

Interval and Sign Patterns

Will Roe, Ian Kim

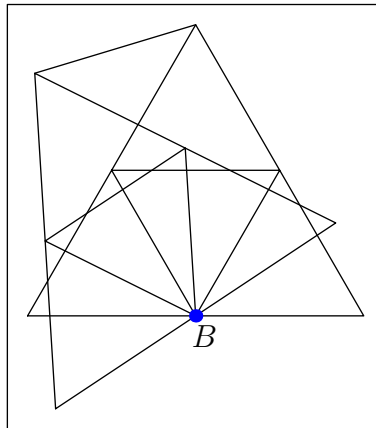
October 16, 2025

1 Abstract

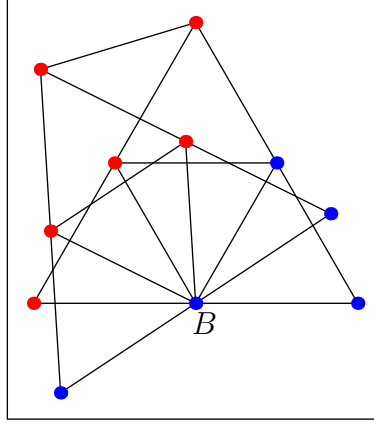
In Euclidean Ramsey Theory, a *line* is a sequence of collinear points with distance 1. Such a line is notated as an l_n where n is the number of points. For k -colorings of $\mathbb{R}^n$, a red l_n is said to be an l_n whose points are all colored red. $\mathbb{E}^n \rightarrow (l_r, l_b)$ is written to mean that in every coloring of $\mathbb{R}^n$, $\exists$ a red l_r or a blue l_b . Below is an example of a red l_3 .



It is known for a few small m that $\mathbb{E}^n \rightarrow (l_3, l_m)$. The simplest example is when $m = 2$. It is trivial to prove that $\mathbb{E}^n \rightarrow (l_3, l_2)$. For example, if there are no blue points in a coloring then every l_3 is red, and if there is a blue point, the figure below shows a set of 11 points, including the point known to be blue (labeled B) for which it is impossible to color them all blue or red without getting a red l_3 or a blue l_2 .



An example coloring that fails is shown below:



It has been proven by Conlon and Wu [1] that $\mathbb{E}^n \not\rightarrow (l_3, l_m)$ for some number m . Their paper used two smaller lemmas to determine that $m \leq 10^{50}$. The first lemma discussed bounds on the number of sign patterns of some quantity of polynomials with given maximum degree that accept a given number of inputs. The second lemma discussed bounds on the number of specific intervals that any quadratic falls into.

2 Introduction

Keywords Euclidean Ramsey Theory; Ramsey Theory; Sign Patterns; Intervals; Colorings

Def 2.1 The sign pattern of a series of functions, calculated at a point, is defined to be the sequence of the signs of the functions at the point.

Here is an example of a sign pattern being calculated:

Say we want to find the sign pattern of functions $(x + 1)^2 y, \sin(y + x), 2x - xy$ when $x = 0$ and $y = 3\pi$. It would be calculated as such:

Function	Value	Sign
$x^2 y$	$(0 + 1)^2 \cdot 3\pi = 3\pi$	+
$x^2 y$	$\sin(3\pi + 0) = -1$	-
$x^2 y$	$2 \cdot 0 - 0 \cdot 5\pi = 0$	0

Reading from top to bottom the Sign column, we find the sign pattern to be the string "+ - 0".

Def 2.2 The sign pattern of a series of functions, calculated at a point, is defined to be the sequence of the signs of the functions at the point.

Here is an example of what is meant: Given functions $x + 2 + 3y, 3x - 1 -$

	Function	Value	Sign
$8y, 5x + 2y + 7$	x^2y	$(0 + 1)^2 \cdot 3\pi = 3\pi$	+
	x^2y	$\sin(3\pi + 0) = -1$	-
	x^2y	$2 \cdot 0 - 0 \cdot 5\pi = 0$	0

Reading from top to bottom the Sign column, we find the sign pattern to be the string "+ - 0".

Notation 2.3 $\mathbb{E}^n$ is written to mean n -dimensional Euclidean space. A k -coloring of $\mathbb{R}^n$ is some $f : \mathbb{R}^n \rightarrow [k]$. In this paper we focus on 2-colorings of $\mathbb{R}^2$.

