# BILL, RECORD LECTURE!!!!

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# Gen 2-letter Sub and Matrix Codes

October 9, 2021

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Need bijection of  $\{0,\ldots,25\}\times\{0,\ldots,25\}$  that is easy to use.

**Def** Matrix Cipher. Pick M a 2  $\times$  2 matrix.

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Do you recognize the expression ad - bc? Determinant!

# Inverse Matrix in $\ensuremath{\mathbb{C}}$ and in Mods

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#### Inverse Matrix in $\ensuremath{\mathbb{C}}$ and in Mods

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- 1. Matrix *M* over  $\mathbb{C}$  has an inverse iff  $ad bc \neq 0$ .
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- Matrix M over Mod 26 has an inverse iff ad bc is rel prime to 26 iff ad - bc has no factors of 2 or 13 iff has an inverse in Mod 26.

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So what to do?

**Def** Pick  $n \in \mathbb{N}$  and M an  $\mathbf{n} \times \mathbf{n}$  matrix with det rel prime to 26.

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We'll take n = 8.

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- 3. Eve cannot use brute force. Key Space is  $\sim 26^{64} \sim 10^{90}$ , Number of protons is  $\sim 10^{79}$ . (the number of non-invertible matrices is very small so  $26^{64}$  is a good approximation).

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2.2 Decode T into T' using M.

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2. No and we can PROVE we can't do better with ciphertext-only.

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- **Yes**. We can crack in time  $8 \times 26^8$ .

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is every third letter. Can do IS-ENGLISH on it.

Eve knows that Alice and Bob decode with  $8 \times 8$  Matrix *M*. Ciphertext is

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Takes  $8 \times 26^8$  steps.

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- 4. After Alice and Bob learn the technique they have to up their parameters. This is a mild win for Eve in that A and B have to work harder.

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Proofs rely on limiting what Eve can do, and hence do not work if Eve does something else.

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Eve will have old messages and what they decoded to.

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# Example of What Eve Might Know

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5. Eve knows that (3,9) = M(13,24).

Example using  $2 \times 2$  Matrix Cipher. Eve learns that (13,24) encrypts to (3,9). Hence:

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- 3. We need to better refine our notion of **attack**.
- 4. We will do this in a later slide packet.