BILL, RECORD LECTURE!!!!

BILL RECORD LECTURE!!!



Grid Colorings that Avoid Rectangles

January 23, 2025

Credit Where Credit is Due

This talk is based on a paper by Stephen Fenner William Gasarch Charles Glover Semmy Purewal

▲□▶ ▲□▶ ▲三▶ ▲三▶ 三三 - のへぐ

Ramsey Theory

This talk is an example of Ramsey Theory.



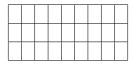
This talk is an example of Ramsey Theory.

I will be teaching CMSC 752: **Ramsey Theory and its "Applications"** in the Spring of 2025.

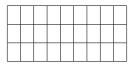
This talk is an example of Ramsey Theory.

I will be teaching CMSC 752: **Ramsey Theory and its "Applications"** in the Spring of 2025.

So this talk is an advertisement for the course.

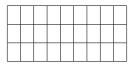






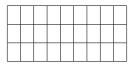
Is there a 2-coloring of 3×9 with no mono rectangles?





▲□▶ ▲□▶ ▲目▶ ▲目▶ 三日 - のへで

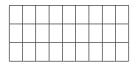
Is there a 2-coloring of 3×9 with no mono rectangles? What is a mono rectangle? Here is a an example:



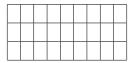
Is there a 2-coloring of 3×9 with no mono rectangles? What is a mono rectangle? Here is a an example:

R			R	
R			R	

▲□▶ ▲□▶ ▲目▶ ▲目▶ 三日 - のへで

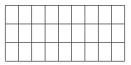






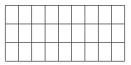
<□▶ <□▶ < □▶ < □▶ < □▶ < □▶ < □ > ○ < ○

Vote



Vote

1. There is a 2-coloring of 3×9 with NO mono rectangles.

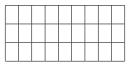


Vote

1. There is a 2-coloring of 3×9 with NO mono rectangles.

▲□▶ ▲□▶ ▲目▶ ▲目▶ 三日 - のへの

2. All 2-colorings of 3×9 have a mono rectangle.

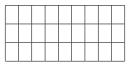


Vote

1. There is a 2-coloring of 3×9 with NO mono rectangles.

▲ロ ▶ ▲周 ▶ ▲ ヨ ▶ ▲ ヨ ▶ → ヨ → の Q @

- 2. All 2-colorings of 3×9 have a mono rectangle.
- 3. The problem is **UNKNOWN TO SCIENCE**.



Vote

1. There is a 2-coloring of 3×9 with NO mono rectangles.

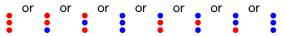
- 2. All 2-colorings of 3×9 have a mono rectangle.
- 3. The problem is **UNKNOWN TO SCIENCE**.

Answer on the next slide.

Given a 2-coloring of 3×9 look at each column.

Given a 2-coloring of 3×9 look at each column.

Each column is either



Given a 2-coloring of 3×9 look at each column.

Each column is either



Key: A 2-coloring of 3×9 is an 8-coloring of the 9 columns.

Given a 2-coloring of 3×9 look at each column.

Each column is either



Key: A 2-coloring of 3×9 is an 8-coloring of the 9 columns. So some column-color appears twice.

Given a 2-coloring of 3×9 look at each column.

Each column is either

or or or or or or or

Key: A 2-coloring of 3×9 is an 8-coloring of the 9 columns.

So some column-color appears twice. Example:

R		R	
B		B	
R		R	

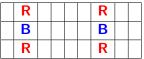
Given a 2-coloring of 3×9 look at each column.

Each column is either

or or or or or or or

Key: A 2-coloring of 3×9 is an 8-coloring of the 9 columns.

So some column-color appears twice. Example:



Can easily show that the two repeat-columns lead to a mono rectangle.

Work in groups:



Work in groups:

1. Is there a 2-coloring of 3×8 with no mono rectangles?

▲□▶ ▲□▶ ▲ 臣▶ ▲ 臣▶ ― 臣 … のへぐ

Work in groups:

- 1. Is there a 2-coloring of 3×8 with no mono rectangles?
- 2. Is there a 2-coloring of 3×7 with no mono rectangles?

▲□▶ ▲□▶ ▲三▶ ▲三▶ 三三 - のへぐ

Work in groups:

- 1. Is there a 2-coloring of 3×8 with no mono rectangles?
- 2. Is there a 2-coloring of 3×7 with no mono rectangles?
- 3. Is there a 2-coloring of 3×6 with no mono rectangles?

Work in groups:

- 1. Is there a 2-coloring of 3×8 with no mono rectangles?
- 2. Is there a 2-coloring of 3×7 with no mono rectangles?
- 3. Is there a 2-coloring of 3×6 with no mono rectangles?
- 4. Is there a 2-coloring of 3×5 with no mono rectangles?

Work in groups:

- 1. Is there a 2-coloring of 3×8 with no mono rectangles?
- 2. Is there a 2-coloring of 3×7 with no mono rectangles?
- 3. Is there a 2-coloring of 3×6 with no mono rectangles?
- 4. Is there a 2-coloring of 3×5 with no mono rectangles?
- 5. Is there a 2-coloring of 3×4 with no mono rectangles?

Work in groups:

- 1. Is there a 2-coloring of 3×8 with no mono rectangles?
- 2. Is there a 2-coloring of 3×7 with no mono rectangles?
- 3. Is there a 2-coloring of 3×6 with no mono rectangles?
- 4. Is there a 2-coloring of 3×5 with no mono rectangles?
- 5. Is there a 2-coloring of 3×4 with no mono rectangles?
- 6. Is there a 2-coloring of 3×3 with no mono rectangles? YES: Example:

R	В	R
R	В	В
R	R	В

・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・
 ・

▲□▶▲圖▶▲≣▶▲≣▶ ≣ の�?

1. Is there a 2-coloring of 3×8 with no mono rectangles?

1. Is there a 2-coloring of 3×8 with no mono rectangles? NO: to avoid a repeat col must have col Easily get mono

rectangle.



1. Is there a 2-coloring of 3×8 with no mono rectangles? NO: to avoid a repeat col must have col Easily get mono

rectangle.

2. Is there a 2-coloring of 3×7 with no mono rectangles?

Is there a 2-coloring of 3 × 8 with no mono rectangles?
 NO: to avoid a repeat col must have col Easily get mono

rectangle.

2. Is there a 2-coloring of 3×7 with no mono rectangles? NO: to avoid a repeat col must have col OR

▲ロ ▶ ▲ □ ▶ ▲ □ ▶ ▲ □ ▶ ▲ □ ▶ ● ○ ○ ○

Easily get mono rectangle.

Is there a 2-coloring of 3 × 8 with no mono rectangles?
 NO: to avoid a repeat col must have col Easily get mono

rectangle.

Is there a 2-coloring of 3 × 7 with no mono rectangles?
 NO: to avoid a repeat col must have col OR

Easily get mono rectangle.