Def 2.4 Let $\text{Reg}_{\max}(d, n, p)$ be the *maximum number of regions n polynomials with maximum degree p with d inputs.*

Def 2.5 Let $\text{Reg}_{\max\text{lin}}(d, n, p)$ be the *maximum number of regions n linear polynomials with maximum degree p with d inputs.*

Def 2.6 Let l_n be a set of points $a_1, a_2, \dots, a_n$ where $a_{i+1} - a_i = u$ for some $|u| = 1$.

Def 2.7 A polynomial $P(a_1, a_2, \dots, a_d) = \sum_{i_1, i_2, \dots, i_d=0}^p C_{i_1, i_2, \dots, i_d} \cdot a_1^{i_1} \cdot a_2^{i_2} \cdot \dots \cdot a_d^{i_d}$ for constant c .

Def 2.8 A linear polynomial $O(a_1, a_2, \dots, a_d) = \sum_{i=1}^d c_i a_i$ for constant c .

Notation 2.9 Let the set $\mathbb{P}_{d,p}$ consist of all P with d inputs and maximum degree p . Let the set $\mathbb{P}_{d,n,p}$ be the power set of $\mathbb{P}_{d,p}$.

It will later be useful to consider trivial cases of $\text{Reg}_{\max}$ and of $\text{Reg}_{\max\text{lin}}$. Some are listed here.

An obvious bound on $\text{Reg}_{\max}(d, n, p)$ is 3^n . For n polynomials, each may return one of three signs: $+$, $-$, or 0 . Therefore, there exist 3^n sign patterns.

$\text{Reg}_{\max\text{lin}}(0, n)$ represents a set of constants. Since the sign of a constant never changes, the sign pattern is always the same, so $\text{Reg}_{\max\text{lin}}(0, n) = 1$.

$\text{Reg}_{\max\text{lin}}(d, 0)$ represents exactly zero polynomials. Since there are no polynomials, the sign pattern is of length 0, for which only one pattern exists. Therefore, $\text{Reg}_{\max\text{lin}}(d, 0) = 1$.

Def 2.10 A function $F(n) = \text{mod}(n^2 + \alpha n + \beta, p)$ for real α, β and prime p . $F(n)$ then takes on values in $[0, p)$.

A function $R(n)$ is taken to be to be the number of regions $[0, 1), [1, 2), \dots, [p-1, p)$ that contain at least one value in $F(1), F(2), \dots, F(n)$. For the sake of this paper, F will be defined in advance.

Def 2.11 $R(n, \alpha, \beta) = |\{i | |\{j | j \in [n], i \leq F(j) < i+1\}| \neq 0\}|$ for $F(n) = \text{mod}(n^2 + \alpha n + \beta, p)$.

It will also be useful to consider the inverse of R , which we will call Q .

Notation 2.12 Let $Q(\alpha, \beta)$ be the smallest number such that $R(Q(\alpha, \beta), \alpha, \beta) = p$, if such a number exists at all.

A particularly powerful lower bound for the number of sign patterns of polynomials was found by Olenik et. al. A special case of their theorem was used by Conlon et. al. to obtain the following corollary:

Theorem 1 n polynomials of degree p with d variable will form, at most, $\left(\frac{50pn}{d}\right)^d$ sign patterns, for $n > d$.

Our proof focuses on a geometric interpretation of $\text{Reg}_{\max}$, where Olenik's mainly relied on matrix manipulation.

In Section 1 we present the problem and a basic overview of Conlon and Wu's paper. In Section 2 we state definitions regarding the two lemmas given in Conlon and Wu's paper. In Section 3 we discuss and attempt to improve the result given by Olenik. In Section 4 we use a Monte Carlo simulation to give a better bound on the number of regions a polynomial falls into. In Section 5 we use our results to improve the bounds on m and present open questions.

3 Signs

It is useful to imagine the sign of a function P with d inputs as a surface cutting $\mathbb{R}^d$ into at most parts: where $P > 0$, $P < 0$, and where $P = 0$. Note that some polynomials, like x^2 or $x^2 + 1$, do not divide space into three pieces.

