  $T, F, R$  which will form a triangle. In any coloring they have different colors which we call  $T, F, R$ . This is 3 vertices and 3 edges. (We won't count these in the end since our crude upper bounds on the vertices and edges in  $G'$  will clearly be over by at least 3.)
2. For  $1 \leq i \leq n$  and  $1 \leq j \leq k$  vertex  $v_{ij}$ . All of these will be connected by an edge to vertex  $R$ . This will be  $kn$  vertices and  $kn$  edges. Here is our intent and how we achieve it:
  - (a) For all  $1 \leq i \leq n$  our intent is:  $v_{ij}$  is colored  $T$  means that vertex  $v_i$  in  $G$  is colored  $j$ ;  $v_{ij}$  is colored  $F$  means that vertex  $v_i$  in  $G$  is NOT colored  $j$ .
  - (b) For all  $1 \leq i \leq n$  we need that *at least one* of  $v_{i1}, \dots, v_{in}$  is colored  $T$ . Hence we need it to NOT be the case that  $v_{i1}, v_{i2}, \dots, v_{in}$  are all colored  $F$ . We place the gadget  $G(v_{i1}, \dots, v_{in}, T)$  in the graph. If  $v_{i1}, \dots, v_{in}$  are all colored  $F$  then this gadget will not be 3-colorable. This is  $\leq 3kn$  vertices and  $\leq 5kn$  edges.
  - (c) For all  $1 \leq i \leq n$  we need that *at most one* of  $v_{i1}, \dots, v_{ik}$  is colored  $T$ . Hence we need that for each pair at most one is colored  $T$ . For each  $1 \leq j_1 < j_2 \leq k$  we place the gadget  $GAD(v_{ij_1}, v_{ij_2}, F)$ . This is  $n \binom{k}{2} \times 2 \leq k^2n$  vertices and  $n \binom{k}{2} \times 5 \leq 2.5k^2n$  edges.
3. For each edge  $(v_i, v_j)$  in the original graph we want to make sure that  $v_i$  and  $v_j$  are not the same color. Place the gadgets  $GAD(v_{i1}, v_{j1}, F), GAD(v_{i2}, v_{j2}, F), \dots, GAD(v_{ik}, v_{jk}, F)$ . This is  $2ke$  vertices and  $5ke$  edges.

Note that the number of vertices in  $G'$  is  $\leq kn + 3kn + k^2n + 2ke \leq 2k^2n + 2ke$  vertices and  $\leq kn + 5kn + 2.5k^2n + 2ke \leq 3k^2n + 2ke$  edges.

Clearly  $G$  is  $k$ -colorable iff  $G'$  is 3-colorable.

■