# **Rev For Mid1: Proofs**

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# Review of Mods and GCD

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1.  $a \equiv b \pmod{m}$  means *m* divides b - a.

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- 2. The remainder when you divide a by m or b by m is the same.

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We usually think of  $a \equiv b \pmod{m}$  to mean that *a* is large and  $0 \le b \le m - 1$  (so small).

### **Do Examples of Mods**

I ask random people in the class what a is congruent to mod m.

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# **Theorem** Assume $a_1 \equiv b_1 \pmod{m}$ and $a_2 \equiv b_2 \pmod{m}$ .

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Then:

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Then:

1.  $a_1 + a_2 \equiv b_1 + b_2 \pmod{m}$ .

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2.  $a_1 a_2 \equiv b_1 b_2 \pmod{m}$ .

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**Theorem**  $(\forall a \in \mathbb{N})[a^7 \equiv a \pmod{7}].$ **Note** The theorem is about ALL  $a \in \mathbb{N}$ . Do we have to consider all  $a \in \mathbb{N}$ .

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Can compute  $a^n \pmod{m}$  in  $\leq 2 \log n$  steps.

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3^{100} \pmod{13}.
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Can compute  $a^n \pmod{m}$  in  $\leq 2 \log n$  steps. 3<sup>100</sup> (mod 13). DO WITH YOUR NEIGHBOR.

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Can compute  $a^n \pmod{m}$  in  $\le 2 \log n$  steps.  $3^{100} \pmod{13}$ . DO WITH YOUR NEIGHBOR. **Step One**  $100 = 2^6 + 2^5 + 2^2$ . So  $3^{100} = 3^{2^6} \times 3^{2^5} \times 3^{2^2}$ .

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#### **Powering Fast**

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#### Greatest Common Divisor (GCD)

**Definition** The **Greatest Common Divisor** of x, y is the largest number that divides both x and y. We denote this GCD(x, y).



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Do Examples with the class.

Assume x < y. Then



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$$\operatorname{GCD}(x, y) = \operatorname{GCD}(x, y - x)$$

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Assume x < y. Then

$$\operatorname{GCD}(x, y) = \operatorname{GCD}(x, y - x)$$

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**Better** Remove the largest multiple of x that is  $\leq y$ .

Assume x < y. Then

$$\operatorname{GCD}(x, y) = \operatorname{GCD}(x, y - x)$$

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**Better** Remove the largest multiple of x that is  $\leq y$ . Have class do an example.

# **Proof by Example**

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**Theorem** 
$$(\exists x)[\neg(\exists a, b, c)[x = a^2 + b^2 + c^2]]$$

**Theorem**  $(\exists x)[\neg(\exists a, b, c)[x = a^2 + b^2 + c^2]]$ To prove a  $\exists x$  give x and prove the thm for x.

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**Theorem**  $(\exists x)[\neg(\exists a, b, c)[x = a^2 + b^2 + c^2]]$ To prove a  $\exists x$  give x and prove the thm for x. x = 7. We show 7 is not the sum of 3 squares. Cases.

**Theorem**  $(\exists x)[\neg(\exists a, b, c)[x = a^2 + b^2 + c^2]]$ To prove a  $\exists x$  give x and prove the thm for x. x = 7. We show 7 is not the sum of 3 squares. Cases. **Case 1** At least one of a, b, c is  $\geq 3$ . Then  $a^2 + b^2 + c^2 \geq 9 > 7$ .

**Theorem**  $(\exists x)[\neg(\exists a, b, c)[x = a^2 + b^2 + c^2]]$ To prove a  $\exists x$  give x and prove the thm for x. x = 7. We show 7 is not the sum of 3 squares. Cases. **Case 1** At least one of a, b, c is  $\geq 3$ . Then  $a^2 + b^2 + c^2 \geq 9 > 7$ . **Case 2** At least two of a, b, c are > 2. Then  $a^2 + b^2 + c^2 > 8 > 7$ .

