

# **CMSC 474, Introduction to Game Theory**

## **11. Evolutionary Stability**

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# Outline

- Chapter 2 discussed two solution concepts:
  - Pareto optimality and Nash equilibrium
- Chapter 3 discusses several more:
  - Maxmin and Minmax
  - Dominant strategies
  - Correlated equilibrium
  - Trembling-hand perfect equilibrium (complicated definition)
  - $\epsilon$ -Nash equilibrium
  - Evolutionarily stable strategies

# Evolutionary Stability

- This concept comes from evolutionary biology
- Start with a population of some species
  - For us species are those agents playing a particular strategy  $s$
- Add a small population of “invaders” species
  - For us invaders are those agents playing a different strategy  $t$
  - Assume  $t$  **invades**  $s$  at level  $p$ , i.e.,  $p$  is the fraction that uses  $t$
  - $(1-p)$  = the fraction that uses  $s$
- If  $s$ 's fitness against the mixture of both species is higher than  $t$ 's, then  $t$ 's proportion will shrink and  $s$ 's will grow (thus  $s$  is “**stable**”)
  - Fitness for species is the ability to both survive and reproduce
  - For us, fitness of a species = its expected payoff from interacting with a random member of the population, namely with species  $t$  with probability  $p$  and with species  $s$  with probability  $1-p$

# Evolutionary Stability

- Write a payoff matrix for the two species against each other
  - Symmetric 2-player game, so we only need to look at agent 1's payoffs
- A strategy's **fitness** is its expected payoff against a randomly chosen agent
  - $\text{fitness}(s) = (1-p)a + pb$
  - $\text{fitness}(t) = (1-p)c + pd$
- $s$  is **evolutionarily stable against**  $t$  if there is an  $\varepsilon > 0$  such that for every  $p < \varepsilon$ ,  $\text{fitness}(s) > \text{fitness}(t)$ 
  - i.e.,  $(1-p)a + pb > (1-p)c + pd$
- As  $p \rightarrow 0$ ,  $(1-p)a + pb \rightarrow a$  and  $(1-p)c + pd \rightarrow c$ 
  - For sufficiently small  $p$ , the inequality holds if  $a > c$ , or if  $a = c$  and  $b > d$
- Thus  $s$  is evolutionarily stable against  $t$  iff either of the following holds:
  - $a > c$
  - $a = c$  and  $b > d$

	$s$	$t$
$s$	$a, a$	$b, c$
$t$	$c, b$	$d, d$

# Example: the Body-Size Game

- Consider two different sizes of beetles competing for food
  - When beetles of the same size compete, they get equal shares
  - When large competes with small, large gets most of the food
  - Large beetles get less fitness benefit from any given amount of food
    - Some of it is diverted into maintaining their expensive metabolism

	<i>small</i>	<i>large</i>
<i>small</i>	5, 5	1, 8
<i>large</i>	8, 1	3, 3

- Is a population of small beetles evolutionarily stable against large beetles?

	<i>large</i>	<i>small</i>
<i>large</i>	3, 3	8, 1
<i>small</i>	1, 8	5, 5

- Is a population of large beetles evolutionarily stable against small ones?

- Source:

- <http://www.cs.cornell.edu/home/kleinber/networks-book>

# Evolutionary Stability

- More generally, suppose  $s$  is a mixed strategy
  - Represents a population composed of several species
  - We'll talk about  $s$ 's evolutionary stability against all other mixed strategies
  - $s$  is an **evolutionarily stable strategy (ESS)** iff for *every* mixed strategy  $t \neq s$ , either of the following holds:
    - $u(s,s) > u(t,s)$
    - $u(s,s) = u(t,s)$  and  $u(s,t) > u(t,t)$(note that  $u_1 = u_2$  since the game is symmetric)
  - $s$  is **weakly evolutionarily stable** iff for every mixed strategy  $t \neq s$ , either of the following stability conditions holds:
    1.  $u(s,s) > u(t,s)$
    2.  $u(s,s) = u(t,s)$  and  $u(s,t) \geq u(t,t)$
- Includes cases where  $s$  and  $t$  have the same fitness
- So the population that uses  $t$  neither grows nor shrinks