For a set of polynomials $P_1, P_2, \dots, P_n$, the $\mathbb{R}^{d-1}$ surface $P_i = 0$ splits $\mathbb{R}^d$ into some number of regions. (For the sake of this paper, we take $P = 0$ as a region as well as $P > 0$ and $P < 0$.) Due to the continuity of polynomials, every region contains only points with the same sign pattern. Therefore, the number of sign patterns $\leq \text{Reg}_{\max}(d, n, p)$.

Lemma 3.1 *The intersection of any two hyperplanes in $\mathbb{R}^n$ is a hyperplane in $\mathbb{R}^{n-1}$.*

Proof: Let H be the hyperplane defined by the set of points $\vec{x}$ where $\vec{h}_c \cdot \vec{x} + h_i = 0$, and J be the set of points $\vec{x}$ where $\vec{j}_c \cdot \vec{x} + j_i = 0$. Then let $h_j = \sum_{i=2}^n h_{c,i}x_i + h_i$ and $j_j = \sum_{i=0}^{n-2} j_{c,i}x_i + j_i$. Then $h_j + h_{c,0}x_0 + h_{c,1}x_1 = 0$ and $j_j + j_{c,0}x_0 + j_{c,1}x_1 = 0$. If $\begin{bmatrix} h_{c,0} & h_{c,1} \\ j_{c,0} & j_{c,1} \end{bmatrix}$ is invertible, then there is one solution for x_0 and x_1 . ■

Lemma 3.2 $\text{Reg}_{\max\text{lin}}(d, n + 1) \leq \text{Reg}_{\max\text{lin}}(d, n) + 2 \text{Reg}_{\max\text{lin}}(d - 1, n) \forall d, n > 1$.

Proof: Assume you have n hyperplanes, with d inputs, that divide space into the maximum number of regions, e.g. $\text{Reg}_{\max\text{lin}}(d, n)$ regions. Suppose you add another hyperplane H , defined by $P = 0$. This hyperplane may intersect any number of other planes. The intersections of each hyperplane with the new hyperplane form a hyperplane with one fewer dimension. At worst, you will get $\text{Reg}_{\max\text{lin}}(d - 1, n)$ regions formed by the intersection. Notice that these regions exist in the higher level of space, where $P = 0$. Every region that existed before H that is split by H becomes either two or three new regions. All points in this region will have the same sign pattern, except for the sign of P , which might be 1, 0, or -1 for those points. We may assume the cut makes three new regions to calculate a bound for $\text{Reg}_{\max\text{lin}}$. Therefore, you get 2 more regions than there were before for those divided. Recall that since $\text{Reg}_{\max\text{lin}}(d - 1, n)$ regions were formed, $\text{Reg}_{\max\text{lin}}(d - 1, n)$ regions must have been cut. Therefore, you

have $\text{Reg}_{\max\text{lin}}(d-1, n) \cdot 2$ new regions. There were $\text{Reg}_{\max\text{lin}}(d, n)$ regions before this new hyperplane was added, so after adding another hyperplane you have $\text{Reg}_{\max\text{lin}}(d, n) + 2\text{Reg}_{\max\text{lin}}(d-1, n)$ regions. However, the new hyperplane represents another polynomial with degree 1 and d dimensions, so the number of new regions is, at worst, $\text{Reg}_{\max\text{lin}}(d, n+1)$. Therefore, $\text{Reg}_{\max\text{lin}}(d, n+1) \leq \text{Reg}_{\max\text{lin}}(d, n) + 2\text{Reg}_{\max\text{lin}}(d-1, n)$. ■

Lemma 3.3 $\text{Reg}_{\max\text{lin}}(d, n) \leq \sum_{i=0}^d 2^i \binom{n}{i}$.

Proof: Let $R(d, n) = \sum_{i=0}^d 2^i \binom{n}{i}$. R has the unique property that $R(0, n) = R(d, 0) = 1$, and that $R(d, n) = R(d, n-1) + 2R(d-1, n-1)$. Thus for $d = 0$ or $n = 0$, $\text{Reg}_{\max\text{lin}}(d, n) = R(d, n)$. Assume it is known that $\text{Reg}_{\max\text{lin}}(d-1, n) \leq \sum_{i=0}^{d-1} 2^i \binom{n}{i}$ and that $\text{Reg}_{\max\text{lin}}(d-1, n-1) \leq \sum_{i=0}^{d-1} 2^i \binom{n-1}{i}$. Then $\text{Reg}_{\max\text{lin}}(d-1, n) + 2\text{Reg}_{\max\text{lin}}(d-1, n-1) \leq R(d-1, n) + 2R(d-1, n-1) = R(d, n)$. But by the previous lemma $\text{Reg}_{\max\text{lin}}(d, n) \leq \text{Reg}_{\max\text{lin}}(d-1, n) + 2\text{Reg}_{\max\text{lin}}(d-1, n-1)$. Then $\text{Reg}_{\max\text{lin}}(d, n) \leq R(d, n)$. ■