3. Is there a 2-coloring of 3×6 with no mono rectangles?

▲ロ ▶ ▲ □ ▶ ▲ □ ▶ ▲ □ ▶ ▲ □ ▶ ● ○ ○ ○

Is there a 2-coloring of 3 × 8 with no mono rectangles?
 NO: to avoid a repeat col must have col Easily get mono

rectangle.

Is there a 2-coloring of 3 × 7 with no mono rectangles?
 NO: to avoid a repeat col must have col
 OR

Easily get mono rectangle.

3. Is there a 2-coloring of 3×6 with no mono rectangles? YES

▲ロ ▶ ▲ □ ▶ ▲ □ ▶ ▲ □ ▶ ▲ □ ▶ ● ○ ○ ○

2-Coloring 3 \times 8, 3 \times 7, ...

Is there a 2-coloring of 3 × 8 with no mono rectangles?
 NO: to avoid a repeat col must have col Easily get mono

rectangle.

2. Is there a 2-coloring of 3×7 with no mono rectangles? NO: to avoid a repeat col must have col OR

Easily get mono rectangle.

3. Is there a 2-coloring of 3 × 6 with no mono rectangles? YES

R	R	R	В	В	В
R	В	В	R	В	R
В	R	В	В	R	R

▲ロ ▶ ▲ □ ▶ ▲ □ ▶ ▲ □ ▶ ▲ □ ▶ ● ○ ○ ○

2-Coloring 3 \times 8, 3 \times 7, ...

Is there a 2-coloring of 3 × 8 with no mono rectangles?
 NO: to avoid a repeat col must have col Easily get mono

rectangle.

2. Is there a 2-coloring of 3×7 with no mono rectangles? NO: to avoid a repeat col must have col OR

Easily get mono rectangle.

3. Is there a 2-coloring of 3 × 6 with no mono rectangles? YES



ション ふゆ アメビア メロア しょうくり

4. Hence there is a 2-coloring of 3 \times 5, 3 \times 4, 3 \times 3 with no mono rectangles.

Diff proof that all 2-col of 3×7 have mono rectangle.

▲□▶ ▲□▶ ▲ 臣▶ ▲ 臣▶ ― 臣 … のへぐ

Diff proof that all 2-col of 3×7 have mono rectangle. Let COL be a 2-coloring of 3×7 . There are $3 \times 7 = 21$ squares so there must be ≥ 11 that are the same color, say **R**.

Diff proof that all 2-col of 3×7 have mono rectangle. Let COL be a 2-coloring of 3×7 . There are $3 \times 7 = 21$ squares so there must be ≥ 11 that are the same color, say **R**. **Case 1** Some col is Then the other columns have to have ≤ 1 **R** in them (or else you get a mono Rectangle). Total: 3 + 1 + 1 + 1 + 1 + 1 = 9 < 11.

Diff proof that all 2-col of 3×7 have mono rectangle. Let COL be a 2-coloring of 3×7 . There are $3 \times 7 = 21$ squares so there must be ≥ 11 that are the same color, say **R**. **Case 1** Some col is Then the other columns have to have ≤ 1 **R** in them (or else you get a mono Rectangle). Total: 3 + 1 + 1 + 1 + 1 + 1 + 1 = 9 < 11. **Case 2** \leq 4 cols have two **R** in them. Total:

 $\leq 2 + 2 + 2 + 2 + 1 + 1 = 10 < 11.$

Diff proof that all 2-col of 3×7 have mono rectangle. Let COL be a 2-coloring of 3×7 . There are $3 \times 7 = 21$ squares so there must be ≥ 11 that are the same color, say **R**.

Case 1 Some col is Then the other columns have to have $\leq 1 R$

in them (or else you get a mono Rectangle). Total: 3+1+1+1+1+1+1=9<11.

Case 2 \leq 4 cols have two R in them. Total: \leq 2 + 2 + 2 + 2 + 1 + 1 = 10 < 11.

Case 3 \geq 5 cols have two **R** in them. Map each col to the $\{i, j\}$ such that it has **R** in the *i*th and *j*th spot. Domain \geq 5, range $\binom{3}{2} = 3$ so two cols map to the same $\{i, j\}$. Get mono Rectangle.

 $a \times b$ is 2-colorable if there is a 2-coloring with no mono rectangles.

 $a \times b$ is 2-colorable if there is a 2-coloring with no mono rectangles. What we know

 $a \times b$ is 2-colorable if there is a 2-coloring with no mono rectangles. What we know

▲□▶ ▲□▶ ▲ 臣▶ ▲ 臣▶ ― 臣 … のへぐ

 \blacktriangleright 2 × *b* is always 2-colorable

 $a \times b$ is 2-colorable if there is a 2-coloring with no mono rectangles. What we know

- \blacktriangleright 2 × *b* is always 2-colorable
- ▶ 3×3 , ..., 3×6 2-colorable.

 $a \times b$ is 2-colorable if there is a 2-coloring with no mono rectangles. What we know

- \blacktriangleright 2 × *b* is always 2-colorable
- ▶ 3×3 , ..., 3×6 2-colorable.
- ▶ $3 \times b$ where $b \ge 7$ NOT 2-colorable.

 $a \times b$ is 2-colorable if there is a 2-coloring with no mono rectangles. What we know

- \blacktriangleright 2 × *b* is always 2-colorable
- ▶ 3×3 , ..., 3×6 2-colorable.
- ▶ $3 \times b$ where $b \ge 7$ NOT 2-colorable.
- 4×4 , 4×5 , 4×6 unknown so far.

 $a \times b$ is 2-colorable if there is a 2-coloring with no mono rectangles. What we know

- \blacktriangleright 2 × *b* is always 2-colorable
- ▶ 3×3 , ..., 3×6 2-colorable.
- ▶ $3 \times b$ where $b \ge 7$ NOT 2-colorable.
- 4×4 , 4×5 , 4×6 unknown so far.
- $4 \times b$ where $b \ge 7$ NOT 2-colorable.

 $a \times b$ is 2-colorable if there is a 2-coloring with no mono rectangles. What we know

- \blacktriangleright 2 × *b* is always 2-colorable
- ▶ 3×3 , ..., 3×6 2-colorable.
- ▶ $3 \times b$ where $b \ge 7$ NOT 2-colorable.
- 4×4 , 4×5 , 4×6 unknown so far.
- $4 \times b$ where $b \ge 7$ NOT 2-colorable.
- ▶ 5×5 , 5×6 unknown so far.

 $a \times b$ is 2-colorable if there is a 2-coloring with no mono rectangles. What we know

▲□▶ ▲□▶ ▲ 三▶ ▲ 三▶ ▲ 三 ● ● ●

- \blacktriangleright 2 × *b* is always 2-colorable
- ▶ 3×3 , ..., 3×6 2-colorable.
- ▶ $3 \times b$ where $b \ge 7$ NOT 2-colorable.
- 4×4 , 4×5 , 4×6 unknown so far.
- $4 \times b$ where $b \ge 7$ NOT 2-colorable.
- ▶ 5×5 , 5×6 unknown so far.
- ▶ $5 \times b$ where $b \ge 7$ NOT 2-colorable.