**Theorem**  $(\exists x)[\neg(\exists a, b, c)[x = a^2 + b^2 + c^2]]$ To prove a  $\exists x$  give x and prove the thm for x. x = 7. We show 7 is not the sum of 3 squares. Cases. **Case 1** At least one of a, b, c is  $\geq 3$ . Then  $a^2 + b^2 + c^2 \geq 9 > 7$ . **Case 2** At least two of a, b, c are  $\geq 2$ . Then  $a^2 + b^2 + c^2 \geq 8 > 7$ . **Case 3** The only case left: at most 1 of a, b, c is 2. Then  $a^2 + b^2 + c^2 \leq 4 + 1 + 1 = 6 < 7$ .

**Theorem**  $(\exists x)[\neg(\exists a, b, c)[x = a^2 + b^2 + c^2]]$ To prove a  $\exists x$  give x and prove the thm for x. x = 7. We show 7 is not the sum of 3 squares. Cases. **Case 1** At least one of a, b, c is  $\geq 3$ . Then  $a^2 + b^2 + c^2 \geq 9 > 7$ . **Case 2** At least two of a, b, c are  $\geq 2$ . Then  $a^2 + b^2 + c^2 \geq 9 > 7$ . **Case 3** The only case left: at most 1 of a, b, c is 2. Then  $a^2 + b^2 + c^2 \leq 4 + 1 + 1 = 6 < 7$ .

**Upshot** For  $\exists x$  Theorems SHOW THE x. (Nonconstructive proofs are possible though rare for this course.)

# Irrationals

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Want  $7^{1/3} \notin \mathbb{Q}$ . We need the Lemma. All  $\equiv$  is mod 7.

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Want  $7^{1/3} \notin \mathbb{Q}$ . We need the Lemma. All  $\equiv$  is mod 7. Lemma  $(\forall n)[n^3 \equiv 0 \pmod{7}) \rightarrow n \equiv 0 \pmod{7}]$ .

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Want  $7^{1/3} \notin \mathbb{Q}$ . We need the Lemma. All  $\equiv$  is mod 7. Lemma  $(\forall n)[n^3 \equiv 0 \pmod{7} \rightarrow n \equiv 0 \pmod{7}]$ . Take Contrapositive:

Want  $7^{1/3} \notin \mathbb{Q}$ . We need the Lemma. All  $\equiv$  is mod 7. **Lemma**  $(\forall n)[n^3 \equiv 0 \pmod{7}) \rightarrow n \equiv 0 \pmod{7}]$ . Take Contrapositive:  $(\forall n)[n \neq 0 \pmod{7}) \rightarrow n^3 \neq 0 \pmod{7}]$ . 7 cases

Want  $7^{1/3} \notin \mathbb{Q}$ . We need the Lemma. All  $\equiv$  is mod 7. Lemma  $(\forall n)[n^3 \equiv 0 \pmod{7}) \rightarrow n \equiv 0 \pmod{7}]$ . Take Contrapositive:  $(\forall n)[n \not\equiv 0 \pmod{7}) \rightarrow n^3 \not\equiv 0 \pmod{7}]$ . 7 cases  $n \equiv 1 \rightarrow n^3 \equiv 1 \neq 0$ .

Want 
$$7^{1/3} \notin \mathbb{Q}$$
. We need the Lemma. All  $\equiv$  is mod 7.  
**Lemma**  $(\forall n)[n^3 \equiv 0 \pmod{7} \rightarrow n \equiv 0 \pmod{7}]$ .  
Take Contrapositive:  
 $(\forall n)[n \not\equiv 0 \pmod{7} \rightarrow n^3 \not\equiv 0 \pmod{7}]$ . 7 cases  
 $n \equiv 1 \rightarrow n^3 \equiv 1 \neq 0$ .  
 $n \equiv 2 \rightarrow n^3 \equiv 8 \equiv 1 \neq 0$ .

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Take Contrapositive:  
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 $n \equiv 1 \rightarrow n^3 \equiv 1 \neq 0$ .  
 $n \equiv 2 \rightarrow n^3 \equiv 8 \equiv 1 \neq 0$ .  
 $n \equiv 3 \rightarrow n^3 \equiv 27 \equiv 6 \neq 0$ .

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 $n \equiv 2 \rightarrow n^3 \equiv 8 \equiv 1 \neq 0$ .  
 $n \equiv 3 \rightarrow n^3 \equiv 27 \equiv 6 \neq 0$ .  
 $n \equiv 4 \rightarrow n^3 \equiv (-3)^3 \equiv -3^3 \equiv -6 \equiv 1 \not\equiv 0$ .