Lemma 3.4 $\text{Reg}_{\max}(d, n, p) \leq \text{Reg}_{\max\text{lin}}(d^p, n)$.

Proof: Every polynomial in $\mathbb{P}_{(d, n, p)}$ is equal to $\sum_{i_1, i_2, \dots, i_d=0}^p C_{i_1, i_2, \dots, i_d} \cdot a_1^{i_1} \cdot a_2^{i_2} \cdot \dots \cdot a_d^{i_d}$ for some $C_{i_1, i_2, \dots, i_d}$ and inputs $a_1, a_2, \dots, a_d$. Then view the polynomial as a linear polynomial with inputs $a_1^{i_1} \cdot a_2^{i_2} \cdot \dots \cdot a_d^{i_d}$. Assuming all values of $a_1^{i_1} \cdot a_2^{i_2} \cdot \dots \cdot a_d^{i_d}$ are independent, we can treat them as another set of inputs, $a'_1, a'_2, \dots, a'_{d^p}$, and treat C as another list of constants $C'_{1, 2, \dots, d^p}$. Therefore every polynomial can be written as $\sum_{i=0}^{d^p} C'_i \cdot a'_i$. However, $a'_1, a'_2, \dots, a'_{d^p}$ may or may not take on any set of values, since they are all functions of a smaller set of variables, $a_1, a_2, \dots, a_d$. So it is indeed the case that $\text{Reg}_{\max}(d, n, p) \leq \text{Reg}_{\max\text{lin}}(d^p, n)$. ■

These lemmas give an upper bound of $\text{Reg}_{\max}(d, n, p) \leq \sum_{i=0}^{d^p} 2^i \binom{n}{i}$. In the context of Conlon and Wu's proof, it was useful to consider the case when $n = 4m^3, p = 1, d = 2$, where $\mathbb{E}^n \not\rightarrow (l_3, l_m)$. The form provided by Olenik et. al. gives $\text{Reg}_{\max}(2, 4m^3, 1) \leq 10^4 m^6$, but the upper bound discussed here gives $\text{Reg}_{\max}(2, 4m^3, 1) \leq \sum_{i=0}^{2^1} 2^i \text{nCr}(4m^3, i) = 1 + 32m^6$.

4 Intervals

It is known $F(1), F(2), \dots, F(p^3) \bmod p$ must fall into at least $\frac{p}{6}$ regions for $F = x^2 + \alpha x + \beta$ with $\alpha, \beta \in \mathbb{R}$. However, it often falls into a greater number of regions. One can make a graph of the smallest number of values of F required to have each region contain some value of F . Below is such a graph, with $\alpha \in [0, 1]$ along the x-axis and $\beta \in [0, 1]$ along the y-axis. In other words, the graph is of $Q(\alpha, \beta)$ for $p = 7$. Interestingly, there seem to be

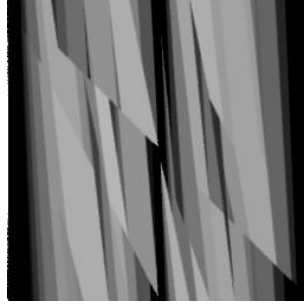


Figure 1: Graph of interval count for $p=7$. Darker values indicate a larger Q , while bright white indicates $Q=0$.

polygonal regions with widely varying values instead of something smooth and [not straight]. A similar graph can be made for larger values of p .

A Monte Carlo simulation was run to evaluate $Q(\alpha, \beta)$ over $0 \leq \alpha < 1, 0 \leq \beta < 1$. Specifically, the region $0 \leq \alpha < 1, 0 \leq \beta < 1$ was divided into 128^2 equal square regions. Then, for each smaller region, 64 values of α and β were picked in the region, and Q was evaluated. The distribution of Q is shown below.