 $a \times b$ is 2-colorable if there is a 2-coloring with no mono rectangles. What we know

▲ロ ▶ ▲周 ▶ ▲ ヨ ▶ ▲ ヨ ▶ → 目 → の Q @

- \blacktriangleright 2 × *b* is always 2-colorable
- ▶ 3×3 , ..., 3×6 2-colorable.
- ▶ $3 \times b$ where $b \ge 7$ NOT 2-colorable.
- 4×4 , 4×5 , 4×6 unknown so far.
- $4 \times b$ where $b \ge 7$ NOT 2-colorable.
- ▶ 5×5 , 5×6 unknown so far.
- ▶ $5 \times b$ where $b \ge 7$ NOT 2-colorable.
- 6×6 unknown so far.

 $a \times b$ is 2-colorable if there is a 2-coloring with no mono rectangles. What we know

- \blacktriangleright 2 × *b* is always 2-colorable
- ▶ 3×3 , ..., 3×6 2-colorable.
- ▶ $3 \times b$ where $b \ge 7$ NOT 2-colorable.
- 4×4 , 4×5 , 4×6 unknown so far.
- $4 \times b$ where $b \ge 7$ NOT 2-colorable.
- ▶ 5×5 , 5×6 unknown so far.
- ▶ $5 \times b$ where $b \ge 7$ NOT 2-colorable.
- 6×6 unknown so far.
- $6 \times b$ where $b \ge 7$ NOT 2-colorable.

 $a \times b$ is 2-colorable if there is a 2-coloring with no mono rectangles. What we know

- \blacktriangleright 2 × *b* is always 2-colorable
- ▶ 3×3 , ..., 3×6 2-colorable.
- ▶ $3 \times b$ where $b \ge 7$ NOT 2-colorable.
- 4×4 , 4×5 , 4×6 unknown so far.
- $4 \times b$ where $b \ge 7$ NOT 2-colorable.
- ▶ 5×5 , 5×6 unknown so far.
- ▶ $5 \times b$ where $b \ge 7$ NOT 2-colorable.
- 6×6 unknown so far.
- $6 \times b$ where $b \ge 7$ NOT 2-colorable.

Work on the 4 \times 4, 4 \times 5 4 \times 6.

 4×6 IS 2-Colorable

R R R В В В В R R R В В Β R BRBR В В RBRR

▲□▶ ▲圖▶ ▲国▶ ▲国▶ - 国 - のへで

▲□▶▲圖▶▲圖▶▲圖▶ 圖 のへで

What we know

What we know

 \blacktriangleright 2 × *b* is always 2-colorable

▲□▶ ▲□▶ ▲ 臣▶ ▲ 臣▶ ― 臣 … のへぐ

What we know

- \blacktriangleright 2 × *b* is always 2-colorable
- ▶ $3 \times 3, \ldots, 3 \times 6$ 2-colorable.

▲□▶ ▲□▶ ▲ 臣▶ ▲ 臣▶ ― 臣 … のへぐ

What we know

- \blacktriangleright 2 × *b* is always 2-colorable
- ▶ 3×3 , ..., 3×6 2-colorable.
- ▶ $3 \times b$ where $b \ge 7$ NOT 2-colorable.

▲□▶ ▲□▶ ▲三▶ ▲三▶ 三三 - のへぐ

What we know

- \blacktriangleright 2 × b is always 2-colorable
- ▶ 3×3 , ..., 3×6 2-colorable.
- ▶ $3 \times b$ where $b \ge 7$ NOT 2-colorable.

▲ロ ▶ ▲周 ▶ ▲ ヨ ▶ ▲ ヨ ▶ → ヨ → の Q @

• 4 × 4, 4 × 5, 4 × 6 are 2-colorable

What we know

- \blacktriangleright 2 × b is always 2-colorable
- ▶ 3×3 , ..., 3×6 2-colorable.
- $3 \times b$ where $b \ge 7$ NOT 2-colorable.
- 4 × 4, 4 × 5, 4 × 6 are 2-colorable
- $4 \times b$ where $b \ge 7$ NOT 2-colorable.

▲ロ ▶ ▲周 ▶ ▲ ヨ ▶ ▲ ヨ ▶ → ヨ → の Q @

What we know

- \blacktriangleright 2 × b is always 2-colorable
- ▶ 3×3 , ..., 3×6 2-colorable.
- $3 \times b$ where $b \ge 7$ NOT 2-colorable.
- 4 × 4, 4 × 5, 4 × 6 are 2-colorable
- $4 \times b$ where $b \ge 7$ NOT 2-colorable.

▲ロ ▶ ▲周 ▶ ▲ ヨ ▶ ▲ ヨ ▶ → ヨ → の Q @

▶ 5×5 , 5×6 unknown so far.

What we know

- \blacktriangleright 2 × b is always 2-colorable
- $3 \times 3, \ldots, 3 \times 6$ 2-colorable.
- $3 \times b$ where $b \ge 7$ NOT 2-colorable.
- 4 × 4, 4 × 5, 4 × 6 are 2-colorable
- $4 \times b$ where $b \ge 7$ NOT 2-colorable.
- ▶ 5×5 , 5×6 unknown so far.
- ▶ $5 \times b$ where $b \ge 7$ NOT 2-colorable.

▲ロ ▶ ▲周 ▶ ▲ ヨ ▶ ▲ ヨ ▶ → ヨ → の Q @

What we know

- \blacktriangleright 2 × b is always 2-colorable
- ▶ 3×3 , ..., 3×6 2-colorable.
- $3 \times b$ where $b \ge 7$ NOT 2-colorable.
- 4 × 4, 4 × 5, 4 × 6 are 2-colorable
- $4 \times b$ where $b \ge 7$ NOT 2-colorable.
- ▶ 5×5 , 5×6 unknown so far.
- ▶ $5 \times b$ where $b \ge 7$ NOT 2-colorable.

▲ロ ▶ ▲周 ▶ ▲ ヨ ▶ ▲ ヨ ▶ → ヨ → の Q @

• 6×6 unknown so far.

What we know

- \blacktriangleright 2 × *b* is always 2-colorable
- ▶ 3×3 , ..., 3×6 2-colorable.
- $3 \times b$ where $b \ge 7$ NOT 2-colorable.
- 4 × 4, 4 × 5, 4 × 6 are 2-colorable
- $4 \times b$ where $b \ge 7$ NOT 2-colorable.
- ▶ 5×5 , 5×6 unknown so far.
- ▶ $5 \times b$ where $b \ge 7$ NOT 2-colorable.
- 6×6 unknown so far.
- $6 \times b$ where $b \ge 7$ NOT 2-colorable.