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Lemma  $(\forall n)[n^3 \equiv 0 \pmod{7} \rightarrow n \equiv 0 \pmod{7}]$ .  
Take Contrapositive:  
 $(\forall n)[n \not\equiv 0 \pmod{7} \rightarrow n^3 \not\equiv 0 \pmod{7}]$ . 7 cases  
 $n \equiv 1 \rightarrow n^3 \equiv 1 \neq 0$ .  
 $n \equiv 2 \rightarrow n^3 \equiv 8 \equiv 1 \neq 0$ .  
 $n \equiv 3 \rightarrow n^3 \equiv 27 \equiv 6 \neq 0$ .  
 $n \equiv 4 \rightarrow n^3 \equiv (-3)^3 \equiv -3^3 \equiv -6 \equiv 1 \neq 0$ .  
 $n \equiv 5 \rightarrow n^3 \equiv (-2)^3 \equiv -2^3 \equiv -1 \equiv 6 \neq 0$ .

Want 
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Take Contrapositive:  
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 $n \equiv 1 \rightarrow n^3 \equiv 1 \neq 0$ .  
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 $n \equiv 5 \rightarrow n^3 \equiv (-2)^3 \equiv -2^3 \equiv -1 \equiv 6 \neq 0$ .  
 $n \equiv 6 \rightarrow n^3 \equiv (-1)^3 \equiv -1 \equiv 6 \neq 0$ .

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Want  $7^{1/3} \notin \mathbb{Q}$ . We need the Lemma. All  $\equiv$  is mod 7. Lemma  $(\forall n)[n^3 \equiv 0 \pmod{7}) \rightarrow n \equiv 0 \pmod{7}]$ . Take Contrapositive:  $(\forall n)[n \not\equiv 0 \pmod{7} \rightarrow n^3 \not\equiv 0 \pmod{7}]$ . 7 cases  $n \equiv 1 \rightarrow n^3 \equiv 1 \neq 0.$  $n \equiv 2 \rightarrow n^3 \equiv 8 \equiv 1 \neq 0.$  $n \equiv 3 \rightarrow n^3 \equiv 27 \equiv 6 \neq 0.$  $n \equiv 4 \rightarrow n^3 \equiv (-3)^3 \equiv -3^3 \equiv -6 \equiv 1 \neq 0.$  $n \equiv 5 \rightarrow n^3 \equiv (-2)^3 \equiv -2^3 \equiv -1 \equiv 6 \neq 0.$  $n \equiv 6 \rightarrow n^3 \equiv (-1)^3 \equiv -1 \equiv 6 \not\equiv 0.$ Proof of Lemma is done. Next slide is proof of irrationality.

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Want  $7^{1/3} \notin \mathbb{Q}$ . Assume BWOC that  $7^{1/3} \in \mathbb{Q}$ .

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Want  $7^{1/3} \notin \mathbb{Q}$ . Assume BWOC that  $7^{1/3} \in \mathbb{Q}$ . So there exists *a*, *b* in lowest terms such that  $7^{1/3} = \frac{a}{b}$  $b7^{1/3} = a$ 

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Want  $7^{1/3} \notin \mathbb{Q}$ . Assume BWOC that  $7^{1/3} \in \mathbb{Q}$ . So there exists *a*, *b* in lowest terms such that  $7^{1/3} = \frac{a}{b}$  $b7^{1/3} = a$  $7b^3 = a^3$  $a^3 \equiv 0$ .

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Want  $7^{1/3} \notin \mathbb{Q}$ . Assume BWOC that  $7^{1/3} \in \mathbb{Q}$ . So there exists *a*, *b* in lowest terms such that  $7^{1/3} = \frac{a}{b}$  $b7^{1/3} = a$  $7b^3 = a^3$  $a^3 \equiv 0$ . By Lemma  $a \equiv 0$ .

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Want  $7^{1/3} \notin \mathbb{Q}$ . Assume BWOC that  $7^{1/3} \in \mathbb{Q}$ . So there exists a, b in lowest terms such that  $7^{1/3} = \frac{a}{b}$  $b7^{1/3} = a$  $7b^3 = a^3$  $a^3 \equiv 0$ . By Lemma  $a \equiv 0$ . a = 7c.  $7b^3 = a^3 = (7c)^3 = 7^3c^3$ .

