# Help On the Classification HW

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### **Classification Problem**

For 
$$n = 14, ..., 20$$
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## **Classification Problem**

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For n = 14, ..., 20.

1) Find the pattern of

10^0 \pmod{n},

10^1 \pmod{n},

10^2 \pmod{n},

:
```

(You should write a program to help you with this.)

## **Classification Problem**

```
For n = 14, ..., 20.

1) Find the pattern of

10^0 \pmod{n},

10^1 \pmod{n},

10^2 \pmod{n},

:

(You should write a program to help you with this.)

2) Find the size of a DFA to classify mod n.
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 $\mathsf{AII} \equiv \mathsf{is} \ \mathsf{mod} \ \mathsf{7}.$ 



 $\text{AII} \equiv \text{is mod 7.} \\ 10^0 \equiv 1$ 



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 $AII \equiv \text{is mod } 7.$  $10^0 \equiv 1$  $10^1 \equiv 10 \equiv 3$ 

 $\begin{aligned} \text{AII} &\equiv \text{is mod 7.} \\ 10^0 &\equiv 1 \\ 10^1 &\equiv 10 &\equiv 3 \\ 10^2 &\equiv 10 \times 10 &\equiv 3 \times 3 &\equiv 9 &\equiv 2. \end{aligned}$ 

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 $\begin{aligned} \text{AII} &\equiv \text{is mod 7.} \\ 10^0 &\equiv 1 \\ 10^1 &\equiv 10 \equiv 3 \\ 10^2 &\equiv 10 \times 10 \equiv 3 \times 3 \equiv 9 \equiv 2. \\ 10^3 &\equiv 10^2 \times 10 \equiv 2 \times 3 \equiv 6. \end{aligned}$ 