▲ロ ▶ ▲ □ ▶ ▲ □ ▶ ▲ □ ▶ ▲ □ ▶ ● ○ ○ ○

What we know

- \blacktriangleright 2 × *b* is always 2-colorable
- ▶ 3×3 , ..., 3×6 2-colorable.
- $3 \times b$ where $b \ge 7$ NOT 2-colorable.
- 4 × 4, 4 × 5, 4 × 6 are 2-colorable
- $4 \times b$ where $b \ge 7$ NOT 2-colorable.
- ▶ 5×5 , 5×6 unknown so far.
- ▶ $5 \times b$ where $b \ge 7$ NOT 2-colorable.
- 6×6 unknown so far.
- $6 \times b$ where $b \ge 7$ NOT 2-colorable.

▲□▶ ▲□▶ ▲ 三▶ ▲ 三▶ ▲ 三 ● ● ●

Work on 5×5 , 5×6 .

5×5 IS NOT 2-Colorable!

Let ${\rm COL}$ be a 2-coloring of 5 \times 5.



5×5 IS NOT 2-Colorable!

Let COL be a 2-coloring of 5×5 . Some color must occur ≥ 13 times.

▲□▶ ▲□▶ ▲目▶ ▲目▶ 三日 - のへで

Case 1: There is a column with 5 R's

Case 1: There is a column with 5 **R**'s.

R	0	0	0	0
R	0	0	0	0
R	0	0	0	0
R	0	0	0	0
R	0	0	0	0

Remaining columns have $\leq 1 \ \mathbf{R}$ so

Number of **R**'s $\leq 5 + 1 + 1 + 1 + 1 = 9 < 13$.

<□ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ >

Case 2: There is a column with 4 *R*'s

Case 2: There is a column with 4 *R*'s.

R	0	0	0	0
R	0	0	0	0
R	0	0	0	0
R	0	0	0	0
0	0	0	0	0

Remaining columns have $\leq 2 \mathbb{R}$'s

Number of R's $\leq 4+2+2+2+2\leq 12<13$

▲□▶ ▲□▶ ▲目▶ ▲目▶ 三日 - のへの

Case 3: Max in a column is 3 R's

Case 3: Max in a column is 3 **R**'s. **Case 3a:** There are ≤ 2 columns with 3 **R**'s.

Number of **R**'s $\leq 3 + 3 + 2 + 2 + 2 \leq 12 < 13$.

Case 3b: There are \geq 3 columns with 3 **R**'s.

R	0	0	0	0
R	0	0	0	0
R	R	0	0	0
0	R	0	0	0
0	R	0	0	0

Can't put in a third column with 3 R's!

Case 4: Max in a column is $\leq 2R$'s

Case 4: Max in a column is $\leq 2\mathbf{R}$'s.

Number of **R**'s $\leq 2 + 2 + 2 + 2 + 2 \leq 10 < 13$.

*ロ * * @ * * ミ * ミ * ・ ミ * の < や

No more cases. We are Done! Q.E.D.

What we know

What we know

 \blacktriangleright 2 × *b* is always 2-colorable

▲□▶ ▲□▶ ▲ 臣▶ ▲ 臣▶ ― 臣 … のへぐ

What we know

- \blacktriangleright 2 × *b* is always 2-colorable
- ▶ 3×3 , ..., 3×6 2-colorable.

▲□▶ ▲□▶ ▲ 臣▶ ▲ 臣▶ ― 臣 … のへぐ

What we know

- \blacktriangleright 2 × *b* is always 2-colorable
- ▶ 3×3 , ..., 3×6 2-colorable.
- ▶ $3 \times b$ where $b \ge 7$ NOT 2-colorable.

▲□▶ ▲□▶ ▲三▶ ▲三▶ 三三 - のへぐ

What we know

- \blacktriangleright 2 × *b* is always 2-colorable
- ▶ 3×3 , ..., 3×6 2-colorable.
- ▶ $3 \times b$ where $b \ge 7$ NOT 2-colorable.

▲ロ ▶ ▲周 ▶ ▲ ヨ ▶ ▲ ヨ ▶ → 目 → の Q @

• 4 × 4, 4 × 5, 4 × 6 are 2-colorable

What we know

- \blacktriangleright 2 × b is always 2-colorable
- ▶ 3×3 , ..., 3×6 2-colorable.
- $3 \times b$ where $b \ge 7$ NOT 2-colorable.
- 4 × 4, 4 × 5, 4 × 6 are 2-colorable
- $4 \times b$ where $b \ge 7$ NOT 2-colorable.

▲ロ ▶ ▲周 ▶ ▲ ヨ ▶ ▲ ヨ ▶ → ヨ → の Q @

What we know

- \blacktriangleright 2 × b is always 2-colorable
- ▶ 3×3 , ..., 3×6 2-colorable.
- $3 \times b$ where $b \ge 7$ NOT 2-colorable.
- 4 × 4, 4 × 5, 4 × 6 are 2-colorable
- $4 \times b$ where $b \ge 7$ NOT 2-colorable.

▲ロ ▶ ▲周 ▶ ▲ ヨ ▶ ▲ ヨ ▶ → ヨ → の Q @

▶ 5×5 , 5×6 NOT 2-colorable.

What we know

- \blacktriangleright 2 × b is always 2-colorable
- ▶ 3×3 , ..., 3×6 2-colorable.
- $3 \times b$ where $b \ge 7$ NOT 2-colorable.
- 4 × 4, 4 × 5, 4 × 6 are 2-colorable
- $4 \times b$ where $b \ge 7$ NOT 2-colorable.
- ▶ 5×5 , 5×6 NOT 2-colorable.
- $5 \times b$ where $b \ge 7$ NOT 2-colorable.

▲ロ ▶ ▲周 ▶ ▲ ヨ ▶ ▲ ヨ ▶ → ヨ → の Q @

What we know

- \blacktriangleright 2 × b is always 2-colorable
- ▶ 3×3 , ..., 3×6 2-colorable.
- $3 \times b$ where $b \ge 7$ NOT 2-colorable.
- 4 × 4, 4 × 5, 4 × 6 are 2-colorable
- $4 \times b$ where $b \ge 7$ NOT 2-colorable.
- ▶ 5×5 , 5×6 NOT 2-colorable.
- ▶ $5 \times b$ where $b \ge 7$ NOT 2-colorable.

• 6×6 NOT 2-colorable.

What we know

- \blacktriangleright 2 × *b* is always 2-colorable
- ▶ 3×3 , ..., 3×6 2-colorable.
- $3 \times b$ where $b \ge 7$ NOT 2-colorable.
- 4 × 4, 4 × 5, 4 × 6 are 2-colorable
- $4 \times b$ where $b \ge 7$ NOT 2-colorable.
- ▶ 5×5 , 5×6 NOT 2-colorable.
- ▶ $5 \times b$ where $b \ge 7$ NOT 2-colorable.
- 6×6 NOT 2-colorable.