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Want  $7^{1/3} \notin \mathbb{Q}$ . Assume BWOC that  $7^{1/3} \in \mathbb{Q}$ . So there exists *a*, *b* in lowest terms such that  $7^{1/3} = \frac{a}{b}$  $b7^{1/3} = a$  $7b^3 = a^3$  $a^3 \equiv 0$ . By Lemma  $a \equiv 0$ . a = 7c.  $7b^3 = a^3 = (7c)^3 = 7^3c^3$ .  $b^3 = 7^2c^3$ .

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Want  $7^{1/3} \notin \mathbb{Q}$ . Assume BWOC that  $7^{1/3} \in \mathbb{Q}$ . So there exists a, b in lowest terms such that  $7^{1/3} = \frac{a}{b}$  $b7^{1/3} = a$  $7b^3 = a^3$  $a^3 \equiv 0$ . By Lemma  $a \equiv 0$ . a = 7c.  $7b^3 = a^3 = (7c)^3 = 7^3c^3$ .  $b^3 = 7^2c^3$ . By Lemma  $b \equiv 0$ .

Want  $7^{1/3} \notin \mathbb{Q}$ . Assume BWOC that  $7^{1/3} \in \mathbb{Q}$ . So there exists *a*, *b* in lowest terms such that  $7^{1/3} = \frac{a}{b}$  $b7^{1/3} = a$  $7b^3 = a^3$  $a^3 \equiv 0$ . By Lemma  $a \equiv 0$ . a = 7c.  $7b^3 = a^3 = (7c)^3 = 7^3c^3$ .  $b^3 = 7^2c^3$ . By Lemma  $b \equiv 0$ . AH-HA! 7 divides both *a* and *b*. So *a*, *b* not in lowest terms.

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Want  $7^{1/3} \notin \mathbb{O}$ . Assume BWOC that  $7^{1/3} \in \mathbb{O}$ . So there exists a, b in lowest terms such that  $7^{1/3} = \frac{a}{b}$  $b7^{1/3} = a$  $7b^3 = a^3$  $a^3 \equiv 0$ . By Lemma  $a \equiv 0$ . a = 7c.  $7b^3 = a^3 = (7c)^3 = 7^3c^3$ .  $b^3 = 7^2 c^3$ . By Lemma  $b \equiv 0$ . AH-HA! 7 divides both a and b. So a, b not in lowest terms. Contradiction!

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The above proof is a template for these kinds of proofs.

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The above proof is a template for these kinds of proofs. To show  $x^{1/z} \notin \mathbb{Q}$ .

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The above proof is a template for these kinds of proofs. To show  $x^{1/z} \notin \mathbb{Q}$ . Step 1 Prove  $(\forall n)[n^z \equiv 0 \pmod{x} \rightarrow n \equiv 0 \pmod{x}]$ .

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The above proof is a template for these kinds of proofs. To show  $x^{1/z} \notin \mathbb{Q}$ . **Step 1** Prove  $(\forall n)[n^z \equiv 0 \pmod{x} \rightarrow n \equiv 0 \pmod{x}]$ . Take Contrapositive  $(\forall n)[n \not\equiv 0 \pmod{x} \rightarrow n^z \not\equiv 0 \pmod{x}]$ .

The above proof is a template for these kinds of proofs. To show  $x^{1/z} \notin \mathbb{Q}$ . **Step 1** Prove  $(\forall n)[n^z \equiv 0 \pmod{x} \rightarrow n \equiv 0 \pmod{x}]$ . Take Contrapositive  $(\forall n)[n \not\equiv 0 \pmod{x} \rightarrow n^z \not\equiv 0 \pmod{x}]$ . Prove by x - 1 cases.

The above proof is a template for these kinds of proofs. To show  $x^{1/z} \notin \mathbb{Q}$ . **Step 1** Prove  $(\forall n)[n^z \equiv 0 \pmod{x} \rightarrow n \equiv 0 \pmod{x}]$ . Take Contrapositive  $(\forall n)[n \not\equiv 0 \pmod{x} \rightarrow n^z \not\equiv 0 \pmod{x}]$ . Prove by x - 1 cases.