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```

```
All \equiv is mod 7.
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10^1 = 10 = 3
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Pattern is  $\overline{1, 3, 2, 6, 4, 5}$ .

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Pattern is  $\overline{1, 3, 2, 6, 4, 5}$ .

The DFA to classify mod 7 has to keep track of

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Pattern is  $\overline{1, 3, 2, 6, 4, 5}$ .

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Pattern is  $\overline{1, 3, 2, 6, 4, 5}$ .

The DFA to classify mod 7 has to keep track of 1) The weighted sum mod 7.

2) The position of the digit mod 6.

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Pattern is 1, 3, 2, 6, 4, 5.

The DFA to classify mod 7 has to keep track of 1) The weighted sum mod 7. 2) The position of the digit mod 6.

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So the number of states is  $7 \times 6 = 42$ .

(This is not required for the HW.)  $Q = \{0, 1, 2, 3, 4, 5, 6\} \times \{0, 1, 2, 3, 4, 5\}.$ 

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(This is not required for the HW.)  

$$Q = \{0, 1, 2, 3, 4, 5, 6\} \times \{0, 1, 2, 3, 4, 5\}.$$
  
 $s = (0, 0).$ 

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 $\delta((q, 0), \sigma) = (q + 1 \times \sigma \pmod{7}, 1)$ 

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(This is not required for the HW.)  $Q = \{0, 1, 2, 3, 4, 5, 6\} \times \{0, 1, 2, 3, 4, 5\}.$  s = (0, 0).Weights are (1, 3, 2, 6, 4, 5).  $\delta((q, 0), \sigma) = (q + 1 \times \sigma \pmod{7}, 1)$   $\delta((q, 1), \sigma) = (q + 3 \times \sigma \pmod{7}, 2)$   $\delta((q, 2), \sigma) = (q + 2 \times \sigma \pmod{7}, 3)$  $\delta((q, 3), \sigma) = (q + 6 \times \sigma \pmod{7}, 4)$ 

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Mod 37



 $\begin{array}{l} \mbox{Mod 37} \\ \mbox{All} \equiv \mbox{are mod 37}. \end{array}$ 



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 $10^0 \equiv 1$ 

 $\begin{array}{l} \mbox{Mod 37} \\ \mbox{All} \equiv \mbox{are mod 37}. \end{array}$ 

 $\begin{array}{l} 10^0 \equiv 1 \\ 10^1 \equiv 10 \end{array}$ 



 $\begin{array}{l} \mbox{Mod 37} \\ \mbox{All} \equiv \mbox{are mod 37}. \end{array}$ 

 $\begin{array}{l} 10^0 \equiv 1 \\ 10^1 \equiv 10 \\ 10^2 \equiv 26 \end{array}$ 

Mod 37 All  $\equiv$  are mod 37.  $10^0 \equiv 1$  $10^1 \equiv 10$  $10^2 \equiv 26$  $10^3 = 26 \times 10 = 260 \equiv 1.$ 

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Mod 37 All  $\equiv$  are mod 37.  $10^0 \equiv 1$   $10^1 \equiv 10$   $10^2 \equiv 26$   $10^3 = 26 \times 10 = 260 \equiv 1$ . Pattern is (1, 10, 26).

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Mod 37 All  $\equiv$  are mod 37.  $10^0 = 1$  $10^1 \equiv 10$  $10^2 \equiv 26$  $10^3 = 26 \times 10 = 260 \equiv 1.$ Pattern is (1, 10, 26). Pattern length 3, mod is 37, so DFA has  $37 \times 3 = 111$  states.

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Mod 37 All  $\equiv$  are mod 37.  $10^0 = 1$  $10^{1} \equiv 10$  $10^2 \equiv 26$  $10^3 = 26 \times 10 = 260 \equiv 1$ . Pattern is (1, 10, 26). Pattern length 3, mod is 37, so DFA has  $37 \times 3 = 111$  states. Drawing the DFA would only be 3 cases.

# Pattern $\overline{a_1, \ldots, a_m}$

Mod n.

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Mod *n*. If pattern is  $\overline{a_1, a_2, \ldots, a_m}$  then the DFA will need to keep track of

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1) The weighted sum mod *n* 

Mod n.

If pattern is  $\overline{a_1, a_2, \ldots, a_m}$  then the DFA will need to keep track of

- 1) The weighted sum mod n
- 2) The position of the digit mod m.

Mod n.

If pattern is  $\overline{a_1, a_2, \ldots, a_m}$  then the DFA will need to keep track of

- 1) The weighted sum mod n
- 2) The position of the digit mod *m*.

So 
$$Q = \{0, \dots, n-1\} \times \{0, \dots, m-1\}.$$

Mod n.

If pattern is  $\overline{a_1, a_2, \ldots, a_m}$  then the DFA will need to keep track of

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- 1) The weighted sum mod n
- 2) The position of the digit mod *m*.

So 
$$Q = \{0, ..., n-1\} \times \{0, ..., m-1\}.$$

The number of states is *nm*.

Mod n.

If pattern is  $\overline{a_1, a_2, \ldots, a_m}$  then the DFA will need to keep track of

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- 1) The weighted sum mod n
- 2) The position of the digit mod *m*.

So 
$$Q = \{0, \dots, n-1\} \times \{0, \dots, m-1\}.$$

The number of states is nm.

What If The patterns might not be of this form?

Mod 6.



Mod 6.

 $\equiv$  is mod 6.



Mod 6.  $\equiv$  is mod 6.  $10^0 \equiv 1$ 



 $10^1=10\equiv 4$ 



Mod 6.  $\equiv$  is mod 6.  $10^0 \equiv 1$   $10^1 = 10 \equiv 4$  $10^2 = 10 \times 10 \equiv 4 \times 4 \equiv 16 \equiv 4$ 

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Mod 6.  $\equiv$  is mod 6.  $10^0 \equiv 1$   $10^1 = 10 \equiv 4$   $10^2 = 10 \times 10 \equiv 4 \times 4 \equiv 16 \equiv 4$ Pattern is  $(1,\overline{4})$ .

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Mod 6.  $\equiv$  is mod 6.  $10^0 \equiv 1$   $10^1 = 10 \equiv 4$   $10^2 = 10 \times 10 \equiv 4 \times 4 \equiv 16 \equiv 4$ Pattern is  $(1, \overline{4})$ . So have the pattern. How big is the DFA?