We now know exactly what grids are 2-colorable.

▲ロ ▶ ▲ □ ▶ ▲ □ ▶ ▲ □ ▶ ▲ □ ▶ ● ○ ○ ○

What we know

- $2 \times b$ is always 2-colorable
- ▶ 3×3 , ..., 3×6 2-colorable.
- $3 \times b$ where $b \ge 7$ NOT 2-colorable.
- 4 × 4, 4 × 5, 4 × 6 are 2-colorable
- $4 \times b$ where $b \ge 7$ NOT 2-colorable.
- ▶ 5×5 , 5×6 NOT 2-colorable.
- $5 \times b$ where $b \ge 7$ NOT 2-colorable.
- 6×6 NOT 2-colorable.

We now know **exactly** what grids are 2-colorable. Can we say it more succinctly?

▲ロ ▶ ▲ □ ▶ ▲ □ ▶ ▲ □ ▶ ▲ □ ▶ ● ○ ○ ○

Def $n \times m$ contains $a \times b$ if $a \le n$ and $b \le m$. **Thm** For all c there exists a unique finite set of grids OBS_c such that

- $n \times m$ is *c*-colorable iff
- $n \times m$ does not contain any element of OBS_c .

Def $n \times m$ contains $a \times b$ if $a \le n$ and $b \le m$. **Thm** For all c there exists a unique finite set of grids OBS_c such that

- $n \times m$ is *c*-colorable iff
- $n \times m$ does not contain any element of OBS_c .
 - 1. $OBS_2 = \{3 \times 7, 7 \times 3, 5 \times 5\}.$

Def $n \times m$ contains $a \times b$ if $a \le n$ and $b \le m$. **Thm** For all c there exists a unique finite set of grids OBS_c such that

- $n \times m$ is *c*-colorable iff
- $n \times m$ does not contain any element of OBS_c .
 - 1. $OBS_2 = \{3 \times 7, 7 \times 3, 5 \times 5\}.$
 - 2. Can prove Thm using well-quasi-orderings. No bound on $|OBS_c|$.

ション ふゆ アメビア メロア しょうくり

Def $n \times m$ contains $a \times b$ if $a \le n$ and $b \le m$. **Thm** For all c there exists a unique finite set of grids OBS_c such that

- $n \times m$ is *c*-colorable iff
- $n \times m$ does not contain any element of OBS_c .
 - 1. $OBS_2 = \{3 \times 7, 7 \times 3, 5 \times 5\}.$
 - 2. Can prove Thm using well-quasi-orderings. No bound on |OBS_c|.

3. We showed $2\sqrt{c}(1-o(1)) \leq |OBS_c| \leq 2c^2$.

The theorem $a \times b$ is 2-colorable iff no elements of $OBS_2 = \{3 \times 7, 7 \times 3, 5 \times 5\}$ fits into $a \times b$

was proven by cleverness.

The theorem $a \times b$ is 2-colorable iff no elements of $OBS_2 = \{3 \times 7, 7 \times 3, 5 \times 5\}$ fits into $a \times b$

was proven by cleverness.

Could it have been proven by

The theorem $a \times b$ is 2-colorable iff no elements of $OBS_2 = \{3 \times 7, 7 \times 3, 5 \times 5\}$ fits into $a \times b$

was proven by cleverness.

Could it have been proven by

1. Quantum?

The theorem $a \times b$ is 2-colorable iff no elements of $OBS_2 = \{3 \times 7, 7 \times 3, 5 \times 5\}$ fits into $a \times b$

▲ロ ▶ ▲周 ▶ ▲ ヨ ▶ ▲ ヨ ▶ → ヨ → の Q @

was proven by cleverness.

Could it have been proven by

- 1. Quantum?
- 2. Machine Learning?

The theorem $a \times b$ is 2-colorable iff no elements of $OBS_2 = \{3 \times 7, 7 \times 3, 5 \times 5\}$ fits into $a \times b$

▲ロ ▶ ▲ □ ▶ ▲ □ ▶ ▲ □ ▶ ▲ □ ▶ ● ○ ○ ○

was proven by cleverness.

Could it have been proven by

- 1. Quantum?
- 2. Machine Learning?
- 3. Quantum Machine Learning?

The theorem $a \times b$ is 2-colorable iff no elements of $OBS_2 = \{3 \times 7, 7 \times 3, 5 \times 5\}$ fits into $a \times b$

was proven by cleverness.

Could it have been proven by

- 1. Quantum?
- 2. Machine Learning?
- 3. Quantum Machine Learning? (If so then great for funding!)

▲ロ ▶ ▲ □ ▶ ▲ □ ▶ ▲ □ ▶ ▲ □ ▶ ● ○ ○ ○

The theorem $a \times b$ is 2-colorable iff no elements of $OBS_2 = \{3 \times 7, 7 \times 3, 5 \times 5\}$ fits into $a \times b$

was proven by cleverness.

Could it have been proven by

- 1. Quantum?
- 2. Machine Learning?
- 3. Quantum Machine Learning? (If so then great for funding!)

▲ロ ▶ ▲ □ ▶ ▲ □ ▶ ▲ □ ▶ ▲ □ ▶ ● ○ ○ ○

4. (more seriously)

The theorem $a \times b$ is 2-colorable iff no elements of $OBS_2 = \{3 \times 7, 7 \times 3, 5 \times 5\}$ fits into $a \times b$

was proven by cleverness.

Could it have been proven by

- 1. Quantum?
- 2. Machine Learning?
- 3. Quantum Machine Learning? (If so then great for funding!)

▲ロ ▶ ▲ □ ▶ ▲ □ ▶ ▲ □ ▶ ▲ □ ▶ ● ○ ○ ○

- 4. (more seriously)
- 5. SAT Solvers?

The theorem $a \times b$ is 2-colorable iff no elements of $OBS_2 = \{3 \times 7, 7 \times 3, 5 \times 5\}$ fits into $a \times b$

was proven by cleverness.

Could it have been proven by

- 1. Quantum?
- 2. Machine Learning?
- 3. Quantum Machine Learning? (If so then great for funding!)

▲ロ ▶ ▲ □ ▶ ▲ □ ▶ ▲ □ ▶ ▲ □ ▶ ● ○ ○ ○

- 4. (more seriously)
- 5. SAT Solvers?
- 6. ChatGPT?

The theorem $a \times b$ is 2-colorable iff no elements of $OBS_2 = \{3 \times 7, 7 \times 3, 5 \times 5\}$ fits into $a \times b$

was proven by cleverness.

Could it have been proven by

- 1. Quantum?
- 2. Machine Learning?
- 3. Quantum Machine Learning? (If so then great for funding!)

- 4. (more seriously)
- 5. SAT Solvers?
- 6. ChatGPT?

Since 2-coloring has been solved why ask this question?