# Example

	H	D
H	-2, -2	6, 0
D	0, 6	3, 3

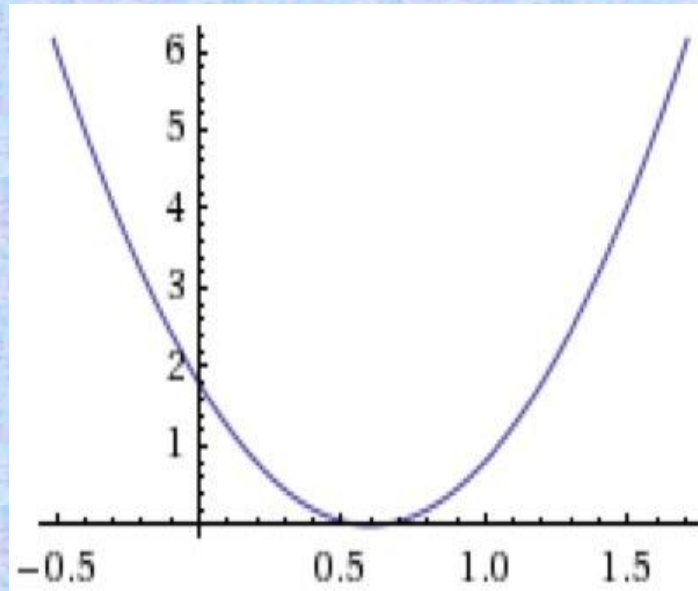
- **The Hawk-Dove game**

- 2 animals contend for a piece of food
- Each animal may be either a hawk (H) or a dove (D)
  - The prize is worth 6 to each
  - Fighting costs each 5
- When a hawk meets a dove, the hawk gets the prize without a fight: payoffs 6, 0
- When 2 doves meet, they split the prize without a fight: payoffs 3, 3
- When 2 hawks meet,
  - They fight, and each has a 50% chance of getting the prize
  - For each, the payoff is  $-5 + 0.5 \cdot 6 = -2$
- Unique Nash equilibrium  $(s, s)$ , where  $s = \{(3/5, H), (2/5, D)\}$ 
  - i.e., 60% hawks, 40% doves

# Example

	H	D
H	-2, -2	6, 0
D	0, 6	3, 3

- To confirm that  $s$  is also an ESS, show that for all  $t \neq s$ ,
  - $u_1(s, s) > u_1(t, s)$  OR
  - $u_1(s, s) = u_1(t, s)$  and  $u_1(s, t) > u_1(t, t)$ 
    - where  $s = \{(3/5, H), (2/5, D)\}$  and  $t = \{(p, H), (1-p, D)\}$
- For every fully-mixed strategy  $s$ , if  $(s, s)$  is a Nash equilibrium then  $u_1(s, s) = u_1(t, s)$
- Next, show  $u_1(s, t) > u_1(t, t)$ :
  - $u_1(s, t) = (3/5)(-2p + 6(1-p)) + (2/5)(0p + 3(1-p))$
  - $u_1(t, t) = p(-2p + 6(1-p)) + (1-p)(0p + 3(1-p))$
- Let  $v = u_1(s, t) - u_1(t, t)$
- Easy to solve using <http://wolframalpha.com>
  - Simplifies to  $v = 5p^2 - 6p + 9/5$
  - Unique minimum  $v = 0$  when  $p = 3/5$ , i.e.,  $t = s$
  - If  $p \neq 3/5$  then  $v > 0$ , i.e.,  $u_1(s, t) > u_1(t, t)$





# Evolutionary Stability and Nash Equilibria

- Recall that  $s$  is **evolutionarily stable** iff for every mixed strategy  $t \neq s$ , either of the following holds:
  - $u(s,s) > u(t,s)$  (1)
  - $u(s,s) = u(t,s)$  and  $u(s,t) > u(t,t)$  (2)

**Theorem.** Let  $G$  be a symmetric 2-player game, and  $s$  be a mixed strategy. If  $s$  is an evolutionarily stable strategy, then  $(s, s)$  is a Nash equilibrium of  $G$ .

**Proof.** By definition, an ESS  $s$  must satisfy  $u(s,s) \geq u(t,s)$ , i.e.,  $s$  is a best response to itself, so it must be a Nash equilibrium.

**Theorem.** Let  $G$  be a symmetric 2-player game, and  $s$  be a mixed strategy. If  $(s,s)$  is a strict Nash equilibrium of  $G$ , then  $s$  is evolutionarily stable.

**Proof.** If  $(s,s)$  is a strict Nash equilibrium, then  $u(s,s) > u(t,s)$ .

- This satisfies (1) above

# Summary

- Maxmin and minmax strategies, and the Minimax Theorem
  - Matching Pennies, Two-Finger Morra
- dominant strategies
  - Prisoner's Dilemma, Which Side of the Road, Matching Pennies
  - Iterated elimination of dominated strategies (IESDS)
- rationalizability
  - the  $p$ -Beauty Contest
- correlated equilibrium
  - Battle of the Sexes
- epsilon-Nash equilibria
- evolutionarily stable strategies
  - Body-Size game, Hawk-Dove game