**Step 2** Assume, BWOC, that  $x^{1/z} = \frac{a}{b}$ : *a*, *b* in lowest terms.

The above proof is a template for these kinds of proofs. To show  $x^{1/z} \notin \mathbb{Q}$ . **Step 1** Prove  $(\forall n)[n^z \equiv 0 \pmod{x} \rightarrow n \equiv 0 \pmod{x}]$ . Take Contrapositive  $(\forall n)[n \not\equiv 0 \pmod{x} \rightarrow n^z \not\equiv 0 \pmod{x}]$ . Prove by x - 1 cases.

**Step 2** Assume, BWOC, that  $x^{1/z} = \frac{a}{b}$ : a, b in lowest terms.  $bx^{1/z} = a$ 

The above proof is a template for these kinds of proofs. To show  $x^{1/z} \notin \mathbb{Q}$ . **Step 1** Prove  $(\forall n)[n^z \equiv 0 \pmod{x} \rightarrow n \equiv 0 \pmod{x}]$ . Take Contrapositive  $(\forall n)[n \not\equiv 0 \pmod{x} \rightarrow n^z \not\equiv 0 \pmod{x}]$ . Prove by x - 1 cases.

**Step 2** Assume, BWOC, that  $x^{1/z} = \frac{a}{b}$ : a, b in lowest terms.  $bx^{1/z} = a$  $b^z x = a^z$ 

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For proofs of irrationality using mods:



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For proofs of irrationality using mods:

- 1. The lemma is the only part that is not a template.
- 2. The lemma may have a lot of cases.
- 3. If you are trying to prove a rational is irrational, the proof will fall apart in the lemma.

Want  $7^{1/3} \notin \mathbb{Q}$ . Assume, BWOC that

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# Want $7^{1/3} \notin \mathbb{Q}$ . Assume, BWOC that $7^{1/3} = \frac{a}{b}$

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Want  $7^{1/3} \notin \mathbb{Q}$ . Assume, BWOC that  $7^{1/3} = \frac{a}{b}$   $b7^{1/3} = a$   $7b^3 = a^3$ Factor both sides  $p_{ab} = p_{ab}$  is the set of  $p_{ab}$ 

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# Proving $7^{1/3} \notin \mathbb{Q}$ Using UFT

Want  $7^{1/3} \notin \mathbb{Q}$ . Assume, BWOC that  $7^{1/3} = \frac{a}{b}$   $b7^{1/3} = a$   $7b^3 = a^3$ Factor both sides.  $p_1, \ldots, p_L$  is the set of primes that divide a or b.  $b = p_1^{b_1} \cdots p_L^{b_L}$   $a = p_1^{a_1} \cdots p_L^{a_L}$   $7p_1^{3b_1} \cdots p_L^{3b_L} = p_1^{3a_1} \cdots p_L^{3a_L}$ The number of 7's on the LHS is  $\equiv 1 \pmod{3}$ .

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## Proving $7^{1/3} \notin \mathbb{Q}$ Using UFT

Want  $7^{1/3} \notin \mathbb{O}$ . Assume, BWOC that  $7^{1/3} = \frac{a}{L}$  $b7^{1/3} = a$  $7h^3 = a^3$ Factor both sides.  $p_1, \ldots, p_L$  is the set of primes that divide *a* or *b*.  $b = p_1^{b_1} \cdots p_r^{b_r}$  $a = p_1^{a_1} \cdots p_l^{a_l}$  $7p_1^{3b_1}\cdots p_l^{3b_l} = p_1^{3a_1}\cdots p_l^{3a_l}$ The number of 7's on the LHS is  $\equiv 1 \pmod{3}$ . The number of 7's on the RHS is  $\equiv 0 \pmod{3}$ .

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## **Proving Irrationality Using UFT**

These proofs also have a very definite template.



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On HW05 you will do this proof for  $\sqrt{p}$ .

# **Primes**

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**Theorem** The number of primes is infinite. Assume, BWOC, that the number of primes is finite.

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**Case 2** N is not prime. Let p be a prime factor of N.

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