The theorem $a \times b$ is 2-colorable iff no elements of $OBS_2 = \{3 \times 7, 7 \times 3, 5 \times 5\}$ fits into $a \times b$

was proven by cleverness.

Could it have been proven by

- 1. Quantum?
- 2. Machine Learning?
- 3. Quantum Machine Learning? (If so then great for funding!)
- 4. (more seriously)
- 5. SAT Solvers?
- 6. ChatGPT?

Since 2-coloring has been solved why ask this question? 3-coloring is known. 4-coloring is known. 5-coloring is open!

Main Question

Fix *c* What is OBS_c



Main Question

Fix *c* What is OBS_c

We developed tools to get us both colorings and non-colorings. They helped us get some of our results, but (alas) to many had to be done ad-hoc.

3-COLORABILITY

We will EXACTLY Characterize which $n \times m$ are 3-colorable!

▲□▶ ▲□▶ ▲ 臣▶ ▲ 臣▶ ― 臣 … のへぐ

Easy 3-Colorable Results

Thm

1. The following grids **are not** 3-colorable. $4 \times 19, 19 \times 4, 5 \times 16, 16 \times 5, 7 \times 13, 13 \times 7, 10 \times 12, 12 \times 10, 11 \times 11.$

Easy 3-Colorable Results

Thm

- 1. The following grids are not 3-colorable. 4×19 , 19×4 , 5×16 , 16×5 , 7×13 , 13×7 , 10×12 , 12×10 , 11×11 .
- The following grids are 3-colorable.
 3 × 19, 19 × 3, 4 × 18, 18, 6 × 15, 15 × 6, 9 × 12, 12 × 9.

▲ロ ▶ ▲ □ ▶ ▲ □ ▶ ▲ □ ▶ ▲ □ ▶ ● ○ ○ ○

Easy 3-Colorable Results

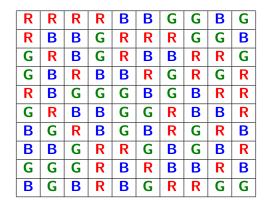
Thm

- 1. The following grids are not 3-colorable. 4×19 , 19×4 , 5×16 , 16×5 , 7×13 , 13×7 , 10×12 , 12×10 , 11×11 .
- 2. The following grids are 3-colorable. 3×19 , 19×3 , 4×18 , 18, 6×15 , 15×6 , 9×12 , 12×9 . Follows from tools.

▲ロ ▶ ▲ □ ▶ ▲ □ ▶ ▲ □ ▶ ▲ □ ▶ ● ○ ○ ○

10×10 is 3-colorable

Thm 10×10 is 3-colorable. UGLY! TOOLS DID NOT HELP AT ALL!!



◆□▶ ◆□▶ ◆三▶ ◆三▶ ・三 ・ つへぐ

10×11 is not 3-colorable

Thm 10×11 is not 3-colorable. You don't want to see this. UGLY case hacking.

▲□▶ ▲□▶ ▲ 臣▶ ▲ 臣▶ ― 臣 … のへぐ

Complete Char of 3-colorability

Thm $OBS_3 =$

 $\{4\times19,5\times16,7\times13,10\times11,11\times10,13\times7,16\times5,19\times4\}$

Complete Char of 3-colorability

Thm $OBS_3 =$

$\{4\times19,5\times16,7\times13,10\times11,11\times10,13\times7,16\times5,19\times4\}$

*ロ * * @ * * ミ * ミ * ・ ミ * の < や

Follows from our tools and the ad-hoc results.

```
The theorem

a \times b is 3-colorable iff no elements of

OBS_3 =

\{4 \times 19, 5 \times 16, 7 \times 13, 10 \times 11\} \cup

\{11 \times 10, 13 \times 7, 16 \times 5, 19 \times 4\}

fits into a \times b
```

was proven by cleverness.

```
The theorem

a \times b is 3-colorable iff no elements of

OBS_3 =

\{4 \times 19, 5 \times 16, 7 \times 13, 10 \times 11\} \cup

\{11 \times 10, 13 \times 7, 16 \times 5, 19 \times 4\}

fits into a \times b
```

▲ロ ▶ ▲周 ▶ ▲ ヨ ▶ ▲ ヨ ▶ → ヨ → の Q @

was proven by cleverness.

Can it be proven by

```
The theorem

a \times b is 3-colorable iff no elements of

OBS_3 =

\{4 \times 19, 5 \times 16, 7 \times 13, 10 \times 11\} \cup

\{11 \times 10, 13 \times 7, 16 \times 5, 19 \times 4\}

fits into a \times b
```

▲ロ ▶ ▲周 ▶ ▲ ヨ ▶ ▲ ヨ ▶ → ヨ → の Q @

was proven by cleverness.

Can it be proven by

1. Machine Learning?

```
The theorem

a \times b is 3-colorable iff no elements of

OBS_3 =

\{4 \times 19, 5 \times 16, 7 \times 13, 10 \times 11\} \cup

\{11 \times 10, 13 \times 7, 16 \times 5, 19 \times 4\}

fits into a \times b
```

▲ロ ▶ ▲周 ▶ ▲ ヨ ▶ ▲ ヨ ▶ → ヨ → の Q @

was proven by cleverness.

Can it be proven by

- 1. Machine Learning?
- 2. SAT Solvers?

```
The theorem

a \times b is 3-colorable iff no elements of

OBS_3 =

\{4 \times 19, 5 \times 16, 7 \times 13, 10 \times 11\} \cup

\{11 \times 10, 13 \times 7, 16 \times 5, 19 \times 4\}

fits into a \times b
```

▲ロ ▶ ▲周 ▶ ▲ ヨ ▶ ▲ ヨ ▶ → ヨ → の Q @

was proven by cleverness.

Can it be proven by

- 1. Machine Learning?
- 2. SAT Solvers?
- 3. ChatGPT?

```
The theorem

a \times b is 3-colorable iff no elements of

OBS_3 =

\{4 \times 19, 5 \times 16, 7 \times 13, 10 \times 11\} \cup

\{11 \times 10, 13 \times 7, 16 \times 5, 19 \times 4\}

fits into a \times b
```

was proven by cleverness.

Can it be proven by

- 1. Machine Learning?
- 2. SAT Solvers?
- 3. ChatGPT?

Since 2-coloring has been solved why ask this question?

```
The theorem

a \times b is 3-colorable iff no elements of

OBS_3 =

\{4 \times 19, 5 \times 16, 7 \times 13, 10 \times 11\} \cup

\{11 \times 10, 13 \times 7, 16 \times 5, 19 \times 4\}

fits into a \times b
```

was proven by cleverness.

Can it be proven by

- 1. Machine Learning?
- 2. SAT Solvers?
- 3. ChatGPT?

Since 2-coloring has been solved why ask this question? 4-coloring is known. 5-coloring is open!

```
The theorem

a \times b is 3-colorable iff no elements of

OBS_3 =

\{4 \times 19, 5 \times 16, 7 \times 13, 10 \times 11\} \cup

\{11 \times 10, 13 \times 7, 16 \times 5, 19 \times 4\}

fits into a \times b
```

was proven by cleverness.

