## HW 5 CMSC 452. Morally DUE Mar 7 THIS HOMEWORK IS THREE PAGES

- 1. (5 points) What is your name? Write it clearly. When is the midterm? Write that clearly too. Staple your HW. WHAT IS THE DAY/TIME OF THE MIDTERM? (HINT: The Midterm is March 30 IN CLASS at 11:00.)
- 2. (40 points)
  - (a) (Use our usual convention for pairs of numbers, so the alphabet is  $\{00, 01, 10, 11\}$ , which we usually write vertically.) Write a DFA for  $\{(x, y) : x = y + 3\}$ . Label each state
    - vince a DFA for  $\{(x, y) : x y + 5\}$ . Labe
    - A for accept,
    - R for reject, or
    - B for Bad Format.

(NOTE - the comma after "reject" is called an Oxford comma.)

- (b) Write a DFA for  $\{(x, y) : x \neq y + 3\}$ . Label each state
  - A for accept,
  - R for reject, or
  - B for Bad Format.

3. (25 points)  $A \subseteq \{0, 1\}^{\omega}$  is *D*-regular if there is a DFA  $M = (Q, \Sigma, \delta, s, F)$  such that:

 $x \in A \rightarrow$  if you run M on x you hit a state in F infinitely often

- $x \notin A \rightarrow$  if you run M on x you DO NOT hit a state in F infinitely often
- (a) Write a D-DFA for

 $L_1 = \{ X : X \text{ is infinite } \}.$ 

- (b) Take the *D*-DFA you wrote for  $L_1$  in part a. Swap the final and non-final states. Let  $L'_1$  be the language (subset of  $\{0, 1\}^{\omega}$ ) that your new automta accepts. Describe  $L'_1$ . Is it  $\overline{L_1}$  (also called the *compliment* of  $L_1$  or  $\{0, 1\}^{\omega} - L_1$ )?
- (c) Write a D-DFA for

 $\{X : X \text{ is infinite and } \mathsf{N} - X \text{ is infinite } \}.$ 

(d) Write an D-DFA for

 $\{X : \text{ there is an infinite number of } x \text{ such that } x \in X \text{ and } x + 1 \in X\}$ 

(e) Show this is NO *D*-DFA for  $\{X : X \text{ is finite}\}$ .

- 4. (35 points) A *J*-automata M is a tuple  $(Q, \Sigma, \delta, s, F)$  such that
  - Q is a set of states just like in a DFA.
  - $\Sigma$  is an alphabet just like in a DFA.
  - $\delta: Q \times \Sigma \to Q$  just like a DFA.
  - $s \in S$ , the start state just like a DFA.
  - F is a NOT a subset of Q. F is a set of subsets of F. For example, if  $Q = \{1, 2, 3, 4, 5, 6\}$  F could be  $\{\{1, 2, 5\}, \{1, 5\}, \{2, 3, 6\}\}$

Let  $x \in \{0,1\}^{\omega}$ . We say that *J*-automta *A* accepts *x* if, when you run *x* through *A*, the set of states that are visiting infinitely often is a set in *F*. For example, in the above example, if the set of states visiting infinitely often was  $\{1,2,5\}$  then ACCEPT, but if its  $\{1,2\}$  then REJECT.

A subset of  $\{0,1\}^{\omega}$  which is accepted by an *J*-automta is called *J*-regular.

- (a) (0 points) How would you complement an *J*-regular set?
- (b) (7 points) Show that if L is J-regular then  $\overline{L}$  is J-regular.
- (c) (7 points) Give a *J*-automata for  $\{X : X \text{ is infinite }\}$ .
- (d) (7 points) Give a *J*-automata for  $\{X : X \text{ is finite }\}$ .
- (e) (7 points) Give a *J*-automata for  $\{X : X \text{ is infinite and } N X \text{ is infinite}\}$ .
- (f) (7 points) Let  $T = \{x : x \equiv 0 \pmod{3}\}$ . Give a *J*-automata for  $\{X : X \cap T \text{ is infinite}\}$ .