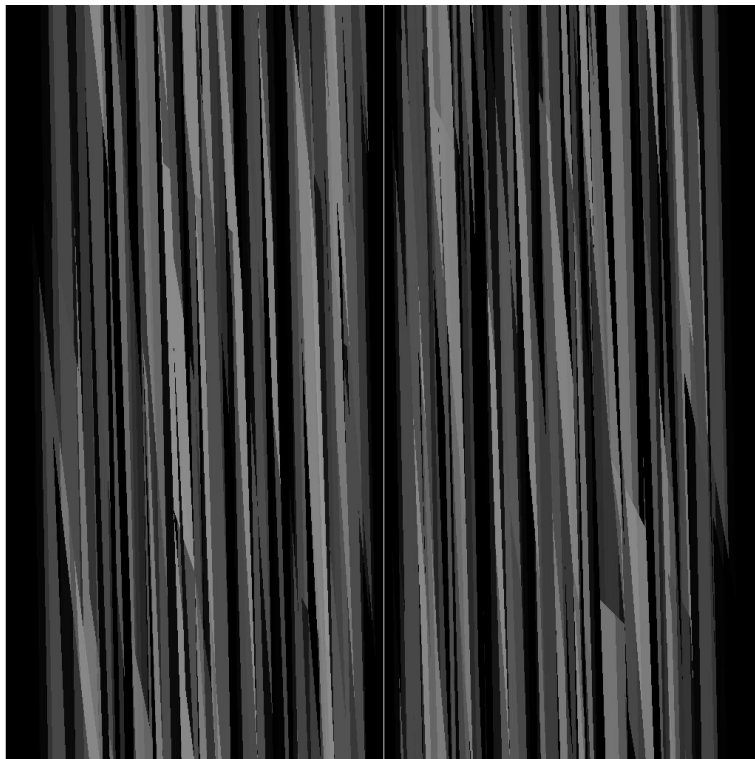


Figure 2: Graph of interval count for $p=17$. Darker values indicate a larger Q .

p	Average	Min	Lower Quartile	Mode	Upper Quartile	Max
7	22.77	9	14.92	14.01	22.29	124.54
11	51.52	18.96	35.13	31.33	50.89	310.13
13	72.72	29.91	55.45	45	71.53	406.2
17	78.34	42.77	50.91	54.69	61.37	399.62
19	112.05	46.57	58.96	72.4	125.96	561.72
23	134.19	63.23	75.96	88.43	153.32	722.8
29	147.86	80.98	93.78	111.02	167.65	733.69
31	167.13	88.76	108.88	115.78	175.66	839.29
37	212.45	114.3	141.97	159.77	228.24	1053.36
41	245	127.08	159.8	177.3	240.42	1194.09
43	272.51	142.88	170.54	192.66	283.93	1334.18
47	263.07	167.74	188.4	201.49	285.29	1253.69

From this data we see that the maximum value of Q grows roughly linearly with p . In fact, for this set of data, $Q \approx 27p$. So, for $m = 27p$, roughly p intervals are hit.

The proof itself uses the polynomial $a + d(i - 1) + i^2 - 3i + 2$ for some a , d , and prime p .

5 Conclusions and Open Problems

Conlon and Wu's result required $12p^{-\frac{1}{4}} < \frac{1}{2}$ and $\text{Reg}_{\max}(2, Q(p), 1) \left(1 - p^{-\frac{3}{4}}\right)^p < \frac{1}{2}$, where $m = Q(p)$. This gives $10^9 < p$, for which there exists at least one prime $\leq 2 \cdot 10^9$. Finally, since $m = 27p$, $\exists m \leq 54 \cdot 10^9 |\mathbb{E}^2 \rightarrow (l_3, l_m)$, and, equivalently, $\mathbb{E}^2 \rightarrow (l_3, l_{10^{10}})$. Without using the results shown here for the interval problem, we get the inequality $(32m^6 + 1) \left(1 - p^{-\frac{3}{4}}\right)^{\frac{p}{6}} \leq \frac{1}{2}$, for $m = p^3$.

6 Acknowledgements

The researcher used Desmos and Wolfram Alpha. They would like to thank Dr. William Gasarch for their mentorship and time.

References

- [1] W. Conlon. More on lines in euclidean ramsey theory. *Arxiv*, 322(10):891–921, 2022.