Can it be proven by

- 1. Machine Learning?
- 2. SAT Solvers?
- 3. ChatGPT?

Since 2-coloring has been solved why ask this question?

4-coloring is known. 5-coloring is open!

If the answer is NO the we have found a problem that AI can't do!

4-COLORABILITY

From now on $G_{a,b}$ is $a \times b$. We will EXACTLY Characterize which $G_{n,m}$ are 4-colorable!

▲□▶ ▲□▶ ▲ 臣▶ ▲ 臣▶ ― 臣 … のへぐ

Easy NOT 4-Colorable Results

Thm The following grids are NOT 4-colorable:

- 1. $G_{5,41}$ and $G_{41,5}$
- 2. $G_{6,31}$ and $G_{31,6}$
- 3. $G_{7,29}$ and $G_{29,7}$
- 4. $G_{9,25}$ and $G_{25,9}$
- 5. $G_{10,23}$ and $G_{23,10}$
- 6. $G_{11,22}$ and $G_{22,11}$
- 7. $G_{13,21}$ and $G_{21,13}$
- 8. $G_{17,20}$ and $G_{20,17}$
- 9. $G_{18,19}$ and $G_{19,18}$

Easy NOT 4-Colorable Results

Thm The following grids are NOT 4-colorable:

- 1. $G_{5,41}$ and $G_{41,5}$
- 2. $G_{6,31}$ and $G_{31,6}$
- 3. $G_{7,29}$ and $G_{29,7}$
- 4. $G_{9,25}$ and $G_{25,9}$
- 5. $G_{10,23}$ and $G_{23,10}$
- 6. $G_{11,22}$ and $G_{22,11}$
- 7. $G_{13,21}$ and $G_{21,13}$
- 8. $G_{17,20}$ and $G_{20,17}$
- 9. $G_{18,19}$ and $G_{19,18}$

Follows from our tools.

Easy IS 4-Colorable Results

Thm The following grids are 4-colorable:

▲□▶ ▲□▶ ▲三▶ ▲三▶ 三三 - のへぐ

- 1. $G_{4,41}$ and $G_{41,4}$.
- 2. $G_{5,40}$ and $G_{40,5}$.
- 3. $G_{6,30}$ and $G_{30,6}$.
- 4. $G_{8,28}$ and $G_{28,8}$.
- 5. $G_{16,20}$ and $G_{20,16}$.

Easy IS 4-Colorable Results

Thm The following grids are 4-colorable:

- 1. $G_{4,41}$ and $G_{41,4}$.
- 2. $G_{5,40}$ and $G_{40,5}$.
- 3. $G_{6,30}$ and $G_{30,6}$.
- 4. $G_{8,28}$ and $G_{28,8}$.
- 5. $G_{16,20}$ and $G_{20,16}$.

Follows from our tools.

Theorems with UGLY Proofs

Thm

1. $G_{17,19}$ is NOT 4-colorable: Some Tools, Some ad-hoc.



Theorems with UGLY Proofs

Thm

1. $G_{17,19}$ is NOT 4-colorable: Some Tools, Some ad-hoc.

<□ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ > < □ >

2. $G_{24,9}$ is 4-colorable: Some Tools, Some ad-hoc.

Results about 4-COL So Far in This Talk

Thm

1. The following grids are in OBS_4 : $G_{5,41}, G_{6,31}, G_{7,29}, G_{9,25}, G_{10,23}, G_{11,22},$ $G_{22,11}, G_{23,10}, G_{25,9}, G_{29,7}, G_{31,6}, G_{41,5}.$

Results about 4-COL So Far in This Talk

Thm

- 1. The following grids are in OBS₄: $G_{5,41}, G_{6,31}, G_{7,29}, G_{9,25}, G_{10,23}, G_{11,22},$ $G_{22,11}, G_{23,10}, G_{25,9}, G_{29,7}, G_{31,6}, G_{41,5}.$
- The following grids status is unknown: G_{17,17}, G_{17,18}, G_{18,17}, G_{18,18}. G_{21,12}, G_{22,10}.

ション ふゆ アメリア メリア しょうくしゃ

Rectangle Free Conjecture

The following is obvious:

Lemma Let $n, m, c \in \mathbb{N}$. If $G_{n,m}$ is *c*-colorable then some color occurs $\geq \lceil nm/c \rceil$ times. Hence there is a rectangle free subset of $G_{n,m}$ with $\geq \lceil nm/c \rceil$ elements.

Rectangle Free Conjecture

The following is obvious:

Lemma Let $n, m, c \in \mathbb{N}$. If $G_{n,m}$ is *c*-colorable then some color occurs $\geq \lceil nm/c \rceil$ times. Hence there is a rectangle free subset of $G_{n,m}$ with $\geq \lceil nm/c \rceil$ elements.

Rectangle-Free Conjecture (RFC) is the converse:

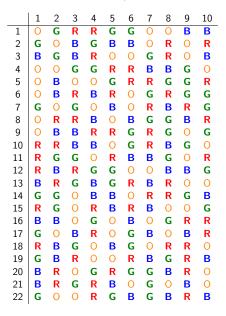
Let $n, m, c \ge 2$. If there is a rectangle free subset of size of $G_{n,m}$ which is $\ge \lceil nm/c \rceil$ then $G_{n,m}$ is c-colorable.

Rect-Free Subset of $G_{22,10}$ of size $55 = \left\lceil \frac{22 \cdot 10}{4} \right\rceil$

	01	02	03	04	05	06	07	08	09	10
1	•						•			
2		٠					•			
3			٠				•			
4				•			•			
5					•		•			
6						•	•			
7	•	•						•		
8			٠	•				•		
9					•	•		•		
10		•	٠						•	
11				•	•				•	
12	•					•			•	
13	•			•						•
14		•				•				•
15			٠		•					•
16		•			•					
17	•		٠							
18				•		•				
19			•			•				
20		٠		•						
21	•				•					
22							•	•	•	•

▲□▶ ▲□▶ ▲三▶ ▲三▶ 三三 - のへぐ

4-coloring of $G_{22,10}$ Due to Brad Loren

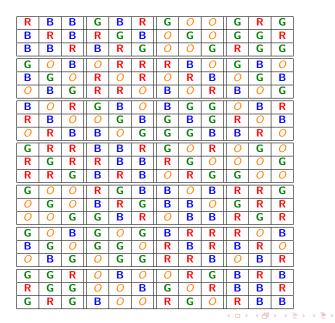


Rect-Free subset of $G_{21,12}$ of size $63 = \left\lceil \frac{21 \cdot 12}{4} \right\rceil$