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Mod 6.  $\equiv$  is mod 6.  $10^0 \equiv 1$   $10^1 = 10 \equiv 4$   $10^2 = 10 \times 10 \equiv 4 \times 4 \equiv 16 \equiv 4$ Pattern is  $(1, \overline{4})$ . So have the pattern. How big is the DFA? In groups try to design the DFA.

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The pattern is  $(1, \overline{4})$ .



The pattern is  $(1, \overline{4})$ . The start state is DIFFERENT from the rest.

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The pattern is  $(1, \overline{4})$ . The start state is DIFFERENT from the rest. Any edge coming out of the start state is multiplied by 1.

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The pattern is  $(1, \overline{4})$ . The start state is DIFFERENT from the rest. Any edge coming out of the start state is multiplied by 1. After that always multiplied to 4.

The pattern is  $(1, \overline{4})$ . The start state is DIFFERENT from the rest. Any edge coming out of the start state is multiplied by 1. After that always multiplied to 4.

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 $Q = \{s\} \cup \{0, 1, 2, 3, 4, 5\}.$ 

The pattern is  $(1, \overline{4})$ .

The start state is DIFFERENT from the rest.

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 $egin{aligned} \mathcal{Q} &= \{s\} \cup \{0,1,2,3,4,5\}. \ \delta(s,\sigma) &= 1 imes \sigma \pmod{6} \end{aligned}$ 

The pattern is  $(1, \overline{4})$ .

The start state is DIFFERENT from the rest.

Any edge coming out of the start state is multiplied by 1. After that always multiplied to 4.

$$Q = \{s\} \cup \{0, 1, 2, 3, 4, 5\}.$$
  

$$\delta(s, \sigma) = 1 \times \sigma \pmod{6}$$
  
If  $q \in \{0, 1, 2, 3, 4, 5\}, \ \delta(q, \sigma) = q + 4\sigma \pmod{6}$ 

The pattern is  $(1, \overline{4})$ .

The start state is DIFFERENT from the rest.

Any edge coming out of the start state is multiplied by 1. After that always multiplied to 4.

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 $\begin{aligned} &Q = \{s\} \cup \{0, 1, 2, 3, 4, 5\}.\\ &\delta(s, \sigma) = -1 \times \sigma \pmod{6}\\ &\text{If } q \in \{0, 1, 2, 3, 4, 5\}, \ \delta(q, \sigma) = q + 4\sigma \pmod{6}\\ &\text{Number of states: } 1 + 6 = 7. \end{aligned}$ 

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Mod 375.

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Mod 375. Pattern is  $(1, 10, 100, \overline{250})$ .

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Mod 375. Pattern is  $(1, 10, 100, \overline{250})$ . Edges out of the start state have weight 1.

Mod 375. Pattern is  $(1, 10, 100, \overline{250})$ . Edges out of the start state have weight 1. Edges out of the second set of states have weight 10.

Mod 375. Pattern is  $(1, 10, 100, \overline{250})$ . Edges out of the start state have weight 1. Edges out of the second set of states have weight 10. Edges out of the third set of states have weight 100.

Mod 375. Pattern is  $(1, 10, 100, \overline{250})$ .

Edges out of the start state have weight 1.

Edges out of the second set of states have weight 10.

Edges out of the third set of states have weight 100.

Edges out of the fourth set of states have weight 250 and only go to the fourth set of states.

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Mod 375.





#### Mod 375. Weights are $(1, 10, 100, \overline{250})$ .



#### Mod 375. Weights are $(1, 10, 100, \overline{250})$ . $Q = \{s\} \cup \{0, \dots, 3742\} \times \{1, 2, 3\}$

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Mod 375. Weights are 
$$(1, 10, 100, \overline{250})$$
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 $Q = \{s\} \cup \{0, \dots, 3742\} \times \{1, 2, 3\}$   
 $\delta(s, \sigma) = (1 \times \sigma \pmod{375}, 1)$ .

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$$\begin{array}{ll} \text{Mod 375. Weights are } (1, 10, 100, 250). \\ Q = \{s\} \cup \{0, \dots, 3742\} \times \{1, 2, 3\} \\ \delta(s, \sigma) = & (1 \times \sigma \pmod{375}, 1). \\ \delta((q, 1), \sigma) = & (q + 10 \times \sigma \pmod{375}, 2). \end{array}$$

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## The DFA

Mod 375. Weights are 
$$(1, 10, 100, 250)$$
.  
 $Q = \{s\} \cup \{0, \dots, 3742\} \times \{1, 2, 3\}$   
 $\delta(s, \sigma) = (1 \times \sigma \pmod{375}, 1)$ .  
 $\delta((q, 1), \sigma) = (q + 10 \times \sigma \pmod{375}, 2)$ .  
 $\delta((q, 2), \sigma) = (q + 100 \times \sigma \pmod{375}, 3)$ .

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## The DFA

Mod 375. Weights are 
$$(1, 10, 100, \overline{250})$$
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 $\delta((q, 2), \sigma) = (q + 100 \times \sigma \pmod{375}, 3)$ .  
 $\delta((q, 3), \sigma) = (q + 250 \times \sigma \pmod{375}, 3)$ .

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Mod n.



Mod n.

Pattern  $(a_1, \ldots, a_{m-1}, \overline{a_m})$ 



Mod n.

- Pattern  $(a_1, \ldots, a_{m-1}, \overline{a_m})$
- Uses 1 + (m-1)n states.

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Mod n.

- Pattern  $(a_1, \ldots, a_{m-1}, \overline{a_m})$
- Uses 1 + (m-1)n states.

Leave it to you to work this out.

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Identical to previous slides. The DFA is easier since the last set of states have transition to themselves.

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