Solutions to HW09 Problems

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A & B do PRIV-LWE with $\vec{k} = (11, 100, 39, 4)$, p = 1009, $\gamma = 2$. All \equiv are mod 1009.

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A computes

 $C = (11, 100, 39, 4) \cdot (1, 2, 3, 4) = 11 + 200 + 117 + 16 = 344 \equiv 344.$

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A sends (1, 2, 3, 4; 598).

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A sends (5, 10, 41, 3; 647).

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 $(11, 100, 39, 4) \cdot (12, 39, 44, 19) = 5824 \equiv 779$

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779 is 0 away from 779 and 0 < 2. So the bit is 0.

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If $|D - C| \le 4$ then output **A probably sent a 0.**

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If $|D - C| \le 4$ then output **A probably sent a 0.**

If $|D - (C + \frac{p}{4})| \le 4$ then output **A probably sent a 1.**

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If $|D - C| \le 4$ then output **A probably sent a 0.**

If $|D - (C + \frac{p}{4})| \le 4$ then output **A probably sent a 1**.

If NEITHER then output *E* tampered with the message.

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 $C \equiv (1, 2, 3, 4) \cdot (10, 201, 89, 8) \equiv 711.$

This is NOT close to 5, nor is 711 + 500 \equiv 1211, so TAMPERED WITH.

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(11, 40, 99, 101; 245).

 $C \equiv (11, 40, 99, 101) \cdot (10, 201, 89, 8) \equiv 1745.$

1745 is NOT 245.

But $1745 + 500 \equiv 242$ IS close to 245. (It needs to be within 4 and it is) So A probably sent 1.

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 $7k_1 + 13k_2 + 22k_3 + 100k_4 \in \{618 - 2, 618 - 1, 618, 618 + 1, 618 + 2\}$

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SO

 $7k_1 + 13k_2 + 22k_3 + 100k_4 \in \{616, 617, 618, 619, 620\}$