	01	02	03	04	05	06	07	08	09	10	11	12
1	•	•										
2	•		•									
3		•	•									
4			•	•	•							
5		•		•		•						
6	•				•	•						
7						•	•	•				
8					•		•		•			
9				•				•	•			
10						•				٠	•	
11					•					•		•
12				•							•	•
13			•			•			•			•
14			•					•		٠		
15			•				•				•	
16		•							•	•		
17		•			•			•			•	
18		•					•					•
19	٠								٠		٠	
20	•							•				•
21	•			•			•			•		

◆□▶ ◆□▶ ◆目▶ ◆目▶ 目 のへの

Tom Sirgedas's 4-coloring of 21×12



= 900

Rectangle Free subset of $G_{18,18}$ of size $81 = \left\lceil \frac{18 \cdot 18}{4} \right\rceil$

	01	02	03	04	05	06	07	08	09	10	11	12	13	14	15	16	17	18
1		٠		•										•		•	•	
2	•	•								•	•		•					
3	٠								•						•	•		•
4						•			•			•	•	•				
5		•	•			•												•
6	•			•		•	•											
7							•	•		•				•				•
8			•				•		•		•						•	
9		٠			•		•					•			•			
10				•							•	•						•
11	٠		•		•									•				
12			•	•				•					•		•			
13					•	•		•			•					•		
14	•							•				•					•	
15				•	•				•	•								
16						•				•					•		•	
17			•							•		•				•		
18					•								•				•	•

If RFC is true then $G_{18,18}$ is 4-colorable. NOTE: If delete 2nd column and 5th Row have 74-sized RFC of $G_{17,17}$.

◆□ ▶ ◆□ ▶ ◆ □ ▶ ◆ □ ▶ ● □ ● ● ● ●

On Nov 30, 2009 I posted a blog with the following offer: The first person to email me both (1) plaintext, and (2) LaTeX, of a 4-coloring of the 17 \times 17 grid that has no monochromatic rectangles will receive \$289.00.

On Nov 30, 2009 I posted a blog with the following offer: The first person to email me both (1) plaintext, and (2) LaTeX, of a 4-coloring of the 17 \times 17 grid that has no monochromatic rectangles will receive \$289.00.

Also want to know about 18×18 .

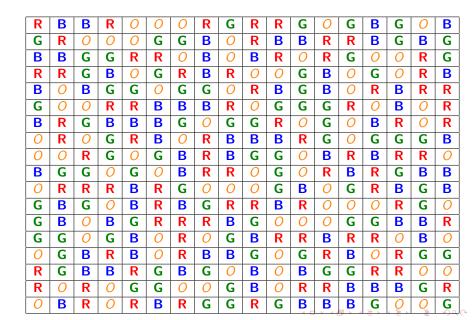
On Nov 30, 2009 I posted a blog with the following offer: The first person to email me both (1) plaintext, and (2) LaTeX, of a 4-coloring of the 17×17 grid that has no monochromatic rectangles will receive \$289.00.

Also want to know about 18×18 .

Bernd Steinbach and Christian Postoff showed both $G_{18,18}$ is 4-colorable and they are \$289 richer!

ション ふゆ アメリア メリア しょうくしゃ

Steinbach & Postoff's 4-Coloring of G_{18,18}





Given the 4-colorings of ${\it G}_{18,18},~{\it G}_{21,12},~{\it G}_{22,10}$ we now have that ${\rm OBS}_4$ is

◆□▶ ◆□▶ ◆目▶ ◆目▶ 目 のへの

Given the 4-colorings of ${\it G}_{18,18},~{\it G}_{21,12},~{\it G}_{22,10}$ we now have that ${\rm OBS}_4$ is

▲ロ ▶ ▲周 ▶ ▲ ヨ ▶ ▲ ヨ ▶ → ヨ → の Q @

 $\{ \textit{G}_{5,41}, \textit{G}_{6,31}, \textit{G}_{7,29}, \textit{G}_{9,25}, \textit{G}_{10,23}, \textit{G}_{11,22} \} \cup \\ \{ \textit{G}_{22,11}, \textit{G}_{23,10}, \textit{G}_{25,9}, \textit{G}_{29,7}, \textit{G}_{31,6}, \textit{G}_{41,5} \}.$

Given the 4-colorings of ${\it G}_{18,18},~{\it G}_{21,12},~{\it G}_{22,10}$ we now have that ${\rm OBS}_4$ is

 $\{G_{5,41}, G_{6,31}, G_{7,29}, G_{9,25}, G_{10,23}, G_{11,22}\} \cup \\\{G_{22,11}, G_{23,10}, G_{25,9}, G_{29,7}, G_{31,6}, G_{41,5}\}.$

The usual **Research Question:** Can we get OBS₄ with Al?

What About Square-Free Colorings?

The following is known but much harder: Thm

▲□▶ ▲□▶ ▲ 臣▶ ▲ 臣▶ ― 臣 … のへぐ

What About Square-Free Colorings?

The following is known but much harder: **Thm**

- 1. There exists N such that, for all 2-colorings of $N \times N$, there exists a monochromatic square. (The proof gives an enormous value of N though by a computer search its known that N = 15 is the min value that works.)
- 2. For all c there exists N = N(c) such that, for all c-colorings of $N \times N$, there exists a monochromatic square. (N(2) = 15 is known. Beyond that I believe nothing is known.)

What About Square-Free Colorings?

The following is known but much harder: **Thm**

- 1. There exists N such that, for all 2-colorings of $N \times N$, there exists a monochromatic square. (The proof gives an enormous value of N though by a computer search its known that N = 15 is the min value that works.)
- 2. For all c there exists N = N(c) such that, for all c-colorings of $N \times N$, there exists a monochromatic square. (N(2) = 15 is known. Beyond that I believe nothing is known.)

If you want to see the proof that for all c, N(c) exists then Take CMSC 752 in Spring 2025

- ・ロト・日本・モト・モト・ ヨー のへぐ

1. What is OBS_5 ?



- 1. What is OBS_5 ?
- 2. Prove or disprove Rectangle Free Conjecture.

▲□▶ ▲□▶ ▲目▶ ▲目▶ 三日 - のへで

- 1. What is OBS_5 ?
- 2. Prove or disprove Rectangle Free Conjecture.
- 3. Have $\Omega(\sqrt{c}) \le |OBS_c| \le O(c^2)$. Get better bounds!

- 1. What is OBS₅?
- 2. Prove or disprove Rectangle Free Conjecture.
- 3. Have $\Omega(\sqrt{c}) \leq |OBS_c| \leq O(c^2)$. Get better bounds!

4. Refine tools so can prove ugly results cleanly.

- 1. What is OBS₅?
- 2. Prove or disprove Rectangle Free Conjecture.
- 3. Have $\Omega(\sqrt{c}) \leq |OBS_c| \leq O(c^2)$. Get better bounds!

▲ロ ▶ ▲周 ▶ ▲ ヨ ▶ ▲ ヨ ▶ → 目 → の Q @

- 4. Refine tools so can prove **ugly** results **cleanly**.
- 5. Unleash AI on these problems!